

MINI PROJECT – FUNDAMENTALS OF OPTIMIZATION

SEMESTER EXAM SCHEDULE

Group 18

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1. INTRODUCTION AND DESCRIPTION

Scheduling semester exam is one of the most popular task at school. To help teachers do this work more conveniently, we give some algorithms that will be explained in the next section to optimize this problem.

PROBLEM DESCRIPTION

- N : Number of courses to be scheduled
- $d(i)$: Number of students participate in course $n(i)$
- M : Number of rooms
- $c(j)$: Number of seats of room $m(j)$
- K : Number of pairs of courses cannot be grouped together (conflicting courses)
- Pairs of conflicting courses

1	20
2	45 39 54 58 23 51 27 47 52 28 48 32 48 53 37 50 27 30 54 25
3	6
4	29 44 24 48 60 26
5	10
6	1 15
7	7 10
8	1 6
9	13 10
10	19 7
11	15 9
12	16 17
13	20 7
14	16 14
15	15 12

We have four periods in a day and the problem is minimizing the number of days scheduled for examination.

2. MODELLING AND ALGORITHMS IMPLEMENTATION

**MIXED INTEGER
PROGRAMMING**

BACKTRACKING

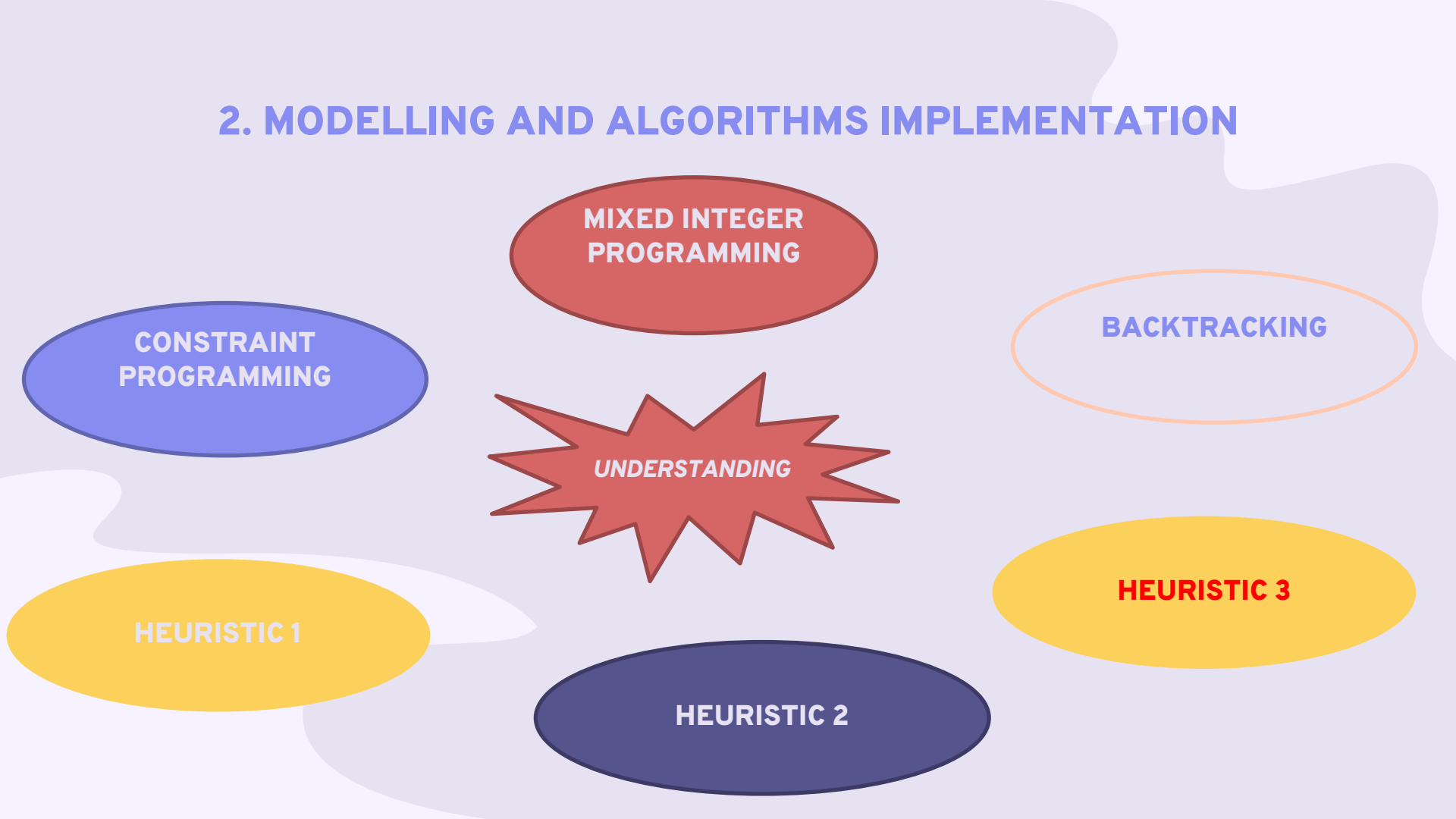
**CONSTRAINT
PROGRAMMING**

UNDERSTANDING

HEURISTIC 1

HEURISTIC 3

HEURISTIC 2



- **MIXED INTEGER PROGRAMMING**

1. **Define variables**

- $X[i][j][k]$: variable that put course $n(i)$ to room $m(j)$ at period k
 $0 \leq i, k \leq N, 0 \leq j \leq M, X[i][j][k] \in \{0; 1\}$
- y : number of periods
- C : list of conflicting courses

2. Constraints

- **Constraint 1: Pairs of conflicting courses cannot be put in the same period**

$$0 \leq X[i1][j][k] + X[i2][j][k] \leq 1, \forall (i1, i2) \in C; j = 0, \dots, M-1; k = 0, \dots, N-1$$

- **Constraint 2: An course room be putted at most one course in a period**

$$0 \leq \sum_{i=0}^{N-1} X[i][j][k] \leq 1, \forall j = 0, \dots, M-1; k = 0, \dots, N-1$$

- **Constraint 3: The number of periods**

$$-\infty \leq k \times X[i][j][k] - y \leq 0, \forall j = 0, \dots, M-1; k = 0, \dots, N-1$$

- **Constraint 4: A course be conducted at most one time in an room**

