

Synthetic Algebraic Geometry

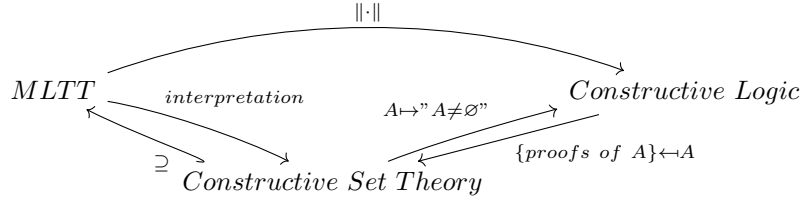
Constructive working over a field with nonzero nilpotents

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1 HoTT revisited

1.1 Curry Howard Correspondence



Martin Löff Type theory	Sets as Types	Propositions as Types
A Type	A is Set	A is Prop
$a : A$	$a \in A$	a is a proof that A holds
$A \rightarrow B$	functions $A \rightarrow B$	A implies B
$\neg A \equiv A \rightarrow \perp$	A is empty	A is false
$A \simeq B$	{ Bijections $A \cong B$ }	$A \leftrightarrow B$
$A \times B$	$A \times B$	$A \wedge B$
$A + B$	$A \sqcup B$	$A \vee B$
$\sum_{x:X} Bx$	$\bigsqcup_{x \in A} Bx$	$\exists x \in X, Bx$ (X a Set)
$\prod_{x:X} Bx$	Sections of $(\bigsqcup_{x \in A} Bx) \rightarrow X$	$\forall x \in X, Bx$ (X a Set)
$x =_A y$	{ proofs of the proposition that $x = y$ holds }	automatic

Here:

$$A \simeq B \equiv \sum_{f:A \rightarrow B} \text{isEquiv}(f),$$

where

$$\text{isEquiv}(f) \equiv \text{sec}(f) \times \text{ret}(f)$$

We have an internal notion, when a type A is a proposition:

$$\begin{aligned} \text{isProp}(A) &\equiv \prod_{x,y:A} x =_A y \\ &\simeq (A \rightarrow \text{isContr}(A)) \end{aligned}$$

where

$$\text{isContr}(A) \equiv \sum_{x:A} \prod_{y:A} x =_A y.$$

Slogan: Once a proposition is true I dont care anymore, why its true.

Proposition 1.1. $\text{isProp}(A), \text{isContr}(A), \text{isEquiv}(f)$ are propositions.

One realizes, that propositions are not stable under $+$ and Σ :

Example. • If 1 denotes the unit type, then $1 + 1$ is not a proposition!

- $\sum_{x:A} Bx$ is the type of all terms x such that Bx holds, so different x such that Bx give different terms!

We see in action where Martin Löff Type theory is not rich enough, when we try to define the image of a map:

Example. Let $f: A \rightarrow B$. If one defines $\text{im } f \equiv \sum_{y:B} \sum_{x:A} fx =_B y$ then we would get

$$\begin{aligned} \text{im } f &\equiv \sum_{y:B} \sum_{x:A} fx =_B y \\ &\simeq \sum_{x:A} \underbrace{\sum_{y:B} fx =_B y}_{\text{contractible}} \\ &\simeq \sum_{x:A} 1 \\ &\simeq A \end{aligned}$$

Similar calculations would show that for the naive definition of surjectivity

$$\text{isSurj}(f) \equiv \prod_{y:B} \sum_{x:A} fx = y \simeq \sum_{g:B \rightarrow A} \prod_{y:B} f(gy) = y \equiv \text{sec}(f).$$

1.2 Propositional Truncation

We need to add another type former to Martin-Löff-Type theory, called the propositional truncation

Definition 1.2. Given a type A , we add a type $\|A\|$, called the propositional truncation, equipped with a map $\eta_A : A \rightarrow \|A\|$, such that

- $\|A\|$ is a proposition
- for any proposition P the map

$$\begin{aligned} (\|A\| \rightarrow P) &\rightarrow (A \rightarrow P) \\ f &\mapsto f \circ \eta_A \end{aligned}$$

is an equivalence.

$$\begin{array}{ccc} A & \xrightarrow{\forall} & P \\ \eta_A \downarrow & \nearrow \exists! & \\ \|A\| & & \end{array}$$

