

# COMP3400

## Assignment 2 Written

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### Tail Recursion

**Question 1.** [2 MARKS]

Produce an equation which computes  $S_{n+1}^2$  from *only*  $(n, S_n^2, \bar{x}_n, \bar{x}_{n+1}, x_{n+1})$

*Answer:* We have

$$S_n^2 = \left( \frac{1}{n} \sum_{k=1}^n x_k^2 \right) - \bar{x}_n^2 \quad (1)$$

$$\implies n(S_n^2 + \bar{x}_n^2) = \sum_{k=1}^n x_k^2 \quad (2)$$

$$\implies (n+1)(S_{n+1}^2 + \bar{x}_{n+1}^2) = \sum_{k=1}^{n+1} x_k^2 \quad (3)$$

Taking Equation (3) – Equation (2) yields

$$(n+1)(S_{n+1}^2 + \bar{x}_{n+1}^2) - n(S_n^2 + \bar{x}_n^2) = x_{n+1}^2$$

and rearranging gives

$$S_{n+1}^2 = \frac{x_{n+1}^2 + n(S_n^2 + \bar{x}_n^2)}{n+1} - \bar{x}_{n+1}^2$$

1 mark for obtaining an expression for  $S_{n+1}^2$

1 mark for showing work

**Question 2.** [2 MARKS]

```
lcVariance :: [Float] -> Float
lcVariance (x:[]) = 0
lcVariance (x:xs) = (x**2 + n*(var' + xb**2))/(n+1) - xb'**2
  where
    xb  = (/) (sum xs) n
    xb' = (/) (sum (x:xs)) (n+1)
    var' = lcVariance xs
    n    = (fromIntegral $ length xs)
```

**Note.** Solution *must* be linear recursive in order to receive *any* marks.

1 mark for base case

1 mark for recursive step

**Question 3.** [4 MARKS]

```
1 trVariance :: Float -> Float -> Float -> [Float] -> Float
2 trVariance _ s _ [] = s
3 trVariance n s xb (x:xs) = trVariance (n+1) s' xb' xs
4   where
5     s'  = (x**2 + n*(s + xb**2)) / (n+1) - xb'**2
6     xb' = (x + n*xb)/(n+1)
```

**Note.** Solution *must* be tail recursive in order to receive *any* marks.

1 mark for base case

1 mark for updating variance

1 mark for updating mean

0.5 mark for updating  $n$

0.5 mark for recursing on tail of input

**Question 4.** [1 MARK]

```
7 variance :: [Float] -> Float
8 variance [] = undefined
9 variance xs = trVariance 0 0 0 xs
```

**Note:** Students *do not* need line 8.

1 mark for line 9

**Question 5.** [1 MARK]

$$\begin{aligned} & \text{Invariant}(\text{trVariance } n \text{ } s \text{ } xb \text{ } xs) \\ &= \left( \frac{n \cdot s + n \cdot xb^2 + \text{sqsum } xs}{n + \text{length } xs} \right) - \left( \frac{n \cdot xb + \text{sum } xs}{n + \text{length } xs} \right)^2 \end{aligned}$$

where  $\text{sqsum} = \text{sum } \$ \text{ map } (**2)$ .

1 mark for line 9 (or anything equivalent)

Notice this implies

$$\begin{aligned} & \text{trVariance } 0 \text{ } 0 \text{ } 0 \text{ } xs \\ &= \left( \frac{0 \cdot s + 0 \cdot xb^2 + \text{sqsum } xs}{0 + \text{length } xs} \right) - \left( \frac{0 \cdot xb + \text{sum } xs}{0 + \text{length } xs} \right)^2 && \text{By invariant} \\ &= \frac{\text{sqsum } xs}{\text{length } xs} - \left( \frac{\text{sum } xs}{\text{length } xs} \right)^2 \\ &= \text{variance } xs && \text{By (1)} \end{aligned}$$

If student's invariant doesn't have this property it is wrong.

**Question 6.** [5 MARKS]

Prove  $\text{trVariance}$  satisfies your iteration invariant.

$$\begin{aligned} & \text{LHS}(\text{line 2}) \\ &= \text{trVariance } n \text{ } s \text{ } xb \text{ } [] \\ &= \left( \frac{n \cdot s + n \cdot xb^2 + \text{sqsum } []}{n + \text{length } []} \right) - \left( \frac{n \cdot xb + \text{sum } []}{n + \text{length } []} \right)^2 && \text{By invariant} \\ &= \left( \frac{n \cdot s + n \cdot xb^2 + 0}{n + 0} \right) - \left( \frac{n \cdot xb + 0}{n + 0} \right)^2 \\ &= s + xb^2 - xb^2 \\ &= s \\ &= \text{RHS}(\text{line 2}). \end{aligned}$$

2 marks

Let

$$s' = \frac{x^2 + n \cdot (s + xb^2)}{n + 1} - xb'^2 \implies (n + 1) \cdot s' = x^2 + n \cdot (s + xb^2) - (n + 1) \cdot xb'^2$$

and

$$xb' = \frac{x + n \cdot xb}{n + 1} \implies (n + 1) \cdot xb' = x + n \cdot xb$$

in

RHS(line 3)

$$= \text{trVariance } (n+1) \ s' \ x b' \ x s$$

$$= \left( \frac{(n+1) \cdot s' + (n+1) \cdot x b'^2 + \text{sqsum } x s}{(n+1) + \text{length } x s} \right) - \left( \frac{(n+1) \cdot x b' + \text{sum } x s}{(n+1) + \text{length } x s} \right)^2$$

