# ต้นไม้แบบทวิภาค

(Binary Trees)

สมชาย ประสิทธิ์จูตระกูล Translated to English by Nuttapong Chentanez

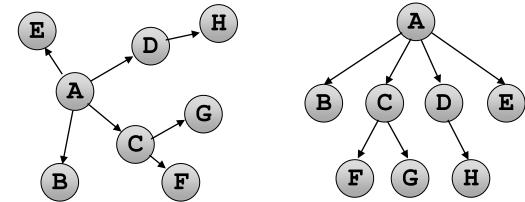
### **Topics**

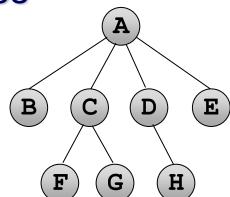
- > Tree Definition
- > Tree Implementation
- ➤ Binary Tree
  - ➤ Huffman Tree
  - > Expression Tree
  - Tree traversal
  - Expression Tree Evaluation
  - ➤ Expression Tree Differentiation

#### Tree

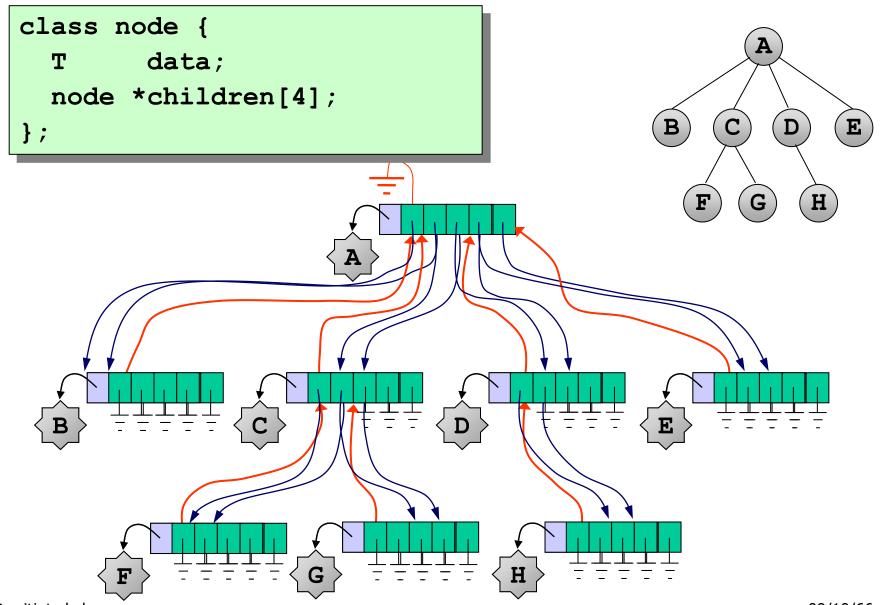
- Tree consists of nodes and edges
- Edge has direction (Directed Edge)
- A is the parent of B when there's an edge from A to B
- Each node has only one parent (Except root, which has no parent)

Tree with v nodes has v – 1 edges





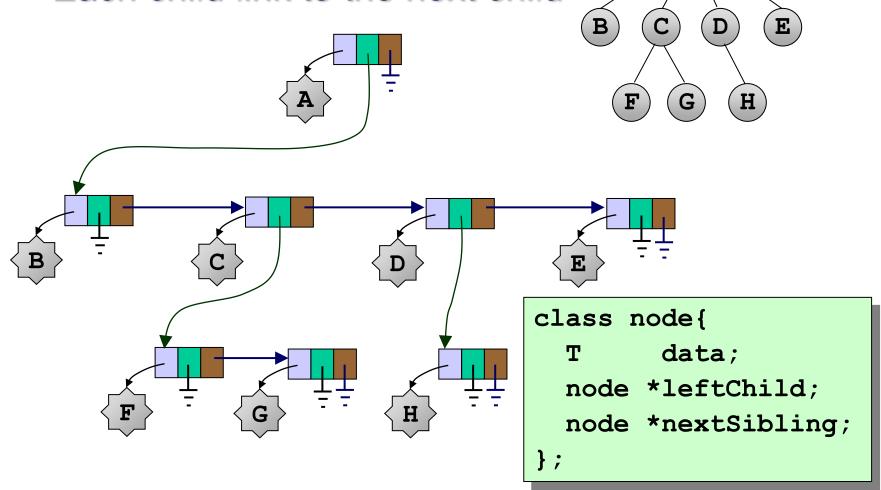
#### Tree Implementation: Use array to store children