$$0 \leq \sum_{j=0}^{M-1} \sum_{k=0}^{N-1} X[i][j][k] \leq 1, \forall i = 0, \dots, N-1$$

- **Constraint 5: A course n(i) must be put into a room m(j) with capacity c(j)**

$$0 \leq \sum_{k=0}^{N-1} X[i][j][k] \times d[i] \leq c[j], \forall i = 0, \dots, N-1; j = 0, \dots, M-1$$

3. Objective

$\Rightarrow \textit{Minimize}(y)$

```
mip_solver = pywraplp.Solver.CreateSolver('SCIP')
INF = mip_solver.infinity()

# Define variables
x = [[[mip_solver.IntVar(0, 1, f'x[{i}][{j}][{k}])] for i in range(N)] for j in range(M)] for k in range(N)]
y = mip_solver.IntVar(0, N - 1, 'y')

# Define constraints

# Constraint 1: Pairs of conflicting courses may not be put in the same time slot
for i in range(K):
    u, v = p[i][0], p[i][1]
    for k in range(N):
        constraint = mip_solver.Constraint(0, 1)
        for j1 in range(M):
            for j2 in range(M):
                if j1 != j2:
                    constraint.SetCoefficient(x[u][j1][k], 1)
                    constraint.SetCoefficient(x[v][j2][k], 1)

# Constraint 2: An course room may be assigned at most one course in a time slot
for j in range(M):
    for k in range(N):
        constraint = mip_solver.Constraint(0, 1)
        for i in range(N):
            constraint.SetCoefficient(x[i][j][k], 1)
```

```
# Constraint 3: The number of time slots ( $k.x[i,j,k] - y \leq 0$ )
for i in range(N):
    for j in range(M):
        for k in range(N):
            constraint = mip_solver.Constraint(-INF, 0)
            constraint.SetCoefficient(y, -1)
            constraint.SetCoefficient(x[i][j][k], k)

# Constraint 4: A course may be conducted at most one time in an course room
for i in range(N):
    constraint = mip_solver.Constraint(1, 1)
    for j in range(M):
        for k in range(N):
            constraint.SetCoefficient(x[i][j][k], 1)

# Constraint 5: A course  $n_i$  must be put into a room  $m_j$  with capacity  $c(j)$ 
for i in range(N):
    for j in range(M):
        constraint = mip_solver.Constraint(0, c[j])
        for k in range(N):
            constraint.SetCoefficient(x[i][j][k], d[i])
```


