

1. *Multiple access channel* Consider a multiple access channel with inputs $\mathcal{X}_1 = \{0, 1, 2, 3\}$ and $\mathcal{X}_2 = \{0, 1\}$ The channel is given by:

$$Y = X_1 + X_2 \bmod 4 \tag{1}$$

Find the capacity region for this channel.

2. *Slepian-Wolf* Let X_i and Z_i be independent random variables with $p_X(x) = [1-p, p]$ and $p_Z(z) = [1-r, r]$ for $0 \leq p, r \leq 1$. Let $Y_i = X_i \oplus Z_i$ where \oplus denotes addition modulo 2. Let the source vector $\mathbf{X} = (X_1, \dots, X_n)$ be encoded at rate R_1 and let $\mathbf{Y} = (Y_1, \dots, Y_n)$ be encoded at rate R_2 .
- What is $p_Y(y)$?
 - What region of rates allows recovering of \mathbf{X}, \mathbf{Y} with probability of error tending to zero as $n \rightarrow \infty$? Draw a pentagon-shaped region and label the key points.
 - On the same figure, draw the region of achievable rates, assuming the correlation between \mathbf{X} and \mathbf{Y} is ignored.

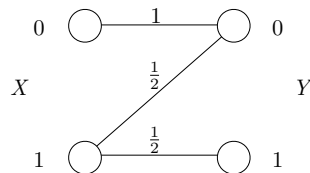
3. Suppose that (X, Y, Z) are jointly Gaussian and that $X \rightarrow Y \rightarrow Z$ forms a Markov chain. Let X and Y have correlation coefficient ρ_1 and let Y and Z have correlation coefficient ρ_2 .
1. Find $E[XZ]$. (Hint: $E[X|Y = y] = \rho_1 \frac{\sigma_X^2}{\sigma_Y^2} y$.)
 2. Find $I(X; Z)$.

4. Let $\Pr(X = 1) = p$, $\Pr(X = 0) = 1 - p$, and let $Y = X + Z$, where Z is uniform over the interval $[0, a]$, $a > 1$, and Z is independent of X .
1. Calculate $I(X; Y) = H(X) - H(X|Y)$.
 2. Now calculate $I(X; Y)$ the other way by $I(X; Y) = H(Y) - H(Y|X)$.
 3. Calculate the capacity of this channel by maximizing over p .

5. *Joint AEP for the binary Z channel* The binary Z channel is a DMC with binary inputs $\mathcal{X} = \{0, 1\}$ and binary outputs $\mathcal{Y} = \{0, 1\}$ and conditional probability distribution $p_{\mathbf{Y}|\mathbf{X}}(y|x)$ given by matrix:

$$\begin{bmatrix} 1 & 0 \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix}. \quad (32)$$

The channel diagram looks like a “Z”:



The channel is memoryless:

$$p_{\mathbf{Y}|\mathbf{X}}(\mathbf{y}|\mathbf{x}) = \prod_{i=1}^n p_{Y_i|X_i}(y_i|x_i)$$

Assume the input distribution is uniform $p_{\mathbf{X}}(x) = [\frac{1}{2}, \frac{1}{2}]$.

- (a) Let m be the number of one's in \mathbf{y} . Find $p_{\mathbf{Y}}(\mathbf{y})$.

- (b) For what values of m and n does the following hold (equivalent to $\mathcal{T}_{\epsilon}^{(n)}$ with $\epsilon = 0$):

$$-\frac{1}{n} \log p_{\mathbf{Y}}(\mathbf{y}) = H(\mathbf{Y}).$$

- (c) Using the memoryless property, compute the following quantities:

- $p_{Y|X}(000|001) =$
- $p_{Y|X}(001|001) =$
- $p_{Y|X}(010|001) =$
- $p_{Y|X}(011|001) =$

Continued on next page

Let $\mathbf{x} = (x_1, \dots, x_n)$ be a sequence with k ones, with $k \leq n$. Consider any sequence \mathbf{x} and \mathbf{y} . If $x_i = 0$ and $y_i = 1$, then $p_{\mathbf{x}\mathbf{y}}(\mathbf{x}, \mathbf{y}) = 0$ (for any $i = 1, 2, \dots, n$). Let \mathcal{V} be the valid sequences:

$$\mathcal{V} = \{(\mathbf{x}, \mathbf{y}) : x_i \neq 0 \text{ or } y_i \neq 1, \text{ all } i = 1, 2, \dots, n\}$$

