

# Secure Chat Application Design

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# Notation

- $A, B$  = Clients
- $S$  = Server of the system
- $PUZ$  = The Puzzle given by Server.
- $SOL$  = Solution to Puzzle
- $M$  = message
- $N$ 's = Nonces
- $t$  = timestamp
- $P$  = Public Key



# Features

- Authentication
- Confidentiality
- Integrity
- Non Repudiation
- Identity Hiding
- DoS Resistance
- Perfect Forward Secrecy
- Password Cracking Resistance
- End - End Encryption

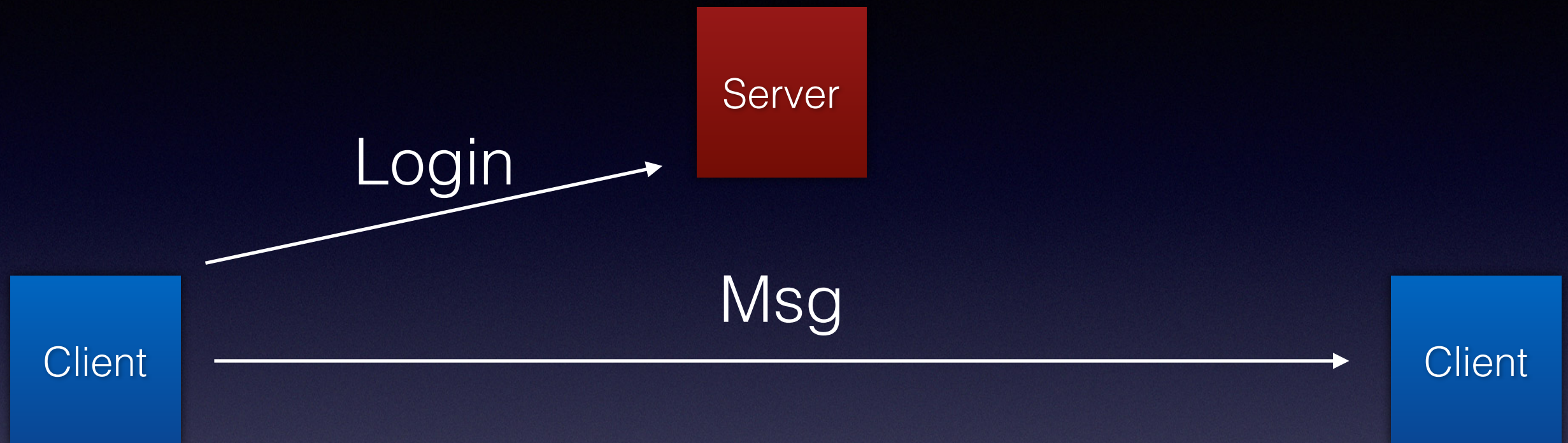


# Changes

- Added an extra assumption
- Added an extra accept/reject message in login which should be obvious, but forgot to add last time.
- Added some more additional details like broadcasting logout by server which was also missed last time.
- Added Diffie Hellman Algorithm in Crypto



# Architecture



- We have a Central Server that handles logins of clients.
- The Clients Communicate with each other directly.



# Crypto Algorithms

- Symmetric Encryption = AES-GCM 256 bit
- RSA Key Size = 2048 bits
- Diffie Hellman = Elliptic Curve SECP256R1
- Password Storage = PBKDF2
- KDF = PBKDF2 with SHA 512



# Assumptions

- Each Client has the correct public key of Server.
- The clocks of all clients and server is synchronized and this synchronization cannot be interrupted by adversary.
- The Server is trusted by all clients after the login.
- The Logout Broadcast Reaches Clients.

