## Homework 9 (2023) — SOLUTIONS

JAIST — School of Information Science — I232 Information Theory

## 1. Packing spheres

- (a) In two dimensions, consider packing circles of radius N=1 inside a large circle of radius S=10. Find an upper bound on M, the number of circles, by dividing the area of the large circle, by the area of small circle.
- (b) Now consider n-dimensional spheres. The volume of a sphere with radius r in n dimensions is Vol(n,r):

$$Vol(n,r) = \frac{\pi^{n/2}}{\Gamma(\frac{n}{2}+1)}r^n,$$

where  $\Gamma$  is the gamma function. Let the radius of the large sphere be  $\sqrt{P+N}$ , and let the radius of the small sphere be  $\sqrt{N}$ . As before, find an upper bound on the number of spheres M

(c) Now let  $R = \frac{1}{n} \log M$ . Using your answer to part (b), what is an upper bound on R?

Solution:



(a)

$$\begin{split} M &\leq \frac{\text{Area of large circle}}{\text{Area of small circle}} \\ &= \frac{\pi S^2}{\pi N^2} = \frac{100}{1} = 100 \end{split}$$

(b) With  $C_n = \frac{\pi^{n/2}}{\Gamma(\frac{n}{2}+1)}$ :

$$\operatorname{Vol}_{\operatorname{big}} = C_n(\sqrt{n(P+N)})^n$$
  
 $\operatorname{Vol}_{\operatorname{small}} = C_n(\sqrt{n(N)})^n$ 

Maximum number of messages M is the maximum number of non-intersecting spheres:

$$M \le \frac{\text{Vol}_{\text{big}}}{\text{Vol}_{\text{small}}} = \frac{C_n(\sqrt{n(P+N)})^n}{C_n(\sqrt{n(N)})^n}$$

$$M \le \left(1 + \frac{P}{N}\right)^{n/2}$$

(c)

$$\begin{split} R &= \frac{1}{n} \log M &\leq \frac{1}{n} \log \left( 1 + \frac{P}{N} \right)^{n/2} \\ R &\leq \frac{1}{2} \log \left( 1 + \frac{P}{N} \right) \end{split}$$

2. Triangular probability distribution Consider two uniformly distributed random variables X and Y:

$$p_{\mathsf{X}}(x) = \begin{cases} \frac{1}{2} & -1 \le x \le 1\\ 0 & \text{otherwise} \end{cases} \quad \text{and} \quad p_{\mathsf{Y}}(y) = \begin{cases} \frac{1}{2} & -1 \le y \le 1\\ 0 & \text{otherwise} \end{cases}$$

(a) Find the probability distribution of X + Y, using convolution.

**Solution:** Let Z = X + Y.

$$p_{\mathsf{Z}}(z) = \int_{-\infty}^{\infty} p_{\mathsf{X}}(z - w) p_{\mathsf{Y}}(w) dw \qquad -1 \le z - w \le 1, -1 \le w \le 1$$

$$= \int_{\max(z - 1, -1)}^{\min(z + 1, 1)} \frac{1}{2} \frac{1}{2} dw \qquad z - 1 \le w \le z + 1$$

$$= \frac{1}{4} \left( \min(z + 1, 1) - \max(z - 1, -1) \right)$$

$$= \begin{cases} \frac{1}{2} + \frac{1}{4}z & -2 \le z \le 0 \\ \frac{1}{2} - \frac{1}{4}z & 0 < z \le 2 \\ 0 & \text{otherwise} \end{cases}$$

(b) Let Z be distributed as a triangular distribution:

$$p_{\mathsf{Z}}(z) = \begin{cases} c + c^2 z & -\frac{1}{c} \le z \le 0\\ c - c^2 z & 0 < z \le \frac{1}{c}\\ 0 & \text{otherwise} \end{cases}$$

for c > 0. Find differential entropy of Z in nats, using natural log ln.