So $\|A\|$ is the proposition, that says that A is inhabited. To emphasize I only have a term of type $\|A\|$, we say A **has merely a term**. Now we can fix the issues we add earlier:

Definition 1.3. • For propositions A, B we define

$$A \vee B \equiv \|A + B\|$$

- For family of propositions Bx indexed over a type $x : A$ we define

$$\exists x : A, Bx \equiv \|\sum_{x:A} Bx\|$$

In words : " There merely exists an $x : A$ such that Bx " , and

$$\forall x : A, Bx \equiv \prod_{x:A} Bx$$

Suggestively: Assuming $\exists x : A, Bx$, only if we want to show a proposition, we are allowed to extract a witness $x : A$ that satisfies Bx . That the dependent product of propositions is a proposition, has to be added as an axiom¹ to Martin-Löf-Type theory.

Remark 1. One can realize propositional truncation as a higher inductive type, but they will not play a role today.

1.3 Subtypes

We fix a universe \mathcal{U} .

We write $\text{Prop} \equiv \sum_{A:\mathcal{U}} \text{isProp}(A)$

Definition 1.4 (Subtypes). Given a type A , a subtype of A is a map

$$U : A \rightarrow \text{Prop}$$

written $U \subset A$. For $x : A$, we write the respective proposition as

$$x \in U \equiv U(x)$$

And we may use the notation

$$U \equiv (x : A \mapsto Ux) \equiv \{x : A \mid Ux\}$$

For subtypes U, V we write $U \subseteq V \equiv \forall x : A, x \in U \rightarrow x \in V$.

Example. *The empty subtype is defined as*

$$\emptyset = \{x : A \mid \perp\}$$

Definition 1.5. Let $(U_i)_i : I \rightarrow (A \rightarrow \text{Prop})$ be a family of subtypes of A . Then

•

$$\bigcup_{i \in I} U_i \equiv \{x : A \mid \exists i : I, x \in U_i\}$$

•

$$\bigcap_{i \in I} U_i \equiv \{x : A \mid \forall i : I, x \in U_i\}$$

Definition 1.6. A map $f : A \rightarrow B$ is an embedding if for all $x, y : A$ the map

$$\text{ap}_f : x =_A y \rightarrow fx =_B fy$$

induced by path-induction, is an equivalence. We write $A \hookrightarrow B$ for the type of maps $A \rightarrow B$ that are embeddings.

¹up to equivalence its called function extensionality

Lemma 1.7. *A map $f : A \rightarrow B$ is an embedding iff for any $y : B$, the fiber $\text{fib}_f(y) \equiv \sum_{x:A} f x =_B y$ is a proposition.*

Proposition 1.8. *Assuming the univalence axiom, there are mutual inverse equivalences*

$$\begin{aligned} (A \rightarrow \text{Prop}) &\simeq \sum_{B:\mathcal{U}} B \hookrightarrow A \\ U &\mapsto (\sum_{x:A} x \in U, \text{proj}_1) \\ \{x:A \mid \text{fib}_f(x)\} &\leftrightarrow (f : B \hookrightarrow A) \end{aligned}$$

2 Synthetic Algebraic Geometry

2.1 Algebra

Definition 2.1. A type A is a set if $\prod_{x,y:A} \text{isProp}(x =_A y)$. We write Set for the subtype of \mathcal{U} that consists of the sets.

We can do algebra in HoTT. The type of (commutative unital) Rings can be written as

$$\text{Ring} \equiv \sum_{R:\text{Set}} \sum_{+, \cdot : R \times R \rightarrow R} \sum_{0, 1 : R} \text{isAbelianGroup}(R, +, 0) \times \text{isAbelianMonoid}(R, \cdot, 1) \times \text{isDistributive}(R, +, \cdot)$$

where we omit the definitions of the algebraic-flavoured propositions.

Remark 2 (Lean). In Lean, this can be formalized by structures. Conversely a structure in Lean can be encoded as an iterated Sigma-type in HoTT.

Definition 2.2. Let $(R, +, \cdot, 0, 1) : \text{Ring}$. Given $x : R$, define

$$\text{isInv}(x) \equiv \sum_{y:R} x \cdot y =_R 1$$

One can show, that this is indeed a proposition, because inverses are unique if they exist.