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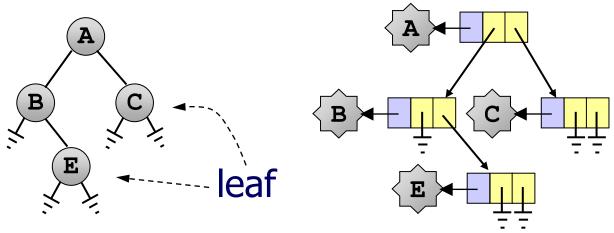
#### Tree Implementation: Use list to store children

- Store link to the left most child
- Each child link to the next child



### Binary Tree

Each node has two children: left and right



```
class node {
   T    data;
   node *left, *right;

   node(T data, node *left, node *right) :
      data(data), left(left), right(right)
   {}

   bool isLeaf() {
      return left==NULL && right==NULL;
};
```

#### **Huffman Code**

### งึกงึกงักงักมันเป็นงึกงึกงักงัก งึกงึกงักงักมันเป็นงึกงึกงักงัก มันเป็นกะอีกกะอักมันเป็นจึกจึกจั๊กจั๊ก

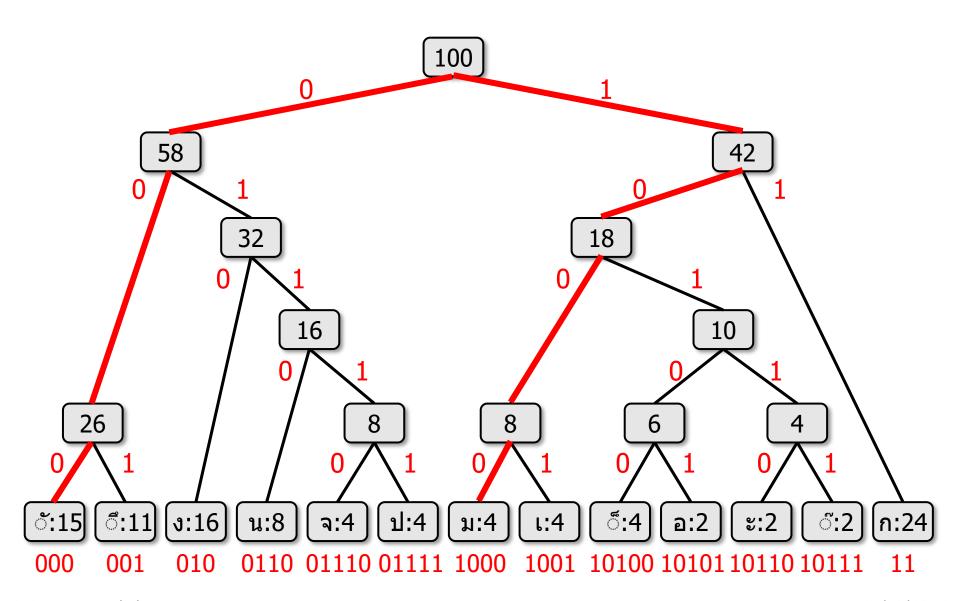
ก	J	े	ឺ	น	ব	ป	ม	l	េរា	อ	66	<b>్</b>
24	16	15	11	8	4	4	4	4	4	2	2	2
0000	0001	0010	1010	0011	0100	0101	0110	0111	1000	1001	1011	1100
11	010	000	001	0110	01110	01111	1000	1001	10100	10101	10110	10111

$$4 \times (24 + 16 + 15 + 11 + 8 + 4 + 4 + 4 + 4 + 4 + 2 + 2 + 2) = 400 \text{ bits}$$

$$2 \times 24 + 
3 \times (16 + 15 + 11) + 
4 \times (8 + 4 + 4) + 
5 \times (4 + 4 + 4 + 2 + 2 + 2)$$
= 328 bits

