# Week 4: STRIPS and Heuristic

COMP90054 – Al Planning for Autonomy

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# Key concepts

- STRIPS problem
- Heuristic functions

Consider a m x m manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible coordinates (walls) W

a state = <current coordinate, a set of remaining coordinates>

Initial state 
$$s_0 = <(0,0), G' \setminus \{(0,0)\}>$$

Goal state 
$$S_G = \{ \langle (x, y), \{ \} \rangle \mid x, y \in \{0, ..., m-1 \} \}$$

State 
$$S = \{ \langle (x, y), V' \rangle | x, y \in \{0, ..., m-1\} \land V' \subseteq G' \}$$

Action A(
$$<(x,y), V'>$$
) =  $\{(dx,dy) \mid dx,dy \in \{-1,0,1\}$   
  $\land |dx| + |dy| = 1$   
  $\land x + dx, y + dy \in \{0,..., m-1\}$   
  $\land (x + dx, y + dy) \notin W\}$ 

Transition 
$$f(<(x,y), V'>, (dx,dy)) = <(x+dx,y+dy), V'\setminus\{(x+dx,y+dy)\}>$$

$$\mathsf{Cost}\;\mathsf{c}(a)=\mathbf{1}$$

### **State-space model**

$$P = \langle S, S_0, S_G, A, T, c \rangle$$
  
 $S = State space$   
 $S_0 = initial state$ 

$$S_0 = \text{minimization}$$

$$c = costs$$

$$P = \langle S, S_0, S_G, A, T, c \rangle$$
  
 $S = State space$   
 $S_0 = initial state$   
 $S_G = goal states$   
 $A = actions$   
 $T = transition functions$   
 $C = costs$ 

Consider a m x m manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible

coordinates (walls) W

$$I = \{at(0,0), visited(0,0)\}$$

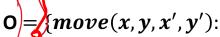
$$\begin{array}{c|c} 2 & & \\ \hline & &$$

$$G = \{visited(x,y)|x,y \in G'\} \leftarrow$$

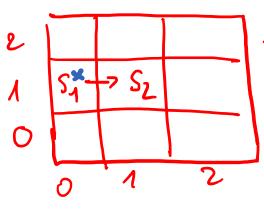
$$F = \{at(x,y), visited(x,y)|x, y \in \{0, ..., m-1\}\} \rightarrow pred$$

- A **problem** in **STRIPS** is a tuple  $P = \langle F, O, I, G \rangle$ :
  - F stands for set of all atoms (boolean vars)
- O stands for set of all operators (actions)
  - $I \subseteq F$  stands for initial situation  $\longleftarrow$
  - $G \subseteq F$  stands for goal situation
- Operators  $o \in O$  represented by

  - the Delete list  $Del(o) \subseteq F$  False
  - the Precondition list  $Pre(o) \subseteq F$



