

Week 4: STRIPS and Heuristic

COMP90054 – AI Planning for Autonomy

Key concepts

- STRIPS problem
- Heuristic functions

Problem 1

Consider a $m \times m$ manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible coordinates (walls) W

a state = <current coordinate, a set of remaining coordinates>

Initial state $s_0 = \langle (0, 0), G' \setminus \{(0, 0)\} \rangle$

Goal state $S_G = \{ \langle (x, y), \{\} \rangle \mid x, y \in \{0, \dots, m-1\} \}$

State $S = \{ \langle (x, y), V' \rangle \mid x, y \in \{0, \dots, m-1\} \wedge V' \subseteq G' \}$

Action $A(\langle (x, y), V' \rangle) = \{ (dx, dy) \mid dx, dy \in \{-1, 0, 1\} \wedge |dx| + |dy| = 1 \wedge x + dx, y + dy \in \{0, \dots, m-1\} \wedge (x + dx, y + dy) \notin W \}$

Transition $f(\langle (x, y), V' \rangle, (dx, dy)) = \langle (x + dx, y + dy), V' \setminus \{(x + dx, y + dy)\} \rangle$

Cost $c(a) = 1$

State-space model

$P = \langle S, s_0, S_G, A, T, c \rangle$

S = State space

s_0 = initial state

S_G = goal states

A = actions

T = transition functions

c = costs

Problem 1

Consider a $m \times m$ manhattan grid, and a set of coordinates G' to visit in any order, and a set of inaccessible coordinates (walls) W

$$I = \{at(0, 0), visited(0, 0)\}$$

$$G = \{visited(x, y) | x, y \in G'\}$$

$$F = \{at(x, y), visited(x, y) | x, y \in \{0, \dots, m - 1\}\}$$

$$O = \{move(x, y, x', y') : \begin{aligned} &\bullet \text{ Prec: } at(x, y) \\ &\bullet \text{ Add: } at(x', y'), visited(x', y') \\ &\bullet \text{ Del: } at(x, y) / \text{ for each adjacent } (x, y), (x', y'), \text{ and } (x', y') \notin W \end{aligned}\}$$

■ A **problem** in **STRIPS** is a tuple $P = \langle F, O, I, G \rangle$:

■ F stands for set of all **atoms** (boolean vars)

■ O stands for set of all **operators** (actions)

■ $I \subseteq F$ stands for **initial situation**

■ $G \subseteq F$ stands for **goal situation**

■ Operators $o \in O$ **represented** by

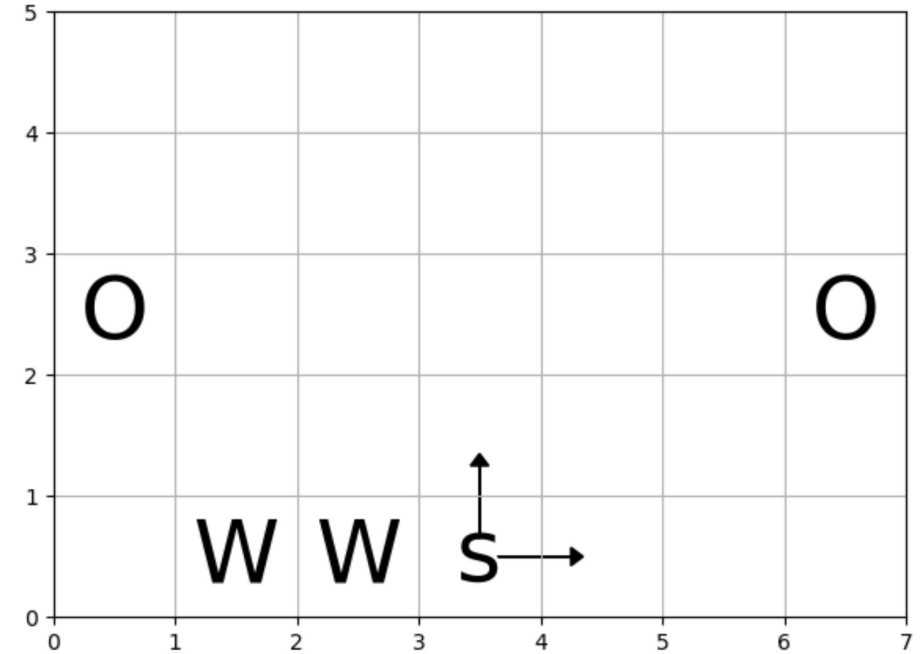
■ the **Add** list $Add(o) \subseteq F$

■ the **Delete** list $Del(o) \subseteq F$

■ the **Precondition** list $Pre(o) \subseteq F$

Problem 2

1. Zero heuristic
2. Goal-counting heuristic
3. Manhattan Distance to Closest Goal heuristic
4. Manhattan Distance to Furthest Goal heuristic
5. Sum of manhattan distances of all goals
6. Minimum spanning tree