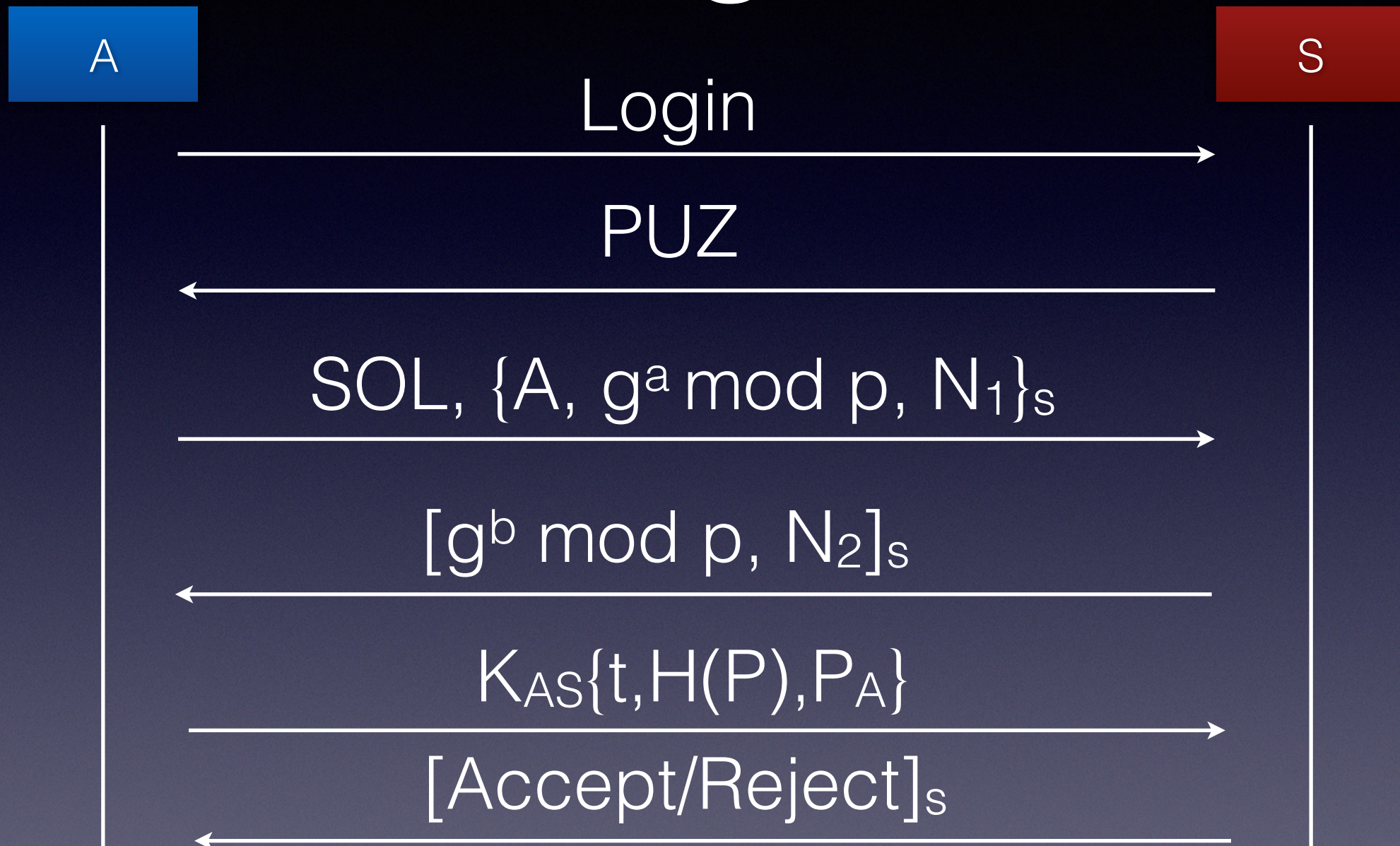


# Important Info

- Every Packet has a Timestamp that the Server or Client uses to prevent replay.
- Most of the Packets have a Signature that is calculated by using the timestamp and other fields.
- All Sockets are UDP
- Active Clients send Heartbeats every 30s.



# Login





# Login contd...

- $H(P)$  is the output of KDF in Password Hashing.
- $N_1$  and  $N_2$  are used for  $K_{AS}$  generation.
- $t$  should be equal to timestamp in last packet.