No code is a prefix of another code

### How to find Huffman code?



### huffman\_tree

```
class huffman tree {
  class node {
                                                  21
                                               ม
   public:
      char c;
      int freq;
      node *left, *right;
      node(char c, int freq, node *left, node *right) :
        c(c), freq(freq), left(left), right(right)
      { }
  };
  node *root;
```

### huffman\_tree: ctor

```
class huffman tree {
  class node { ... }
                                                 21
  node *root;
public:
  huffman tree(char c[], int f[], size t n) {
    class freq comp {
      public:
        bool operator()(node *a, node *b) {
            return a->freq > b->freq;
    };
    priority queue<node*, vector<node*>, freq comp > h;
    for (size t i=0; i<n; ++i) {
       h.push(new node(c[i], f[i], NULL, NULL));
    for (size t i=0; i<n-1; ++i) {
       node *n1 = h.top(); h.pop();
       node *n2 = h.top(); h.pop();
       h.push(new node('*', n1-)freq+n2-)freq, n1, n2);
    root = h.top();
```

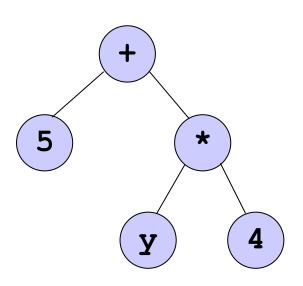
### huffman\_tree: dtor

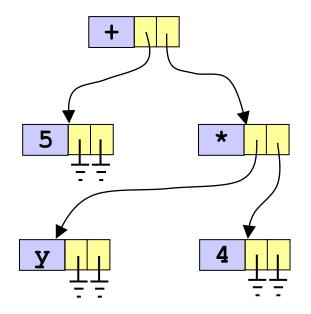
```
class huffman tree {
  node *root;
public:
  huffman tree(char c[], int f[], int n) { ... }
 ~huffman tree() {
    delete all node( root );
  void delete all nodes(hnode *r) {
    if (r == NULL) return;
    delete all nodes(r->left);
    delete all nodes(r->right);
    delete r;
```

## **Expression Tree**

- Represent expression as tree
- Leaf: operand
- Inner Node : operator

$$5 + y * 4$$



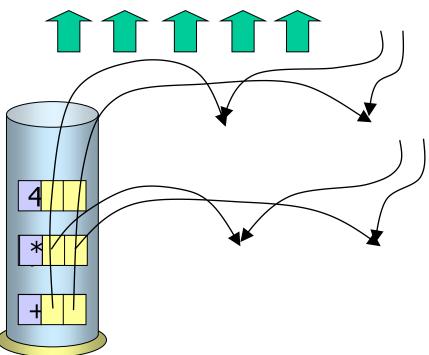


### **Expression Tree Construction**

- Operand: push to stack
- Operator: pop two nodes from stack to be the children of the new node, then push to stack