Number of node expansion + Calculation time of the heuristic function = Total running time

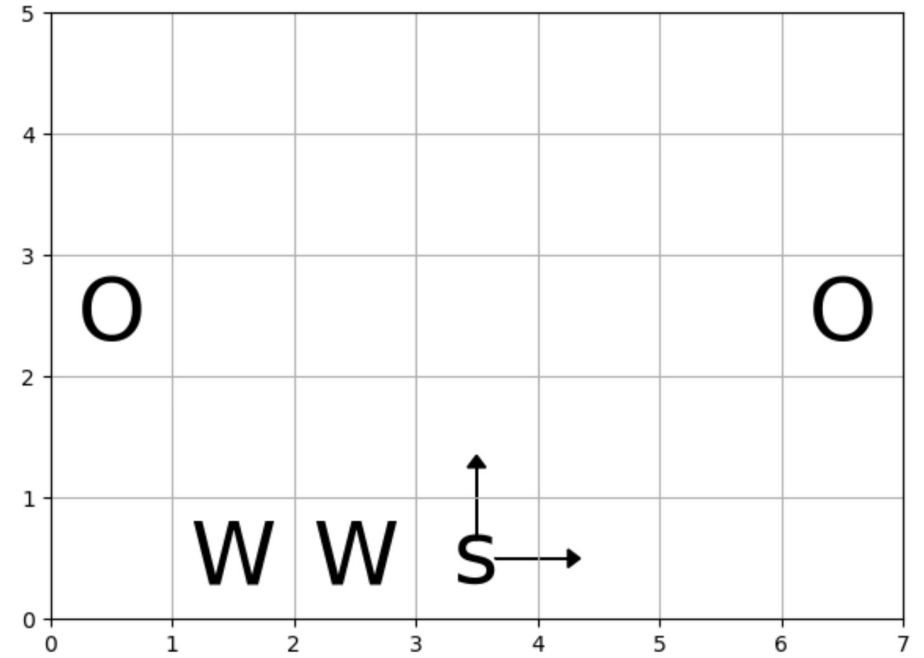
Problem 2

1. Zero heuristic $h = 0$

Admissible: Yes

Consistent: Yes

Time to calculate h : None



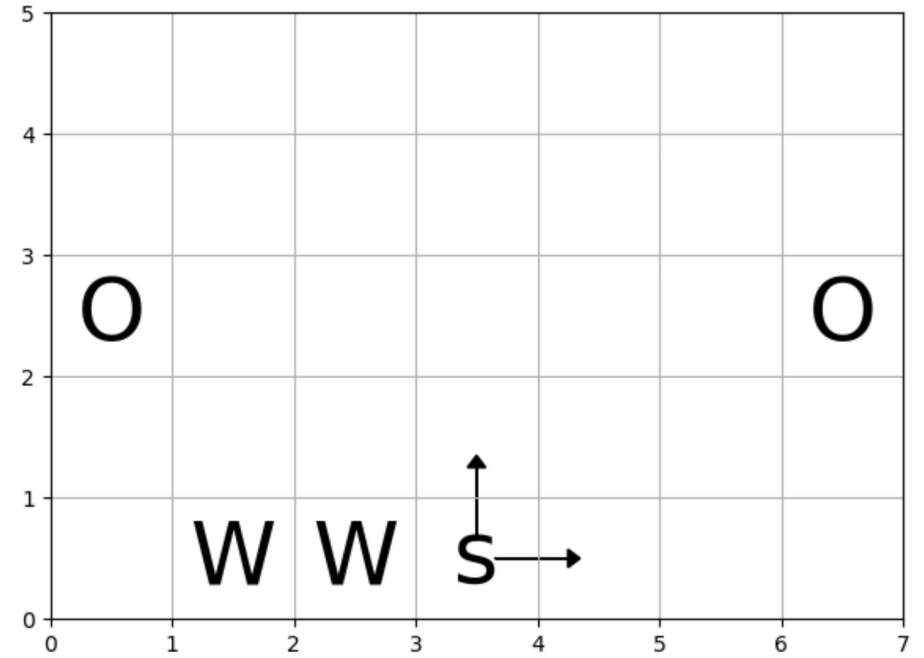
Problem 2

2. Goal-counting heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy



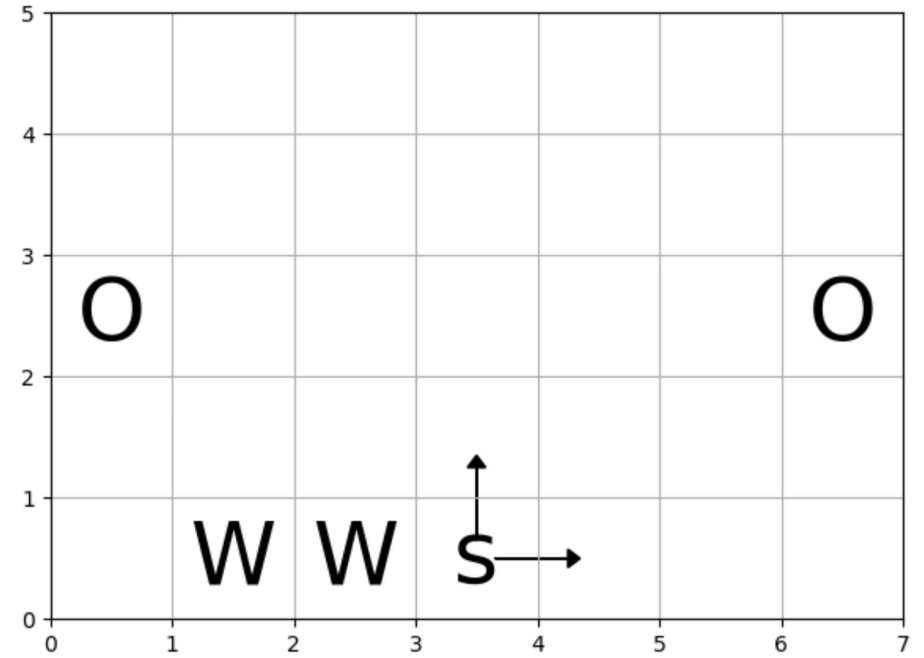
Problem 2

3. Manhattan Distance to Closest Goal heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy



