

DSA Interview Question

1) $\underline{\text{count} = 0}$ \hookrightarrow Bit Manipulation \longrightarrow Divide two Integers
Dividend = 9
Divisor = 3
 $9 - 3 = \underline{6}$
 $\underline{\text{count} = 1}$

\hookrightarrow Left shift
Right shift
No we \longrightarrow {
 1) Multiplication
 2) Division
 3) Mod

$3 - 3 = \underline{0}$
 $6 - 3 = \underline{3}$
 $\underline{\text{count} = 2}$

$\text{count} = 3$ *

$$\begin{array}{c}
 \xrightarrow{\hspace{1cm}} \\
 \left(\begin{array}{cc} a & b \\ c & d \end{array} \right) \downarrow \left(\begin{array}{cc} e & f \\ g & h \end{array} \right) = \left(\begin{array}{cc} c_{11} & c_{12} \\ c_{21} & c_{22} \end{array} \right) \xrightarrow{\hspace{1cm}}
 \end{array}$$

Strassen's Matrix Multiplication

Matrix Multiplication

$i \times k$ $k \times j$ 2×2

Small Problem

$\hookrightarrow \Theta(1)$
(constant)

$$\begin{aligned}
 \checkmark c_{11} &= ae + bg & (1) \\
 \checkmark c_{12} &= af + bh & (2) \\
 \checkmark c_{21} &= ce + dg & (3) \\
 \checkmark c_{22} &= cf + dh & (4)
 \end{aligned}$$

$\Theta(n^3)$

for $i = 0$ to n :
 for $j = 0$ to n :
 $c(i, j) = 0$
 for $k = 0$ to n :
 $c(i, j) += A(i, k) * B(k, j)$

Matrix addition $\rightarrow \Theta(n^2)$

$$\left[\begin{array}{cc} a & e \\ 1 & 2 \\ 3 & 4 \end{array} \right] + \left[\begin{array}{cc} b & g \\ 5 & 6 \\ 7 & 8 \end{array} \right] = \left[\begin{array}{cc} 6 & 8 \\ 10 & 12 \end{array} \right]$$

← ←

$$\left(\begin{array}{cc} 1 & 2 \\ 3 & 4 \end{array} \right) * \left(\begin{array}{cc} 5 & 6 \\ 7 & 8 \end{array} \right) = \left(\begin{array}{cc} 5 & 12 \\ 21 & 32 \end{array} \right)$$

$A_{21} \rightarrow 2 \rightarrow 2, 3$
 $C \rightarrow 0, 1$

$A_{22} \rightarrow 2, 3 \rightarrow 8$
 $2, 3 \rightarrow C$

$$\begin{array}{c|cc|cc}
 & A_{11} & A_{12} & A_{13} & A_{14} \\
 \hline
 0 & a_{11} & a_{12} & a_{13} & b \\
 1 & a_{21} & a_{22} & a_{23} & a_{24} \\
 \hline
 2 & a_{31} & a_{32} & a_{33} & a_{34} \\
 3 & a_{41} & a_{42} & a_{43} & a_{44} \\
 \hline
 & A_{21} & & A_{22} &
 \end{array} \times
 \begin{array}{c|cc|cc}
 & B_{11} & B_{12} & B_{13} & B_{14} \\
 \hline
 0 & b_{11} & e & b_{13} & f \\
 1 & b_{21} & b_{22} & b_{23} & b_{24} \\
 \hline
 2 & b_{31} & b_{32} & b_{33} & b_{34} \\
 3 & b_{41} & g & b_{43} & h \\
 \hline
 & B_{21} & B_{22} & &
 \end{array} \quad 4 \times 4 \quad 4 \times 4$$

$A_{11} \rightarrow$

rows $\rightarrow 0$ to 1

$$\text{cols } \rightarrow 0 \text{ to } 1, \quad C_{11} = \underline{ae} + \underbrace{bg}_{f} \xrightarrow{\text{matrix multiplication}}$$

$A_{12} \rightarrow$

$$\text{rows } \rightarrow 0, 1, \quad C_{12} = ce + dg \xrightarrow{\text{matrix addition}}$$

cols $\rightarrow 2, 3$

$$C_{22} = cf + dh$$

Pseudocode

matrixMul(A, B, m): small problem

if $m \leq 2$:

