

Chiang Mai University

PramorKangDamCPP

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Contest (1)

template.cpp

```
86 lines
```

```
#include <bits/stdc++.h>
using namespace std;
void print(int x) {cerr << x;}</pre>
void __print(long x) {cerr << x;}</pre>
void __print(long long x) {cerr << x;}</pre>
void __print(unsigned x) {cerr << x;}</pre>
void __print(unsigned long x) {cerr << x;}</pre>
void __print(unsigned long long x) {cerr << x;}</pre>
void __print(float x) {cerr << x;}</pre>
void __print(double x) {cerr << x;}</pre>
void __print(long double x) {cerr << x;}</pre>
void __print(char x) {cerr << '\'' << x << '\'';}</pre>
void __print(const char *x) {cerr << '\"' << x << '\"';}</pre>
void __print(const string &x) {cerr << '\"' << x << '\"';}</pre>
void __print(bool x) {cerr << (x ? "true" : "false");}</pre>
template<typename T, typename V>
void __print(const pair<T, V> &x);
template<typename T>
void __print(const T &x) {int f = 0; cerr << '{'; for (auto &i:</pre>
     x) cerr << (f++ ? ", " : ""), __print(i); cerr << "}";}
template<typename T, typename V>
void __print(const pair<T, V> &x) {cerr << '{'; __print(x.first</pre>
    ); cerr << ", "; print(x.second); cerr << '}';}
void print() {cerr << "]\n";}</pre>
template <typename T, typename... V>
void _print(T t, V... v) {__print(t); if (sizeof...(v)) cerr <</pre>
      ", "; _print(v...);}
//#ifdef DEBUG
#define dbg(x...) cerr << "\e[91m"<< func <<":"<< LINE <<"
    [" << #x << "] = ["; _print(x); cerr << "\e[39m" << endl;
//#else
//\#define\ dbg(x...)
//#endif
typedef long long 11;
typedef long double ld;
typedef complex<ld> cd;
typedef pair<int, int> pi;
typedef pair<11,11> pl;
typedef pair<ld, ld> pd;
typedef vector<int> vi;
typedef vector<ld> vd;
typedef vector<11> v1;
typedef vector<pi> vpi;
typedef vector<pl> vpl;
typedef vector<cd> vcd;
template<class T> using pq = priority_queue<T>;
template < class T > using pqg = priority_queue < T, vector < T > ,
    greater<T>>;
#define rep(i, a) for(int i=0;i<a;++i)
#define FOR(i, a, b) for (int i=a; i<(b); i++)
#define FOR(i, a) for (int i=0; i<(a); i++)
#define FORd(i,a,b) for (int i = (b)-1; i \ge a; i--)
#define FORd(i,a) for (int i = (a)-1; i >= 0; i--)
#define trav(a,x) for (auto& a : x)
#define uid(a, b) uniform_int_distribution<int>(a, b) (rng)
#define sz(x) (int)(x).size()
#define mp make_pair
#define pb push_back
```

```
//#define f first
//#define s second
#define lb lower bound
#define ub upper bound
#define all(x) x.begin(), x.end()
#define ins insert
template < class T > bool ckmin(T& a, const T& b) { return b < a ?
     a = b, 1 : 0; }
template < class T > bool ckmax (T& a, const T& b) { return a < b ?
     a = b, 1 : 0; }
mt19937 rng(chrono::steady_clock::now().time_since_epoch().
    count());
const char nl = '\n';
const int N = 2e5+1;
const int INF = 1e9+7;
const long long LINF = 1e18+7;
void solve(){
int main(){
    ios::sync_with_stdio(false);cin.tie(nullptr);
    int t = 1;
    while (t--) solve();
.vimrc
                                                           21 lines
"General editor settings
set tabstop=4
set nocompatible
set shiftwidth=4
set expandtab
set autoindent
set smartindent
set ruler
set showcmd
set incsearch
set shellslash
set number
set relativenumber
set cino+=L0
"keybindings for { completion, "jk" for escape, ctrl-a to
    select all
inoremap {<CR> {<CR>}<Esc>0
inoremap {}
               {}
                <Esc>
imap jk
map <C-a> <esc>ggVG<CR>
set belloff=all
.bashrc
                                                            1 lines
export PATH=$PATH:~/scripts/
build.sh
q++ -static -DLOCAL -lm -s -x c++ -Wall -Wextra -O2 -std=c++17
    -o $1 $1.cpp
stress.sh
                                                           22 lines
#!/usr/bin/env bash
```

for ((testNum=0;testNum<\$4;testNum++))

```
./$3 > input
    ./$2 < input > outSlow
    ./$1 < input > outWrong
    H1='md5sum outWrong'
    H2= 'md5sum outSlow'
    if !(cmp -s "outWrong" "outSlow")
        echo "Error found!"
        echo "Input:"
        cat input
        echo "Wrong Output:"
        cat outWrong
        echo "Slow Output:"
        cat outSlow
        exit
    fi
done
echo Passed $4 tests
```