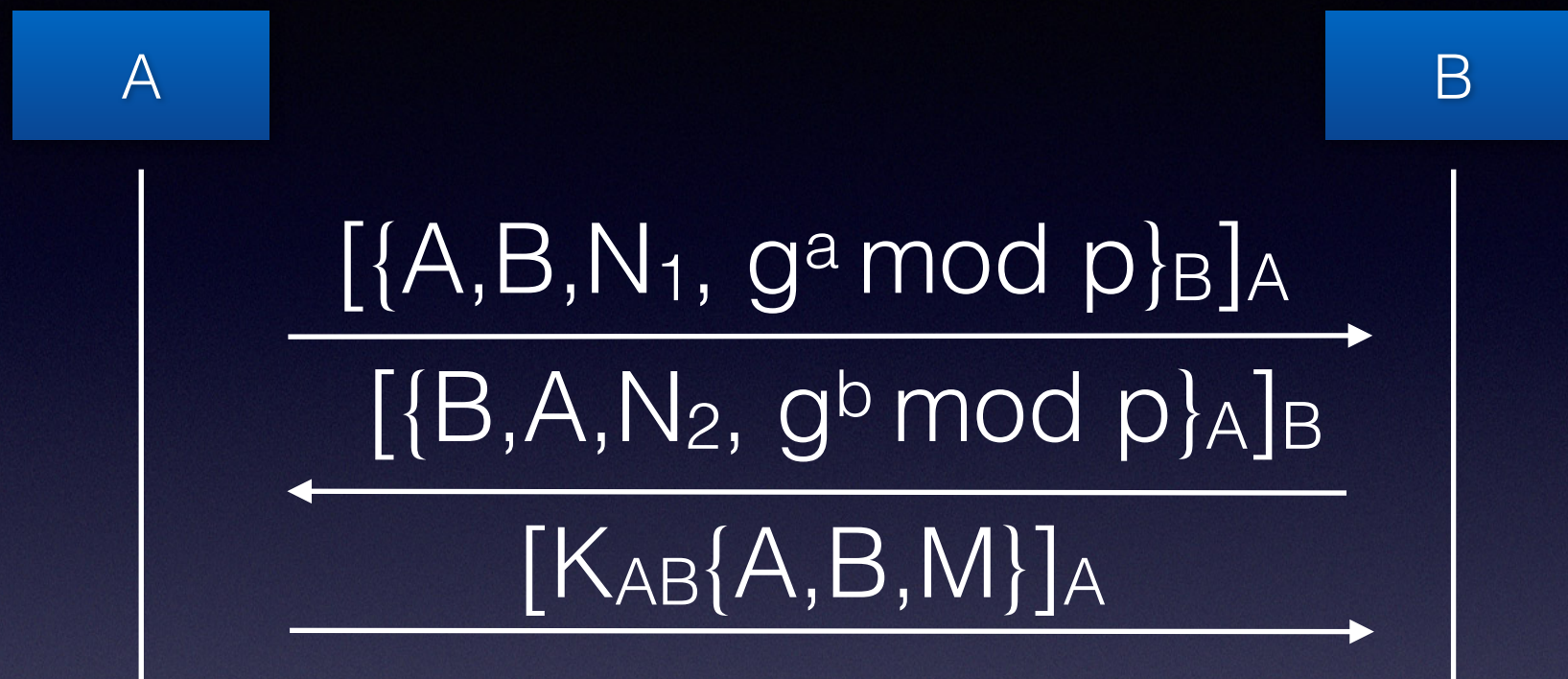


# Client Puzzle

- The Server prepares a Puzzle Certificate with  $[t, D, N_s]_s$ 
  - $t$  is the expiry timestamp of certificate.
  - $D$  is the difficulty of the puzzle
  - $N_s$  is the Server's Nonce.
- When Client requests to login, the server sends the certificate which is precomputed for  $D$  and  $N_s$
- Then the Client verifies the certificate and computes  $H(N_s, N_c, X)$  where  $N_c$  is the client Nonce and  $X$  is the solution. The client sends  $N_c, X, H(N_s, N_c, X)$  to the Server.
- The Server checks if  $N_c$  is not repeat for the current  $N_s$  and then computes  $H(N_s, N_c, X)$ , if the first  $D$  bytes are 0 then the solution is correct and the Server continues with Authentication else it rejects it.
- $N_s$  is changed after a certain amount of time after which the certificate is recomputed. If the Server is attacked then the certificate is remade for every change to  $D$ .
- When  $N_s$  is changed then all  $N_c$ 's which were recorded will be dropped.
- We modified the puzzle from this paper <http://www.tcs.hut.fi/old/papers/aura/aura-nikander-leiwo-protocols00.ps>



# Message



- A sends  $N_1$  and its DH contribution which is signed by A and encrypted using  $P_B$
- B sends back similar packet if initial packet is authentic.
- $K_{AB}$  is derived from DH,  $N_1$  and  $N_2$  using Key generation algorithm.
- We will use timestamps or sequence numbers to keep track of duplicates or replays.
- Accept is sent by B if message is received successfully from A



# List



- A can also ask for specific clients details.
- A uses this to get public keys of other clients.



# Logout



- The Server uses IP:PORT to identify user in cases like this.
- After send an OK, the server broadcasts a message stating that A has logged out
- The Broadcast message looks like  $[\text{IP:PORT, LOGOUT}]_S$



