A Cunning Plan

The Lambda Calculus, Lisp, Scheme, and R

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Outline

Disclaimer

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Motivation

What is this?

```
Y(\(r) \(n)
    `zerop?`(n)(C0)
    (`IF`(eq(n)(C1))
        (C1)
        (\(x) add (r( sub(n)(C1) )) (r( sub(n)(C2) )) (x))))
```

It’s R!

```r
Y(\(r) \(n)
    `zerop?`(n)(C0)
    (`IF`(eq(n)(C1))
        (C1)
        (\(x) add (r( sub(n)(C1) )) (r( sub(n)(C2) )) (x)))) (C10) |> .to.integer ()
```

```
[1] 55
```

Fibonacci: R

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    (`IF`(eq(n)(C1))
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lapply(list(C0, C1, C2, C3, C4, C5, C6, C7, C8, C9, C10),
       \(n) cat(sprintf("%+2s: %+3s\n", n |> .to.integer(), fibonacci(n) |> .to.integer())))
           -> dev.null
```

```
 0:   0
 1:   1
 2:   1
 3:   2
 4:   3
 5:   5
 6:   8
 7:  13
 8:  21
 9:  34
10:  55
```

Fibonacci: Lisp

```lisp
(defun fibonacci (n)
    (if (or (= n 0) (= n 1)) n
        (+ (fibonacci (- n 1)) (fibonacci (- n 2)))))

(let ((x '(0 1 2 3 4 5 6 7 8 9 10)))
(mapcar (lambda (x) (princ (format "%2s:%+3s\n" x (fibonacci x)))) x))
```

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- "no side effects..."

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- "use functions, instead of objects..."
- "no side effects..."
- "mumble, mumble..."

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  - recursion (no “loops”)
- scalability

R: A Language for Data Analysis and Graphics

Ross IHAKA and Robert GENTLEMAN

In this article we discuss our experience designing and implementing a statistical computing language. In developing this new language, we sought to combine what we felt were useful features from two existing computer languages. We feel that the new language provides advantages in the areas of portability, computational efficiency, memory management, and scoping.

Key Words: Computer language; Statistical computing.

1. INTRODUCTION

This article discusses some issues involved in the design and implementation of a computer language for statistical data analysis. Our experience with these issues occurred while developing such a language. The work has been heavily influenced by two existing languages—Becker, Chambers, and Wilks' S (1985) and Steel and Sussman's Scheme (1975). We felt that there were strong points in each of these languages and that it would be interesting to see if the strengths could be combined. The resulting language is very similar in appearance to S, but the underlying implementation and semantics are derived from Scheme. In fact, we implemented the language by first writing an interpreter for a Scheme subset and then progressively mutating it to resemble S.

Source: [Ihaka and Gentleman, 1996]

Scheme

MASSACHUSETTS INSTITUTE OF TECHNOLOGY
ARTIFICIAL INTELLIGENCE LABORATORY

AI Memo No. 349 December 1975

SCHEME

AN INTERPRETER FOR EXTENDED LAMBDA CALCULUS

by

Gerald Jay Sussman and Guy Lewis Steele Jr.

Abstract:
Inspired by ACTORS [Greif and Hewitt] [Smith and Hewitt], we have implemented an interpreter for a LISP-like language, SCHEME, based on the lambda calculus [Church], but extended for side effects, multiprocessing, and process synchronization. The purpose of this implementation is tutorial. We wish to:

Source: [Sussman and Steele, 1975]

Functional languages

Functional languages

Haskell

An advanced, purely functional programming language

Functional languages

Haskell

An advanced, purely functional programming language

Scala

A programming language that scales with you: from small scripts to large multiplatform applications.

Lambda Calculus: Theory

Functions

- Mathematics

$$x \mapsto x^2$$

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$$x \mapsto x^2$$

- Haskell

```
(\x -> x^2) 3
```

9

Functions

- Mathematics

$$x \mapsto x^2$$

- Haskell

```haskell
(\x -> x^2) 3
```

9

- R

```r
(\(x) x^2) (3)
```

```
[1] 9
```

Functions

- Mathematics

$$x \mapsto x^2$$

- Haskell

```
(\x -> x^2) 3
```

9

- R

```
(function(x) x^2) (3)
```

[1] 9

Haskell: `Int`

- Mathematics: constrain x

$$f : x \mapsto x^2, \quad \forall x \in \mathbb{Z}$$

Haskell: `Int`

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$$f : x \mapsto x^2, \quad \forall x \in \mathbb{Z}$$

- type is metadata

```
:{
  f :: Int -> Int
  f = \x -> x^2
:}
f 3
```

9

Haskell: `Double`

- Mathematics: constrain x

$$f : x \mapsto x^2, \quad \forall x \in \mathbb{R}$$

Haskell: `Double`

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```haskell
:{
  f :: Double -> Double
  f = \x -> x^2
:}
f 3.0
```

```
9.0
```

Haskell: `Num`

- Mathematics: what if there is no constraint on x?

$$x \mapsto x^2, \quad \forall x \in ???$$

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```
:type \x -> x^2
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```
\x -> x^2 :: Num a => a -> a
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- Haskell infers type

```
:type \x -> x^2
```

`\x -> x^2 :: Num a => a -> a`

- type is metadata

```
:{
  f :: Num a => a -> a
  f x = x^2
:}
f 3
```

9

Haskell: What is `Num`?

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- Num is a *type class*

```
:info Num
```

```
class Num a where
  (+) :: a -> a -> a
  (-) :: a -> a -> a
  (*) :: a -> a -> a
  negate :: a -> a
  abs :: a -> a
  signum :: a -> a
  fromInteger :: Integer -> a
  {-# MINIMAL (+), (*), abs, signum, fromInteger, (negate | (-)) #-}
-- Defined in ‘GHC.Num’
instance Num Word -- Defined in ‘GHC.Num’
instance Num Integer -- Defined in ‘GHC.Num’
instance Num Int -- Defined in ‘GHC.Num’
instance Num Float -- Defined in ‘GHC.Float’
instance Num Double -- Defined in ‘GHC.Float’
```

Haskell: What is `Int`?

