Inter-Process Communication (IPC)

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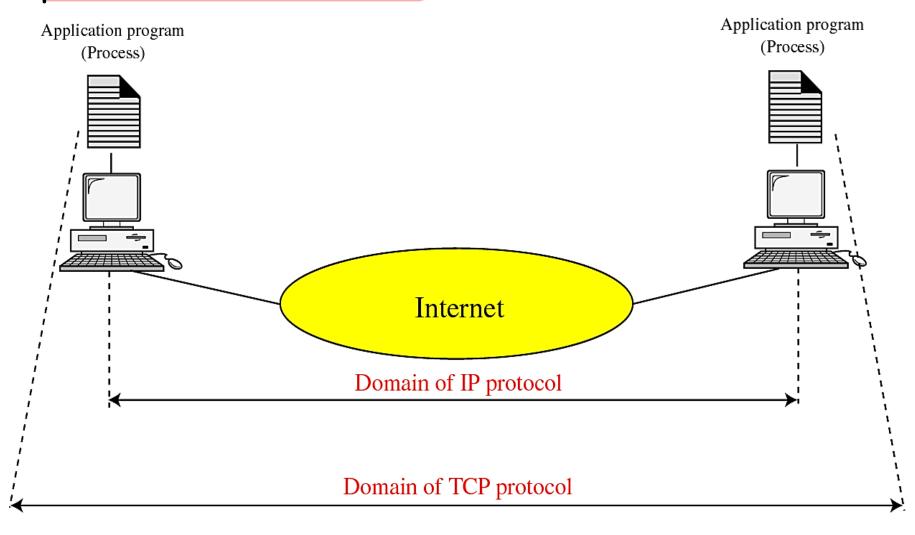
Outlines

- □ What is inter-process communication?
- pipe()

Process Communication

巴芬 IDel 华丽到安皇 郑和.

Host-to-host communication and process-toprocess communication

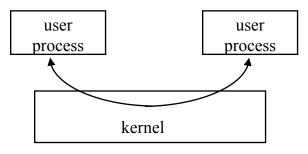


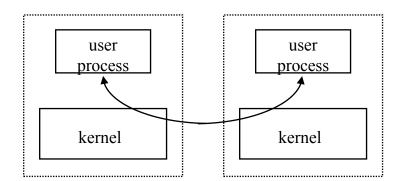
Interprocess Communication (IPC) 學吃咖啡

- The general, processes cannot influence each other
 - Each process is executed in isolation from the others (isolated memory): 沙型型地 电弧分数
 - For example, one process cannot write into the memory of another (建學性先型性的 可學問告生)
 - ☐ IPC between Two Processes

☐ IPC between Two Systems

at the same system





Classifications of IPC

- But you can let several processes communicate with each other via the following methods:
 - Signals: Send a signal (SIGHUP, SIGINT, SIGKILL, etc.)
 to another process
 - OPIPE: a communication channel to transfer data
 - FIFOs, Message Queues, Shared memory (the same memory area is accessible to multiple processes),
 Semaphores (e.g., to regulate access to shared memory)
- Generally, all these IPCs work only between processes of the same computer

PIPE CH11

- A specific type of file that IPC is supported by OS
 - Temporal file that is managed by OS, unlike general file
 - It's used to transfer data between processes not to store data
- IPC using a PIPE THEORY PIPEON WITE TO WHITE THE MENT WITH THE MENT WITH THE MENT WITH THE MENT WITH THE MENT WHITE THE MENT WHITE THE MENT WHITE THE MENT WITH THE MENT WITH THE MENT WHITE THE MENT
 - Sender writes to a PIPE and receiver reads from that
 - Supports a stream channel
 - Sent data has a sending order

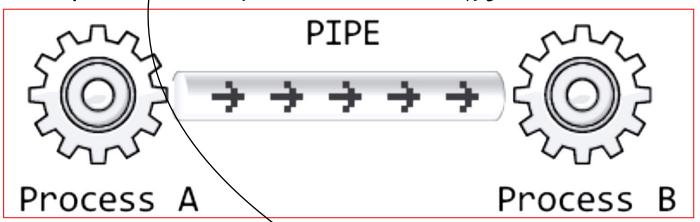
PIPE

□ Acts as a conduit/channel allowing two processes to communicate 平型加度等的 如鄉 零

- □ Some questions?
 - · Is communication unidirectional or bidirectional?
 - Must there exist a relationship (i.e., parent-child) between the communicating processes?

