Week 4: STRIPS and Heuristic

COMP90054 – Al Planning for Autonomy

Key concepts

- STRIPS problem
- Heuristic functions

Consider a m x m manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible coordinates (walls) W

a state = <current coordinate, a set of remaining coordinates>

Initial state
$$s_0 = <(0,0), G' \setminus \{(0,0)\}>$$

Goal state
$$S_G = \{ \langle (x, y), \{ \} \rangle \mid x, y \in \{0, ..., m-1 \} \}$$

State
$$S = \{ \langle (x, y), V' \rangle | x, y \in \{0, ..., m-1\} \land V' \subseteq G' \}$$

Action A(
$$<(x,y), V'>$$
) = $\{(dx,dy) \mid dx,dy \in \{-1,0,1\}$
 $\land |dx| + |dy| = 1$
 $\land x + dx, y + dy \in \{0,...,m-1\}$
 $\land (x + dx, y + dy) \notin W\}$

Transition
$$f(<(x,y), V'>, (dx,dy)) = <(x+dx,y+dy), V'\setminus\{(x+dx,y+dy)\}>$$

$$\mathsf{Cost}\,\mathsf{c}(a)=\mathbf{1}$$

State-space model

$$P = \langle S, S_0, S_G, A, T, c \rangle$$

P = $\langle S, S_0, S_G, A, T, c \rangle$ S = State space S_0 = initial state S_G = goal states A = actions T = transition functions

c = costs

Consider a m x m manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible coordinates (walls) W

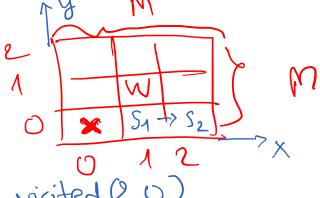
$$(I) = \{at(0,0), visited(0,0)\}$$

$$G = \{visited(x, y) | x, y \in G'\}$$

$$F = \{at(x, y), visited(x, y) | x, y \in \{0, ..., m-1\}\}$$

 $O = \{move(x, y, x', y'):$

- A **problem** in STRIPS is a tuple $P = \langle F, O, I, G \rangle$:
 - F stands for set of all atoms (boolean vars)
 - O stands for set of all operators (actions)
 - $I \subseteq F$ stands for initial situation \longleftarrow
 - \blacksquare $G \subseteq F$ stands for goal situation
- lacktriangle Operators $o \in O$ represented by
 - the Add list $Add(o) \subseteq F$
 - the Delete list $Del(o) \subseteq F$
 - the Precondition list $Pre(o) \subseteq F$



Prec: at(x,y)

Add: at(x',y'), visited(x',y')

Del: at(x,y) | for each adjacent (x,y),(x',y'), and

(x',y') \neq W \}

The action

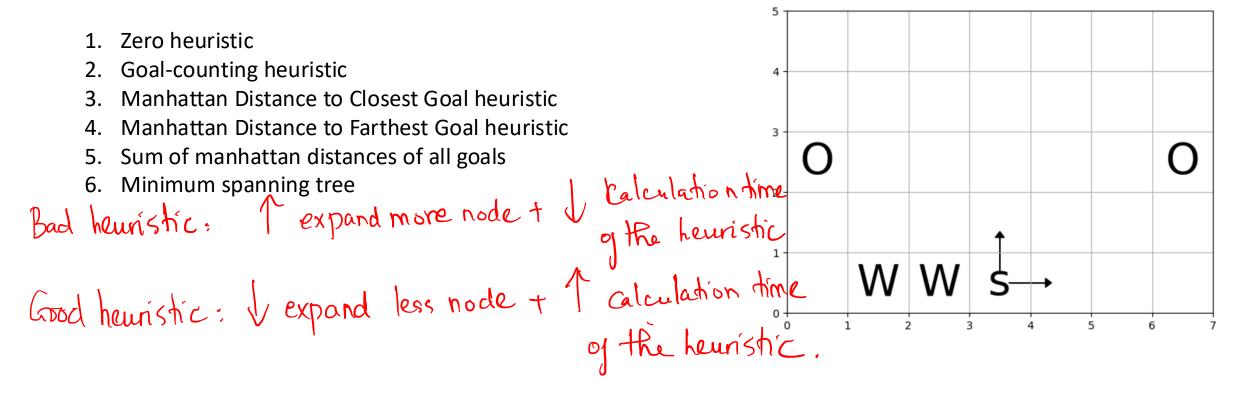
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Number of node expansion + Calculation time of the heuristic function = Total running time

(1) (2)

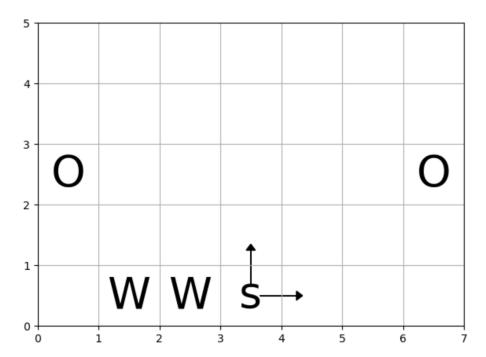
Tip: If h1 dominates h2 then A* with h1 will expand less or equal node to h2