- `Int` is a *data constructor*: it is an *instance* of type class `Num`

```
:info Int
```

```
data Int = GHC.Types.I# GHC.Prim.Int#  -- Defined in ‘GHC.Types’
instance Eq Int -- Defined in ‘GHC.Classes’
instance Ord Int -- Defined in ‘GHC.Classes’
instance Enum Int -- Defined in ‘GHC.Enum’
instance Num Int -- Defined in ‘GHC.Num’
instance Real Int -- Defined in ‘GHC.Real’
instance Show Int -- Defined in ‘GHC.Show’
instance Integral Int -- Defined in ‘GHC.Real’
instance Bounded Int -- Defined in ‘GHC.Enum’
instance Read Int -- Defined in ‘GHC.Read’
```

Haskell: What is Double?

- Double is a *data constructor*: it is an *instance* of type class Num

```
:info Double
```

```
data Double = GHC.Types.D# GHC.Prim.Double#
-- Defined in ‘GHC.Types’
instance Eq Double -- Defined in ‘GHC.Classes’
instance Ord Double -- Defined in ‘GHC.Classes’
instance Enum Double -- Defined in ‘GHC.Float’
instance Floating Double -- Defined in ‘GHC.Float’
instance Fractional Double -- Defined in ‘GHC.Float’
instance Num Double -- Defined in ‘GHC.Float’
instance Real Double -- Defined in ‘GHC.Float’
instance RealFloat Double -- Defined in ‘GHC.Float’
instance RealFrac Double -- Defined in ‘GHC.Float’
instance Show Double -- Defined in ‘GHC.Float’
instance Read Double -- Defined in ‘GHC.Read’
```

Back to functions...

- Mathematics

$$f : x \mapsto x^2$$

- R

```
f<- \(x) x^2
f(3)
```

```
[1] 9
```

What if we *only* had anonymous functions? (no *function names*)

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```r
(\(x) x^2) (3)
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\(x) x
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$$\underbrace{1}_{\text{reduced}}$$

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  1. application: substitute the input λ expression for all instances of bound variables within the body of the abstraction
  2. eliminate the head of the abstraction (its only purpose is to bind a variable)

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$$(\lambda z.z)(\lambda y.y) \quad \overset{\alpha}{=} \quad \underbrace{(\lambda x.x)}_{\alpha \text{ equivalence}} \ (\lambda y.y)$$

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$$
\begin{aligned}
(\lambda z.z)(\lambda y.y) &\overset{\alpha}{=} \underbrace{(\lambda x.x)}_{\alpha \text{ equivalence}} (\lambda y.y) \\
&= \lambda[x := \lambda y.y].x
\end{aligned}
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&\overset{\beta}{\rightarrow} (\lambda y.y)\, z \\
&= \lambda[y := z].y \\
&\overset{\beta}{\rightarrow} z
\end{aligned}$$

Free variables

- here, z is a *free variable*
  - $\big((\lambda x.x)(\lambda y.y)\big)\,z$
  - β reduce the λ expression:

$$
\begin{aligned}
\big((\lambda x.x)(\lambda y.y)\big)\,z &= \big(\lambda[x := \lambda y.y].x\big)\,z \\
&\overset{\beta}{\rightarrow} (\lambda y.y)\,z \\
&= \lambda[y := z].y \\
&\overset{\beta}{\rightarrow} z
\end{aligned}
$$

- here, y is a *free variable*

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\big((\lambda x.x)(\lambda y.y)\big)\, z &= \big(\lambda[x := \lambda y.y].x\big)\, z \\
&\overset{\beta}{\rightarrow} (\lambda y.y)\, z \\
&= \lambda[y := z].y \\
&\overset{\beta}{\rightarrow} z
\end{aligned}
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- here, y is a *free variable*
  - $\lambda x.xy$

Combinators

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    - *not* a combinator (y is a free variable: it does not occur in the head):

$$\lambda x.y$$

Combinators in use: select *left*

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$$(\lambda xy.x)(ab) \quad = \quad (\lambda x.\lambda y.x)(ab)$$

Combinators in use: select *left*

$$
\begin{aligned}
(\lambda xy.x)(ab) &= (\lambda x.\lambda y.x)(ab) \\
&= \big((\lambda x.\lambda y.x)(a)\big)(b)
\end{aligned}
$$

Combinators in use: select *left*

$$
\begin{aligned}
(\lambda xy.x)(ab) &= (\lambda x.\lambda y.x)(ab) \\
&= \big((\lambda x.\lambda y.x)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.x)(b)
\end{aligned}
$$

Combinators in use: select *left*

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(\lambda xy.x)(ab) &= (\lambda x.\lambda y.x)(ab) \\
&= \big((\lambda x.\lambda y.x)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.x)(b) \\
&\overset{\beta}{\rightarrow} (\lambda y.a)(b)
\end{aligned}$$

Combinators in use: select *left*

$$
\begin{aligned}
(\lambda xy.x)(ab) &= (\lambda x.\lambda y.x)(ab) \\
&= \big((\lambda x.\lambda y.x)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.x)(b) \\
&\overset{\beta}{\to} (\lambda y.a)(b) \\
&= \lambda[y := b].a
\end{aligned}
$$

Combinators in use: select *left*

$$
\begin{aligned}
(\lambda xy.x)(ab) &= (\lambda x.\lambda y.x)(ab) \\
&= \big((\lambda x.\lambda y.x)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.x)(b) \\
&\overset{\beta}{\rightarrow} (\lambda y.a)(b) \\
&= \lambda[y := b].a \\
&\overset{\beta}{\rightarrow} a
\end{aligned}
$$

