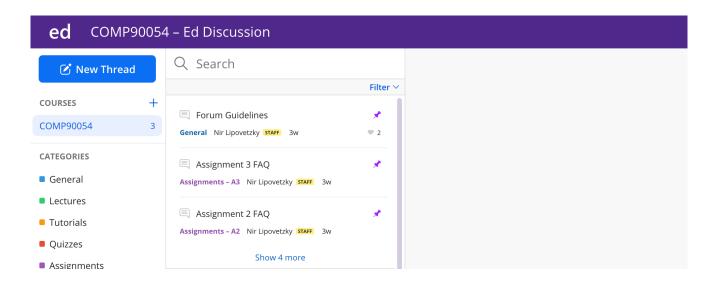
Week 2: Blind Search

COMP90054 – Al Planning for Autonomy

• Tutor Name: Thao Le

• Questions? Ed Discussion



Key concepts

- State-space model
- Blind search algorithms:
 - Breadth First Search (BFS)
 - Depth First Search (DFS)
 - Iterative Deepening (ID)

State space model:

$$P = \langle s_0, S, S_G, A, f, c \rangle$$

• s_0 is an initial state

*s*1

• S is a set that includes all states in the state space

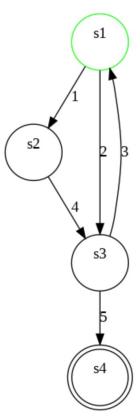
$$S = \{s1, s2, s3, s4\}$$

• S_G is a goal set with all goal states

$$S_G = \{s4\}$$

▼ Problem 1:

Following the above example, define the state-space model of the graph:



State space model:

$$P = \langle s_0, S, S_G, A, f, c \rangle$$

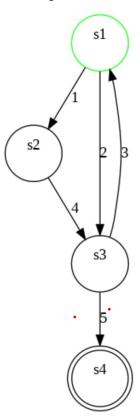
• A is an action set that includes all possible actions you can take from a state

A(s1) = {(s1,s2), (s1,s3)} A(s4) = ?
A(s4) =
4

- f is a transition function: f(s, a) = s'f(s1, (s1, s2)) = s2
- c is a cost function between two states c(s1, s2) = 1

▼ Problem 1:

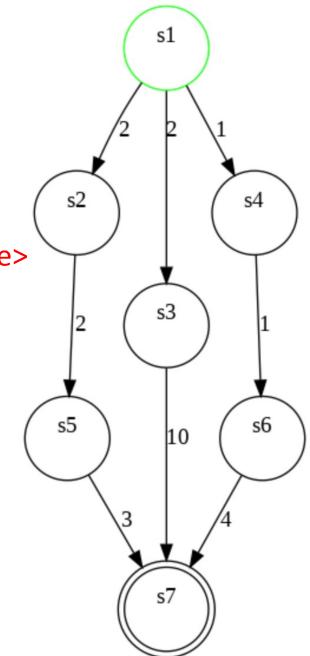
Following the above example, define the state-space model of the graph:



State vs Search node?

Search node n = <state, accumulated cost g(n), id of parent node>

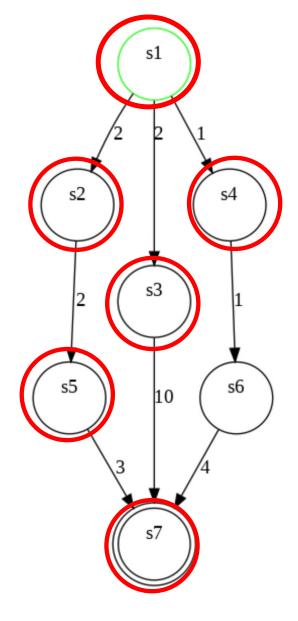
n0 = <s1, 0, None>



Breadth First Search (BFS) Queue: FIFO

Search node n = (state, accumulated cost g(n), id of parent node)

Solution: s1 -> s3 -> s7



Depth First Search (DFS) Stack: LIFO

Search node n = (state, accumulated cost g(n), id of parent node)