1. Zero heuristic h = 0

Admissible: Yes

Consistent: Yes

Time to calculate h: None



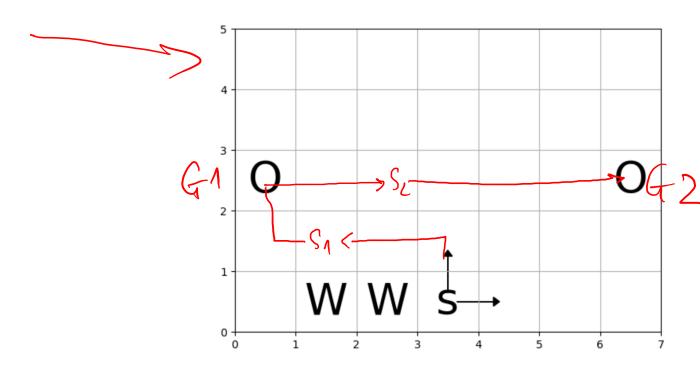
2. Goal-counting heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy

$$h(s_1) = 2$$
 $h(s_2) = 1$



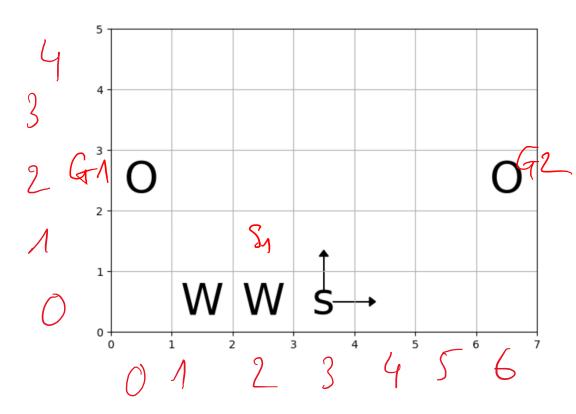
3. Manhattan Distance to Closest Goal heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy

$$h(s_1) = d(G_1, s_1) = 3$$



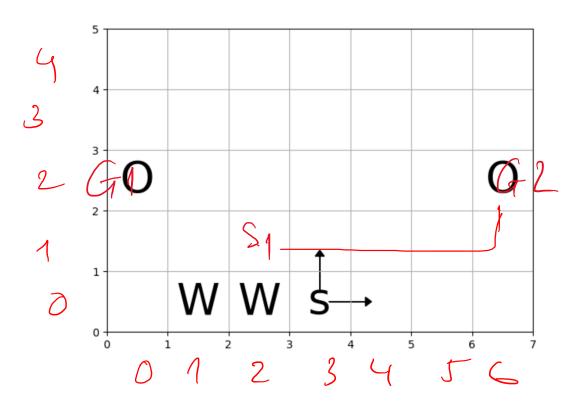
4. Manhattan Distance to Farthest Goal heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy

$$\chi(s_1) = d(s_1, G_2) = 5$$



5. Sum of Manhattan distances of all goals

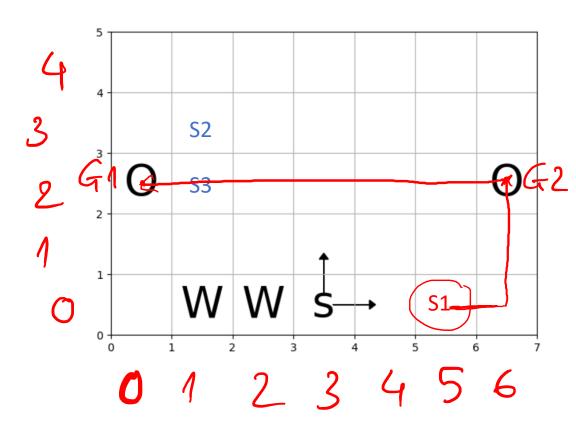
Admissible: No Consistent: No

Time to calculate h: Easy

$$h(s_1) = d(s_1, G_1) + d(s_1, G_2)$$

$$= 7 + 3 = 10$$
 $h*(s_1) = d(s_1, G_2) + d(G_2, G_1)$

$$= 3 + 6 = 9$$
 $h(s_1) > h*(s_1) = Not admissible$

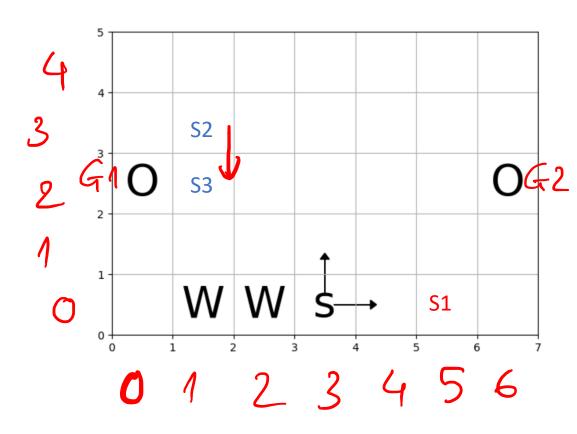


5. Sum of Manhattan distances of all goals

Admissible: No Consistent: No

Time to calculate h: Easy

$$h(s_2) = d(s_2, G_1) + h d(s_2, G_2)$$
 $h(s_3) = d(s_3, G_1) + d(s_3, G_2)$
 $h(s_3) = d(s_3, G_1) + d(s_3, G_2)$
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 $h(s_3) = d(s_3, G_1) + d(s_3, G_2)$