Combinators in use: select *right*

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$$(\lambda xy.y)(ab) \quad = \quad (\lambda x.\lambda y.y)(ab)$$

Combinators in use: select *right*

$$\begin{aligned}(\lambda xy.y)(ab) &= (\lambda x.\lambda y.y)(ab)\\ &= \big((\lambda x.\lambda y.y)(a)\big)(b)\end{aligned}$$

Combinators in use: select *right*

$$
\begin{aligned}
(\lambda xy.y)(ab) \quad &= \quad (\lambda x.\lambda y.y)(ab) \\
&= \quad \big((\lambda x.\lambda y.y)(a)\big)(b) \\
&= \quad (\lambda[x := a].\lambda y.y)(b)
\end{aligned}
$$

Combinators in use: select *right*

$$
\begin{aligned}
(\lambda xy.y)(ab) &= (\lambda x.\lambda y.y)(ab) \\
&= \big((\lambda x.\lambda y.y)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.y)(b) \\
&\xrightarrow{\beta} (\lambda y.y)(b)
\end{aligned}
$$

Combinators in use: select *right*

$$
\begin{aligned}
(\lambda xy.y)(ab) &= (\lambda x.\lambda y.y)(ab) \\
&= \big((\lambda x.\lambda y.y)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.y)(b) \\
&\overset{\beta}{\rightarrow} (\lambda y.y)(b) \\
&= \lambda[y := b].b
\end{aligned}
$$

Combinators in use: select *right*

$$
\begin{aligned}
(\lambda xy.y)(ab) &= (\lambda x.\lambda y.y)(ab) \\
&= \big((\lambda x.\lambda y.y)(a)\big)(b) \\
&= (\lambda[x := a].\lambda y.y)(b) \\
&\overset{\beta}{\to} (\lambda y.y)(b) \\
&= \lambda[y := b].b \\
&\overset{\beta}{\to} b
\end{aligned}
$$

Divergence: non-terminating λ expressions

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- consider

$$(\lambda x.xx)(\lambda x.xx)$$

Divergence: non-terminating λ expressions

- consider

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$$\lambda[x := (\lambda x.xx)].xx$$

Divergence: non-terminating λ expressions

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Divergence: non-terminating λ expressions

- consider

$$(\lambda x.xx)(\lambda x.xx)$$

$$\lambda[x := (\lambda x.xx)].xx$$

$$(\lambda x.xx)(\lambda x.xx)$$

- some λ expressions clearly disallow reduction to normal form: such λ expressions *diverge*

Divergence: non-terminating λ expressions

- consider

$$(\lambda x.xx)(\lambda x.xx)$$

$$\lambda[x := (\lambda x.xx)].xx$$

$$(\lambda x.xx)(\lambda x.xx)$$

- some λ expressions clearly disallow reduction to normal form: such λ expressions *diverge*
- the λ expression $(\lambda x.xx)(\lambda x.xx)$ is called "omega"

The BNF form summarizes the grammar for the Lambda Calculus:

$$
\begin{aligned}
\langle \lambda\ \textit{expression} \rangle \quad & ::= \ <\text{variable}> \\
& \quad\, | \ \ <\text{application}> \\
& \quad\, | \ \ <\text{abstraction}>
\end{aligned}
$$

$$\langle \textit{application} \rangle \quad ::= \ (<\lambda\ \text{expression}>)\ <\lambda\ \text{expression}>$$

$$\langle \textit{abstraction} \rangle \quad ::= \ \lambda\ <\text{variable}>.\ <\lambda\ \text{expression}>$$

Source: [Révész, 1988]

Lambda Calculus: Implementation

R package: `al`

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Title: An implementation of the Lambda Calculus in R

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`https://github.com/tharte/al`

Title: An implementation of the Lambda Calculus in R

Description: R is based on S and Scheme: syntactically, R has the appearance of S, but its implementation is similar to Scheme. Scheme is "An Interpreter for Extended Lambda Calculus". Insight into the Lambda Calculus from within R therefore gives us a fresh understanding of R's roots in functional programming. The Lambda Calculus was created by the American logician Alonzo Church; `al` is named in his honor.

R package: `al`

Source: https://en.wikipedia.org/wiki/Alonzo_Church

Booleans & branching

Booleans & branching

- Truth and Falsity

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$$\texttt{true} \stackrel{\text{def}}{=} T \stackrel{\alpha}{=} \lambda x.\lambda y.x$$

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$$\texttt{true} \stackrel{\text{def}}{=} T \stackrel{\alpha}{=} \lambda x.\lambda y.x$$

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Booleans & branching

- Truth and Falsity

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- if-then-else (the Branching Combinator)

$$\texttt{IF} \stackrel{\text{def}}{=} p(ab)$$

$$p(ab) \stackrel{\beta}{\to} T(ab) \stackrel{\alpha}{=} (\lambda x.\lambda y.x)(ab) \stackrel{\beta}{\to} a$$

Booleans & branching

- Truth and Falsity

$$\texttt{true} \stackrel{\text{def}}{=} T \stackrel{\alpha}{=} \lambda x.\lambda y.x$$

$$\texttt{false} \stackrel{\text{def}}{=} F \stackrel{\alpha}{=} \lambda x.\lambda y.y$$

- if-then-else (the Branching Combinator)

$$\texttt{IF} \stackrel{\text{def}}{=} p(ab)$$

$$p(ab) \xrightarrow{\beta} T(ab) \stackrel{\alpha}{=} (\lambda x.\lambda y.x)(ab) \xrightarrow{\beta} a$$

$$p(ab) \xrightarrow{\beta} F(ab) \stackrel{\alpha}{=} (\lambda x.\lambda y.y)(ab) \xrightarrow{\beta} b$$

