

I217E: Functional Programming

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Schedule

10/12	introduction	11/9	interpreters
10/14	algebraic data types I	11/11	compilers
10/19	algebraic data types II	11/16	termination
10/21	program reasoning	11/18	confluence
10/26	applications	11/25	verification
10/28	data structures I	11/30	review
11/2	data structures II	12/5	exam
11/4	computational models		

Evaluation

exam (60) + reports (40)

Lambda Expressions and Infix Operators

Lambda Expressions

$(\lambda x \rightarrow x + 1) 2 = 2 + 1 = 3$
 $(\lambda x y \rightarrow x * y) 2 3 = 2 * 3 = 6$

Infix Operators and Partial Applications

$(+) 10 2 = 10 + 2 = 12$
 $(/) 10 2 = 10 / 2 = 5.0$
 $(/ 2) 10 = 10 / 2 = 5.0$
 $(10 /) 2 = 10 / 2 = 5.0$

Higher-Order Functions

Definition

higher-order functions take/return functions

$\text{twice } f \ x = f (f \ x)$

$\text{twice } (*2) \ 1 = (*2) ((*2) \ 1)$

Map

$\text{map } f [x_1, \dots, x_n] = [f x_1, \dots, f x_n]$

```
map (* 10) [] = []
map (* 10) [-4] = [-40]
map (* 10) [3, -4] = [30, -40]
map (* 10) [-2, 3, -4] = [-20, 30, -40]
map (* 10) [1, -2, 3, -4] = [10, -20, 30, -40]
```

```
myMap :: (a -> b) -> [a] -> [b]
myMap f [] = ...
myMap f (x : xs) = ...
```

Filter

$\text{filter } p [x_1, \dots, x_n]$ only keeps x_i s with $p x_i == \text{True}$

```
filter (> 0) [] = []
filter (> 0) [-4] = []
filter (> 0) [3, -4] = [3]
filter (> 0) [-2, 3, -4] = [3]
filter (> 0) [1, -2, 3, -4] = [1, 3]
```

```
myFilter :: (a -> Bool) -> [a] -> [a]
myFilter p [] = ...
myFilter p (x : xs) | ... = ...
                    | ... = ...
```

Partition

```
partition (> 0) [] = ([], [])
partition (> 0) [-4] = ([], [-4])
partition (> 0) [3, -4] = ([3], [-4])
partition (> 0) [-2, 3, -4] = ([3], [-2, -4])
partition (> 0) [1, -2, 3, -4] = ([1, 3], [-2, -4])
```

```
partition p [] = ...
partition p (x : xs)
  | ... = ...
  | ... = ...
  where (ys, zs) = partition p ...
```

Fold Functions

$\text{foldl } (\oplus) e [x_1, \dots, x_n] = ((e \oplus x_1) \oplus x_2) \oplus \dots \oplus x_n$
 $\text{foldr } (\oplus) e [x_1, \dots, x_n] = x_1 \oplus \dots \oplus (x_{n-1} \oplus (x_n \oplus e))$

```
foldl (-) 10 [1, 2, 3] = ((10 - 1) - 2) - 3 = 4
foldr (-) 10 [1, 2, 3] = 1 - (2 - (3 - 10)) = -8
```

```
myFoldl :: (a -> b -> a) -> a -> [b] -> a
myFoldl f e [] = ...
myFoldl f e (x : xs) = ...
```

List Comprehension

Set Comprehension (in Mathematics)

$$\begin{aligned}\{x \times 10 \mid x \in \{1, -2, 3, -4\}\} &= \{10, -20, 30, -40\} \\ \{x + 1 \mid x \in \{1, -2, 3, -4\} \text{ and } x > 0\} &= \{2, 4\} \\ \{x + y \mid x \in \{10, 20\} \text{ and } y \in \{1, 2\}\} &= \{11, 12, 21, 22\}\end{aligned}$$

List Comprehension (in Haskell)

$$\begin{aligned}[x * 10 \mid x \leftarrow [1, -2, 3, -4]] &= [10, -20, 30, -40] \\ [x \mid x \leftarrow [1, -2, 3, -4], x > 0] &= [1, 3] \\ [x + y \mid x \leftarrow [10, 20], y \leftarrow [1, 2]] &= [11, 12, 21, 22]\end{aligned}$$

Quick Sort (Hoare 1960)

```
qsort [] = []
qsort (x : xs) = qsort [y | y <- xs, y < x] ++ [x] ++
                 qsort [y | y <- xs, y >= x]
```

Note

- beautiful and practically efficient algorithm
- $O(n^2)$ in worst case (why?)

Divide and Conquer

```
qsort [3, 4, 2, 5, 1]
=
qsort [2, 1] ++ [3] ++ qsort [4, 5]
=
(qsort [1] ++ [2] ++ qsort []) ++ [3] ++ (qsort [] ++ [4] ++ qsort [5])
=
([1] ++ [2] ++ []) ++ [3] ++ ([] ++ [4] ++ [5])
=
[1, 2] ++ [3] ++ [4, 5]
=
[1, 2, 3, 4, 5]
```

Merge Sort (van Neumann 1945)

$$\text{msort } [x_1, \dots, x_n] = \begin{cases} [] & \text{if } n = 0 \\ [x_1] & \text{if } n = 1 \\ \text{merge } (\text{msort } ys) (\text{msort } zs) & \text{otherwise} \end{cases}$$

where

- $(ys, zs) = ([x_1, x_3, x_5, \dots], [x_2, x_4, x_6, \dots])$
- **merge** *xs ys* merges two **sorted** lists

Note

- problem is divided into subproblems of **same size**
- $O(n \log n)$

Divide and Conquer

let $xs \otimes ys = \text{merge } xs \text{ } ys$

```
msort [3,4,2,5,1]
=
msort [3,2,1]  $\otimes$  msort [4,5]
=
(msort [3,1]  $\otimes$  msort [2])  $\otimes$  (msort [4]  $\otimes$  msort [5])
=
(( $[3] \otimes [1]$ )  $\otimes$  [2])  $\otimes$  ( $[4] \otimes [5]$ )
=
( $[1,3] \otimes [2]$ )  $\otimes$  [4,5]
=
 $[1,2,3] \otimes [4,5] = [1,2,3,4,5]$ 
```

Homework

1 Use `foldr` to implement `sumList`:

$$\text{sumList } [x_1, \dots, x_n] = x_1 + \dots + x_n$$

2 Use list comprehension to re-implement `oddplus1`:

```
oddplus1 xs =
  map (+ 1) (filter (\x -> mod x 2 == 1) xs)
```

3 Implement `merge :: [Int] → [Int] → [Int]`.

4 Implement `split :: [Int] → ([Int], [Int])`:

$$\text{split } [x_1, \dots, x_n] = ([x_1, x_3, \dots], [x_2, x_4, \dots])$$

5 Implement `msort :: [Int] → [Int]`.