C —

$C_{11} = ae + bg$	(1)
$C_{12} = af + bh$	(2)
$C_{21} = ce + dg$	(3)
$C_{22} = cf + dh$	(4)

else :

$$m \text{id} = n/2$$

matrix addition

$$c_{11} = \underbrace{\text{matrixMul}(a, e, n/2)}_{T(n/2)} + \underbrace{\text{matrixMul}(b, g, n/2)}_{T(n/2)}$$

$$c_{12} = \underbrace{\text{matrixMul}(a, f, n/2)}_{T(n/2)} + \underbrace{\text{matrixMul}(b, h, n/2)}_{T(n/2)}$$

$$c_{21} = \underbrace{\text{matrixMul}(c, e, n/2)}_{T(n/2)} + \underbrace{\text{matrixMul}(d, g, n/2)}_{T(n/2)}$$

$$c_{22} = \underbrace{\text{matrixMul}(c, f, n/2)}_{T(n/2)} + \underbrace{\text{matrixMul}(d, h, n/2)}_{T(n/2)}$$

Divide & conquer approach

$$T(n) = \begin{cases} c & n \leq 2 \\ 8T(n/2) + n^2 & n > 2 \end{cases}$$

Recurrence Relation

$$a = 8 \quad k = 2$$

$$b = 2 \quad P = 0$$

$$\log_b a = \log_2 8 = \underline{\quad}$$

$$\log_b a > k \rightarrow \frac{\text{case 1}}{\Theta(n^3)}$$

Strassen's approach

+ - Matrix

addition

$$P = \underline{(A_{11} + A_{12})} * \underline{(B_{11} + B_{12})}$$

$$Q = \underline{(A_{21} + A_{22})} * \underline{B_{11}}$$

$$R = \underline{A_{11}} * \underline{(B_{12} - B_{22})}$$

$$S = \underline{A_{22}} * \underline{(B_{21} - B_{11})}$$

$$T = \underline{(A_{11} + A_{12})} * \underline{B_{22}}$$

$$U = \underline{(A_{21} - A_{11})} * \underline{(B_{11} + B_{12})}$$

$$V = \underline{(A_{12} - A_{22})} * \underline{(B_{21} + B_{22})}$$

$$C_{11} = P + S - T + V$$

$$C_{12} = R + T$$

$$C_{21} = Q + S$$

$$C_{22} = P + R - Q + U$$

$$T(n) = \begin{cases} C & n \leq 2 \\ 7T(n/2) + n^2 & n > 2 \end{cases}$$

$$a = 7, b = 2$$

$$k = 2, p = 0$$

→ Case 1

$$\log \frac{a}{b} = \log \frac{7}{2} = 2.81 > k$$

$$\underline{\Theta(n^{2.81})}$$

Power of an element

$$\begin{array}{c}
 n < 0 \\
 \xrightarrow{\quad} \\
 \text{oval containing } 2^{-2} \\
 \text{with } 2 \text{ circled in green} \\
 \text{and } -2 \text{ circled in green} \\
 \downarrow \\
 x^n \\
 = \frac{1}{2^2} = \frac{1}{4} = 0.25
 \end{array}$$

$$\begin{array}{c}
 x = 2 \\
 \hline
 n = -2 \\
 \hline
 x^n
 \end{array}$$

$$\begin{array}{c}
 \text{Divide \& conquer} \rightarrow \Theta(\log n) \\
 \hline
 n < 0 \\
 \xrightarrow{\quad} \\
 x = \frac{1}{x} \\
 \hline
 n = -n
 \end{array}$$

$$n = 32 \quad x = 2$$

$$2^{32}$$

$$2^{16} \quad 2^{16}$$

$$\text{mid} = 32/2 = 16$$

Small problem

$$\begin{array}{c}
 n = 1 \\
 \xrightarrow{\quad} \\
 x
 \end{array}$$

$$\begin{array}{c}
 2^{256} \\
 / \\
 16 \\
 / \\
 4 \\
 / \\
 2 \\
 / \\
 2 \\
 / \\
 2 \\
 / \\
 2 \\
 / \\
 2 \\
 / \\
 2 \\
 / \\
 2
 \end{array}$$

Median of $\underline{\text{nums1}}$ & $\underline{\text{nums2}}$

$\Theta(\log(m+n))$

$\underline{\text{nums1}} \rightarrow \underline{\left[\frac{0}{1}, 2 \right]}$ after merging

$\underline{\text{nums2}} \rightarrow \underline{\left[3, 4 \right]} \Rightarrow \underline{\left[0, 1, 2, 3 \right]}$

$$\begin{array}{c} \text{1} \rightarrow \text{float} \\ \text{1} \rightarrow \text{Lower} \end{array} \quad \boxed{m=4} \quad \frac{2+3}{2} = 2.5$$

Bound

Approach 1 → Merge Procedure

$$\underline{\text{mid} = 4//2 = 2} \quad \underline{\text{sorted}(\text{nums1} + \text{num2})}$$

$$\begin{array}{cccc} 0 & 1 & 2 & 3 \\ (1 & 2 & 3 & 4) \end{array} \quad \frac{(2+3)/2}{2}$$

$$\underline{\frac{\text{nums}(\text{mid}) + \text{nums}(\text{mid}-1)}{2}}$$