Tanya (2)

SegmentTree.h

Description: Segment tree with point update for range sum

Time: $\mathcal{O}(\log N)$ efd738, 29 lines

```
//TODO: use 0 base indexing
vector<long long>tree;
void update(int node,int n_l,int n_r,int q_i,long long value) {
    if (n_r<q_i || q_i<n_l) return;</pre>
    if(q_i==n_l && n_r==q_i){
        tree[node] = value;
        return:
    int mid = (n r+n 1)/2;
    update(2*node, n_l, mid, q_i, value);
    update(2*node+1, mid+1, n_r, q_i, value);
    tree[node] = tree[2*node] + tree[2*node+1];
long long f(int node,int n_l,int n_r,int q_l,int q_r) {
    if(n_r<q_l || q_r<n_l)return 0;</pre>
    if(q_1<=n_1 && n_r<=q_r) return tree[node];</pre>
    int mid = (n_1+n_r)/2;
    return f(2*node,n_1,mid,q_1,q_r) + f(2*node+1,mid+1,n_r,q_1
         , a r);
void build tree(vi &a,int n){
    tree.clear();
    int m=n;
    while (__builtin_popcount (m) !=1) ++m;
    tree.resize(2*m+1,0);
    for(int i=0;i<n;++i)tree[i+m]=a[i];</pre>
    for(int i=m-1;i>=1;--i)tree[i]=tree[2*i]+tree[2*i+1];
```

LazySegmentTree.h

Description: Segment tree with lazy propagation update for range sum **Time:** $\mathcal{O}(\log N)$.

```
if(n 1!=n r){
            lazy[2*node] +=lazy[node];
            lazy[2*node+1]+=lazy[node];
        lazy[node] = 0;
    if (n_r<q_l || q_r<n_l) return;</pre>
    if(q_l<=n_l && n_r<=q_r){
        tree[node] += (long long) (n_r-n_l+1) *value;
        // for range + update
        if (n_l!=n_r) {
            lazy[2*node]+=value;
            lazy[2*node+1]+=value;
        return;
    int mid = (n_r+n_1)/2;
    update(2*node,n_l,mid,q_l,q_r,value);
    update(2*node+1, mid+1, n_r, q_1, q_r, value);
    tree[node] = tree[2*node] + tree[2*node+1];
long long f(int node,int n_l,int n_r,int q_l,int q_r) {
    if(lazy[node]!=0){
        tree[node] += (long long) (n_r-n_l+1) *lazy[node];
        if (n_l!=n_r) {
            lazy[2*node] += lazy[node];
            lazy[2*node+1] += lazy[node];
        lazy[node] = 0;
    if (n_r<q_1 || q_r<n_1) return 0;</pre>
    if(q_1<=n_1 && n_r<=q_r)return tree[node];</pre>
    int mid = (n_l+n_r)/2;
    return f(2*node,n_1,mid,q_1,q_r) + f(2*node+1,mid+1,n_r,q_1
void build tree(vi &a,int n){
    tree.clear(); lazy.clear();
    while( builtin popcount(m)!=1)++m;
    tree.resize(2*m+1,0); lazy.resize(2*m+1,0);
    for(int i=0;i<n;++i)tree[i+m]=a[i];</pre>
    for(int i=m-1;i>=1;--i)tree[i]=tree[2*i]+tree[2*i+1];
OrderStatisticTree.h
Description: find nth largest element, count elements strictly less than x
Time: \mathcal{O}(\log N)
<ext/pb_ds/assoc_container.hpp>, <ext/pb_ds/tree_policy.hpp>
typedef __qnu_pbds::tree<int, __qnu_pbds::null_type, less<int>,
     __gnu_pbds::rb_tree_tag, __gnu_pbds::
     tree order statistics node update> ordered set;
```