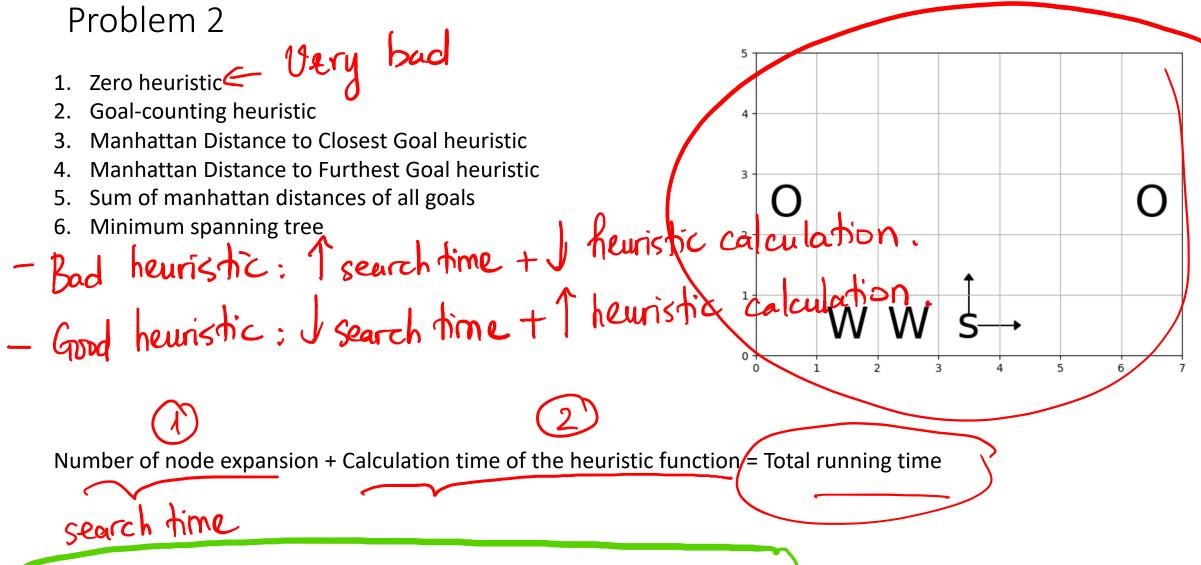
- Prec: at(x,y) Add: at(x', y'), visited(x',y')
- Del: at(x, y) | for each adjacent (x,y), (x',y'), and  $(x',y') \notin W \mathcal{V}$



move ((0,1), (1,1)): +pre: at (0,1)

tadd: at (1, 1), visited (1,1)

+ del. at (0,1)



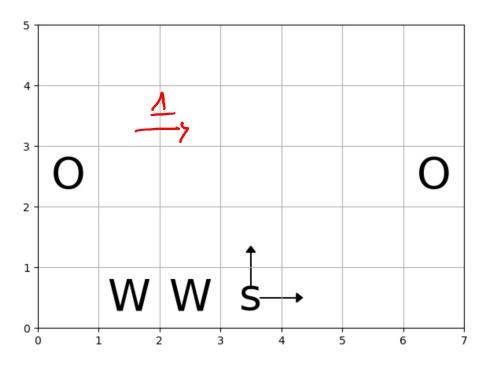
Tip: If h1 dominates h2 then A\* with h1 will expand less or equal node to h2

#### 1. Zero heuristic h = 0

Admissible: Yes Consistent: Yes

Time to calculate h: None

parent child  $h(s) - h(s') \leq c(s, s')$   $0 - 0 \leq 1$ 



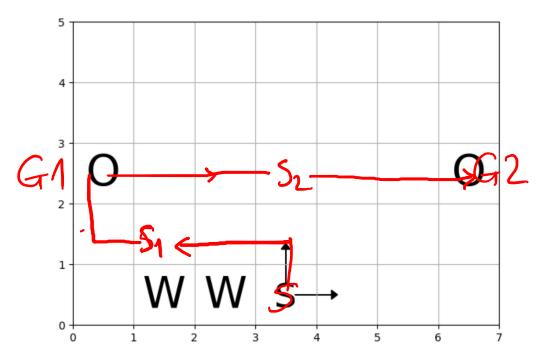
? admissible

- ? consistent

### 2. Goal-counting heuristic

Admissible: Yes Consistent: Yes

Time to calculate h: Easy



$$k(s_1) = 2$$
  
 $k(s_2) = 1$ 

$$f_{s}(s_{s}) = 1$$

#### 3. Manhattan Distance to Closest Goal heuristic

Admissible: Yes Consistent: Yes

Time to calculate h: Easy

$$h(s_1) = min (d(s_1, 41), d(s_1, 42))$$

$$= min (3, 5) = 3$$

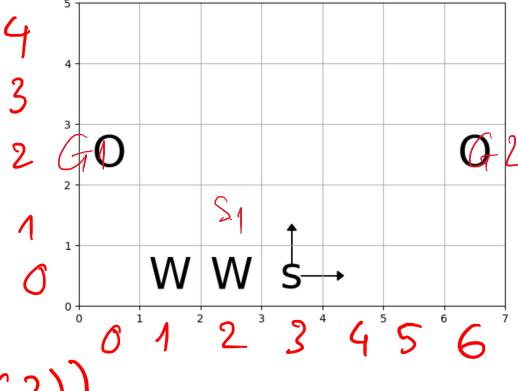
$$c(s, s') = 1$$

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#### 4. Manhattan Distance to Furthest Goal heuristic