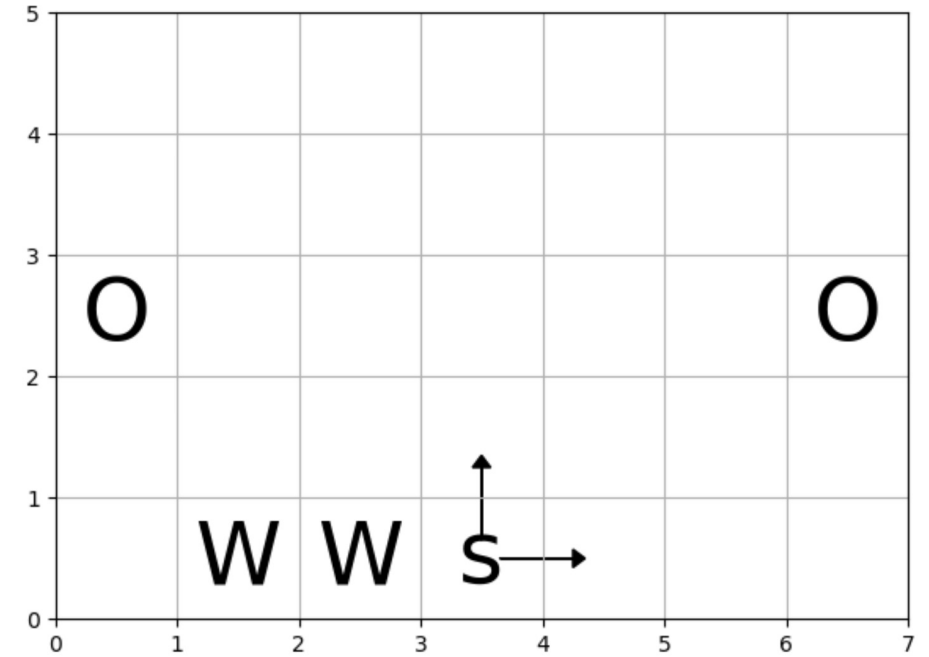
Problem 2

4. Manhattan Distance to Furthest Goal heuristic

Admissible: Yes

Consistent: Yes

Time to calculate h: Easy



Problem 2

5. Sum of Manhattan distances of all goals

Admissible: No

Consistent: No

Time to calculate h: Easy

$$h(s1) = 3 + 7 = 10$$

$$h^*(s1) = 3 + 6 = 9$$

$h(s1) > h^*(s1) \Rightarrow$ Not admissible

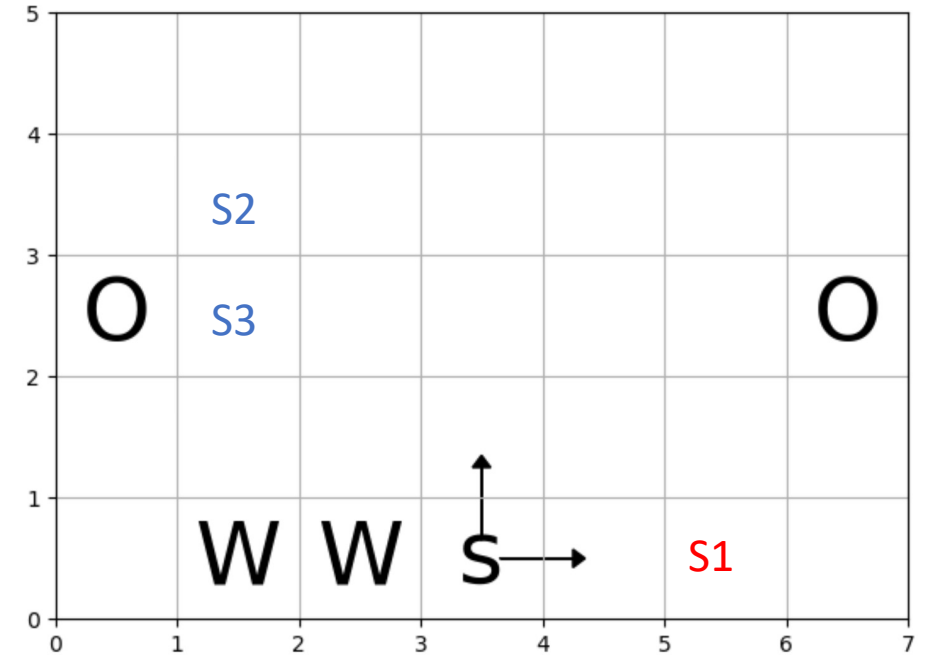
$$h(s2) = 2 + 6 = 8$$

$$h(s3) = 1 + 5 = 6$$

$$h(s2) - h(s3) = 2$$

$$c(s2, s3) = 1$$

$\Rightarrow h(s2) - h(s3) > c(s2, s3) \Rightarrow$ Not consistent



Problem 2

Dominate relation

1. Zero heuristic: Admissible, Consistent
2. Goal-counting heuristic: Admissible, Consistent
3. Manhattan Distance to Closest Goal heuristic: Admissible, Consistent
4. Manhattan Distance to Furthest Goal heuristic: Admissible, Consistent
5. Sum of manhattan distances of all goals: Not admissible, Not consistent

