

Subarray Sums Divisible by K . -

974. Subarray Sums Divisible by K

Solved

Medium Topics Companies

Given an integer array `nums` and an integer `k`, return the number of non-empty **subarrays** that have a sum divisible by `k`.

A **subarray** is a contiguous part of an array.

Example 1:

Input: `nums = [4,5,0,-2,-3,1]`, `k = 5`

Output: 7

Explanation: There are 7 subarrays with a sum divisible by `k = 5`:

`[4, 5, 0, -2, -3, 1]`, `[5, 0]`, `[5, 0, -2, -3]`, `[0]`, `[0, -2, -3]`, `[-2, -3]`

Example 2:

Input: `nums = [5]`, `k = 9`

Output: 0

Constraints:

- $1 \leq \text{nums.length} \leq 3 \times 10^4$
- $-10^4 \leq \text{nums}[i] \leq 10^4$
- $2 \leq k \leq 10^4$

Seen this question in a real interview before? 1/5

Yes No

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eg. $1/p$; `nums = [4, 5, 0, -2, -3, 1]`, `k = 5`

`[4, 9, 9, 7, 4, 5]`

`[4, 0, 0, 3, 2, 1]`

`[4, 4, 4, 2, 4, 0]`

4 - 0
4 - 1
4 - 2
4 - 5

1 + 2 + 3 + 1

$a + x$, $\frac{1}{k} \cdot k - a$
 $\frac{1}{k} \cdot k$, n/k

```
1 class Solution {
2 public:
3     int subarraysDivByK(vector<int>& a, int k) {
4         int n=a.size();
5         int ans=0;
6         int cur=0;
7         unordered_map<int, int> mp;
8         for(int i=0; i<n; i++){
9             cur+=a[i];
10            int p = (cur % k + k) % k;
11            if(mp.count(p)){
12                ans+=mp[p];
13            }
14            if(p==0){ans++;}
15            mp[p]++;
16        }
17        return ans;
18    }
19 }
```

T.C. $O(n)$
S.C. $O(n)$

```
C++ Java Copy
1 class Solution {
2 public:
3     int subarraysDivByK(vector<int>& nums, int k) {
4         int n = nums.size();
5         int prefixMod = 0, result = 0;
6
7         // There are k mod groups 0...k-1.
8         vector<int> modGroups(k);
9         modGroups[0] = 1;
10
11         for (int num : nums) {
12             // Take modulo twice to avoid negative remainders.
13             prefixMod = (prefixMod + num % k + k) % k;
14             // Add the count of subarrays that have the same remainder as the current
15             // one to cancel out the remainders.
16             result += modGroups[prefixMod];
17             modGroups[prefixMod]++;
18         }
19
20         return result;
21     }
22 };
```

Complexity Analysis

Here, `n` is the length of `nums` and `k` is the given integer.

- Time complexity: $O(n + k)$
 - We require $O(k)$ time to initialize the `modGroups` array.
 - We also require $O(n)$ time to iterate over all the elements of the `nums` array. The computation of the `prefixSum` and the calculation of the subarrays divisible by `k` take $O(1)$ time for each index of the array.
- Space complexity: $O(k)$
 - We require $O(k)$ space for the `modGroups` array.

Comments (72)

Sort by: Best