• **CONSTRAINT PROGRAMMING**

1. Define variables

- $X[i]$: period of course $n(i)$
 $i \in \{1, 2, \dots, n\}, X[i] \in \{1, 2, \dots, n\}$
- $Y[i][j]$: course $n(i)$ be putted in room $m(j)$
 $i \in \{1, 2, \dots, N\}, j \in \{1, 2, \dots, M\}, Y[i][j] \in \{0; 1\}$
- m : number of periods that need for scheduling exam
- p : list of pairs conflicting courses

2. Constraint

- **Constraint 1: Pairs of conflicting courses may not be put in the same period**

$$X[i] \neq X[j], \forall (i, j) \in p$$

- **Constraint 2: An course room is assigned at most one course in a period**

$$\sum_{j=1}^M Y[i][j] = 1, \forall i = 1, \dots, N$$

- **Constraint 3: Courses with same period cannot use the same room**

$$X[i] = X[j] \Rightarrow Y[i][k] + Y[j][k] \leq 1, \forall i, j \in \{1, \dots, N\}; k \in \{1, \dots, M\}$$

- **Constraint 4: The attendance of course n(i) must be smaller than capacity of room**

$$d[i] \leq \sum_{j=1}^M Y[i][j] \times c[j], \forall i = 1, \dots, N$$

3. Objective

$\Rightarrow \text{Minimize}(m)$ with $m = \max(X)$

```

# Initiation
model = cp_model.CpModel()

# Variable x[i]: period of course ni
x = [model.NewIntVar(1, N, f'x[{i}]') for i in range(N)]

# Variable y[i][j]: whether course ni takes room j or not
y = [[model.NewIntVar(0, 1, f'y[{i}][{j}]') for j in range(M)] for i in range(N)]

# Define constraints
# Constraint 1: Pairs of conflicting courses may not be put in the same period
for pair in p:
    model.Add(x[pair[0]] != x[pair[1]])

# Constraint 2: An course room is assigned at most one course in a period
for i in range(N):
    model.Add(sum(y[i]) == 1)

# Constraint 3: Courses with same period cannot use the same room
for j in range(M):
    for i1 in range(N - 1):
        for i2 in range(i1 + 1, N):
            b = model.NewBoolVar(f'b[{j}][{i1}][{i2}]')
            model.Add(y[i1][j] + y[i2][j] <= 1).OnlyEnforceIf(b)
            model.Add(x[i1] == x[i2]).OnlyEnforceIf(b)
            model.Add(x[i1] != x[i2]).OnlyEnforceIf(b.Not())

# Constraint 4: The attendance of course n_i must be smaller than capacity of room
for i in range(N):
    model.Add(sum([y[i][j] * c[j] for j in range(M)]) >= d[i])

```

• ***BACKTRACKING***

end := a very very large number

define function: dfs(u,slot):

if u == N:

end = min(slot, end)

return

if slot > end:

return

for each room:

if room is free:

for each course:

if course can be putted and attendance <= capacity:

put course to that room

put course to that slot

dfs(u+1,slot)

free that room

free that slot

dfs(u,slot+1)

return

```

end = 100000000
conflict = [[] for _ in range(N)]

for k in p:
    u, v = k[0], k[1]
    conflict[u].append(v)
    conflict[v].append(u)

# assign period
period = [-1] * N

# room
room = []
for _ in range(N):
    room.append([-1] * M)

def isPlaceable(u, slot):
    if period[u] >= 0:
        return False
    for v in conflict[u]:
        if period[v] == slot:
            return False
    return True

