

# Tower of Hanoi: Recursion and Algorithm Design

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# 1 Concept Explanation: Tower of Hanoi Rules and Logic

The Tower of Hanoi is a classic mathematical puzzle consisting of three rods and a set of disks of different sizes. The puzzle begins with all the disks neatly stacked on one rod in ascending order, forming a conical shape. The objective is to move the entire stack to another rod under the following rules:

- Only one disk may be moved at a time.
- Each move must involve taking the top disk from one of the stacks and placing it on another rod.
- No disk may be placed on top of a smaller disk.

The problem lends itself naturally to a recursive solution. For  $n$  disks:

- Move  $n - 1$  disks to the auxiliary rod (recursive subproblem)
- Move the largest disk to the target rod
- Move the  $n - 1$  disks from the auxiliary rod to the target rod (recursive subproblem)

Let  $T(n)$  be the minimum number of moves required. Then:

$$T(n) = 2T(n - 1) + 1$$

**Closed-form:**  $T(n) = 2^n - 1$

## Proof of Closed-Form Using Induction

**Base Case:**  $T(1) = 2^1 - 1 = 1$

**Inductive Hypothesis:** Assume  $T(k) = 2^k - 1$

**Inductive Step:**

$$\begin{aligned} T(k + 1) &= 2T(k) + 1 \\ &= 2(2^k - 1) + 1 \\ &= 2^{k+1} - 2 + 1 \\ &= 2^{k+1} - 1 \end{aligned}$$

**Conclusion:** By the principle of mathematical induction, the closed-form  $T(n) = 2^n - 1$  holds for all  $n \geq 1$ .

## 2 Algorithm and Pseudocode

### Pseudocode:

Listing 1: Recursive Hanoi Algorithm

```
Algorithm Hanoi(n, source, target, auxiliary):  
    if n == 1:  
        print("Move disk 1 from", source, "to", target)  
    else:  
        Hanoi(n-1, source, auxiliary, target)  
        print("Move disk", n, "from", source, "to", target)  
        Hanoi(n-1, auxiliary, target, source)
```

**Explanation:** The algorithm breaks down the problem recursively. Each recursive call handles one part of the movement:

- Move the top  $n - 1$  disks to the auxiliary rod.
- Move the bottom disk directly to the target.
- Move the  $n - 1$  disks onto the bottom disk.

**Proof of Validity:** By structural induction:

- Base case  $n = 1$ : move one disk directly
- Inductive step: assume it works for  $n - 1$ , then prove it works for  $n$

### 3 Code Implementation

Listing 2: Tower of Hanoi with ASCII Display

```
# Tower of Hanoi - Recursive solution with ASCII visualization
# This script accepts user input for the number of disks and calculates the steps
# needed
# to move the disks from Peg A to Peg C following the Tower of Hanoi rules.
# It displays each step and the current state of the towers in ASCII notation.

def print_towers(towers):
    """
    Displays the current state of the towers in ASCII.

    Parameters:
    towers (dict): A dictionary with keys 'A', 'B', 'C' representing pegs,
                   and values as lists representing disks on each peg.
    """
    max_height = max(len(peg) for peg in towers.values())
    for level in reversed(range(max_height)):
        for peg in ['A', 'B', 'C']:
            if level < len(towers[peg]):
                # Print the disk number at the current level of the peg
                print(f" {towers[peg][level]} ", end="\t")
            else:
                # Print a vertical bar if no disk is present at this level
                print(" | ", end="\t")
        print()
    print(" A \t B \t C ")
    print("-" * 24)

def hanoi(n, source, target, auxiliary, towers):
    """
    Recursive function to solve Tower of Hanoi problem.

    Parameters:
    n (int): Number of disks to move.
    source (str): The peg to move disks from.
    target (str): The peg to move disks to.
    auxiliary (str): The peg used as auxiliary storage.
    towers (dict): The current state of the towers.
    """
    if n == 1:
        # Move the top disk from source to target
        disk = towers[source].pop()
        towers[target].append(disk)
        print(f"Move disk 1 from {source} to {target}")
        print_towers(towers)
    else:
        # Move n-1 disks from source to auxiliary, so they are out of the way
        hanoi(n - 1, source, auxiliary, target, towers)
        # Move the nth disk from source to target
```

```

        disk = towers[source].pop()
        towers[target].append(disk)
        print(f"Move disk {n} from {source} to {target}")
        print_towers(towers)
        # Move the n-1 disks that we left on auxiliary to target
        hanoi(n - 1, auxiliary, target, source, towers)

def main():
    """
    Main function to run the Tower of Hanoi program.
    """
    print("Tower of Hanoi - Recursive with ASCII Visualization")
    n = int(input("Enter number of disks: "))
    # Initialize towers with disks on peg A, and pegs B and C empty
    towers = {
        'A': list(reversed(range(1, n + 1))),
        'B': [],
        'C': []
    }

    print("Initial State:")
    print_towers(towers)
    # Start the recursive solution
    hanoi(n, 'A', 'C', 'B', towers)

if __name__ == "__main__":
    main()

```

*Note:* Code is recursive, accepts input, prints the full move sequence, and uses ASCII visualization.

## 4 Worked-Out Example: $n = 4$

For  $n = 4$ , the sequence of 15 moves is:

1. Move disk 1 from A to C
2. Move disk 2 from A to B
3. Move disk 1 from C to B
4. Move disk 3 from A to C
5. Move disk 1 from B to A
6. Move disk 2 from B to C
7. Move disk 1 from A to C
8. Move disk 4 from A to B
9. Move disk 1 from C to B
10. Move disk 2 from C to A
11. Move disk 1 from B to A
12. Move disk 3 from C to B
13. Move disk 1 from A to C
14. Move disk 2 from A to B
15. Move disk 1 from C to B

## 5 Complexity Analysis and Efficiency

The recurrence relation:

$$T(n) = 2T(n - 1) + 1$$

is exponential. Solving it gives the closed-form:

$$T(n) = 2^n - 1$$

Thus, the time complexity is:

$$\mathcal{O}(2^n)$$

The recursive depth is  $n$ , and each level of the recursion performs two recursive calls. Compared to the recursive Fibonacci algorithm, which redundantly recomputes values, Tower of Hanoi's algorithm is more efficient because it does not revisit the same subproblems.