${\bf MergeSortTree.h}$

 ${\bf Description:}\ \ {\rm do\ the\ same\ with\ orderstatistic\ tree\ but\ now\ over\ interval\ can}\ \ {\rm be\ used/modify\ for\ some\ possible\ interval/subbarray\ queries}$

```
Time: \mathcal{O}(\log N^3)
<ext/pb_ds/assoc_container.hpp>, <ext/pb_ds/tree_policy.hpp>,
<ext/pb_ds/assoc_container.hpp>, <ext/pb_ds/tree_policy.hpp>
                                                       2d0612, 94 lines
//merge sort tree with fenwick tree
typedef __gnu_pbds::tree<pair<int,int>, __gnu_pbds::null_type,
     less<pair<int,int>>, __gnu_pbds::rb_tree_tag, __gnu_pbds::
     tree_order_statistics_node_update> ordered_set;
ordered set st;
//st.order\_of\_key(\{x,-1\}) - find \# of elements in st stricty
     less than x
//st.clear() - clear container
vector<ordered_set> mtree;
ordered_set merge(ordered_set &a, ordered_set &b){
    ordered set result;
    for (auto&p:a) {
        result.ins(p);
    for (auto&p:b) {
        result.ins(p);
    return result;
void update(int node,int n_l,int n_r,int q_i,int id,int old_val
     ,int value) {
    if (n_r<q_i || q_i<n_l) return;</pre>
    if(q_i==n_l && n_r==q_i){
        auto it=mtree[node].find({old_val,id});
        mtree[node].erase(it);
        mtree[node].ins({value,id});
        return;
    int mid = (n_r+n_1)/2;
    update(2*node, n_1, mid, q_i, id, old_val, value);
    update(2*node+1, mid+1, n_r, q_i, id, old_val, value);
    auto it=mtree[node].find({old_val,id});
    if(it!=mtree[node].end()){
        mtree[node].erase(it);
        mtree[node].ins({value,id});
int f(int node,int n_l,int n_r,int q_l,int q_r, int value) {
    if (n r<q 1 || q r<n 1) return 0;</pre>
    if (q_1<=n_1 && n_r<=q_r) {</pre>
        return mtree[node].order_of_key({value,-1});
    int mid = (n_1+n_r)/2;
    return f(2*node,n_l,mid,q_l,q_r,value)+f(2*node+1,mid+1,n_r
         ,q l,q r,value);
void build mtree (vi &a) {
    int n=(int)a.size();
    int m=n; while( builtin popcount(m)!=1)++m;
    //for(int i=0;i<2*m+i)mtree[i].clear();
    mtree.resize(2*m);
    for (int i=0; i < n; ++i) mtree[i+m].ins({a[i],i});</pre>
    for (int i=m-1; i>=1; --i) mtree[i] = merge (mtree[2*i], mtree[2*i]
//merge sort tree with fenwick tree(BIT) (4 times less space)
typedef __gnu_pbds::tree<pair<int,int>, __gnu_pbds::null_type,
     less<pair<int,int>>, __gnu_pbds::rb_tree_tag, __gnu_pbds::
     tree order statistics node update> ordered set;
ordered set st:
```

```
//st.order\_of\_key(x) - find \# of elements in st strictly less
//TODO: use 1 base indexing
vector<ordered_set>bit;
void update(int i,int k,int old value, int new value) {
    while(i<(int)bit.size()){</pre>
        auto it=bit[i].find({old_value,k});
        assert(it!=bit[i].end());
        if(it!=bit[i].end()){
            bit[i].erase(it);
        bit[i].ins({new value,k});
        i+=i&-i;//add last set bit
int F(int i, int k){//culmulative sum to ith data
    int sum=0;
    while(i>0){
        sum+=bit[i].order_of_key({k,-1});
        i-=i&-i;
    return sum;
void build bit(vi &a){
    bit.resize((int)a.size());
    for (int i=1; i < (int) a.size(); ++i) bit[i].ins({a[i],i});</pre>
    for(int i=1; i<(int)bit.size();++i){</pre>
        int p=i+(i&-i);//index to parent
        if (p<(int)bit.size()){</pre>
             for(auto&x:bit[i])bit[p].ins(x);
MoQueries.h
Description: answering offline quries
Time: \mathcal{O}\left((N+Q)\sqrt{N}\right)
                                                       9df2fb, 46 lines
/* TODO: use 0 based indexing*/
void remove(int idx); // TODO: remove value at idx from data
     structure
                        // TODO: add value at idx from data
void add(int idx);
     structure
int get_answer(); // TODO: extract the current answer of the
     data structure
int block size:
struct Ouerv {
    int 1, r, idx;
    bool operator<(Query other) const
        return make_pair(l / block_size, r) <</pre>
                make_pair(other.l / block_size, other.r);
vector<int> mo_s_algorithm(vector<Query>& queries) {
    vector<int> answers(queries.size());
    sort(queries.begin(), queries.end());
    // TODO: initialize data structure
    int cur_1 = 0;
    int cur r = -1;
```

FenwickTree RMQ HashMap DSU Geometry

rmq[i][j]=min(rmq[i][j-1],rmq[i+(1<<(j-1))][j]

} else if(i+(1<<(j-1))<(int)a.size()){</pre>

-11);

}

#pragma GCC optimize("03")

#pragma GCC optimize("unroll-loops")