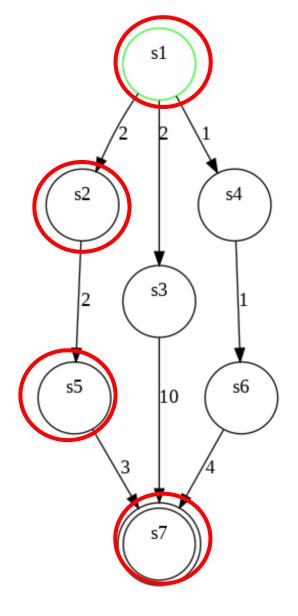
$$n0 = (s1, 0, None)$$

$$n1 = (s2, 2, 0)$$

$$n2 = (s5, 4, 1)$$

$$n3 = (s7, 7, 2)$$

Solution: s1 -> s2 -> s5 -> s7

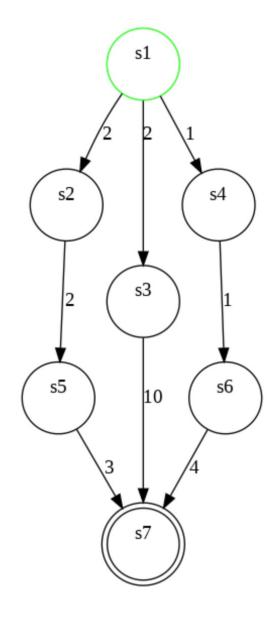


Question 2

Depth First Search (DFS)

	Open (Stack)	Close (Visited)
Iteration 1	n0 = <s1, 0,="" none=""></s1,>	
Iteration 2	n1 = <s2, 0="" 2,=""> n2 = <s3, 0="" 2,=""> n3 = <s4, 0="" 1,=""></s4,></s3,></s2,>	n0
Iteration 3	n4 = <s5, 1="" 4,=""> n2 = <s3, 0="" 2,=""> n3 = <s4, 0="" 1,=""></s4,></s3,></s5,>	n0, n1
Iteration 4	n5 = <s7, 4="" 7,=""> n2 = <s3, 0="" 2,=""> n3 = <s4, 0="" 1,=""></s4,></s3,></s7,>	n0, n1, n4
Iteration 5	n2 = <s3, 0="" 2,=""> n3 = <s4, 0="" 1,=""></s4,></s3,>	n0, n1, n4, n5

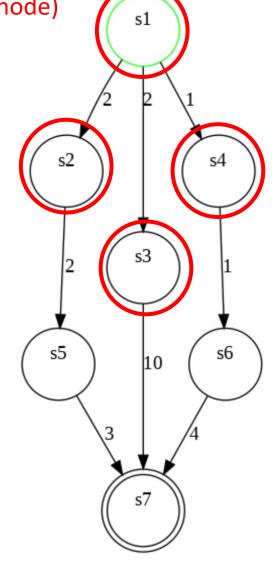
(s1, 0, None), (s2, 2, 0), (s5, 4, 1), (s7, 7, 2)



Search node n = (state, accumulated cost g(n), id of parent node)

Iterative Deepening (ID) DFS with a depth limit

Limit	Step	Open (Stack)	Close
0	1	n0 = (s1,0, None)	
	2		<mark>n0</mark>
1	3	n1 = (s1, 0, None)	
	4	n2 = (s2, 2, 1)	n1
		n3 = (s3, 2, 1)	
		n4 = (s4, 1, 1)	
	5	n3 = (s3, 2, 1)	n1, n2
		n4 = (s4, 1, 1)	
	6	n4 = (s4, 1, 1)	n1, n2, n3
	7		n1, n2, n3, n4

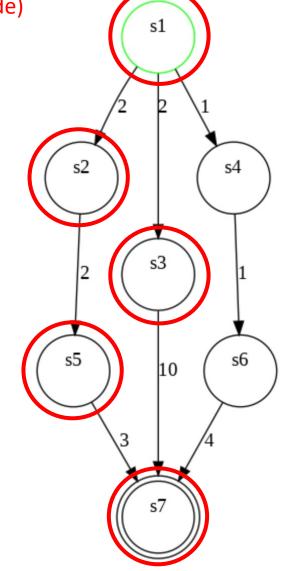


Search node n = (state, accumulated cost g(n), id of parent node)

Iterative Deepening (ID)

DFS with depth limit

Limit	Step	Open (Stack)	Close
2	8	n5= (s1, 0, None)	
	9	n6 = (s2, 2, 5) n7= (s3, 2, 5) n8= (s4, 1, 5)	n5
	10	n9= (s5, 4, 6) n7= (s3, 2, 5) n8= (s4, 1, 5)	n5, n6
	11	n7 = (s3, 2, 5) n8 = (s4, 1, 5)	n5, n6, n9
	12	n10 = (s7, 12, 7) n8 = (s4, 1, 5)	n5, n6, n9, n7
	13	n8 = (s4, 1, 5)	<mark>n5, n6, n9, n7, n</mark> 10



n0, n1, n2, n3, n4, n5, n6, n9, n7, n10

Search node n = (state, accumulated cost g(n), id of parent node)

Iterative Deepening (ID)

