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## Lambda Calculus

Question 1:

Not =  $\lambda a.(a \text{ False}) (True)$   
=  $\lambda a.(a (\lambda x(\lambda y.y)) (\lambda x(\lambda y.x)))$

Question 2:

Lambda expression for  $\neg(\neg p \wedge \neg q)$  when let  $a = p$  and  $b = q$  is:

NOT (NOT p AND NOT q)

Write this expression in prefix position, we have:

```
NOT (AND (NOT p NOT q))
= \a.(a (False) (True)) (\a.(\b.(ab)a) (\a.(a (False) (True)) p \a.(a (False) (True)) q))
= (a (False) (True)) [a := (\a.(\b.(ab)a) (\a.(a (False) (True)) p \a.(a (False) (True)) q)]]
= (\a.(\b.(ab)a) (\a.(a (False) (True)) p \a.(a (False) (True)) q)) (False) (True)
= (\a.(\b.(ab)a) (\a.(a (False) (True)) p \a.(a (False) (True)) q)) (False) (True)
= (\a.(\b.(ab)a) ((a (False) (True))[a:= p] (a (False) (True)) [a:=q])) (False) (True)
= (\a.(\b.(ab)a) ((p (False) (True)) (q (False) (True)))) (False) (True)
= ((\b.(ab)a) [a = (p (False) (True))] (q (False) (True))) (False) (True)
= (\b.((p (False) (True))b)(p (False) (True)) (q (False) (True))) (False) (True)
= (((p (False) (True))b)(p (False) (True)) [b:= (q (False) (True))] ) (False) (True)
= (((p(False)(True)) (q(False)(True))) (p(False)(True))) (False) (True)
```

Let this be (1)

We also have (2):

- True True True  
= (\x.(\y.x)) True True  
= x [x := True] [y := True]  
= True
- False False False  
= (\x.(\y.y)) False False  
= y [x := False] [y := False]  
= False
- False True False  
= (\x.(\y.y)) True False  
= y [x := True] [y := False]  
= False
- False False True  
= (\x.(\y.y)) False True  
= y [x := False] [y := True]  
= True

From (1) there are 4 cases of p and q which is simpler to solve by substitute (2) in:

Case 1: p = False and q = False

```
((False(False)(True)) (False(False)(True))) (False(False)(True)) (False) (True)
= ((True) (True) (True)) (False) (True)
= False
```

Case 2: p = True and q = False

```
((True(False)(True)) (False(False)(True))) (True(False)(True)) (False) (True)
= ((False) (True) (False)) (False) (True)
= True
```

Case 3: p = False and q = True

```
((False(False)(True)) (True(False)(True))) (False(False)(True)) (False) (True)
= ((True) (False) (True)) (False) (True)
```

= True

Case 4:  $p = \text{True}$  and  $q = \text{True}$

= (((True(False)(True)) (True(False)(True))) (True(False)(True))) (False) (True)

= ((False) (False) (False)) (False) (True)

= True

In summary, by grouping these cases and comparing them to OR ( $p \vee q$ ) as the truth table below, we can make the conclusion that both of them are equivalent.

p	q	$\neg(\neg p \wedge \neg q)$	$(p \vee q)$
T	T	T	T
T	F	T	T
F	T	T	T
F	F	F	F

Question 3:

$$\begin{aligned} &= (\lambda xy.x)(\lambda abc.cab)z(\lambda z.zz) \\ &= (\lambda xy.x)[x:=\lambda abc.cab]z(\lambda z.zz) \\ &= (\lambda y.(\lambda abc.cab))z(\lambda z.zz) \\ &= ((\lambda abc.cab))[y:=z](\lambda z.zz) \\ &= (\lambda abc.cab)(\lambda z.zz) \\ &= (\lambda bc.cab)[a:=\lambda z.zz] \\ &= \lambda bc.(c(\lambda z.zz)b) \end{aligned}$$

## Principle Types

(These questions are coded in file PrincipalType.hs)

## Blockus (Question 10)

We can use The Principle of Mathematical Induction theorem to prove a  $2^n \times 2^n$  Blockus board with northwest corner removed can be covered with V3 pieces (3 blocks).