}

```
// invariant: data structure will always reflect the range
          [cur_l, cur_r]
    for (Query q : queries) {
         while (cur_l > q.1) {
             cur_1--;
             add(cur 1);
         while (cur_r < q.r) {</pre>
             cur_r++;
             add(cur r):
         while (cur_1 < q.1) {
             remove(cur_1);
             cur 1++;
         while (cur_r > q.r) {
             remove(cur_r);
             cur_r--;
         answers[q.idx] = get_answer();
    return answers;
FenwickTree.h
Description: find culmulative sum to ith element
Time: \mathcal{O}(\log N)
                                                          a9def9, 27 lines
//TODO: use 1 base indexing
vector<long long>bit;
//range sum point update(k=new_val-old_val)
void add(int i, int k){//add k to ith data and it's parent
     range so on
    while(i<(int)bit.size()){</pre>
        bit[i]+=k;
        i+=i&-i;//add last set bit
11 sum(int i) {//culmulative sum to ith data
    11 sum=0;
    while(i>0){
        sum+=bit[i];
        i-=i&-i;
    return sum:
void build bit(vl &a){
    for (int i=1; i < (int) bit.size(); ++i) {</pre>
        int p=i+(i&-i);//index to parent
        if (p<(int)bit.size())bit[p]+=bit[i];</pre>
RMQ.h
Description: Range Minimum Queries on an array, solving offline queries
Time: build \mathcal{O}(N \log N) query \mathcal{O}(1)
                                                         e37d08, 24 lines
int rmq[N][20];
void build_rmq(vi &a) {
    for(int j=0; j<20;++j) {</pre>
         for (int i=0; i < (int) a.size(); ++i) {</pre>
             if(j==0){
                  rmq[i][0]=a[i];
                                                                         typedef long long int 11;
```

```
int query(int 1, int r){
    int i=1, sub_array_size=r-1+1, ans=INF;
    for(int j=0; j<30;++j) {
        if((1<<j)&(sub_array_size)){</pre>
             ans=min(ans,rmq[i][j]);
             i+=(1<<j);
    return ans;
HashMap.h
Description: Hash map with mostly the same API as unordered_map, but
~3x faster. Uses 1.5x memory. Initial capacity must be a power of 2 (if
provided).
<ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
const int RANDOM = chrono::high resolution clock::now().
     time since epoch().count();
struct chash {
    int operator()(int x) const { return x ^ RANDOM; }
gp_hash_table<int, int, chash> table;
DSU.h
\textbf{Description:} \ \ \text{Disjoint-set data structure}.
Time: \mathcal{O}(\log N)
                                                         7efb0c, 22 lines
//TODO: initialized parent[] and _rank[] array
int find set(int v) {
    if (v == parent[v])
        return v:
    return parent[v] = find_set(parent[v]);
void make_set(int v) {
    parent[v] = v;
    _{rank[v]} = 1;
void union_sets(int a, int b) {
    a = find set(a);
    b = find_set(b);
    if (a != b) {
        if (_rank[a] < _rank[b])
             swap(a, b);
        parent[b] = a;
        _rank[a]+=_rank[b];
Geometry.h
Description: pt, line, polygon, circle
                                                        963fd3, 167 lines
//Geometry
#pragma GCC target("avx2")
```

```
typedef long double ld;
typedef complex<ld> pt;
struct line {
 pt P, D; bool S = false;
 line(pt p, pt q, bool b = false) : P(p), D(q - p), S(b) {}
 line(pt p, ld th) : P(p), D(polar((ld)1, th)) {}
struct circ { pt C; ld R; };
#define X real()
#define Y imag()
#define CRS(a, b) (conj(a) * (b)).Y //scalar cross product
#define DOT(a, b) (conj(a) * (b)).X //dot product
#define U(p) ((p) / abs(p)) //unit vector in direction of p (
     don't use if Z(p) = true
#define Z(x) (abs(x) < EPS)
#define A(a) (a).begin(), (a).end() //shortens sort(),
     upper_bound(), etc. for vectors
//constants (INF and EPS may need to be modified)
1d PI = acosl(-1), INF = 1e20, EPS = 1e-12;
pt I = \{0, 1\};
//true if d1 and d2 parallel (zero vectors considered parallel
     to everything)
bool parallel(pt d1, pt d2) { return Z(d1) || Z(d2) || Z(CRS(U(
    d1), U(d2))); }
//"above" here means if l & p are rotated such that l.D points
     in the +x direction, then p is above l. Returns arbitrary
     boolean if p is on l
bool above_line(pt p, line 1) { return CRS(p - 1.P, 1.D) > 0; }
//true if p is on line l
bool on_line(pt p, line 1) { return parallel(1.P - p, 1.D) &&
    (!1.S | | DOT(1.P - p, 1.P + 1.D - p) \le EPS); }
//returns 0 for no intersection, 2 for infinite intersections,
    1 otherwise. p holds intersection pt
ll intsct(line l1, line l2, pt& p) {
 if (parallel(11.D, 12.D)) //note that two parallel segments
       sharing one endpoint are considered to have infinite
       intersections here
    return 2 * (on line(11.P, 12) || on line(11.P + 11.D, 12)
         || on_line(12.P, 11) || on_line(12.P + 12.D, 11));
 pt q = 11.P + 11.D * CRS(12.D, 12.P - 11.P) / CRS(12.D, 11.D)
 if(on_line(q, 11) && on_line(q, 12)) { p = q; return 1; }
  return 0;
//closest pt on l to p
pt cl_pt_on_l(pt p, line l) {
 pt q = 1.P + DOT(U(1.D), p - 1.P) * U(1.D);
 if(on_line(q, 1)) return q;
 return abs(p - 1.P) < abs(p - 1.P - 1.D) ? 1.P : 1.P + 1.D;
//distance from p to l
ld dist_to(pt p, line l) { return abs(p - cl_pt_on_l(p, l)); }
//p reflected over l
pt refl_pt(pt p, line l) { return conj((p - 1.P) / U(1.D)) * U(
    1.D) + 1.P; }
//ray r reflected off l (if no intersection, returns original
    ray)
line reflect_line(line r, line l) {
 pt p; if(intsct(r, 1, p) - 1) return r;
```