DFS with depth limit

Expansion node order

n0 = (s1, 0, None),

Expansion node order

<mark>n0,</mark> <mark>n1, n2, n3, n4,</mark> n5, n6, n9, n7, n10

$$n1 = (s1, 0, None),$$

$$n2 = (s2, 2, 1),$$

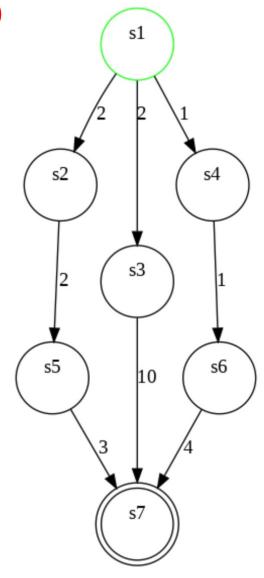
$$n3 = (s3, 2, 1),$$

$$n4 = (s4, 1, 1),$$

$$116 = (52, 2, 5),$$

$$n9 = (s5, 4, 6),$$

$$n10 = (s7, 12, 7)$$



Solution: s1 -> s3 -> s7

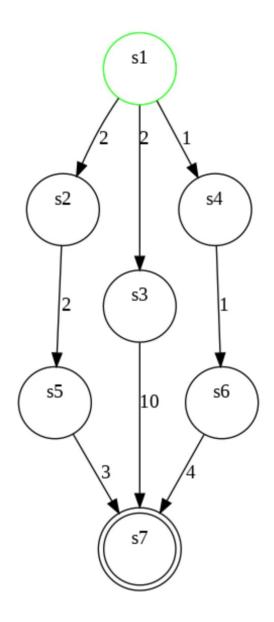
Q2: What is the actual optimal solution?

Q3: Explain under which conditions the algorithms guarantee optimality?

BrFS and ID will be optimal if the costs are equal, such as, all cost are 1

Q4: Can any of the previous algorithms be adapted to account for g(n) in order to make it optimal?

Dijkstra. Expanding the node with lowest accumulated cost, instead of the lowest depth.



Describe a simple example of *Travelling Salesman Problem* along with its corresponding **State Space Model**.

Definition should be brief, clear, and *compact* (*compact* means using mathematical notation to define sets, i.e. $S = \{x | x \in V\}$ to define that there are as many states as elements in the set V, and pseudo-code, i.e. to define the transition function.)

- 1. State space S
- 2. Initial state $s_0 \in S$
- 3. Set of goal states $S_G \subseteq S$
- 4. Applicable actions function A(s) for each state $s \in S$
- 5. Transition function f(s, a) for $s \in S$ and $a \in A(s)$
- 6. Cost of each action c(a) for $a \in A(s)$

Hint: Given

- V = a set of cities
- v_{start} = a starting city location
- E = a set of edges specifying if there is an edge between two cities <v1, v2>
- V' = a set of cities that have been visited

Note for the TSP:

A city can be visited more than once

- You can start at any city

The goal state is to visit all cities
and no need to go back to the starting

city

Hint: Given

- V = a set of cities
- v_{start} = a starting city location
- E = a set of edges specifying if there is an edge between two cities <v1, v2>
- V' = a set of cities that have been visited

a state = <current city, a set of visited cities>

an edge/action= <current city, next city>

Initial state
$$s_0 = \langle v_{start}, \{v_{start}\} \rangle$$

Goal state
$$S_G = \{ \langle v_{current}, V \rangle \mid v_{current} \in V \}$$

State
$$S = \{ \langle v_{current}, V' \rangle | v_{current} \in V \land V' \subseteq V \}$$

$$\textbf{Action A(} < v_{current}, V' > \textbf{)} = \{ < v_{current}, v_{next} > | < v_{current}, v_{next} > \in E \}$$

Transition f(<
$$v_{current}$$
, $V'>$, < $v_{current}$, $v_{next}>$) =< v_{next} , $V'\cup\{v_{next}\}>$

$$c(< v_{current}, v_{next} >) = cost(< v_{current}, v_{next} >)$$

Criteria for Evaluating Search Strategies

Completeness: Is the strategy guaranteed to find a solution when there is one?

Optimality: Are the returned solutions guaranteed to be optimal?

Criteria for Evaluating Search Strategies

Consider when the graph is finite and cyclic

	BFS	DFS	ID
Complete	Yes	No	Yes
Optimal	Yes when costs are all equal	No	Yes when costs are all equal
Time complexity	$O(b^d)$	$O(b^m)$	$O(b^d)$
Space complexity	$O(b^d)$	O(b.m)	O(b.d)

b is the maximal branching factor, d is the depth of the shallowest goal state m is maximal depth