## Tower of Hanoi: Recursion and Algorithm Design

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# 1 Concept Explanation: Tower of Hanoi Rules and Logic

The Tower of Hanoi is a classic mathematical puzzle consisting of three rods and a set of disks of different sizes. The puzzle begins with all the disks neatly stacked on one rod in ascending order, forming a conical shape. The objective is to move the entire stack to another rod under the following rules:

- Only one disk may be moved at a time.
- Each move must involve taking the top disk from one of the stacks and placing it on another rod.
- No disk may be placed on top of a smaller disk.

The problem lends itself naturally to a recursive solution. For n disks:

- Move n-1 disks to the auxiliary rod (recursive subproblem)
- Move the largest disk to the target rod
- Move the n-1 disks from the auxiliary rod to the target rod (recursive subproblem)

Let T(n) be the minimum number of moves required. Then:

$$T(n) = 2T(n-1) + 1$$

Closed-form:  $T(n) = 2^n - 1$ 

#### **Proof of Closed-Form Using Induction**

**Base Case:**  $T(1) = 2^1 - 1 = 1$ 

Inductive Hypothesis: Assume  $T(k) = 2^k - 1$ 

**Inductive Step:** 

$$T(k + 1) = 2T(k) + 1$$

$$= 2(2^{k} - 1) + 1$$

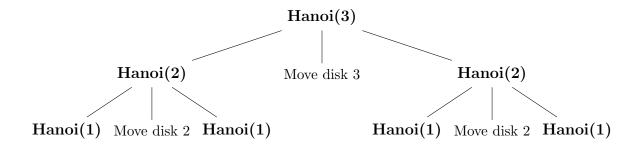
$$= 2^{k+1} - 2 + 1$$

$$= 2^{k+1} - 1$$

**Conclusion:** By the principle of mathematical induction, the closed-form  $T(n) = 2^n - 1$  holds for all  $n \ge 1$ .

## 2 Visualizing the Recursion Tree

Below is a diagram representing the recursive breakdown of the Tower of Hanoi solution for n=3. Each node represents a function call, and the arrows show the sequence of recursive steps taken to solve the subproblems.



## 3 Algorithm and Pseudocode

#### Pseudocode:

```
Listing 1: Recursive Hanoi Algorithm

Algorithm Hanoi(n, source, target, auxiliary):

if n == 1:

print("Move-disk-1-from", source, "to", target)

else:

Hanoi(n-1, source, auxiliary, target)

print("Move-disk", n, "from", source, "to", target)

Hanoi(n-1, auxiliary, target, source)
```

**Explanation:** The algorithm breaks down the problem recursively. Each recursive call handles one part of the movement:

- Move the top n-1 disks to the auxiliary rod.
- Move the bottom disk directly to the target.
- Move the n-1 disks onto the bottom disk.

**Proof of Validity:** By structural induction:

- Base case n = 1: move one disk directly
- Inductive step: assume it works for n-1, then prove it works for n

#### 4 Code Implementation

Listing 2: Tower of Hanoi with ASCII Display

```
# Tower of Hanoi - Recursive solution with ASCII visualization
# This script accepts user input for the number of disks and calculates the steps
# to move the disks from Peg A to Peg C following the Tower of Hanoi rules.
# It displays each step and the current state of the towers in ASCII notation.
def print_towers(towers):
   Displays the current state of the towers in ASCII.
   Parameters:
   towers (dict): A dictionary with keys 'A', 'B', 'C' representing pegs,
                 and values as lists representing disks on each peg.
   max_height = max(len(peg) for peg in towers.values())
   for level in reversed(range(max_height)):
       for peg in ['A', 'B', 'C']:
           if level < len(towers[peg]):</pre>
              # Print the disk number at the current level of the peg
              print(f" {towers[peg][level]} ", end="\t")
              # Print a vertical bar if no disk is present at this level
              print(" | ", end="\t")
       print()
   print(" A \t B \t C ")
   print("-" * 24)
def hanoi(n, source, target, auxiliary, towers):
   Recursive function to solve Tower of Hanoi problem.
   Parameters:
   n (int): Number of disks to move.
   source (str): The peg to move disks from.
   target (str): The peg to move disks to.
   auxiliary (str): The peg used as auxiliary storage.
   towers (dict): The current state of the towers.
   11 11 11
   if n == 1:
       # Move the top disk from source to target
       disk = towers[source].pop()
       towers[target].append(disk)
       print(f"Move disk 1 from {source} to {target}")
       print_towers(towers)
   else:
       # Move n-1 disks from source to auxiliary, so they are out of the way
       hanoi(n - 1, source, auxiliary, target, towers)
       # Move the nth disk from source to target
```

```
disk = towers[source].pop()
       towers[target].append(disk)
       print(f"Move disk {n} from {source} to {target}")
       print_towers(towers)
       \# Move the n-1 disks that we left on auxiliary to target
       hanoi(n - 1, auxiliary, target, source, towers)
def main():
   11 11 11
   Main function to run the Tower of Hanoi program.
   print("Tower of Hanoi - Recursive with ASCII Visualization")
   n = int(input("Enter number of disks: "))
   # Initialize towers with disks on peg A, and pegs B and C empty
   towers = {
       'A': list(reversed(range(1, n + 1))),
       'B': [],
       'C': []
   print("Initial State:")
   print_towers(towers)
   # Start the recursive solution
   hanoi(n, 'A', 'C', 'B', towers)
if __name__ == "__main__":
   main()
```

*Note:* Code is recursive, accepts input, prints the full move sequence, and uses ASCII visualization. If by any chance there are some errors with the code, I have provided a backup version on my GitHub page.

GitHub: github.com/thatrandomasiandev/LaTeX-Projects/tree/main

#### 5 Worked-Out Example: n = 4

For n = 4, the sequence of 15 moves is:

- 1. Move disk 1 from A to C
- 2. Move disk 2 from A to B
- 3. Move disk 1 from C to B
- 4. Move disk 3 from A to C
- 5. Move disk 1 from B to A
- 6. Move disk 2 from B to C
- 7. Move disk 1 from A to C
- 8. Move disk 4 from A to B
- 9. Move disk 1 from C to B
- 10. Move disk 2 from C to A
- 11. Move disk 1 from B to A
- 12. Move disk 3 from C to B
- 13. Move disk 1 from A to C
- 14. Move disk 2 from A to B
- 15. Move disk 1 from C to B

#### 6 Complexity Analysis and Efficiency

My algorithm for the Tower of Hanoi uses a recursive approach with the recurrence:

$$T(n) = 2T(n-1) + 1$$

Solving this gives the closed-form:

$$T(n) = 2^n - 1$$

So the time complexity is:

$$\mathcal{O}(2^n)$$

#### Reflection on Efficiency:

Although the algorithm is exponential, it is still efficient in recursion. In my implementation, each recursive call solves a distinct subproblem exactly once. For example, when calling hanoi(n-1, ...), it completes that entire subproblem before moving on.

In contrast, the recursive Fibonacci algorithm repeatedly recalculates the same values. For example, fib(5) will recompute fib(4) and fib(3) multiple times, leading to a huge number of redundant calls.

#### **Key Difference:**

- 1. My algorithm (Hanoi): No repeated work each move is unique and necessary.
- 2. **Fibonacci**: Many repeated subproblems highly inefficient without memorization.

**Conclusion:** Even though both algorithms are recursive and exponential, my Tower of Hanoi algorithm is more efficient than the recursive Fibonacci method because it avoids redundant computations.