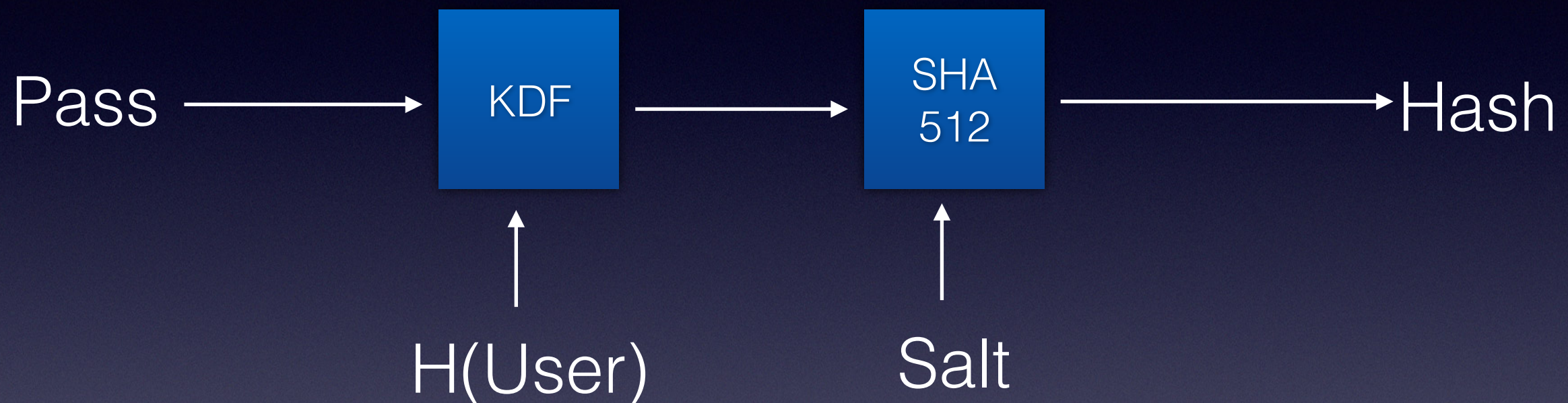
# Heartbeat



- If the Server doesn't get HeartBeat from Client for 60s then it is logged out and a message is broadcasted indicating A has logged out.



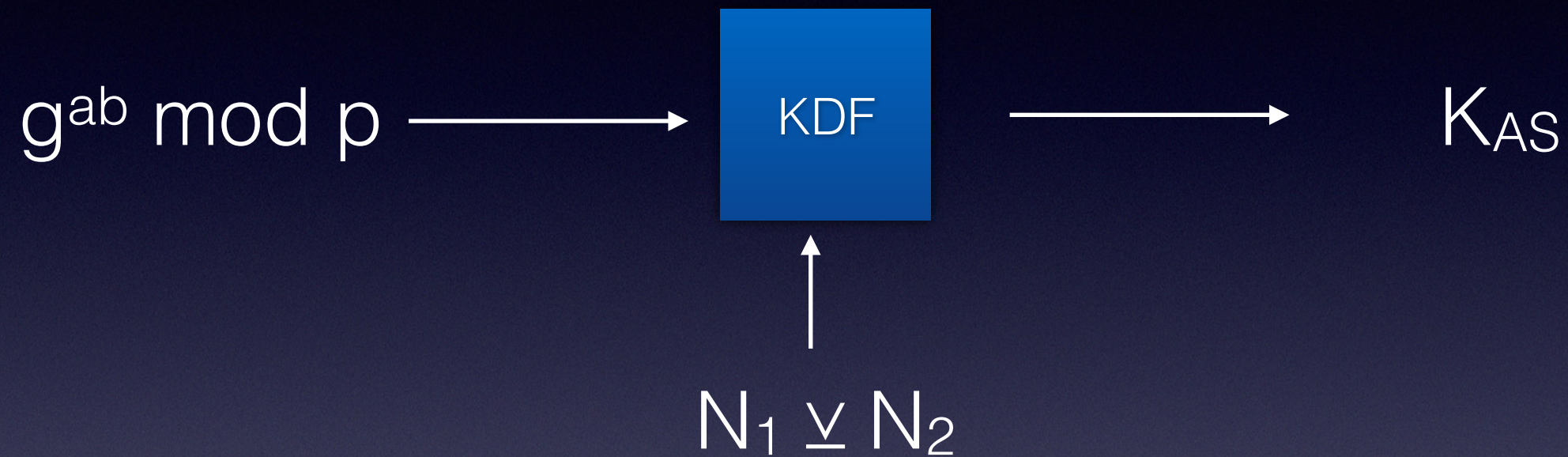
# Password Hash Generation



- The Final Hash is stored on Server along with Salt. Server stores this information in a database.
- The Client computes the KDF ( $H(P)$ ) and sends it to Server where the Server computes the Hash and compares.



# Key Generation



- The output of DH is given to KDF with  $N_1 \neq N_2$  as salt to get  $K_{AS}$



# Packet Format

1      256 / 32    256 / 64      4      16    16      1000 bytes



- The Entire Grey Region is Encrypted by Key.
- Encrypted Key has  $K + IV + \text{Tag}$  , if using RSA, else it has  $IV + \text{Tag}$  only.
- Signature = Sign[timestamp, Grey Region]
- Server or Client Checks Timestamp first after which it verifies the signature. Encrypt then Sign Methodology will be used.
- After the check is passed then the server will decrypt the message.



# Stretch Goals

- Dynamically Change the Difficulty of the Puzzles
- Network Reliability like resending a message if it was rejected by server.