Admissible: Yes Consistent: Yes

Time to calculate h: Easy

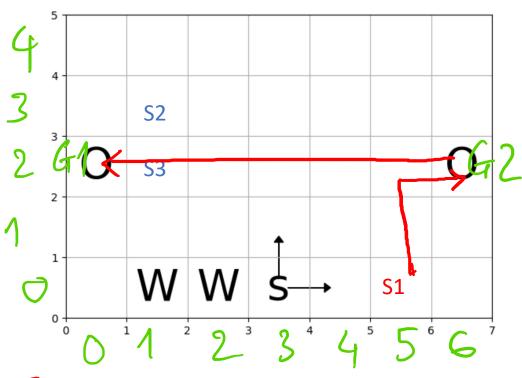


$$h(s_1) = \max(d(s_1,G_1), d(s_1,G_2))$$
 $= \max(3)$ 

### 5. Sum of Manhattan distances of all goals

Admissible: No Consistent: No

Time to calculate h: Easy

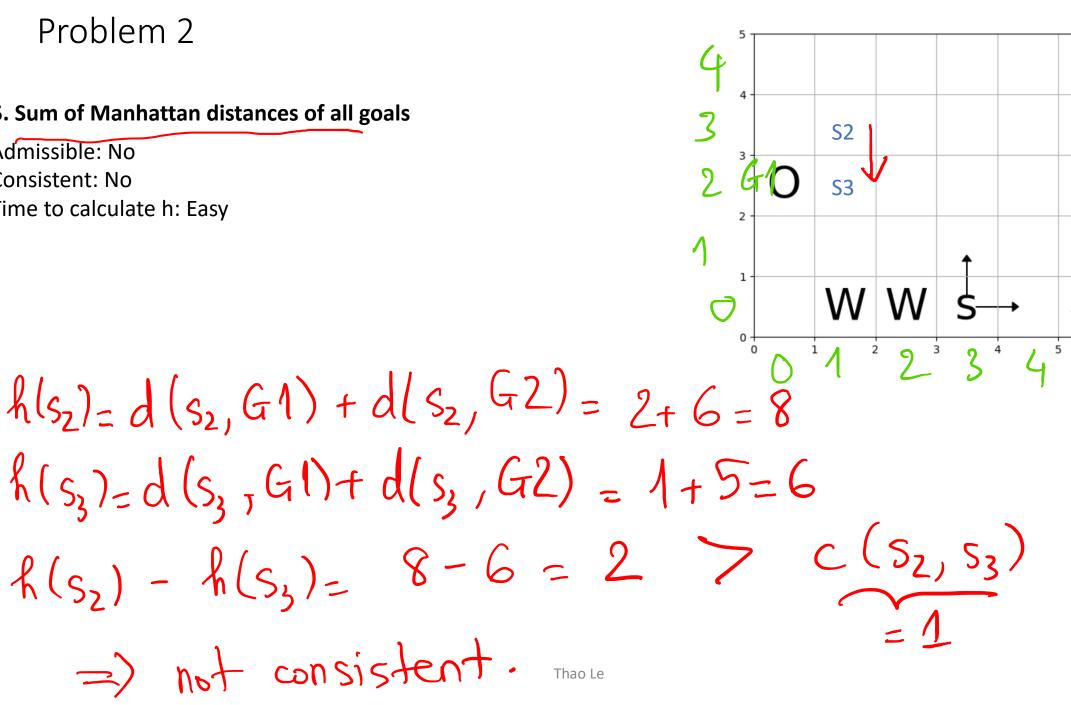


$$h(s_1) = d(s_1, G_1) + d(s_1, G_2) = 7 + 3 = 10$$
  
 $h^*(s_1) = d(s_1, G_2) + d(G_1, G_2) = 3 + 6 = 9$   
 $h(s_1) > h^*(s_1) = not admissible$ 