Verify that branching works

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```
IF(true)(\(a) a)(\(b) b)
```

Verify that branching works

```
IF(true)(\(a) a)(\(b) b)
```

\(a) a

Verify that branching works

```
IF(true)(\(a) a)(\(b) b)
```

```
\(a) a
```

```
IF(false)(\(a) a)(\(b) b)
```

Verify that branching works

```
IF(true)(\(a) a)(\(b) b)
```

```
\(a) a
```

```
IF(false)(\(a) a)(\(b) b)
```

```
\(b) b
```

Switching between worlds: `.to.logical`

Switching between worlds: `.to.logical`

```
.to.logical
```

```
function (b)
(IF(b)(TRUE))(FALSE)
<bytecode: 0x55fc452616e0>
<environment: namespace:al>
```

Switching between worlds: `.to.logical`

```
.to.logical
```

```
function (b)
(IF(b)(TRUE))(FALSE)
<bytecode: 0x55fc452616e0>
<environment: namespace:al>
```

```
true |> .to.logical()
```

Switching between worlds: `.to.logical`

```r
.to.logical
```

```
function (b)
(IF(b)(TRUE))(FALSE)
<bytecode: 0x55fc452616e0>
<environment: namespace:al>
```

```r
true |> .to.logical()
```

```
[1] TRUE
```

Switching between worlds: `.to.logical`

```r
.to.logical
```

```
function (b)
(IF(b)(TRUE))(FALSE)
<bytecode: 0x55fc452616e0>
<environment: namespace:al>
```

```r
true |> .to.logical()
```

```
[1] TRUE
```

```r
false |> .to.logical()
```

Switching between worlds: `.to.logical`

```
.to.logical
```

```
function (b)
(IF(b)(TRUE))(FALSE)
<bytecode: 0x55fc452616e0>
<environment: namespace:al>
```

```
true |> .to.logical()
```

```
[1] TRUE
```

```
false |> .to.logical()
```

```
[1] FALSE
```

Logic: And

Logic: And

$$\texttt{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

Logic: And

$$\texttt{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

$$xyF \xrightarrow{\beta} (FT)F \xrightarrow{\beta} F$$

Logic: And

$$\texttt{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

$$xyF \xrightarrow{\beta} (FT)F \xrightarrow{\beta} F$$

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Logic: And

$$\texttt{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

$$xyF \xrightarrow{\beta} (FT)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (FF)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (TF)F \xrightarrow{\beta} F$$

Logic: And

$$\texttt{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

$$xyF \xrightarrow{\beta} (FT)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (FF)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (TF)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (TT)F \xrightarrow{\beta} T$$

Logic: And

$$\text{and} \stackrel{\text{def}}{=} \lambda x.\lambda y.xyF$$

$$xyF \xrightarrow{\beta} (FT)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (FF)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (TF)F \xrightarrow{\beta} F$$
$$xyF \xrightarrow{\beta} (TT)F \xrightarrow{\beta} T$$

```
and(false)(true) |> .to.logical()
and(false)(false) |> .to.logical()
and(true)(false) |> .to.logical()
and(true)(true) |> .to.logical()
```

```
[1] FALSE
[1] FALSE
[1] FALSE
[1] TRUE
```

Logic: Or

Logic: Or

$$\texttt{or} \stackrel{\text{def}}{=} \lambda x.\lambda y.xTy$$

Logic: Or

$$\texttt{or} \stackrel{\text{def}}{=} \lambda x.\lambda y.xTy$$

$$xTy \xrightarrow{\beta} (F)T(T) \xrightarrow{\beta} T$$

Logic: Or

$$\texttt{or} \stackrel{\text{def}}{=} \lambda x.\lambda y.xTy$$

$$xTy \xrightarrow{\beta} (F)T(T) \xrightarrow{\beta} T$$

$$xTy \xrightarrow{\beta} (F)T(F) \xrightarrow{\beta} F$$

Logic: Or

$$\texttt{or} \stackrel{\text{def}}{=} \lambda x.\lambda y.xTy$$

$$xTy \xrightarrow{\beta} (F)T(T) \xrightarrow{\beta} T$$
$$xTy \xrightarrow{\beta} (F)T(F) \xrightarrow{\beta} F$$
$$xTy \xrightarrow{\beta} (T)T(F) \xrightarrow{\beta} T$$

Logic: Or

$$\texttt{or} \stackrel{\text{def}}{=} \lambda x.\lambda y.xTy$$

$$xTy \xrightarrow{\beta} (F)T(T) \xrightarrow{\beta} T$$
$$xTy \xrightarrow{\beta} (F)T(F) \xrightarrow{\beta} F$$
$$xTy \xrightarrow{\beta} (T)T(F) \xrightarrow{\beta} T$$
$$xTy \xrightarrow{\beta} (T)T(T) \xrightarrow{\beta} T$$

Logic: Or

$$\texttt{or} \overset{\text{def}}{=} \lambda x.\lambda y.xTy$$

$$xTy \xrightarrow{\beta} (F)T(T) \xrightarrow{\beta} T$$

$$xTy \xrightarrow{\beta} (F)T(F) \xrightarrow{\beta} F$$

$$xTy \xrightarrow{\beta} (T)T(F) \xrightarrow{\beta} T$$

$$xTy \xrightarrow{\beta} (T)T(T) \xrightarrow{\beta} T$$