Definition 2.3.

$$\text{isLocal}(R) \equiv 1 \neq 0 \wedge \forall x, y : R, (\text{isInv}(x + y) \rightarrow \text{isInv}(x) \vee \text{isInv}(y))$$

In synthetic algebraic geometry, there is always the datum of a commutative ring R in the context. The first axiom says:

Axiom 1 (Loc). R is a local ring.

Definition 2.4. • The type of algebras is

$$\text{Alg} \equiv \sum_{A:\text{Ring}} \text{Hom}_{\text{Ring}}(R, A)$$

$\text{Hom}_R(A, B)$ is the set of homomorphisms of R -algebras $A \rightarrow B$.

- An algebra A is finitely presented, if

$$\parallel \sum_{m,n:\mathbb{N}} \sum_{f_1,\dots,f_m:R[X_1,\dots,X_n]} A =_{\text{Alg}} R[X_1,\dots,X_n]/(f_1,\dots,f_m) \parallel$$

Remark 3. By the univalence axiom, the latter equality in Alg can be replaced by asking for an isomorphism of R -algebras.

Definition 2.5. For A a (finitely presented) R -algebra, we define the spectrum as the set

$$\text{Spec } A \equiv \text{Hom}_R(A, R)$$

The fundamental observation of algebraic geometry is, that once we pick a representation of A through polynomials f_1, \dots, f_m , the spectrum corresponds to the set of common zeros over R of those polynomials.

Lemma 2.6 (The Fundamental Theorem of AG). *We have an equivalence*

$$\begin{aligned} \text{Spec}(R[X_1, \dots, X_n]/(f_1, \dots, f_m)) &\rightarrow \{(x_1, \dots, x_n) \in R^n \mid f_i(x_1, \dots, x_n) = 0 \forall i\} \\ \phi &\mapsto (\phi(X_1), \dots, \phi(X_n)) \end{aligned}$$

Proof. This is a direct consequence of the universal property of the quotient ring and the universal property of the polynomial algebra. \square

Remark 4 (Classically!). The above equivalence holds if R is an algebraically closed field and one replace Spec by MaxSpec , the set of maximal ideals.

Definition 2.7. Given a type X , we have the algebra R^X of R -valued functions on X , i.e. Its underlying set is the function type $X \rightarrow R$ and it inherits an R -algebra structure from the target by doing everything pointwise. Its not finitely presented in general.

With the next axiom, the spectrum gets new interesting geometric structure:

Axiom 2 (Duality). For any finitely presented algebra A the **duality map**

$$\begin{aligned} A &\rightarrow R^{\text{Spec } A} \\ f &\mapsto (x \mapsto x(f)) \end{aligned}$$

is an equivalence.

We may use this as a coercion, i.e for $f : A, x : \text{Spec } A$ we may write $f(x) \equiv x(f)$.

What does this tell us?

Example. *Every function $R \rightarrow R$ is a polynomial!*

Proof. Indeed, $\text{Spec } R[X] = \text{Hom}_R(R[X], R) \simeq R$ as a set by [The Fundamental Theorem of AG](#), hence we only have to apply [Duality](#) to $A = R[X]$. \square

Let A denote a finitely presented algebra.

Proposition 2.8 (Weak Nullstellensatz). $A = 0$ iff $\text{Spec } A = \emptyset$.

Proof. Both sides are propositions, where we rephrase the left hand side as $1 =_A 0$. If $1 =_A 0$, then there are no R -algebra homomorphisms $A \rightarrow R$ because $0, 1$ have to get preserved but they are not equal in R by [Loc](#).

If $\text{Spec } A = \emptyset$ then

$$A \simeq R^{\text{Spec } A} = R^\emptyset$$

is contractible as a set, hence $1 =_A 0$. \square

We now see some consequences of this:

The next proposition shows, that R is a field in a weak way (from a constructive point of view) and $\neg\neg(x = 0)$ is the same as x being nilpotent.

Proposition 2.9.

- A vector in R^n is non zero iff one of its entries is invertible.

- Let $x \in R$. x is not invertible iff x is nilpotent

Proof.

- $(x_1, \dots, x_n) \neq 0$ iff $\text{Hom}_R(R/(x_1, \dots, x_n), R) = \text{Spec}(R/(x_1, \dots, x_n))$ is empty iff $R/(x_1, \dots, x_n) = 0$ iff $(x_1, \dots, x_n) = R$ iff one of the x_i is invertible by [Loc](#).
- x is not invertible iff $\text{Spec}(R_x)$ is empty iff $R_x = 0$ iff $1 = 0$ in R_x iff x is nilpotent.