```

```

def dfs(u, slot):
    global end
    if u == N:
        end = min(end, slot)
        return
    if slot > end:
        return
    for j in range(M):
        if room[slot][j] == -1:
            for i in range(N):
                if isPlaceable(i, slot) and d[i] <= c[j]:
                    period[i], room[slot][j] = slot, i
                    dfs(u + 1, slot)
                    period[i], room[slot][j] = -1, -1
            dfs(u, slot + 1)
    return

# Solve
start_time = time.process_time()
dfs(0, 0)
end_time = time.process_time()

# Solution
if end != 100000000:
    print(f'Objective value: {end + 1}')
else:
    print('No solution.')
print('-----')
print(f'Used time: {1000*(end_time - start_time)} milliseconds')

```

• *HEURISTIC 1*

sort list of (capacity,room) in ascending order of capacity

for each **Course**:

for each **Period**:

if **Course** test cannot be putted in any existing **Period**:

add new **Period**

if other course in **Period** conflicts with **Course**:

consider the next **Period**

else:

for each sorted **Room**:

if attendants \leq capacity and at that **Room** and **Period** have no test:

put **Course** to that **Room** and **Period**

consider the next **Course**

```
# List of (capacity, room) are sorted by capacity in ascending order
sorted_c = sorted([(c[i], i) for i in range(M)])
```

```
# Conflicts
```

```
conflicts = {} # conflicts[i] = list of courses that cannot be administered in the same period as course i+1
for pair in p:
    conflicts.setdefault(pair[0], []).append(pair[1])
    conflicts.setdefault(pair[1], []).append(pair[0])
```

```
def greedy_2():
    result = [[-1] * M] # initiate with first period
                        # Result[i, k] = course exam administered in period i+1 and room k+1
    for exam in range(N): #sequentially assign a period and a room to each course
        nextCourse = False
        for period in range(len(result) + 1): #consider existing periods first
            if period == len(result):
                #if this exam cannot be held in any existing period, create a new period
                result.append([-1] * M) # new period with M rooms
            not_ThisPeriod = False
            if exam in conflicts:
                for otherCourse in result[period]:
                    if otherCourse in conflicts[exam]:
                        not_ThisPeriod = True
                        break
            if not_ThisPeriod == True:
                continue
            for room in range(M): #consider smaller rooms first to save bigger ones for other courses
                capacity = sorted_c[room][0]
                roomIndex = sorted_c[room][1]
                if result[period][roomIndex] == -1 and capacity >= d[exam]:
                    result[period][roomIndex] = exam
                    nextCourse = True
                    break
            if nextCourse == True:
                break
    return len(result), result
```



• HEURISTIC 2

```
list_of_exam = sorted([(attendant, i)])  
while list_of_exam not empty:  
    allocate a new period for remaining exams  
    for each Room:  
        for exam in list_of_exam:  
            if attendant <= capacity:  
                if have no exam scheduled conflicts in this period:  
                    put exam in this Room and this period  
                    consider the next room
```



```

conflicts = {} #conflicts[i] = list of exams that cannot be administered in the same period as exam i+1
for pair in p:
    conflicts.setdefault(pair[0], []).append(pair[1])
    conflicts.setdefault(pair[1], []).append(pair[0])

print('\nPeriod', 'Room', 'Exam', sep='\t')

sortedExams = sorted([(d[i], i) for i in range(N)], reverse=True) #sort exams in ascending order of capacity

schedule = [] #schedule[i, k] = exam administered in period i+1 and hall k+1
period = 0
startTime = time.process_time()
while sortedExams: #sequentially fill each period with as many exams as possible until all exams have been scheduled
    schedule.append([None] * M)
    for room in range(M):
        for exam in sortedExams: #consider more popular exams first
            if exam[0] <= c[room]: #if a hall has adequate capacity
                #check if any exam already scheduled in this period has common candidates with the one being considered
                noConflict = True
                if exam[1] in conflicts:
                    for scheduledExam in schedule[period]:
                        if scheduledExam in conflicts[exam[1]]:
                            noConflict = False
                            break
                if noConflict: #schedule exam in period and hall and remove from list of exams to schedule
                    schedule[period][room] = exam[1]
                    sortedExams.remove(exam)
                    break
    period += 1