Note that $p(\mathbf{x}, \mathbf{y} | \mathcal{V}^c) = 0$, where \mathcal{V}^c is the complement of \mathcal{V} . For $n = 2$, here is a list of valid sequences \mathcal{V} and invalid sequences \mathcal{V}^c :

Valid \mathcal{V} , (x_1x_2, y_1y_2)	Not Valid \mathcal{V}^c , (x_1, x_2, y_1, y_2)
(00,00)	
(00,01)	(01,00)
(01,01)	(10,00)
(00,10)	(11,00)
(10,10)	(10,01)
(00,11)	(11,01)
(01,11)	(01,10)
(10,11)	(11,10)
(11,11)	

- (d) Find $\Pr(\mathbf{y} | \mathbf{x} \mathcal{V})$ (That is, find $\Pr(\mathbf{y} | \mathbf{x})$, given a valid input/output sequence). Express using k , the number of ones in \mathbf{x} .

- (e) Using your answer to part (d), find $p_{\mathbf{x}\mathbf{y}}(\mathbf{x}, \mathbf{y})$. Then, find $-\frac{1}{n} \log p_{\mathbf{x}\mathbf{y}}(\mathbf{x}, \mathbf{y})$. Express using n and k .

- (f) Find $H(\mathbf{X}, \mathbf{Y})$.

(g) Let \mathcal{T}'_ϵ be:

$$\mathcal{T}'_\epsilon = \left\{ (x^n, y^n) : \left| -\frac{1}{n} \log p(x^n, y^n) - H(X, Y) \right| < \epsilon \right\}.$$

For $n = 20$ and $\epsilon = 0.06$, describe the sequences that are in the set \mathcal{T}'_ϵ .

7. *General information theory* For each question (a)-(d), write an expression

- (a) For the Markov chain $X \rightarrow Y \rightarrow Z$, write the data processing inequality.
- (b) For variables X and Y , write an inequality expressing "independence bound on entropy"
- (c) For variables X and Y , write the entropy chain rule:
- (d) For variables X and Y write an inequality expressing "conditioning reduces entropy":
- (e) **True or False** ? A Markov chain $X_1, X_2, X_3 \dots$, has entropy rate $H(\mathcal{X})$:

$$H(\mathcal{X}) \leq H(X_2).$$

- (f) **True or False?** For a continuous random variable X , the uniform distribution $f(x) = \frac{1}{a}$ for $0 \leq x \leq a$ maximizes the differential entropy $H(X) = \int f(x) \ln f(x) dx$.
- (g) Consider the random variable X :

$$\Pr(X = i) = \begin{cases} \frac{1}{4} & \text{if } i = -1 \\ \frac{1}{2} & \text{if } i = 0 \\ \frac{1}{4} & \text{if } i = 1 \end{cases}$$

Find $H(X)$. If $g(x) = x^2$, find $H(g(X))$ and $H(g(X)|X)$.

9. *Proof of Fano's Inequality* Write a justification (for example, “data processing inequality”) for each step in the proof of Fano's inequality

Fano's Inequality For any estimator \hat{X} such that $X \rightarrow Y \rightarrow \hat{X}$, with event $E = \{X \neq \hat{X}\}$ and with $P_e = \Pr(E)$, we have:

$$h(P_e) + P_e \log |\mathcal{X}| \geq H(X|\hat{X}). \quad (34)$$

Proof:

Justification

$$\begin{aligned}
 H(E, X|\hat{X}) &= H(X|\hat{X}) + H(E|X, \hat{X}) & (a) & \text{_____} \\
 &= H(X|\hat{X}) & (b) & \text{_____} \\
 H(E, X|\hat{X}) &= H(E|\hat{X}) + H(X|E, \hat{X}) & (c) & \text{_____} \\
 H(X|\hat{X}) &= H(E|\hat{X}) + H(X|E, \hat{X}) & & \text{equality of (b) and (c)} \\
 &\leq H(E) + H(X|E, \hat{X}) & (d) & \text{_____} \\
 &= h(P_e) + H(X|E, \hat{X}) & & H(E) = h(P_e) \\
 &= h(P_e) + H(X|\hat{X}, E=1)P_e \\
 &\quad + H(X|\hat{X}, E=0)(1-P_e) & (e) & \text{_____} \\
 &= h(P_e) + H(X|\hat{X}, E=1)P_e & (f) & \text{_____} \\
 &\leq h(P_e) + P_e \log |\mathcal{X}| & (g) & \text{_____}
 \end{aligned}$$