pipe()-unnamed PIPE

- □ A simple, unnamed pipe provides a one-way flow of
- - O Can be thought as a special file that can store a limited amount of data in a first-in-first-out (FIFO) manner, exactly aking to a queue (REMENTATION (FIRST ON)



以相思。由 GIOK的 FIFOSIS 和影片能等现象的影。

□Other variations:

- Stream pipes
- FIFOs

pipe()-unnamed PIPE

- □An unnamed pipe is created by calling pipe(), which returns an array of 2 file descriptors (int)
 - The file descriptors are for reading and writing, respectively
- Unnamed pipes can only be used between related process, such as parent/child, or child/child processes: parent/child child/child to the child/chil
- □ Unnamed pipes can exist only as long as the processes using them :unnamed 如此 妈妈 ?!!

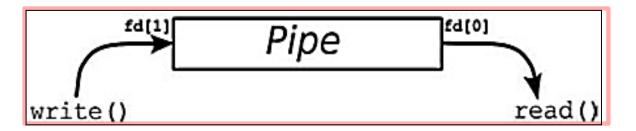
pipe() System Call (unnamed)

Creates a unidirectional pipe.

```
#include < unistd.h>
int pipe (int pipefd[2]);
```

Return 0 on Success; -1 on Failure;

- ☐ If successful, the pipe system call will return two integer file descriptors, pipefd[0] and pipefd[1]
 - opipefd[0] is the read end from the pipe
 - opipefd[1] is the write end to the pipe
- □ Example (pipe1.c)
 - Parent/child processes communicating via unnamed pipe

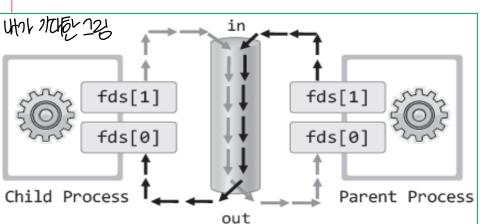


```
Example #1-pipe1.c
int main(int argc, char *argv[])
                                            Junidirectional communication
    int fds[2];
                                                        in
    char str[]="Who are you?";
    char buf[BUF_SIZE];
    pid t pid;
                                                              fds[1]
                                            fds[1]
                      [o]: read
    pipe(fds);
                                            fds[0]
                                                              fds[0]
                      [1]: WINTE
    pid=fork();
    if(pid==0): child process
                                    Child Process
                                                                  Parent Process
                                                       out
        write(fds[1], str, sizeof(str));
    else: Paront process
        read(fds[0], buf, BUF_SIZE);
        puts(buf);
                                  root@my_linux:/tcpip# gcc pipe1.c -o pipe1
                                  root@my linux:/tcpip# ./pipe1
   return 0;
                                  Who are you?
```

Example #2-pipe2.c (bidirectional)

```
int main(int argc, char *argv[])
    int fds[2];
    char str1[]="Who are you?";
    char str2[]="Thank you for your message";
    char buf[BUF SIZE];
   pid t pid;
    pipe(fds);
                              [o]: road
    pid=fork();
                              [1]: WINTE
    if(pid==0)
       write(fds[1], str1, sizeof(str1));
       sleep(2);
       read(fds[0], buf, BUF_SIZE);
       printf("Child proc output: %s \n", buf);
    else
       read(fds[0], buf, BUF SIZE);
       printf("Parent proc output: %s \n", buf);
       write(fds[1], str2, sizeof(str2));
       sleep(3);
   return 0;
```

Bad Example!! Why?



```
char str1[]="Who are you?";
char str2[]="Thank you for your message";
char buf[BUF SIZE];
                                                fds1[1]
                                                                      fds1[0]
pid t pid;
pipe(fds1), pipe(fds2);
                                                fds2[0]
                                                                      fds2[1]
pid=fork();
if(pid==0)
                                         Child Process
                                                                        Parent Process
   write(fds1[1], str1, sizeof(str1));
    read(fds2[0], buf, BUF_SIZE);
    printf("Child proc output: %s \n", buf);
else
    read(fds1[0], buf, BUF_SIZE);
    printf("Parent proc output: %s \n", buf);
    write(fds2[1], str2, sizeof(str2));
    sleep(3);
                                    root@my linux:/tcpip# gcc pipe3.c -o pipe3
                                    root@my linux:/tcpip# ./pipe3
return 0;
                                    Parent proc output: Who are you?
                                    Child proc output: Thank you for your message
```