```

• *HEURISTIC 3*

sort list of rooms in ascending order of capacity

for each **Course**:

for each **Period**:

for each **Room**:

if capacity \geq attendants and no exam scheduled at **Period** and **Room** yet:

if no other course in period conflicts with **Course**:

put **Course** to **Period** and **Room**

consider the next **Course**

if **Course** exam cannot be held in any **Period**:

add new **Period**

```

conflicts = {} #conflicts[i] = list of exams that cannot be administered in the same period as exam i+1
for pairss in p:
    conflicts.setdefault(pairss[0], []).append(pairss[1])
    conflicts.setdefault(pairss[1], []).append(pairss[0])

sortedRooms = sorted([(c[i], i) for i in range(M)]) #sort rooms in ascending order of capacity
result = [[None] * M] #result[i, k] = exam administered in period i+1 and room k+1
print('\nExam', 'Period', 'Room', sep='\t')
for exam in range(N): #sequentially assign a period and a room to each exam
    stop = False
    for period in range(len(result) + 1): #consider existing periods first
        for room in range(M): #consider smaller rooms first to save bigger ones for other exams
            capacity = sortedRooms[room][0]
            roomIndex = sortedRooms[room][1]
            if capacity >= d[exam] and result[period][roomIndex] == None:
                noConflict = True
                if exam in conflicts:
                    for otherExam in result[period]:
                        if otherExam in conflicts[exam]:
                            noConflict = False
                            break
                if noConflict:
                    result[period][roomIndex] = exam
                    print(exam + 1, period + 1, room + 1, sep='\t') #print schedule by exam
                    stop = True
                    break
    if stop:
        break
    if period == len(result) - 1:
        #if this exam cannot be held in any existing period, set up a new period
        result.append([None] * M)

```

3. ANALYSIS AND CONCLUSION



OUR RESULTS

Test 1: With some small datasets ($N/M < 4$):

- MIP, CP work quite well. Meanwhile, BT takes much time to find the result.

N	MIP value	MIP time	CP value	CP time	BT value	BT time
6	3	32,01	3	34,1	3	109,38
9	3	54,02	3	88,13	3	4273000
12	5	57,44	5	170,44	-	INF
16	4	1039,88	4	30147	-	INF
20	7	203,01	7	30312	-	INF

***time unit: millisecond**

value unit: period

Test 1: With some small datasets ($N/M < 4$):

- MIP, CP work quite well. Meanwhile, BT takes much time to find the result.

CP, MIP



Backtracking



OUR RESULTS

Test 1: With some small datasets ($N/M < 4$):

- Heuristic 1, Heuristic 2 and Heuristic 3 outperform MIP, CP, BT. They likely give the same results with an small amount of running time

N	Heu 1 value	Heu 1 time	Heu 2 value	Heu 2 time	Heu 3 value	Heu 3 time
6	3	~0	3	~0	3	~0
9	3	~0	3	~0	3	~0
12	5	~0	5	~0	4	~0
16	4	0,999	4	0,997	4	0,81
20	7	1,004	7	1,002	7	0,9

***time unit: millisecond**

value unit: period

OUR RESULTS

Test 2: With Larger datasets ($N/M > 4$):

- MIP, CP become worse and worse.

N	K	MIP value	MIP time	CP value	CP time
50	20	50	39187	6	31291
100	45	53	30433	-	Timed out
150	60	-	Timed out	-	Timed out
200	100	-	Timed out	-	Timed out
200	200	-	Timed out	-	Timed out

***time unit: millisecond**

value unit: period

time limit = 30000 milisecond

OUR RESULTS

Test 2: With Larger datasets ($N/M > 4$): Heuristic algorithms still run fast with OPTIMAL solutions with nearly the same running time.