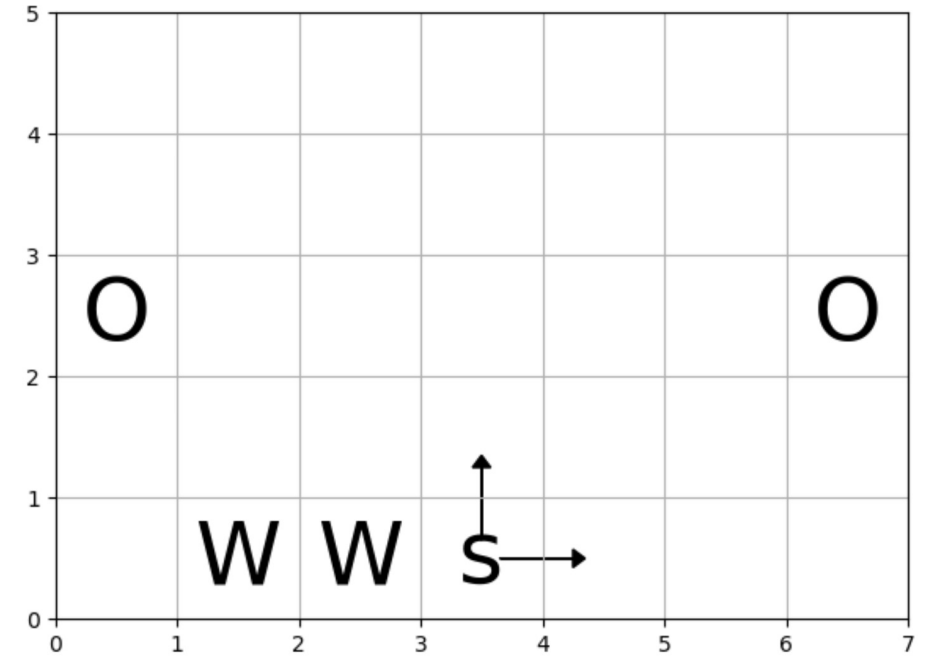
$h(\text{goal counting}) > h(\text{zero})$

$h(\text{closest}) > h(\text{zero})$

$h(\text{furthest}) > h(\text{zero})$

$h(\text{furthest}) > h(\text{closest})$

$h(\text{furthest}) > h(\text{goal counting})$



Problem 2

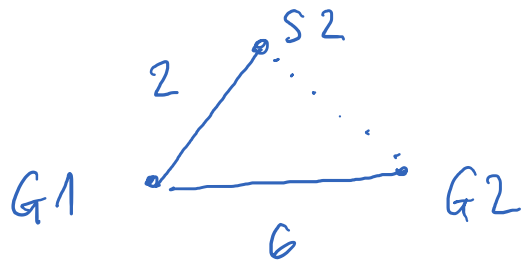
6. Minimum spanning tree

Admissible: Yes

Consistent: No

Time to calculate h: Medium

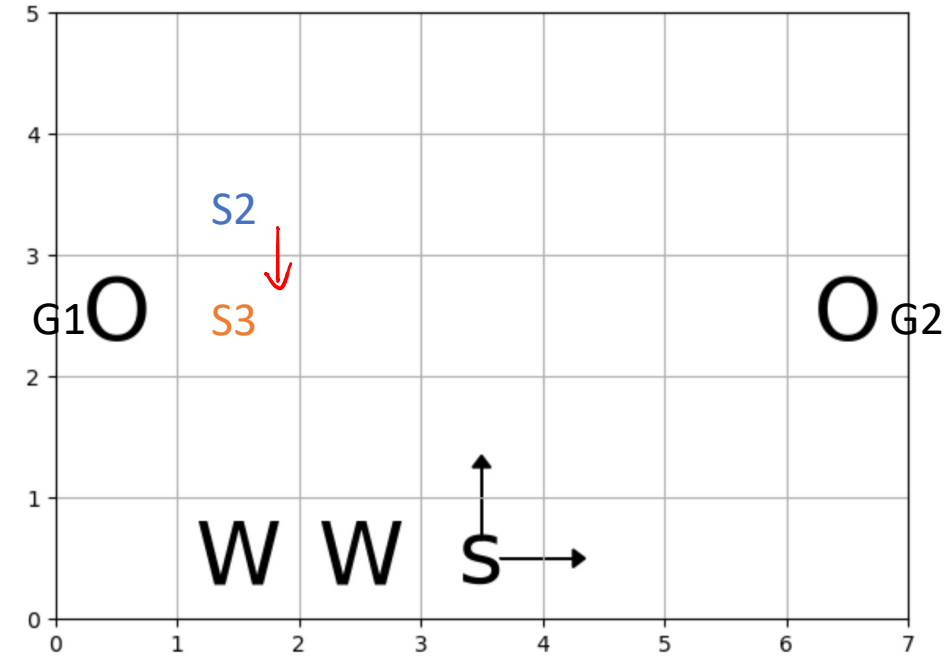
Minimum spanning tree: Only select a subset of edges to connect all vertices and the total cost is minimum



$$h(s_2) = 2 + 6 = 8$$



$$h(s_3) = 1 + 5 = 6$$



$$\Rightarrow h(s_2) - h(s_3) = 8 - 6 = 2$$
$$> c(s_2, s_3) = 1$$

\Rightarrow Not consistent