```
or(false)(true) |> .to.logical()
or(false)(false) |> .to.logical()
or(true)(false) |> .to.logical()
or(true)(true) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
[1] TRUE
[1] TRUE
```

Logic: Not

Logic: Not

$$\texttt{not} \stackrel{\text{def}}{=} pFT$$

Logic: Not

$$\texttt{not} \stackrel{\text{def}}{=} pFT$$

$$pFT \xrightarrow{\beta} (F)FT \xrightarrow{\beta} T$$

Logic: Not

$$\texttt{not} \stackrel{\text{def}}{=} pFT$$

$$pFT \xrightarrow{\beta} (F)FT \xrightarrow{\beta} T$$

$$pFT \xrightarrow{\beta} (T)FT \xrightarrow{\beta} F$$

Logic: Not

$$\texttt{not} \stackrel{\text{def}}{=} pFT$$

$$pFT \xrightarrow{\beta} (F)FT \xrightarrow{\beta} T$$

$$pFT \xrightarrow{\beta} (T)FT \xrightarrow{\beta} F$$

```
not(false) |> .to.logical()
not(true) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
```

Church Numerals

$$\texttt{C0} \quad \stackrel{\text{def}}{=} \quad C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false}$$

Church Numerals

$$\texttt{C0} \stackrel{\text{def}}{=} C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false}$$

$$\texttt{C1} \stackrel{\text{def}}{=} C_1 \stackrel{\alpha}{=} \lambda f.\lambda x.f(x)$$

Church Numerals

$$\texttt{C0} \stackrel{\text{def}}{=} C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false}$$

$$\texttt{C1} \stackrel{\text{def}}{=} C_1 \stackrel{\alpha}{=} \lambda f.\lambda x.f(x)$$

$$\texttt{C2} \stackrel{\text{def}}{=} C_2 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(x))$$

Church Numerals

$$\texttt{C0} \stackrel{\text{def}}{=} C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false}$$

$$\texttt{C1} \stackrel{\text{def}}{=} C_1 \stackrel{\alpha}{=} \lambda f.\lambda x.f(x)$$

$$\texttt{C2} \stackrel{\text{def}}{=} C_2 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(x))$$

$$\texttt{C3} \stackrel{\text{def}}{=} C_3 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(f(x)))$$

Church Numerals

$$
\begin{aligned}
\texttt{C0} &\stackrel{\text{def}}{=} C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false} \\
\texttt{C1} &\stackrel{\text{def}}{=} C_1 \stackrel{\alpha}{=} \lambda f.\lambda x.f(x) \\
\texttt{C2} &\stackrel{\text{def}}{=} C_2 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(x)) \\
\texttt{C3} &\stackrel{\text{def}}{=} C_3 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(f(x))) \\
&\vdots
\end{aligned}
$$

Church Numerals

$$
\begin{aligned}
\texttt{C0} &\stackrel{\text{def}}{=} C_0 \stackrel{\alpha}{=} \lambda f.\lambda x.x \stackrel{\alpha}{=} \lambda x.\lambda y.y = \texttt{false} \\
\texttt{C1} &\stackrel{\text{def}}{=} C_1 \stackrel{\alpha}{=} \lambda f.\lambda x.f(x) \\
\texttt{C2} &\stackrel{\text{def}}{=} C_2 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(x)) \\
\texttt{C3} &\stackrel{\text{def}}{=} C_3 \stackrel{\alpha}{=} \lambda f.\lambda x.f(f(f(x))) \\
&\vdots \\
&\stackrel{\text{def}}{=} C_N \stackrel{\alpha}{=} \lambda f.\lambda x.f^N(x)
\end{aligned}
$$

Switching between worlds: `.to.integer`

Switching between worlds: `.to.integer`

```
.to.integer
```

```
function (n)
n(function(x) x + 1)(0)
<bytecode: 0x55fc442ceef0>
<environment: namespace:al>
```

Switching between worlds: `.to.integer`

```r
.to.integer
```

```
function (n)
n(function(x) x + 1)(0)
<bytecode: 0x55fc442ceef0>
<environment: namespace:al>
```

```r
C0 |> .to.integer()
C1 |> .to.integer()
C2 |> .to.integer()
C3 |> .to.integer()
```

```
[1] 0
[1] 1
[1] 2
[1] 3
```

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \xrightarrow{\beta} 0$$

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \quad \xrightarrow{\beta} \quad 0$$

$$\underbrace{(\lambda f.\lambda x.f(x))}_{n=C_1}(\lambda x.x+1)(0) \quad \xrightarrow{\beta} \quad \underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \xrightarrow{\beta} 0$$

$$\begin{aligned}\underbrace{(\lambda f.\lambda x.f(x))}_{n=C_1}(\lambda x.x+1)(0) &\xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}\\ &\xrightarrow{\beta} 1\end{aligned}$$

Switching between worlds: `.to.integer`

```r
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \xrightarrow{\beta} 0$$

$$\underbrace{(\lambda f.\lambda x.f(x))}_{n=C_1}(\lambda x.x+1)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} 1$$

$$\Big(\underbrace{(\lambda f.\lambda x.f(f(x)))}_{n=C_2}(\lambda x.x+1)\Big)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \xrightarrow{\beta} 0$$

$$\underbrace{(\lambda f.\lambda x.f(x))}_{n=C_1}(\lambda x.x+1)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} 1$$

$$\Big(\underbrace{(\lambda f.\lambda x.f(f(x)))}_{n=C_2}(\lambda x.x+1)\Big)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} \underbrace{(\lambda x.(\lambda x.x+1)+1)}_{f(f(x))}\underbrace{(0)}_{x}$$

Switching between worlds: `.to.integer`

```
`.to.integer`<- \(n) n(\(x) x+1)(0)
```

$$\underbrace{(\lambda f.\lambda x.x)}_{n=C_0}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x} \xrightarrow{\beta} 0$$

$$\underbrace{(\lambda f.\lambda x.f(x))}_{n=C_1}(\lambda x.x+1)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} 1$$

$$\Big(\underbrace{(\lambda f.\lambda x.f(f(x)))}_{n=C_2}(\lambda x.x+1)\Big)(0) \xrightarrow{\beta} \underbrace{(\lambda x.x+1)}_{f}\underbrace{(\lambda x.x+1)}_{f}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} \underbrace{(\lambda x.(\lambda x.x+1)+1)}_{f(f(x))}\underbrace{(0)}_{x}$$

$$\xrightarrow{\beta} 2$$

Arithmetic

Arithmetic

- successor

$$\texttt{succ} \stackrel{\text{def}}{=} \lambda n.\lambda f.\lambda x.f((nf)x)$$