KMP StringHashing DnQDp

```
return line(p, p + INF * (p - refl_pt(r.P, 1)), 1);
//altitude from p to l
line alt(pt p, line 1) { 1.S = 0; return line(p, cl_pt_on_l(p,
    1)); }
//angle bisector of angle abc
line ang_bis(pt a, pt b, pt c) { return line(b, b + INF * (U(a
    - b) + U(c - b)), 1); }
//perpendicular\ bisector\ of\ l\ (assumes\ l.S=1)
line perp_bis(line 1) { return line(l.P + 1.D / (ld)2, arg(l.D
    * I)); }
//orthocenter of triangle abc
pt orthocent(pt a, pt b, pt c) { pt p; intsct(alt(a, line(b, c)
    ), alt(b, line(a, c)), p); return p; }
//incircle of triangle abc
circ incirc(pt a, pt b, pt c) {
  pt cent; intsct(ang_bis(a, b, c), ang_bis(b, a, c), cent);
  return {cent, dist to(cent, line(a, b))};
//circumcircle of triangle abc
circ circumcirc(pt a, pt b, pt c) {
  pt cent; intsct(perp_bis(line(a, b, 1)), perp_bis(line(a, c,
      1)), cent);
  return {cent, abs(cent - a)};
//is pt p inside the (not necessarily convex) polygon given by
    poly
bool in_poly(pt p, vector<pt>& poly) {
  line l = line(p, {INF, INF * PI}, 1);
  bool ans = false;
  pt lst = poly.back(), tmp;
  for(pt q : polv) {
   line s = line(q, lst, 1); lst = q;
    if(on_line(p, s)) return false; //change if border included
    else if(intsct(l, s, tmp)) ans = !ans;
  return ans;
//area of polygon, vertices in order (cw or ccw)
ld area(vector<pt>& poly) {
  1d ans = 0;
  pt lst = poly.back();
  for(pt p : poly) ans += CRS(lst, p), lst = p;
  return abs(ans / 2);
//perimeter of polygon, vertices in order (cw or ccw)
ld perim(vector<pt>& poly) {
 1d ans = 0;
  pt lst = poly.back();
  for(pt p : poly) ans += abs(lst - p), lst = p;
  return ans;
//centroid of polygon, vertices in order (cw or ccw)
pt centroid(vector<pt>& poly) {
 1d area = 0;
  pt lst = poly.back(), ans = \{0, 0\};
  for(pt p : poly) {
   area += CRS(lst, p);
   ans += CRS(lst, p) \star (lst + p) / (ld)3;
```

```
lst = p;
 return ans / area;
//invert a point over a circle (doesn't work for center of
    circle)
pt circInv(pt p, circ c) {
   return c.R * c.R / conj(p - c.C) + c.C;
//vector of intersection pts of two circs (up to 2) (if circles
     same, returns empty vector)
vector<pt> intsctCC(circ c1, circ c2) {
 if(c1.R < c2.R) swap(c1, c2);</pre>
 pt d = c2.C - c1.C;
 if(Z(abs(d) - c1.R - c2.R)) return {c1.C + polar(c1.R, arg(c2
       .C - c1.C));
 if(!Z(d) && Z(abs(d) - c1.R + c2.R)) return {c1.C + c1.R * U(
      d) }:
 if(abs(abs(d) - c1.R) >= c2.R - EPS) return {};
 ld th = acosl((c1.R * c1.R + norm(d) - c2.R * c2.R) / (2 * c1)
       .R * abs(d));
 return {c1.C + polar(c1.R, arg(d) + th), c1.C + polar(c1.R,
      arg(d) - th)};
//vector of intersection pts of a line and a circ (up to 2)
vector<pt> intsctCL(circ c, line 1) {
 vector<pt> v, ans;
 if (parallel(1.D, c.C - 1.P)) v = \{c.C + c.R * U(1.D), c.C - c\}
       .R * U(1.D);
 else v = intsctCC(c, circ{refl_pt(c.C, 1), c.R});
 for(pt p : v) if(on_line(p, l)) ans.push_back(p);
 return ans;
//external tangents of two circles (negate c2.R for internal
    tangents)
vector<line> circTangents(circ c1, circ c2) {
 pt d = c2.C - c1.C;
 1d dr = c1.R - c2.R, d2 = norm(d), h2 = d2 - dr * dr;
 if(Z(d2) || h2 < 0) return {};
 vector<line> ans;
 for(ld sq : {-1, 1}) {
   pt u = (d * dr + d * I * sqrt(h2) * sq) / d2;
   ans.push back(line(c1.C + u * c1.R, c2.C + u * c2.R, 1));
 if(Z(h2)) ans.pop back();
 return ans;
KMP.h
Description: pattern searching
Time: \mathcal{O}(N+M)
                                                     4dfee5, 37 lines
int b[N];
int cnt = 0;
void knp_proc(string t,string p) {
   int i=0, j=-1; b[0] = -1;
   while(i<(int)p.length()){</pre>
       while(j>=0 && p[i]!=p[j]) j =b[j];
       ++i;++j;
       b[i] = j;
```

```
void knp_search(string t, string p) {//count number of occurrence
     of p in t
    int i=0, i=0;
    while(i<(int)t.length()){</pre>
        while(j>=0 && t[i] != p[j]) j = b[j];
        ++i;++j;
        if(j==(int)p.length()){
            ++cnt;
             j = b[j];
vector<int> prefix function(string s) {
    int n = (int)s.length();
    vector<int> pi(n);
    for (int i = 1; i < n; i++) {</pre>
        int j = pi[i-1];
        while (j > 0 \&\& s[i] != s[j])
            j = pi[j-1];
        if (s[i] == s[j])
            j++;
        pi[i] = j;
    return pi;
StringHashing.h
Description: check equality of two substrings
Time: \mathcal{O}(N) preprocessing \mathcal{O}(1) query
                                                        276f01, 25 lines
typedef long long 11;
typedef pair<ll, ll> pl;
#define M 1000000321
#define OP(x, y) pl operator x (pl a, pl b) {return {a.first x b
     .first, (a.second y b.second) % M }; }
OP(+, +) OP(*, *) OP(-, + M -)
//random number generator
mt19937 gen(chrono::steady_clock::now().time_since_epoch().
uniform_int_distribution<11> dist(256, M - 1);
//queries - check \ if \ S[i:i+l] = S[j:j+l](inclusive), \ S \ is \ a
     string, [:] is slice
#define H(i, 1) (h[(i) + (1)] - h[i] * p[1])
#define EQ(i, j, l) (H(i, l) == H(j, l))
//preprocessina
const int N = 2e5:
string s;
pl p[N], h[N];
11 n;
int main(){
    cin >> n >> s;
    p[0] = \{1,1\}, p[1] = \{dist(gen) | 1, dist(gen) \};
    for(ll i=1;i<=(ll)s.length();++i){</pre>
        p[i] = p[i-1] * p[1];
        h[i] = h[i-1] * p[1] + make_pair(s[i-1], s[i-1]);
DnQDp.h
Description: find best k consecutive subbarray partition
Time: \mathcal{O}(N \log N)
//TODO: initialize dp , initialize cost() function of use slide
```