By invariant

$$= \left( \frac{x^2 + n \cdot (s + x b^2) - (n+1) \cdot x b'^2 + (n+1) \cdot x b'^2 + \text{sqsum } x s}{(n+1) + \text{length } x s} \right)$$

$$- \left( \frac{x + n \cdot x b + \text{sum } x s}{(n+1) + \text{length } x s} \right)^2$$

Substitution by definition of  $s'$  and  $x b'$

$$= \left( \frac{n \cdot (s + x b^2) + \text{sqsum } (x : x s)}{n + \text{length } (x : x s)} \right) - \left( \frac{n \cdot x b + \text{sum } (x : x s)}{n + \text{length } (x : x s)} \right)^2$$

Simplification and by definition of sum, sqsum, and length

$$= \left( \frac{n \cdot s + n \cdot x b^2 + \text{sqsum } (x : x s)}{n + \text{length } (x : x s)} \right) - \left( \frac{n \cdot x b + \text{sum } (x : x s)}{n + \text{length } (x : x s)} \right)^2$$

$$= \text{trVariance } n \ s \ x b \ (x : x s)$$

By invariant

$$= \text{LHS}(\text{line 3})$$

2 marks

Therefore  $\text{trVariance}$  satisfies the iteration invariant.

1 mark for justifying lines

**Question 7.** [1 MARK]

$$\text{BV}(\text{trVariance } n \ s \ x b \ x s) = \text{length } x s$$

1 mark

**Question 8.** [2 MARKS]

By definition  $\text{length } x s$  is always nonnegative.

$$\text{BV}(\text{trVariance } n \ s \ x b \ (x : x s))$$

$$= \text{length } (x : x s)$$

$$= 1 + \text{length } x s$$

By definition

$$\begin{aligned}
&> \text{length } xs \\
&= \text{BV}(\text{trVariance } (n + 1) \text{ s}' \text{ x b}' xs)
\end{aligned}$$

Thus the bound value is strictly descending and always nonnegative. 2 mark

**Question 9.** [2 MARKS]

For  $n$  an integer...

```

variance xs = variance $ map (+n) variance
variance xs = n^2 * variance $ map (*n) variance
variance xs = variance $ reverse xs

```

Note: The solution doesn't have to be a valid quick check. It only needs to adequately communicate what property is being checked.

1 mark for Quickcheck 1

1 mark for Quickcheck 2

# Induction

Consider the following definitions for implementing addition on natural numbers.

```
1 data Nat = Zero | Succ Nat deriving Show
2 plus :: Nat -> Nat -> Nat
3 plus m Zero = m
4 plus m (Succ(n)) = Succ (plus m n)
```

## Question 10. [7 MARKS]

Using induction prove:

$$\text{plus Zero } n = n$$

for any  $n :: \text{Nat}$ .

Answer: Let  $P :: \text{Nat} \rightarrow \text{Bool}$  where

1 mark for defining  $P(n)$

$$P(n) \iff \text{plus Zero } n = n$$

Base case:  $\text{plus Zero Zero} = \text{Zero}$  by line 3 and therefore  $P(\text{Zero})$ .

1 mark for base case

Assume  $P(n)$  as an induction hypothesis and notice this implies

1 mark for IH

$$\begin{aligned} &\text{plus Zero (Succ } n) \\ &= \text{Succ (plus Zero } n) \\ &= \text{Succ } n \end{aligned}$$

apply line 4  
By IH

and therefore  $P(\text{Succ } n)$ .

2 marks for induction

By the PSI/PMI we have  $P(n)$  for all  $n :: \text{Nat}$ .

1 mark for conclusion

1 marks for justifying lines

**Question 11.** [8 MARKS]

Using induction prove that plus is commutative, that is,

$$\text{plus } m \ n = \text{plus } n \ m$$

for any  $n, m :: \text{Nat}$ .

*Answer:* Let  $m :: \text{nat}$  be arbitrary and  $P :: \text{Nat} \rightarrow \text{Bool}$  where

$$P(n) \iff \text{plus } m \ n = \text{plus } n \ m$$

1 mark for defining  $P(n)$

*Base case:*  $\text{plus } m \ \text{Zero} = m$  by line 3.  $\text{plus } \text{Zero } m = m$  by Question 10. Therefore  $P(\text{Zero})$ .

1 mark for defining  $P(n)$

Assume  $P(n)$  as an induction hypothesis and notice this implies

1 mark for IH

$\text{plus } m \ (\text{Succ } n)$	
$= \text{Succ } (\text{plus } m \ n)$	Apply line 4
$= \text{Succ } (\text{plus } n \ m)$	By IH
$= \text{plus } n \ (\text{Succ } m)$	Unapply line 4
$= \text{plus } (\text{Succ } n) \ m$	Given

and therefore  $P(\text{succ } n)$ .

3 marks for induction

By the PSI/PMI we have  $P(n)$  for all  $n :: \text{Nat}$ .

1 mark for conclusion

1 marks for justifying lines