Arithmetic

- successor

$$\texttt{succ} \stackrel{\text{def}}{=} \lambda n.\lambda f.\lambda x.f((nf)x)$$

```
succ(C0) |> .to.integer()
succ(C10) |> .to.integer()
```

```
[1] 1
[1] 11
```

Arithmetic

- successor

$$\texttt{succ} \stackrel{\text{def}}{=} \lambda n.\lambda f.\lambda x.f((nf)x)$$

```
succ(C0) |> .to.integer()
succ(C10) |> .to.integer()
```

```
[1] 1
[1] 11
```

- predecessor

$$\texttt{pred} \quad \stackrel{\text{def}}{=} \quad \lambda n.\lambda n(\lambda p.\lambda z.z(\texttt{succ}(p(T)))(p(T)))(\lambda z.z(C_0)(C_0))(F)$$

Arithmetic

- successor

$$\texttt{succ} \overset{\text{def}}{=} \lambda n.\lambda f.\lambda x.f((nf)x)$$

```
succ(C0) |> .to.integer()
succ(C10) |> .to.integer()
```

```
[1] 1
[1] 11
```

- predecessor

$$\texttt{pred} \overset{\text{def}}{=} \lambda n.\lambda n(\lambda p.\lambda z.z(\texttt{succ}(p(T)))(p(T)))(\lambda z.z(C_0)(C_0))(F)$$

```
pred(C1) |> .to.integer()
pred(C10) |> .to.integer()
```

```
[1] 0
[1] 9
```

Arithmetic

Arithmetic

- addition

$$\texttt{add} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\texttt{succ})n$$

Arithmetic

- addition

$$\texttt{add} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\texttt{succ})n$$

```
add(C0)(C1) |> .to.integer()
add(C1)(C2) |> .to.integer()
```

```
[1] 1
[1] 3
```

Arithmetic

- addition

$$\texttt{add} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\texttt{succ})n$$

```
add(C0)(C1) |> .to.integer()
add(C1)(C2) |> .to.integer()
```

```
[1] 1
[1] 3
```

- subtraction

$$\texttt{sub} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\text{pred})n$$

Arithmetic

- addition

$$\mathtt{add} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\mathtt{succ})n$$

```
add(C0)(C1) |> .to.integer()
add(C1)(C2) |> .to.integer()
```

```
[1] 1
[1] 3
```

- subtraction

$$\mathtt{sub} \stackrel{\text{def}}{=} \lambda n.\lambda m.(m\,\text{pred})n$$

```
sub(C1)(C1) |> .to.integer()
sub(C2)(C1) |> .to.integer()
```

```
[1] 0
[1] 1
```

Arithmetic

Arithmetic

- multiplication

$$\texttt{mul} \stackrel{\text{def}}{=} \lambda n.\lambda m.m(\texttt{add}\,n)\,C_0$$

Arithmetic

- multiplication

$$\texttt{mul} \stackrel{\text{def}}{=} \lambda n.\lambda m.m(\texttt{add}n)C_0$$

```
mul(C0)(C1) |> .to.integer()
mul(C1)(C2) |> .to.integer()
```

```
[1] 0
[1] 2
```

Arithmetic

- multiplication

$$\texttt{mul} \stackrel{\text{def}}{=} \lambda n.\lambda m.m(\texttt{add}n)C_0$$

```
mul(C0)(C1) |> .to.integer()
mul(C1)(C2) |> .to.integer()
```

```
[1] 0
[1] 2
```

- exponentiation

$$\texttt{exp} \stackrel{\text{def}}{=} \lambda n.\lambda m.mn$$

Arithmetic

- multiplication

$$\texttt{mul} \stackrel{\text{def}}{=} \lambda n.\lambda m.m(\texttt{add}n)C_0$$

```
mul(C0)(C1) |> .to.integer()
mul(C1)(C2) |> .to.integer()
```

```
[1] 0
[1] 2
```

- exponentiation

$$\texttt{exp} \stackrel{\text{def}}{=} \lambda n.\lambda m.mn$$

```
exp(C2)(C1) |> .to.integer()
exp(C2)(C2) |> .to.integer()
```

```
[1] 2
[1] 4
```

Comparisons

Comparisons

- is-zero

$$\texttt{zerop?} \stackrel{\text{def}}{=} \lambda n.n(\lambda m.F)\,T$$

Comparisons

- is-zero

$$\texttt{zerop?} \stackrel{\text{def}}{=} \lambda n.n(\lambda m.F)\,T$$

```r
`zerop?`(C0) |> .to.logical()
`zerop?`(C1) |> .to.logical()
`zerop?`(C2) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
[1] FALSE
```

Comparisons

- is-zero

$$\texttt{zerop?} \stackrel{\text{def}}{=} \lambda n.n(\lambda m.F)\,T$$

```
`zerop?`(C0) |> .to.logical()
`zerop?`(C1) |> .to.logical()
`zerop?`(C2) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
[1] FALSE
```

- less-than-or-equal

$$\texttt{le} \stackrel{\text{def}}{=} \lambda n.\lambda m.\texttt{zerop?}(\texttt{sub}(n)(m))$$

Comparisons

- is-zero

$$\texttt{zerop?} \stackrel{\text{def}}{=} \lambda n.n(\lambda m.F)\,T$$

```
`zerop?`(C0) |> .to.logical()
`zerop?`(C1) |> .to.logical()
`zerop?`(C2) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
[1] FALSE
```

