

## Coin Change :-

similar to 0-1 Knapsack Problem (inclusion/exclusion principle)

**Coin Change (Minimum Coins)**

Difficulty: Medium Accuracy: 26.74% Submissions: 247K+ Points: 4

You are given an array `coins[]`, where each element represents a coin of a **different** denomination, and a target value `sum`. You have an **unlimited** supply of each coin type {coins1, coins2, ..., coinsm}.

Your task is to determine the **minimum** number of coins needed to obtain the target `sum`. If it is not possible to form the sum using the given coins, return `-1`.

**Examples:**

**Input:** `coins[] = [25, 10, 5], sum = 30`  
**Output:** 2  
**Explanation:** Minimum 2 coins needed, 25 and 5

**Input:** `coins[] = [9, 6, 5, 1], sum = 19`  
**Output:** 3  
**Explanation:**  $19 = 9 + 9 + 1$

**Input:** `coins[] = [5, 1], sum = 0`  
**Output:** 0  
**Explanation:** For 0 sum, we do not need a coin

**Input:** `coins[] = [4, 6, 2], sum = 5`  
**Output:** -1  
**Explanation:** Not possible to make the given sum.

As we know in 0-1 Knapsack problem, only single unit was given. But here, we have unlimited supply of items.

Brute force :- parameters defined :-

$(i, \text{currSum})$

Transition :-  $\text{ans} = \text{INT\_MAX}$

if (`coins[i] ≤ currSum`) { if (`fn(i-1, currSum - coins[i]) != INT_MAX`)  $\text{ans} = 1 + \text{ans}$ ; }

$\text{ans} = \min(\text{ans}, \text{fn}(i-1, \text{currSum}))$

return `ans`;

Base Case :-

if (`currSum == 0`) return 0;

if (`i < 0`) return `INT_MAX`;

```

class Solution {
public:
    int minCoins(vector<int> &coins, int sum) {
        // code here
        int n=coins.size();
        int ans=help(coins, n-1, sum);
        return ans==INT_MAX?-1:ans;
    }
    int help(vector<int> &coins, int i, int sum){
        if(sum==0) return 0;
        if(i<0) return INT_MAX;
        int ans=INT_MAX;
        if(sum>=coins[i]){
            int x=help(coins, i, sum-coins[i]);
            if(x!=INT_MAX) ans=1+x;
        }
        ans=min(ans, help(coins, i-1, sum));
        return ans;
    }
};

```

T.C.  $O(2^n)$  S.C.  $O(n)$  ← recursion stack

Optimized Solution :- Using Recursion + Memoization :-

T.C.  $O(\text{size of coins}[] \cdot \text{sum})$

S.C.  $O(\text{size of coins}[] \cdot \text{sum})$  ← dp array

```

1 // Driver Code Ends
2
3 class Solution {
4 public:
5     int minCoins(vector<int> &coins, int sum) {
6         // code here
7         int n=coins.size();
8         vector<vector<int>> dp(n, vector<int>(sum+1, -1));
9         int ans=help(coins, n-1, sum, dp);
10        return ans==INT_MAX?-1:ans;
11    }
12    int help(vector<int> &coins, int i, int sum, vector<vector<int>> &dp){
13        if(sum==0) return 0;
14        if(i<0) return INT_MAX;
15        if(dp[i][sum]!=-1) return dp[i][sum];
16        int ans=INT_MAX;
17        if(sum>=coins[i]){
18            int x=help(coins, i, sum-coins[i], dp);
19            if(x!=INT_MAX) ans=1+x;
20        }
21        ans=min(ans, help(coins, i-1, sum, dp));
22        return dp[i][sum]=ans;
23    }
24 };
25

```