

Minor Research Report 4

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I. Overview

- Objectives:
 - Translate IMP code into type expressions and generate type constraint equations
 - Perform type inference by using unification algorithm
- Progress: Done
- Source code:
https://github.com/thanhtcptit/imp-typescript/blob/main/imp_interpreter/typing.ts

II. Implementation

1. Type terms

Term	Form
Type Variable	α
Type Arrow	$\alpha \rightarrow \alpha$
Type Application	$\alpha \alpha$
Type Constructor	Int, Bool, Unit

2. Type environment (Γ)

Variable	Type
0, 1, 2, ...	Int
true, false	Bool
+, -	$\text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$
>, >=, <, <=, ==, !=	$\text{Int} \rightarrow \text{Int} \rightarrow \text{Bool}$
&&,	$\text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool}$
!	$\text{Bool} \rightarrow \text{Bool}$
:=	$\alpha \rightarrow \alpha \rightarrow \text{Unit}$
if	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit}$
if-else	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit} \rightarrow \text{Unit}$
while	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit}$
;	$\text{Unit} \rightarrow \alpha \rightarrow \alpha$

Variable	Type
0, 1, 2, ...	Int
true, false	Bool
+, -	$\text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$
>, >=, <, <=, ==, !=	$\text{Int} \rightarrow \text{Int} \rightarrow \text{Bool}$
&&,	$\text{Bool} \rightarrow \text{Bool} \rightarrow \text{Bool}$
!	$\text{Bool} \rightarrow \text{Bool}$
:=	$\alpha \rightarrow \alpha \rightarrow \text{Unit}$
if	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit}$
if-else	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit} \rightarrow \text{Unit}$
while	$\text{Bool} \rightarrow \text{Unit} \rightarrow \text{Unit}$
return	$\alpha \rightarrow \alpha$

3. Type expressions

Expression (e)	Form	Example
Variable	x	x
Integer	0, 1, 2, ...	1
Boolean	true, false	true
λ -abstraction	$\lambda x. e$	$\lambda x. x$
Application	$e1\ e2$	$\lambda x. \lambda y. (+)\ x\ y$
Let	let t = e1 in e2	let t = $\lambda x. (+)\ x\ 1$ in t 1

4. Type inference rules

Expression	Deduction form
Variable	$\frac{\Gamma(x) = \alpha}{\Gamma \vdash x : \alpha}$
Integer	$\frac{}{\Gamma \vdash 1 : \text{Int}}$
Boolean	$\frac{}{\Gamma \vdash \text{true} : \text{Bool}}$

λ -abstraction	$\frac{\Gamma \cup \{x : \alpha\} \vdash e : \beta}{\Gamma \vdash \lambda x. e : \alpha \rightarrow \beta}$
Application	$\frac{\Gamma \vdash e_1 : \alpha \rightarrow \beta \quad \Gamma \vdash e_2 : \alpha}{\Gamma \vdash e_1 e_2 : \beta}$
Let	$\frac{\Gamma \vdash e_1 : \alpha \quad \Gamma \cup \{x : \alpha\} \vdash e_2 : \beta}{\Gamma \vdash \text{let } x = e_1 \text{ in } e_2 : \beta}$

5. Type constraints

Expression	Constraint
Variable	$\frac{\Gamma(x) = \perp}{\Gamma \vdash x : a} [a \approx \alpha]$
	$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : a} [a \approx \tau']$
λ -abstraction	$\frac{\Gamma \cup \{x : b\} \vdash e : c}{\Gamma \vdash \lambda x. e : a} [a \approx b \rightarrow c]$
Application	$\frac{\Gamma \vdash e_1 : b \quad \Gamma \vdash e_2 : c}{\Gamma \vdash e_1 e_2 : a} [b \approx c \rightarrow a]$
Let	$\frac{\Gamma \vdash e_1 : b \quad \Gamma \cup \{x : b\} \vdash e_2 : c}{\Gamma \vdash \text{let } x = e_1 \text{ in } e_2 : a} [a \approx c]$