Solution:

$$\begin{split} H(\mathsf{Z}) &= -\int p_{\mathsf{Z}}(z) \ln p_{\mathsf{Z}}(z) dz \\ &= -\int_{-1/c}^{0} (c+c^2z) \ln(c+c^2z) dz - \int_{0}^{1/c} (c-c^2z) \ln(c-c^2z) dz \end{split}$$

The two integrals are equal by symmetry. The first integral is:

$$-\int_{-1/c}^{0} (c+c^2z) \ln(c+c^2z) dz = \frac{1}{4} + \frac{1}{2} \ln \frac{1}{c}$$

One technique to evaluate the integral is to use Wolfram Alpha or Mathematica:

Integrate[c(1+cz) Log[ c (1+c z)],{z,-1/c,0}, Assumptions 
$$\rightarrow$$
 c > 0]

So the differential entropy is twice that:

$$H(\mathsf{Z}) = \frac{1}{2} + \ln \frac{1}{c}$$

(c) What is the differential entropy of X + Y, in nats?

**Solution:** Since X + Y has a triangular distribution with c = 1/2, the differential entropy is  $\frac{1}{2} + \ln 2$ .

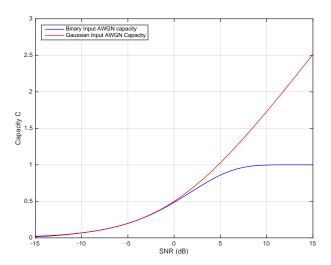
3. Capacity of the binary-input AWGN channel The binary-input AWGN channel achieves capacity with  $p_X^*(x) = [\frac{1}{2}, \frac{1}{2}]$ . The signal-to-noise ratio SNR in decibels (dB) is:

SNR in 
$$dB = 10 \log_{10} \frac{1}{\sigma^2}$$
.

(a) Make a plot of capacity C versus SNR dB for the binary-input AWGN channel. (b) On the same graph, plot the capacity of the AWGN channel with Gaussian input distribution. (c) What happens to achievable rates as SNR  $\rightarrow$  0? What happens to achievable rates as SNR  $\rightarrow$   $\infty$ ? What can you conclude from this?

**Solution:** (a), (b) The figure below shows the plot of capacity of the binary-input AWGN channel and the Gaussian-input AWGN channel, and below that is the source code used to generate them.

(c) Note that SNR  $\to 0$  means SNR dB  $\to -\infty$ . As SNR  $\to 0$ , the capacity goes to 0. As SNR  $\to \infty$ , the capacity of the binary-input goes to 1, since the input is binary, but the capacity of the Gaussian-input AWGN channel goes to infinity. From these two, we can conclude that as SNR  $\to 0$ , it is sufficient to use binary inputs, i.e. there is no benefit to using Gaussian-input signaling. But as SNR  $\to \infty$ , Gaussian inputs are needed to achieve higher capacities.



```
clear all
   close all
   SNRdb = linspace(-15,15);
   SNR
         = 10.^{(SNRdb/10)};
         = 1 ./ SNR;
   IXY = zeros(size(var));
9
   for ii = 1:length(var)
        IXY(ii) = biawqnCapacity(var(ii));
11
   end
12
   plot(SNRdb,IXY,'b-');
   hold on
   C = 0.5 * log2(1 + SNR);
16
   plot(SNRdb,C,'r-');
   xlabel('SNR (dB)');
18
  ylabel('Capacity C');
19
  legend('Binary Input AWGN capacity','Gaussian Input AWGN Capacity','Location','
       Northwest');
```

```
grid on

function IXY = biawgnCapacity(var)

fun = @(y) (1/sqrt(2*pi*var)) .* exp( - (y-1).^2 / (2*var)) .* log2( 2 ./ (1 + exp( - 2*y / var)) );

IXY = integral(fun,-10,10);

>> var = 0.9578;
>> fun = @(y) (1/sqrt(2*pi*var)) .* exp( - (y-1).^2 / (2*var)) .* log2( 2 ./ (1 + exp( - 2*y / var)) );
>> IXY = integral(fun,-10,10)

IXY =

0.5000
```