 $//ll \stackrel{\sim}{dp}[N][M];$

CHT KthAncestor LCA Sieve IntFact

most and rightmost of convex hull so do search on both l

 $<=cur.intersectX(dq[dq.size()-2]))dq.pop_back();$

dq.push_front({0, 0});//cant be put in global, remove this to

//TODO NOTE***: maximum and minimum value exist in bot left

//constructing hull from l to r, maintain correct hull at

while(dq.size()>=266cur.intersectX(dq.back())

local function

to r and r to l

/* ***inserting line (maximum hull)

 $line \ cur = line \{ \dots some \ m, \dots some \ c \}$

rightmost

```
//generic implementation for sliding range technique for logn
//(persistent segtree alternative)
11 ccost = 0;
int c1 = 0, cr = -1;
void slide(int 1, int r){
    while(cr < r){</pre>
        ++cr;
         //add();
         //...
    while (cl > 1) {
        --cl;
         //add();
        //...
    while(cr > r){
        //remove();
        //...
        --cr;
    while(c1 < 1){
        //remove();
        //...
        ++cl;
void compute(int 1, int r, int opt1, int optr, int j){
    if (1>r) return;
    int mid = (1+r) >> 1;
    //pair < ll, int > best = \{0,-1\};
    //pair < ll, int > best = \{LINF, -1\};
    //dp is satisfy quadrangle IE if cost() satisfy quadrangle
    //if \ cost() \ is \ QF \Rightarrow opt() \ is \ nondecreasing
    for(int k=optl;k<=min(mid,optr);++k){</pre>
        slide(k, mid);
        //best = max(best, \{((k>0)?dp[k-1][j-1]:0) + ccost, k\}
        //best = min(best, \{((k>0)?dp[k-1][j-1]:0) + ccost, k\}
    //dp[mid][j] = max(dp[mid][j], best.first);
    //dp[mid][j] = min(dp[mid][j], best.first);
    int opt = best.second;
    if(l!=r){
        compute(1, mid-1, opt1, opt, j);
        compute(mid+1, r, opt, optr, j);
//TODO: set dp to LINF or -LINF
CHT.h
Description: convex-hull trick
Time: \mathcal{O}(N) or \mathcal{O}(N \log N) if sort the slope
                                                         87ad47, 28 lines
struct line {
    long long m, c;
    long long eval(long long x) { return m * x + c; }
    long double intersectX(line 1) { return (long double) (c -
         1.c) / (1.m - m); }
deque<line> dq;
```

```
dq.pb(cur);
 //constructing hull from r to l, maintain correct hull at
               leftmost
 /* inserting line (maximum hull)
              line \ cur = line \{ \dots some \ m, \dots some \ c \}
               while(dq.size() \ge 28 \& cur.intersectX(dq[0]) \ge -cur.intersectX(dq[0]) \ge -cur.intersectX(dq[0])
                              dq[1]))_{\{}
                           dq.pop\_front();
               dq.push_front(cur);
 KthAncestor.h
 Description: find kth-ancestor of a tree-node
 Time: \mathcal{O}(\log N)
                                                                                                                                                                               62aeee, 9 lines
 int kth_ancestor(int node,int k) {
              if(depth[node] < k)return -1;</pre>
              for(int i = 0; i < LOG; ++i){</pre>
                          if(k & (1<<i)){
                                       node = up[node][i];
              return node;
 LCA.h
 Description: find lowest common ancestor of two tree-nodes
 Time: \mathcal{O}(\log N)
                                                                                                                                                                            300c4f, 36 lines
 //TODO: initialize tree(adj list)
const int LOG = 20;
 int depth[N], parent[N];
int up[N][LOG]; // 2^j-th ancestor of n
void dfs(int a,int e){
              for(auto b:adj[a]) {
                          if(b == e)continue;
                          depth[b] = depth[a] + 1;
                          parent[b] = a;
                          up[b][0] = parent[b];
                           for (int i=1; i < LOG; ++i) {</pre>
                                        up[b][i] = up[up[b][i-1]][i-1];
                          dfs(b,a);
int lca(int a, int b) {
              if (depth[a] < depth[b]) swap(a, b);</pre>
              int k = depth[a] - depth[b];
              for (int i=LOG-1; i>=0; --i) {
                          if(k &(1<<i)){
```

```
a = up[a][i];
//if query ask for minimum remove this line after 1st insertion
                                                                           if (a == b) return a;
                                                                           for(int i=LOG-1;i>=0;--i){
                                                                               if(up[a][i] != up[b][i]){
                                                                                   a = up[a][i];
                                                                                   b = up[b][i];
                                                                           return up[a][0];