Example #3-pipe3.c

(bidirectional)-Good!

int main(int argc, char *argv[])

int fds1[2], fds2[2];

```
/*echo storeserv.c*/
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <unistd.h>
#include <signal.h>
#include <sys/wait.h>
#include <arpa/inet.h>
#include <sys/socket.h>
#define BUFSIZE 100
void errorHandling(char* msg);
void readChildProc(int sig);
int main(int argc, char* argv[])
 int servSock, clntSock;
  struct sockaddr in servAddr, clntAddr;
 int fds[2];
  pid t pid;
  struct signation act:
  socklen t addrSz;
  int strLen, state;
 char buf[BUFSIZE];
 if( argc != 2 )
  printf("usage : %s <port>\n", argv[0]);
  exit(1);
  act.sa handler = readChildProc;
  sigemptyset(&act.sa mask);
  act.sa flags = 0;
  state = sigaction(SIGCHLD, &act, 0);
  servSock = socket(PF_INET, SOCK_STREAM, 0);
  memset(&servAddr, 0, sizeof(servAddr));
  servAddr.sin family = AF INET:
  servAddr.sin addr.s addr = htonl(INADDR ANY);
  servAddr.sin port = htons(atoi(argv[1]));
 if( bind(servSock, (struct sockaddr*)&servAddr, sizeof(servAddr)) == -1 )
    errorHandling("bind() error");
 if( listen(servSock, 5) == -1 )
    errorHandling("listen() error");
```

```
Example#4 Echo server
storing message (1/2)
    /*Server stores all strings sent by client*/
```

```
pipe(fds);
                                                              Example#4 Echo
pid = fork();
if( pid == 0 )
 FILE *fp = fopen("echomsg.dat", "wt");
                                                              server storing
 char msqbuf[BUFSIZE];
 int i. len:
 for( i = 0 ; i < 10 ; i++ )
                                                              message (2/2)
 len = read(fds[0], msgbuf, BUFSIZE);
 fwrite((void*)msgbuf, 1, len, fp);
 fclose(fp);
 return 0;
while(1)
 addrSz = sizeof(clntAddr);
 clntSock = accept(servSock, (struct sockaddr*)&clntAddr, &addrSz);
 if( clntSock == -1 )
   continue;
 else
   puts("new client connected...");
 pid = fork();
 if( pid == 0 )
  close(servSock);
 while((strLen = read(clntSock, buf, BUFSIZE)) != 0 )
                                                                    void errorHandling(char* msg)
   write(clntSock, buf, strLen);
   write(fds[1], buf, strLen);
                                                                        fputs(msg, stderr);
                                                                        fputc('\n',stderr);
                                                                        exit(1);
 close(clntSock);
 puts("client disconnected...");
 return 0;
```

else

return 0;

close(clntSock);

close(servSock);

void readChildProc(int sig)

pid = waitpid(-1, &status, WNOHANG);

printf("removed proc id: %d\n", pid);

pid_t pid;
int status;

Example#4 Echo server-additional explanations

```
pipe(fds);
             서버에서 처음으로
pid=fork();
             생성하는 자식프로세스
if(pid==0)
   FILE * fp=fopen("echomsg.txt", "wt");
   char msgbuf[BUF SIZE];
   int i, len;
   for(i=0; i<10; i++)
       len=read(fds[0], msgbuf, BUF_SIZE)
       fwrite((void*)msgbuf, 1, len, fp);
   fclose(fp);
   return 0;
```

accept 함수 호출 후 fork 함수호출을 통해서 파이 프의 디스크립터를 복사하고, 이를 이용해서 이전에 만들어진 자식 프로세스에게 데이터를 전송한다.

파이프를 생성하고 자식 프로세스를 생성해서, 자식 프로세스가 파이프로부터 데이터를 읽어서 저장하도록 구현되어 있다.

```
clnt_sock=accept(serv_sock, (struct sockaddr*)&clnt_adr, &adr_sz);
                    서버에서 연결
pid=fork();
                    허용시마다 생성하는
if(pid==0)
                    자식프로세스
   close(serv_sock);
   while((str_len=read(clnt_sock, buf, BUF_SIZE))!=0)
      write(clnt_sock, buf, str_len);
      write(fds[1], buf, str_len);
   }
   close(clnt_sock);
   puts("client disconnected...");
   return 0;
 }
 else
    close(clnt_sock);
```