infix: 5 + y \* 4

postfix : 5 y 4 \* +

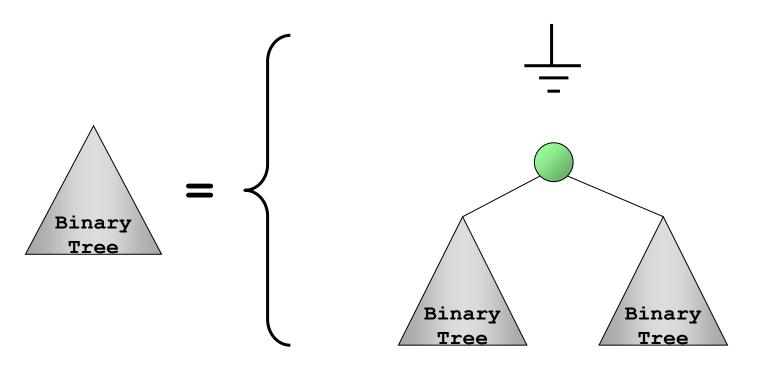


#### expression\_tree

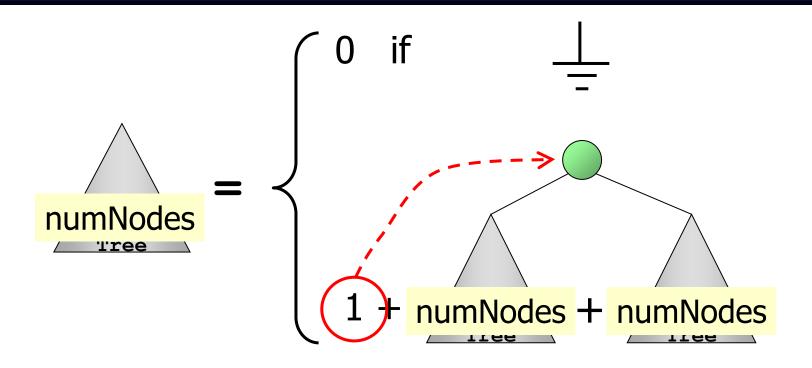
```
node *expression tree(string &infix) {
    string postfix = infix2postfix(infix);
    stack<node*> s:
    for (size t i=0; i<postfix.size(); i++) {</pre>
      char token = postfix[i];
      if (!isOperator(token)) {
        s.push(new node(token, NULL, NULL));
      } else {
        node *right = s.top(); s.pop();
        node *left = s.top(); s.pop();
        s.push(new node(token, left, right));
    return s.top();
```

### Binary tree: recursive view

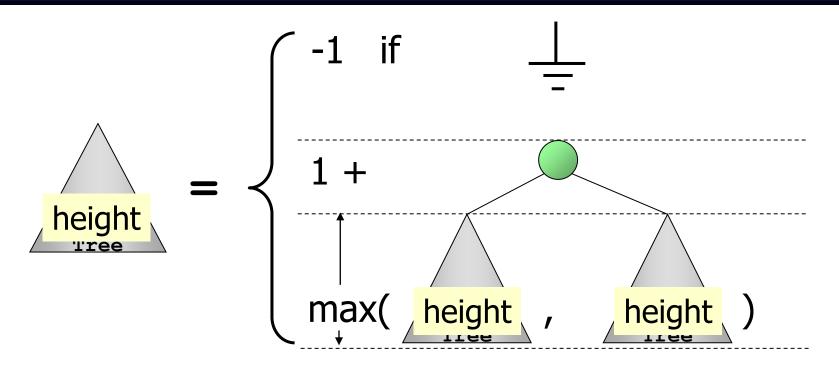
- Binary tree is
  - Empty tree (NULL) or
  - A node with binary trees as left and right children



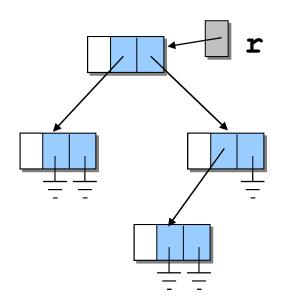
### Counting number of nodes



### Finding height



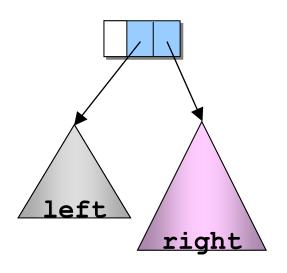
# Copying Tree

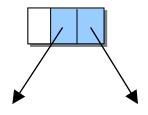




node 
$$*rr = copy(r);$$

### Copying tree (Recursive View)





```
node *copy(node *r) {
   if (r == NULL) return NULL;
   node *rL = copy(r->left);
   node *rR = copy(r->right);
   return new node(r->data, rL, rR);
}
```