### 5. Sum of Manhattan distances of all goals

Admissible: No Consistent: No

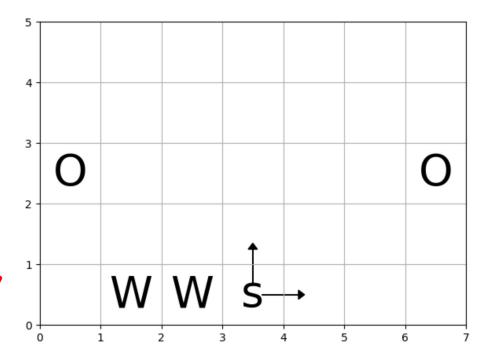
Time to calculate h: Easy

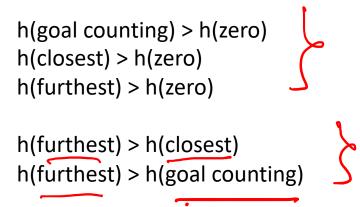


$$h(s_3) = d(s_3, G1) + d(s_3, G2) = 1 + 5 = 6$$
  
 $h(s_2) - h(s_3) = 8 - 6 = 2$ 
  
=) not consistent. That

#### **Dominate relation**

- 1. Zero heuristic: Admissible, Consistent
- 2. Goal-counting heuristic: Admissible, Consistent
- 3. Manhattan Distance to Closest Goal heuristic: Admissible, Consistent
- 4. Manhattan Distance to Furthest Goal heuristic: Admissible, Consistent
- 5. Sum of manhattan distances of all goals: Not admissible, Not consistent



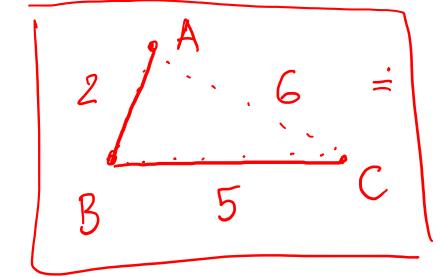


### 6. Minimum spanning tree

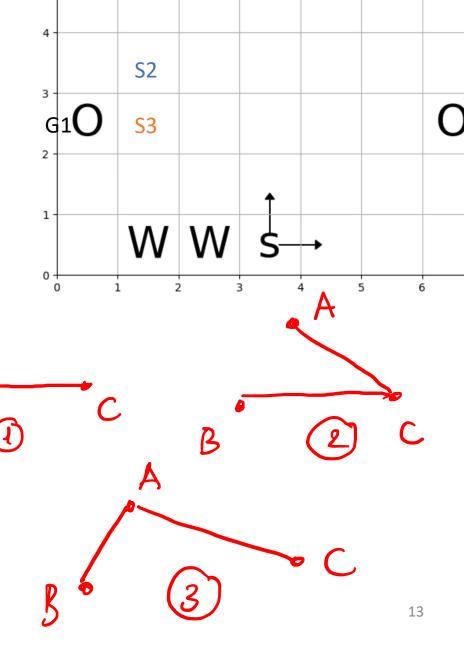
Admissible: Yes Consistent: No

Time to calculate h: Medium

Minimum spanning tree: Only select a subset of edges to connect all vertices and the total cost is minimum



Cost = 7



### 6. Minimum spanning tree

Admissible: Yes Consistent: No

Time to calculate h: Medium

Minimum spanning tree: Only select a subset of edges to connect all vertices and the total cost is minimum

$$h(s_2) = \frac{2}{3}$$

$$h(s_2) = 2+6=8$$