- less-than-or-equal

$$\texttt{le} \stackrel{\text{def}}{=} \lambda n.\lambda m.\texttt{zerop?}(\texttt{sub}(n)(m))$$

```
le(C1)(C1) |> .to.logical()
le(C1)(C2) |> .to.logical()
le(C2)(C1) |> .to.logical()
```

```
[1] TRUE
[1] TRUE
[1] FALSE
```

Comparisons

Comparisons

- is-equal

$$\texttt{eq} \stackrel{\text{def}}{=} \lambda n.\lambda m.\texttt{and}(\texttt{le}(n)(m))(\texttt{le}(n)(m))$$

Comparisons

- is-equal

$$\mathtt{eq} \stackrel{\text{def}}{=} \lambda n.\lambda m.\mathtt{and}(\mathtt{le}(n)(m))(\mathtt{le}(n)(m))$$

```
eq(C1)(C1) |> .to.logical()
eq(C1)(C2) |> .to.logical()
eq(C2)(C1) |> .to.logical()
eq(C2)(C2) |> .to.logical()
```

```
[1] TRUE
[1] FALSE
[1] FALSE
[1] TRUE
```

Recursion with the Y-Combinator

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)$$

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\Big(\lambda x.y(xx)\Big)\Big(\lambda x.y(xx)\Big)$$

apply Y to an arbitrary function F

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the Y-combinator: Curry's paradoxical combinator

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apply Y to an arbitrary function F

$$YF \quad = \quad \bigg(\lambda y.\Big(\lambda x.y(xx)\Big)\Big(\lambda x.y(xx)\Big)\bigg)F$$

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the Y-combinator: Curry's paradoxical combinator

$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)$$

apply Y to an arbitrary function F

$$\begin{aligned} YF &= \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\ &= \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \end{aligned}$$

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$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)$$

apply Y to an arbitrary function F

$$\begin{aligned} YF &= \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\ &= \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\ &\stackrel{\beta}{\to} \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \end{aligned}$$

Recursion with the Y-Combinator

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$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\Big(\lambda x.y(xx)\Big)\Big(\lambda x.y(xx)\Big)$$

apply Y to an arbitrary function F

$$\begin{array}{rcl} YF & = & \Big(\lambda y.\Big(\lambda x.y(xx)\Big)\Big(\lambda x.y(xx)\Big)\Big)F \\ & = & \lambda[y := F].\Big(\lambda x.y(xx)\Big)\Big(\lambda x.y(xx)\Big) \\ & \stackrel{\beta}{\to} & \Big(\lambda x.F(xx)\Big)\Big(\lambda x.F(xx)\Big) \ = YF \end{array}$$

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)$$

apply Y to an arbitrary function F

$$\begin{aligned}
YF &= \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\
&= \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\
&\stackrel{\beta}{\to} \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \ \color{blue}{= YF} \\
&= \lambda[x := \lambda x.F(xx)].F(xx)
\end{aligned}$$

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

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apply Y to an arbitrary function F

$$\begin{aligned}
YF &= \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\
&= \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\
&\xrightarrow{\beta} \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \; = YF \\
&= \lambda[x := \lambda x.F(xx)].F(xx) \\
&\xrightarrow{\beta} F\big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big)
\end{aligned}$$

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

$$\mathtt{Y} \stackrel{\text{def}}{=} \lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)$$

apply Y to an arbitrary function F

$$\begin{array}{rcl}
YF & = & \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\
 & = & \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\
 & \stackrel{\beta}{\to} & \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \;\; = YF \\
 & = & \lambda[x := \lambda x.F(xx)].F(xx) \\
 & \stackrel{\beta}{\to} & F\big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \\
 & = & F(YF)
\end{array}$$

Recursion with the Y-Combinator

the Y-combinator: Curry's paradoxical combinator

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apply Y to an arbitrary function F

$$\begin{aligned}
YF &= \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\
&= \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\
&\stackrel{\beta}{\to} \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \;\color{blue}{= YF} \\
&= \lambda[x := \lambda x.F(xx)].F(xx) \\
&\stackrel{\beta}{\to} F\color{blue}{\big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big)} \\
&= F(YF)
\end{aligned}$$

x in $Fx = x$ is called a *fixed point* of F

Recursion with the Y-Combinator

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apply Y to an arbitrary function F

$$\begin{array}{rcl}
YF & = & \Big(\lambda y.\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big)\Big)F \\
 & = & \lambda[y := F].\big(\lambda x.y(xx)\big)\big(\lambda x.y(xx)\big) \\
 & \stackrel{\beta}{\to} & \big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \;\; = YF \\
 & = & \lambda[x := \lambda x.F(xx)].F(xx) \\
 & \stackrel{\beta}{\to} & F\big(\lambda x.F(xx)\big)\big(\lambda x.F(xx)\big) \\
 & = & F(YF)
\end{array}$$

x in $Fx = x$ is called a *fixed point* of F

x ends up back at x when F is applied to x

Infinite loop

Infinite loop

$$YI = I(YI)$$

Infinite loop

$$
\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\Big(Y(\lambda x.x)\Big)
\end{aligned}
$$

Infinite loop

$$
\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\big(Y(\lambda x.x)\big) \\
& \lambda[x := Y(\lambda x.x)].x
\end{aligned}
$$

Infinite loop

$$
\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\big(Y(\lambda x.x)\big) \\
&\ \lambda[x := Y(\lambda x.x)].x \\
&\xrightarrow{\beta} Y(\lambda x.x)
\end{aligned}
$$

Infinite loop

$$\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\big(Y(\lambda x.x)\big) \\
& \quad \lambda[x := Y(\lambda x.x)].x \\
&\xrightarrow{\beta} Y(\lambda x.x) \\
&= YI
\end{aligned}$$