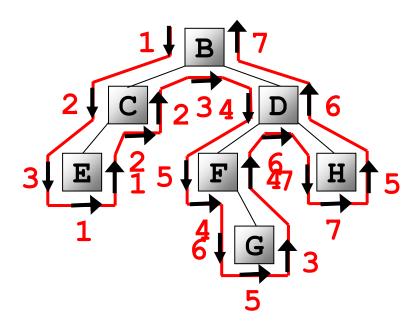
#### Tree traversal

 Tree traversal: mechanism to visit all nodes in the tree once per node

Preorder (แบบก่อนลำดับ)
 B, C, E, D, F, G, H

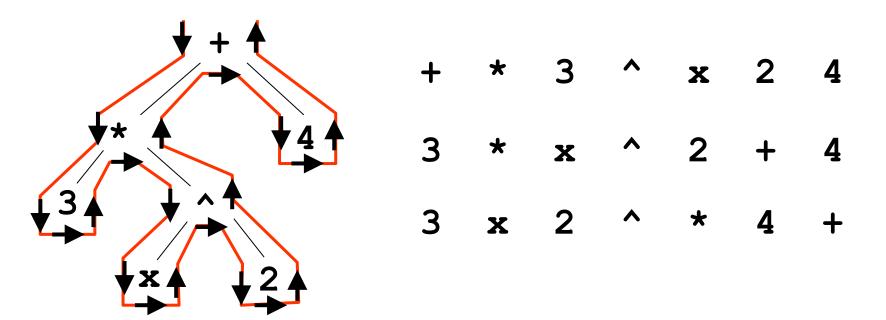
Inorder (แบบตามลำดับ)E, C, B, F, G, D, H

– Postorder (แบบหลังลำดับ)E, C, G, F, H, D, B



### Expression tree traversal

- Preorder traversal => Prefix expression
- Inorder traversal => Infix expression
- Postorder traversal => Postfix expression



#### Tree traversal with node x as root

#### Preorder

- visit x
- traverse left
- traverse right

#### Inorder

- traverse left
- visit x
- traverse right

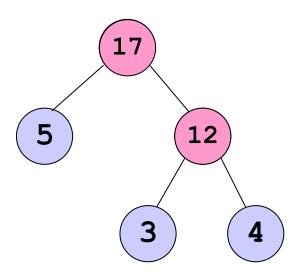
#### Postorder

- traverse left
- traverse right
- visit x

```
void preorder(node *x) {
  if (x == NULL) return;
 visit(x->data);
 preorder(x->left);
 preorder(x->right);
   void inorder(node *x) {
     if (x == NULL) return;
     inorder(x->left);
     visit(x->data);
     inorder(x->right);
        void postorder(node *x) {
          if (x == NULL) return;
          postorder(x->left);
          postorder(x->right);
          visit(x->data);
```

### Expression tree evaluation

- Need to know the values of children first
- Similar to postorder traversal



### **Expression Tree Evaluation**

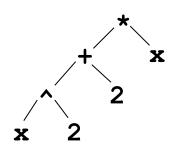
```
float eval(node *r) {
    if (r == NULL) return 0;
    if (r->isLeaf()) return (r->data - '0');
    float vLeft = eval(r->left);
    float vRight = eval(r->right);
    if (r->data == '+') return vLeft + vRight;
    if (r->data == '-') return vLeft - vRight;
    if (r->data == '*') return vLeft * vRight;
    if (r->data == '/') return vLeft / vRight;
    if (r->data == '^') return pow(vLeft, vRight);
    exit(-9); // Should never reach this line
```

#### Differentiation

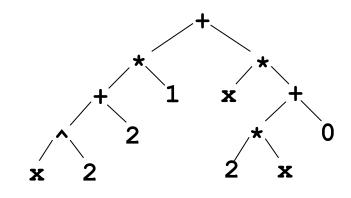
$$f(x) = (x^{2} + 2)x$$

$$f'(x) = (x^{2} + 2) \cdot 1 + x \cdot (2x + 0)$$

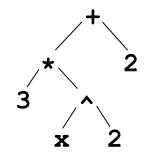
$$= 3x^{2} + 2$$







diff(f)



simplify(f)