N	K	Heu 1 Value	Heu 1 Time	Heu 2 Value	Heu 2 Time	Heu 3 Value	Heu 3 Time
50	20	6	0,997	6	0,1	7	0,996
100	45	8	1,001	7	0,969	10	1,01
150	60	8	4,117	9	0,01	8	2,00
200	100	9	6,228	9	1,757	9	4,98
200	200	10	10,09	10	4,038	10	2,99

***time unit: millisecond
value unit: period**

OUR RESULTS

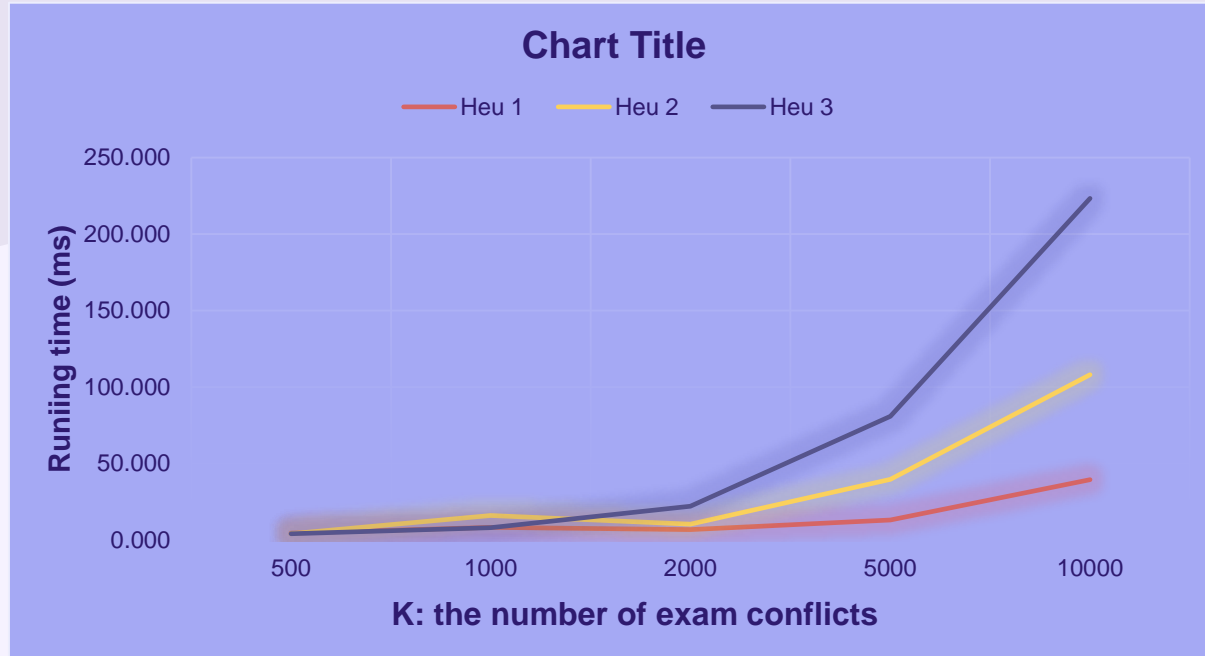
Test 3: Keep N, M at 200, 25; Increase the value of K (number of exams' conflict): There are some small differences between heuristic algorithms

N	K	Heu 1 Value	Heu 1 Time	Heu 2 Value	Heu 2 Time	Heu 3 Value	Heu 3 Time
200	500	9	4,516	10	4,154	9	4.02
200	1000	9	8,238	10	15,999	10	7,98
200	2000	11	6,792	11	10,425	11	21,94
200	5000	19	12,997	20	39,699	20	80,8
200	10000	36	39,476	37	108,015	39	223,4

*time unit: millisecond
value unit: period

OUR RESULTS

Test 3: Keep N, M at 200, 25; Increase the value of K (number of exams' conflict): There are considerable gaps between running time of algorithms



AWESOME WORDS