                                                                      Sieve.h
                                                                      Description: prime sieve
                                                                      Time: \mathcal{O}(N \log N)
                                                                                                                               abb3a3, 21 lines
                                                                      //can use to find all prime factor of a number in O(\log n)
                                                                      const int M = 2e5+1;
                                                                      vector<bool> is_prime(M+1, true);
                                                                      void sieve(){
                                                                           is_prime[0] = is_prime[1] = false;
                                                                           for (int i = 2; i * i <= M; i++) {
                                                                               if (is prime[i]) {
                                                                                   for (int j = i * i; j <= M; j += i)</pre>
                                                                                        is prime[j] = false;
                                                                           /* log n sieve (use sieve to find all prime factors in O(
                                                                           for (int i = 2; i \le M; i++)is_prime[i] = i;
                                                                           for (int i = 2; i * i <= M; i++) {
                                                                               if (is\_prime[i] == i)  {
                                                                                   for (int \ j = i * i; \ j <= M; \ j += i)
                                                                                        is_prime[j] = i;
                                                                      IntFact.h
                                                                      Description: integer factorization algorithm
                                                                      Time: \mathcal{O}\left(\sqrt{N}\right)
                                                                                                                              497201, 40 lines
                                                                       //in general any natural number n has at most n^1/3 divisors in
                                                                             practice
                                                                      bool isPrime(ll x) {
                                                                           for (11 d = 2; d * d <= x; d++) {
                                                                               if (x % d == 0)
                                                                                   return false;
                                                                           return x >= 2;
                                                                      void decompose(11 x) {
                                                                           vl temp;
                                                                           while(x % 2 == 0){
                                                                               temp.pb(2);
                                                                               x/=2;
                                                                           for(11 i=3;i*i <= x;i+=2){</pre>
                                                                               if(x % i == 0){
                                                                                   while(x % i == 0){
                                                                                        x/=i;
                                                                                        temp.pb(i);
```

GcdExtended Matrix Binpow

```
if(x>1) temp.pb(x);
    //do something
void find_all_divisors(ll x){
    vl temp;
    for (int i=1; (ll) i*i<=x;++i) {</pre>
        if(x%i==0){
             if(i==x/i) temp.pb(i);
             else {
                 temp.pb(i); temp.pb(x/i);
    //temp = all \ divisors \ of \ x
```

GcdExtended.h

Description: find Bezout's coefficient

Time: $\mathcal{O}(\log N)$

af07ae, 12 lines

```
int gcd(int a, int b, int& x, int& y) {
    if (b == 0) {
       x = 1;
       v = 0;
        return a;
    int x1, y1;
    int d = gcd(b, a % b, x1, y1);
    y = x1 - y1 * (a / b);
    return d;
```