Infinite loop

$$
\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\big(Y(\lambda x.x)\big) \\
& \quad \lambda[x := Y(\lambda x.x)].x \\
&\xrightarrow{\beta} Y(\lambda x.x) \\
&= YI
\end{aligned}
$$

```
(\(y) (\(x) y(x(x))) (\(x) y(x(x)))) (\(x) x)
```

Infinite loop

$$
\begin{aligned}
YI &= I(YI) \\
&= (\lambda x.x)\big(Y(\lambda x.x)\big) \\
&\quad \lambda[x := Y(\lambda x.x)].x \\
&\xrightarrow{\beta} Y(\lambda x.x) \\
&= YI
\end{aligned}
$$

```
(\(y) (\(x) y(x(x))) (\(x) y(x(x)))) (\(x) x)
```

```
Error: C stack usage  9522372 is too close to the limit
```

Arithmetic series

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i$$

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i \quad = \quad \begin{cases} S_0 & = 0 \\ S_n & = n + S_{n-1} \end{cases}$$

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i \quad = \quad \begin{cases} S_0 & = 0 \\ S_n & = n + S_{n-1} \end{cases}$$

$$F \stackrel{\text{def}}{=} \lambda r.\lambda n.\mathtt{IF}(\grave{}\mathtt{zerop?}\grave{}\, n)(C_0)\Big(\mathtt{add}(n)(r(\mathtt{pred}n))\Big)$$

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i \quad = \quad \begin{cases} S_0 & = 0 \\ S_n & = n + S_{n-1} \end{cases}$$

$$F \stackrel{\text{def}}{=} \lambda r.\lambda n.\mathtt{IF}(\text{`}\mathtt{zerop?}\text{`}\, n)(C_0)\Big(\mathtt{add}(n)(r(\mathtt{pred}n))\Big)$$

$$YFC_3 \xrightarrow{\beta} F(YF)C_3$$

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i \quad = \quad \begin{cases} S_0 & = 0 \\ S_n & = n + S_{n-1} \end{cases}$$

$$F \stackrel{\text{def}}{=} \lambda r.\lambda n.\mathtt{IF}(\mathtt{`zerop?`}\, n)(C_0)\Big(\mathtt{add}(n)(r(\mathtt{pred}\, n))\Big)$$

$$YFC_3 \xrightarrow{\beta} F(YF)C_3$$

```
F<- (\(r) \(n) IF(`zerop?`(n)) (C0) (add(n)(r(pred(n)))))

Y(F)(C3) |> .to.integer()
```

Arithmetic series

$$S_n = \sum_{i=0}^{i=n} i \quad = \quad \begin{cases} S_0 & = 0 \\ S_n & = n + S_{n-1} \end{cases}$$

$$F \stackrel{\text{def}}{=} \lambda r.\lambda n.\mathtt{IF}(\texttt{`zerop?`}\, n)(C_0)\Big(\mathtt{add}(n)(r(\mathtt{pred}\, n))\Big)$$

$$YFC_3 \xrightarrow{\beta} F(YF)C_3$$

```
F<- (\(r) \(n) IF(`zerop?`(n)) (C0) (add(n)(r(pred(n)))))

Y(F)(C3) |> .to.integer()
```

[1] 6

Arithmetic series

```
F<- (\(r) \(n) IF(`zerop?`(n)) (C0) (add(n)(r(pred(n)))))

Y(F)(C3) |> .to.integer()
```

```
[1] 6
```

Arithmetic series

```
Y((\(r) \(n) IF(`zerop?`(n)) (C0) (add(n)(r(pred(n))))))(C3) |> .to.integer()
```

```
[1] 6
```

Arithmetic series

```
(\(y) (\(x) y(x(x))) (\(x) y(x(x)))) (
    \(r) \(n) IF(`zerop?`(n)) (C0) (add(n)(r(pred(n))))) (C3) |> .to.integer()
```

```
[1] 6
```

Arithmetic series

```
(\(y) (\(x) y(x(x))) (\(x) y(x(x)))) (
    (\(r) \(n) (
        \(p) \(x) \(y) p(x)(y)) (
            (\(n) n(\(m) \(x) \(y) y)(\(x) \(y) x))(n)) (
                \(f) \(x) x) ((\(n) \(m) m(\(n) \(f) \(x) f(n(f)(x)))(n))(n)(
                    r((\(n) n(\(p) \(z) z((\(n) \(f) \(x) f(n(f)(x)))(p(\(x) \(y) x)))(
                        p(\(x) \(y) x)))(\(z) z(\(f) \(x) x)(\(f) \(x) x))(
                            \(x) \(y) y))(n)))))) (\(f) \(x) f(f(f(x)))) |> .to.integer()
```

```
[1] 6
```

Fibonacci: R

Fibonacci: R

$$F_n = \begin{cases} F_0 & = 0 \\ F_1 & = 1 \\ F_n & = F_{n-1} + F_{n-2} \end{cases}$$

Fibonacci: R

$$F_n = \begin{cases} F_0 & = 0 \\ F_1 & = 1 \\ F_n & = F_{n-1} + F_{n-2} \end{cases}$$

```
Y(\(r) \(n)
    `zerop?`(n)(C0)
    (`IF`(eq(n)(C1))
        (C1)
        (\(x) add (r( sub(n)(C1) )) (r( sub(n)(C2) )) (x)))) -> fibonacci

lapply(list(C0, C1, C2, C3, C4, C5, C6, C7, C8, C9, C10),
       \(n) cat(sprintf("%+2s: %+3s\n", n |> .to.integer(), fibonacci(n) |> .to.integer())))
           -> dev.null
```

```
 0:   0
 1:   1
 2:   1
 3:   2
 4:   3
 5:   5
 6:   8
 7:  13
 8:  21
 9:  34
10:  55
```

References

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