Logically, by removing 1 block out of  $2^n \times 2^n$  Blockus board we can have:

$$2^n \times 2^n - 1 \text{ (blocks)}$$

Moreover, to make sure that all V3 pieces can cover the whole board, we need the board is not contained a straight line of at least 2 blocks so that it can fit at least 1 V3 piece (fortunately,  $2^n \times 2^n$  is always a square so that not necessary) and the number of remaining blocks is divisible by 3:

$$(2^n \times 2^n - 1) : 3 \text{ (blocks)}$$

$$\text{Let } P(n) \leftrightarrow (2^n \times 2^n - 1) : 3$$

Where base case is:

$$P(0) \leftrightarrow (2^0 \times 2^0 - 1) : 3 \leftrightarrow 0 : 3 \leftrightarrow T.$$

Induction Hypothesis:

Assuming that  $P(n) = T$  in  $P(n) \rightarrow P(n + 1)$  and showing  $P(n + 1)$  is a consequence of  $P(n)$ :

$$P(n + 1) \leftrightarrow (2^{n+1} \times 2^{n+1} - 1) : 3$$

Starting with the left-hand side of the expression:

$$\begin{aligned} & ((2^n \times 2^n) \times (2^1 \times 2^1) - 1) \\ &= 4 \times (2^n \times 2^n) - 1 \text{ (1)} \end{aligned}$$

Using the induction hypothesis, we already know that  $(2^n \times 2^n - 1)$  is divisible by 3. Therefore, the expression can be written as:

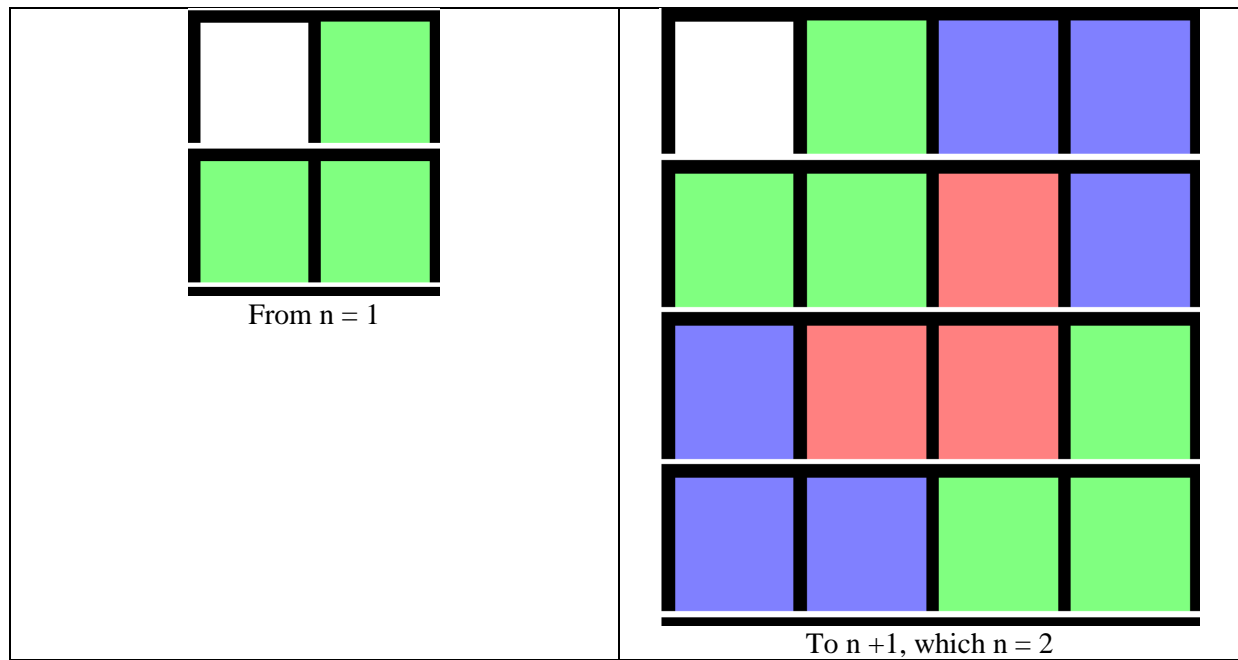
$$\begin{aligned} & (2^n \times 2^n) - 1 = 3 \times m, \text{ where } m \text{ is an arbitrary integer.} \\ & \leftrightarrow (2^n \times 2^n) = 3 \times m + 1 \text{ (2)} \end{aligned}$$

Substituting (2) into the equation (1) gives:

$$\begin{aligned} &= 4 \times (3 \times m + 1) - 1 \\ &= 12 \times m + 3 \end{aligned}$$

Thus,  $(2^{n+1} \times 2^{n+1} - 1)$  is equal to an integer  $m$  multiply by 3 which also divisible by 3.

For instance,



By the principle of mathematical induction, we have proved that  $(2^n \times 2^n - 1)$  is divisible by 3 for all positive integers  $n$ . Finally, there is always true that there is exist at least one order of V3 piece(s) that the square  $2^n \times 2^n$  Blockus board with northwest corner removed can be covered with it.