#### Rules

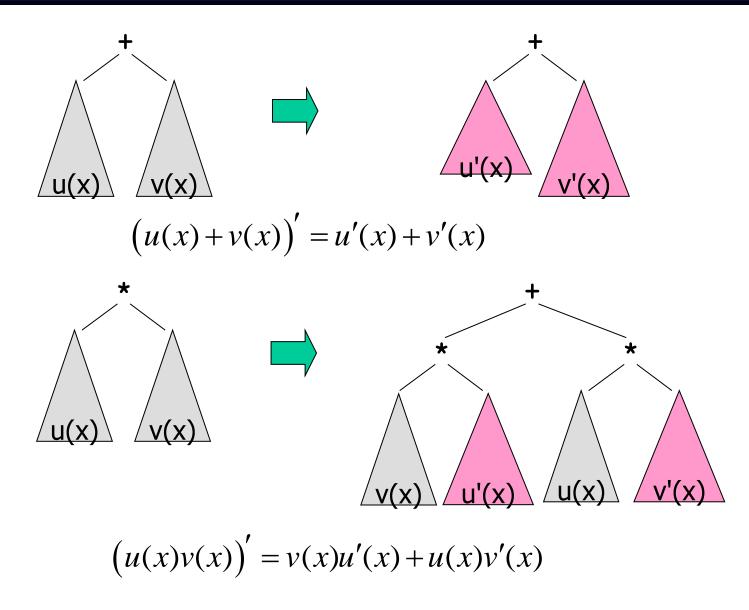
$$(u(x) + v(x))' = u'(x) + v'(x)$$

$$(u(x)v(x))' = v(x)u'(x) + u(x)v'(x)$$

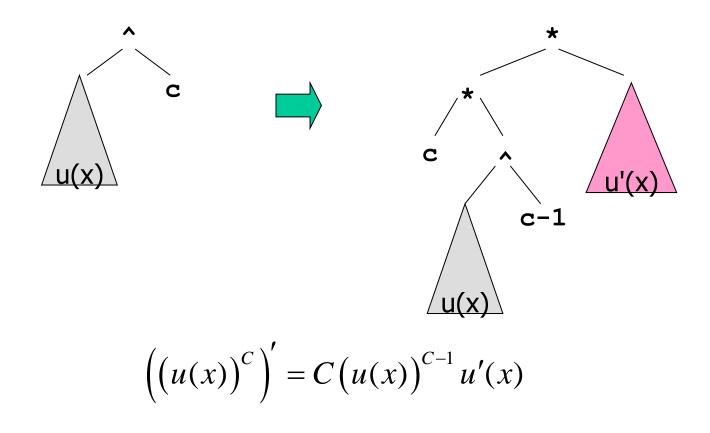
$$\left(\frac{u(x)}{v(x)}\right)' = \frac{v(x)u'(x) - u(x)v'(x)}{(v(x))^2}$$

$$((u(x))^C)' = C(u(x))^{C-1}u'(x)$$

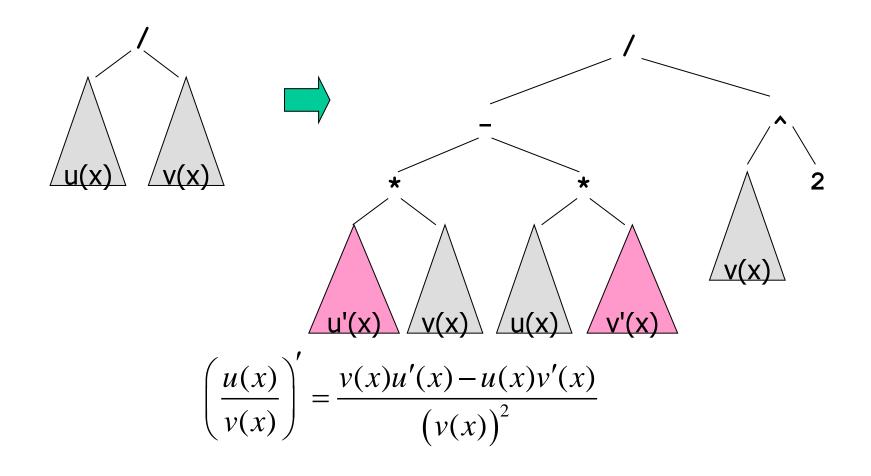
### Derivative of expression trees: +, \*