□

Classically, the next example does not work in a field, but constructively with our notion of a field ($x \neq 0 \rightarrow \text{isInv}(x)$) it does.

Example. In R not every nilpotent is zero.

Proof. We prove something stronger, namely that not every element that squares to zero is itself zero. If it would then $\text{Spec}(R[x]/x^2)$ would be a singleton by [The Fundamental Theorem of AG](#). But the R algebra $R[x]/x^2$ is not isomorphic to R^1 as an R -algebra. □

2.2 Spec is right adjoint and fully faithful when restricted to FP -algebras

Proposition 2.10 (Adjunction). *For any $A : \text{Alg}$, we have a natural equivalence*

$$(X \rightarrow \text{Spec } A) \xrightarrow{\sim} \text{Hom}_R(A, R^X) \\ \text{id}_{\text{Spec } A} \mapsto \text{duality map of } A$$

Proof. We have the natural flip map

$$(X \rightarrow (A \rightarrow R)) \xrightarrow{\sim} (A \rightarrow (X \rightarrow R))$$

which is an equivalence. I claim that this map restricts to a map $(X \rightarrow \text{Spec } A) \simeq \text{Hom}_R(A, R^X)$. One now has to check, that an element on the left hand side has image in algebra homomorphisms $\text{Hom}_R(A, R)$ if and only if the corresponding function on the right hand side will be an algebra homomorphism. □

Now invoking, that by [Duality](#) the counit of the above adjunction is an equivalence, we deduce the following two purely formal statements

Corollary (Fullyfaithfulness of Spec). *Spec enhances to a fully faithful functor*

$$\text{Spec} : \text{Alg}_{FP} \rightarrow \mathcal{U}$$

That means: For A, B finitely presented A -algebras, the map

$$\text{Hom}_R(A, B) \rightarrow (\text{Spec } B \rightarrow \text{Spec } A) \\ \varphi \mapsto (x \mapsto x \circ \varphi)$$

is an equivalence.

Corollary. *A type X is equivalent to $\text{Spec } A$ for some finitely presented A iff the unit map $X \rightarrow \text{Spec}(R^X)$ is an equivalence. In this case we call X **affine**.*

2.3 Topology

A subtype $U : X \rightarrow \text{Prop}$ is open if all the propositions $U(x)$ are merely of the form $(f_1 = 0) \wedge \dots \wedge (f_n = 0)$ for some $f_i : R$.

Definition 2.11. • A proposition P is open iff there merely exists $f_1, \dots, f_n : R$ such that

$$P = \text{isInv}(f_1) \vee \dots \vee \text{isInv}(f_n)$$

.

- A proposition P is closed iff there merely exists $f_1, \dots, f_n : R$ such that

$$P = (f_1 = 0) \wedge \dots \wedge (f_n = 0).$$

- A subtype $U \subset X$ is open (closed) iff the proposition $x \in U$ is open (closed) for all $x : X$.
- for $f : A$, define the *principal open subsets* of $\text{Spec } A$

$$D(f) = \{x : \text{Spec } A \mid \text{isInv}(f(x))\}$$

- For $f_1, \dots, f_n : A$ define the closed subset, the *vanishing locus*

$$V(f_1, \dots, f_n) = \{x : \text{Spec } A \mid f_i(x) = 0 \ \forall i\}$$

Example. For any open subtype $U \subset Y$ and any map $f : X \rightarrow Y$ the preimage $\{x : X \mid f(x) \in U\}$ is open.

Example. We have $D((X - a_1) \dots (X - a_n)) = \text{Spec } R[X] \setminus \{a_1, \dots, a_n\}$

On the other hand one could define for $X = \text{Spec } A$ a subtype $U \subset X$ to be open if its merely of the form $U = \bigcup_{i=1}^n D(f_i)$ for some $f_1, \dots, f_n \in A$. To show that those two topologies coincide we need one last axiom.

2.4 The last axiom

On the other hand, there is also the *Zariski topology* on $\text{Spec } A$, i.e. where the principal opens form a basis of the topology. To show, that this coincide with the pointwise topology above, we need one last axiom.