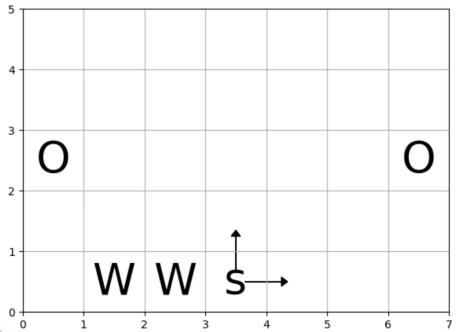


$$-h(s_3) = 8-6=2$$
 > $C(s_2, s_3)=1$. \Rightarrow Not consistent

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Dominate relation

- 1. Zero heuristic: Admissible, Consistent
- 2. Goal-counting heuristic: Admissible, Consistent
- 3. Manhattan Distance to Closest Goal heuristic: Admissible, Consistent
- 4. Manhattan Distance to Farthest Goal heuristic: Admissible, Consistent
- 5. Sum of manhattan distances of all goals Not admissible, Not consistent





h(goal counting) > h(zero) h(closest) > h(zero) h(furthest) > h(zero)

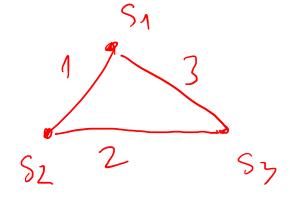
h(furthest) > h(closest) h(furthest) > h(goal counting)

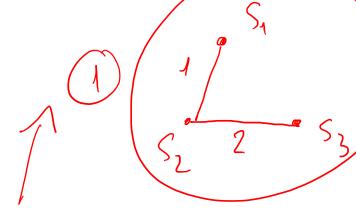
6. Minimum spanning tree

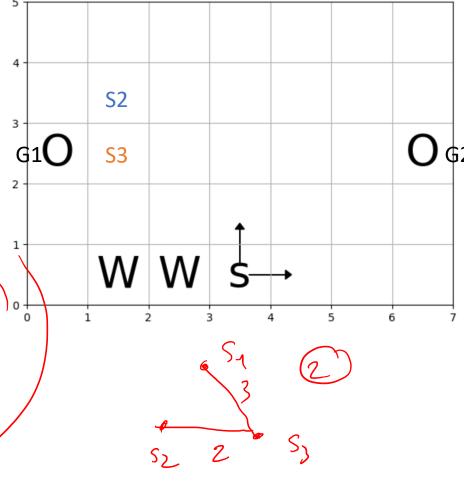
Admissible: Yes Consistent: No

Time to calculate h: Medium

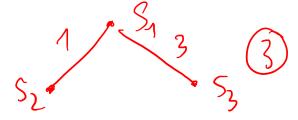
Minimum spanning tree: Only select a subset of edges to connect all vertices and the total cost is minimum







minimum total cost

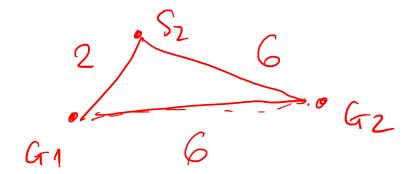


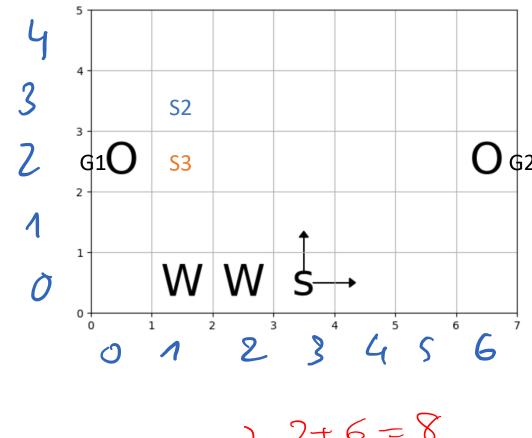
6. Minimum spanning tree

Admissible: Yes Consistent: No

Time to calculate h: Medium

Minimum spanning tree: Only select a subset of edges to connect all vertices and the total cost is minimum





$$\frac{2}{6} = \frac{2}{6} = 8$$

$$\frac{2}{6} = \frac{6}{6} = \frac{8}{6}$$

$$\frac{1}{6} = \frac{8}{6} = \frac{8}{6}$$