Description: mainly for matrix exponentation and find nth recurrence Time: $\mathcal{O}(\log N)$

1e389f, 28 lines

```
struct Matrix{
    double a[2][2] = \{\{0,0\},\{0,0\}\};
    Matrix operator * (const Matrix& other) {
        Matrix product:
        for(int i=0;i<2;++i){</pre>
            for(int j=0; j<2;++j) {
                 for(int k=0; k<2; ++k) {
                     product.a[i][k] += a[i][j] * other.a[j][k];
        return product;
//for calculating nth recurrence of function
//entry Matrix[i][j] is problability/number of way ith state
     change to jth state
Matrix expo_power(Matrix a, int n) {
    Matrix product;
    for(int i=0;i<2;++i)product.a[i][i]=1;</pre>
    while (n>0) {
        if(n&1){
            product = product * a;
        a = a * a;
        n >> = 1;
    return product;
```

```
Binpow.h
```

Description: binary exponentation

Time: $\mathcal{O}(\log N)$

d4debd, 11 lines

```
long long binpow(long long a, int n) {
    long long res = 1;
    while(n>0){
        if(n&1){
            res = res * a;
        a = a * a;
    return res;
```

Mathematics (3)

3.1 Equations

$$ax^{2} + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^{2} - 4ac}}{2a}$$

The extremum is given by x = -b/2a.

$$ax + by = e$$

$$cx + dy = f$$

$$\Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

In general, given an equation Ax = b, the solution to a variable x_i is given by

$$x_i = \frac{\det A_i'}{\det A}$$

where A'_i is A with the *i*'th column replaced by b.

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1 n + d_2)r^n.$

3.2 Trigonometry

$$\sin(v + w) = \sin v \cos w + \cos v \sin w$$
$$\cos(v + w) = \cos v \cos w - \sin v \sin w$$

$$\tan(v+w) = \frac{\tan v + \tan w}{1 - \tan v \tan w}$$
$$\sin v + \sin w = 2\sin\frac{v+w}{2}\cos\frac{v-w}{2}$$
$$\cos v + \cos w = 2\cos\frac{v+w}{2}\cos\frac{v-w}{2}$$

$$(V+W)\tan(v-w)/2 = (V-W)\tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w.

$$a\cos x + b\sin x = r\cos(x - \phi)$$

$$a\sin x + b\cos x = r\sin(x + \phi)$$

where
$$r = \sqrt{a^2 + b^2}$$
, $\phi = \operatorname{atan2}(b, a)$.

3.3 Geometry

3.3.1 Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area:
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{}$

Length of median (divides triangle into two equal-area triangles):

 $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):

$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$ Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

3.3.2 Quadrilaterals $\tan \frac{\alpha + \beta}{2}$

With of itemsengths $\frac{a+b}{a,b,b}$, \overline{a} , $\overline{diagonals}$ e,f, diagonals angle θ , area A and magic flux $F=b^2+t^2$ $\frac{a}{2}$ $\frac{a}{2}$ $\frac{a}{2}$ $\frac{a}{2}$ $\frac{a}{2}$ $\frac{a}{2}$

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

3.3.3 Spherical coordinates

For cyclic quadrilaterals the sum of opposite angles is 180°, ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.



$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \arccos(z/\sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \operatorname{atan2}(y, x) \end{aligned}$$

3.4 Derivatives/Integrals

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \qquad \frac{d}{dx}\arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x\sin ax = \frac{\sin ax - ax\cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \qquad \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

3.5 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c-1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

3.6 Series

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

3.7 Probability theory

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x. It will then have an expected value (mean) $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$ and variance $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent X and Y,

$$V(aX + bY) = a^2V(X) + b^2V(Y).$$

3.7.1 Discrete distributions Binomial distribution

The number of successes in n independent yes/no experiments, each which yields success with probability p is $Bin(n, p), n = 1, 2, ..., 0 \le p \le 1$.

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \, \sigma^2 = np(1-p)$$

Bin(n, p) is approximately Po(np) for small p.

First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability p is Fs(p), $0 \le p \le 1$.

$$p(k) = p(1-p)^{k-1}, k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \ \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $Po(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$

$$\mu = \lambda, \, \sigma^2 = \lambda$$

3.7.2 Continuous distributions Uniform distribution

If the probability density function is constant between a and b and 0 elsewhere it is U(a, b), a < b.

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \, \sigma^2 = \frac{(b-a)^2}{12}$$

Exponential distribution

The time between events in a Poisson process is $\operatorname{Exp}(\lambda)$, $\lambda > 0$.

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$
$$\mu = \frac{1}{\lambda}, \, \sigma^2 = \frac{1}{\lambda^2}$$

Normal distribution

Most real random values with mean μ and variance σ^2 are well described by $\mathcal{N}(\mu, \sigma^2)$, $\sigma > 0$.

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$ and $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$ then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$