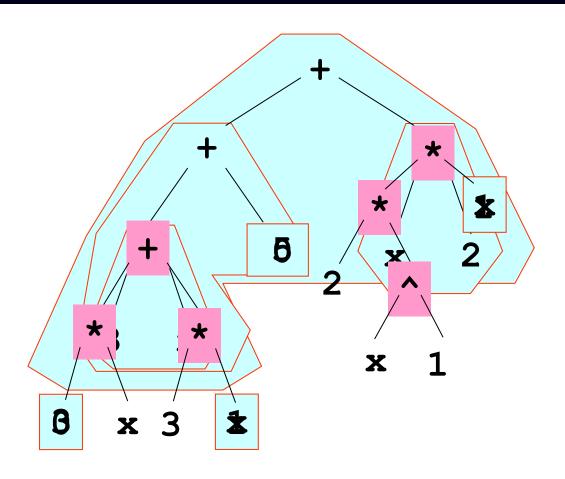
## Derivative of expression tree: ^



# Derivative of expression tree:/



### Example of expression tree differentiation



#### Differentiation

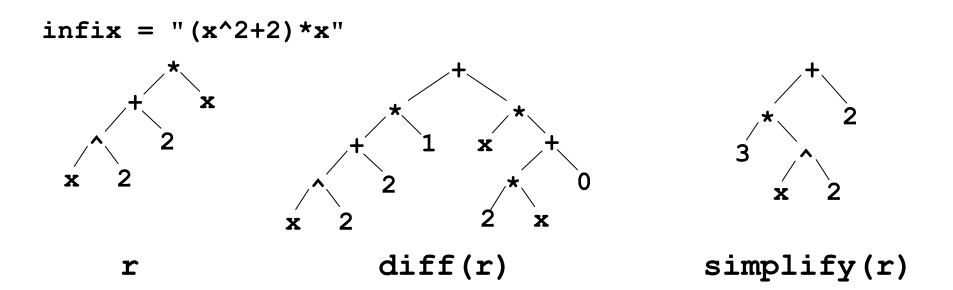
```
node *diff(node *r) {
    if (r == NULL) return NULL;
    char s = r->data;
    if ( r.isLeaf() ) {
      r->data = (s == 'x' ? '1' : '0');
    } else {
     if (s == '+') r = diffSum(r);
     else if (s == '-') r = diffSum(r);
     else if (s == '^') r = diffPow(r);
     else if (s == '*') r = diffMult(r);
     else if (s == '/') r = diffDiv(r);
   return r;
```

### Differentiation: +, \*

```
node *diffSum(node *r) {
                                      \left| \left( u(x) + v(x) \right)' = u'(x) + v'(x) \right|
    r->left = diff(r->left);
    r->right = diff(r->right);
    return r;
node *diffMult(node *r) {
                                                            du
                                                 dv
    node *u = copy(r->left);
    node *v = copy(r->right);
    node *du = diff(r->left);
    node *dv = diff(r->right);
    node *t1 = new node('*', u, dv);
    node *t2 = new node('*', v, du);
    return new node('+', t1, t2);
                                (u(x)v(x))' = v(x)u'(x) + u(x)v'(x)
```

### Put everything into code

```
char infix[100];
scanf("%s", infix);
node *r = newExpressionTree(infix);
printInorder(r); printPostorder(r);
r = diff(r);
r = simplify(r);
printInorder(r); printPostorder(r);
```



### Summary

- > Tree is a kind of data structure
  - ➤ Can be constructed by linking nodes
- ➤ Binary tree is a tree whose nodes has 2 children
  - > Tree consists of small trees
  - Most operations can be naturally expressed with recursive functions
  - Tree operations can be done by preorder, inorder, or postorder traversal