Axiom 3. for any family B of inhabited types over $\text{Spec } A$ we merely find a principal open cover $\text{Spec } A = \bigcup_i D(f_i)$ and local sections $\prod_{x:D(f_i)} Bx$:

$$\left\| \underbrace{\sum_{n:\mathbb{N}} \sum_{(f_1, \dots, f_n): A^n} \left(\bigcup_{i=1}^n D(f_i) = \text{Spec } A \right)}_{\text{principal open cover}} \times \prod_{i=1}^n \underbrace{\prod_{x:D(f_i)} Bx}_{\text{local sections}} \right\|$$

With this we also understand how closed subtypes of $\text{Spec } A$ look like. Also computing cohomology seems not possible without this axiom.

There is this connection of principal opens to algebra:

Lemma 2.12. Given $f_1, \dots, f_n : A$ we have $(f_1, \dots, f_n) = A$ iff $\bigcup_i D(f_i) = \text{Spec } A$

Proof. We prove more generally, that $D(g) \subset \bigcup_i D(f_i)$ iff $g \in \sqrt{(f_1, \dots, f_n)}$. First observe that open subsets satisfy the law of excluded middle by 2.9. Hence we may understand when $(\bigcup D(f_i))^c \cap D(g) = \emptyset$, but by the same proposition

$$\left(\bigcup D(f_i)\right)^c \cap D(g) = V(f_1, \dots, f_n) \cap D(g) = \text{Spec}(A/(f_1, \dots, f_n)_g) = \emptyset$$

iff $(A/(f_1, \dots, f_n))_g = 0$ by the nullstellensatz iff $g \in \sqrt{(f_1, \dots, f_n)}$. □

Definition 2.13. The type of unimodular functions is defined as:

$$\text{Um}(A) = \sum_{n \in \mathbb{N}} \sum_{f_1, \dots, f_n : A} (f_1, \dots, f_n) = (1)$$

In words, For any finitely presented algebra A and any family $B : \text{Spec } A \rightarrow \mathcal{U}$ we have a map

$$(\forall x : X, \|Bx\|) \rightarrow \left\| \sum_{(n, f_1, \dots, f_n) : \text{Um}(A)} \prod_{i=1}^n \prod_{x : D(f_i)} Bx \right\|$$

Now we are finally able to state the definition of a scheme.

Definition 2.14. A type X is a scheme if there merely exists a cover of open subtypes $(U_i)_i$ such that each U_i is affine.

2.5 More on topology

Lemma 2.15. *The Open subsets of X contain \emptyset, X and are stable under finite unions and intersections. Furthermore*

Lemma 2.16. *For $f_1, \dots, f_n : A$, There is an equivalence*

$$V(f_1, \dots, f_n) \rightarrow \text{Spec}(A/(f_1, \dots, f_n))$$

and $D(f) \simeq \text{Spec}(A_f)$

Proof. By the universal property, An algebra homomorphism $A/(f_1, \dots, f_n) \rightarrow R$ is an algebra homomorphism $x : A \rightarrow R$ such that $f_i(x) = x(f_i) = 0$ for $i = 1, \dots, n$. □

Proof. Indeed $(x - a_1) \dots (x - a_n)$ is invertible iff one of the $(x - a_i)$ is invertible, i.e. $x \neq a_i$ by 2.9 □

Lemma 2.17. *We have $V(f_1, \dots, f_n) \subset V(g_1, \dots, g_m)$ iff $(g_1, \dots, g_m) \subset (f_1, \dots, f_n)$.*

Proof. Such an inclusion corresponds to a map $\text{Spec}(A/(f_1, \dots, f_n)) \rightarrow \text{Spec}(A/(g_1, \dots, g_m))$ over $\text{Spec } A$. We conclude by fully faithfulness of Spec . □

Example. Classically, Consider an inclusion of subsets $V(x^2) \subset V(x)$ inside $\text{Spec } R[x]$. Now it does hold $(x) \not\subset (x^2)$. Instead, classically we have an inclusion of subsets of $\text{Spec } A$ $V(I) \subset V(J)$ iff $J \subset \sqrt{I}$ and only if we consider them as closed subschemes we get the condition $J \subset I$. Hence internally the vanishing locus carries more structure than just the underlying set!