

# Floyd Warshall Algorithm:-

## Floyd Warshall

Difficulty: Medium Accuracy: 32.89% Submissions: 182K+ Points: 4 Average Time: 15m

You are given an weighted **directed** graph, represented by an adjacency matrix, **dist[][]** of size **n x n**, where **dist[i][j]** represents the weight of the edge from **node i to node j**. If there is no direct edge, **dist[i][j]** is set to a large value (i.e.,  $10^8$ ) to represent infinity.

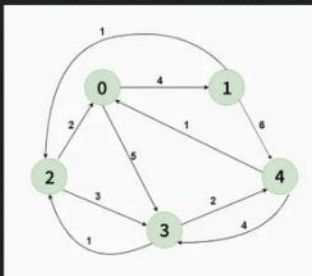
The graph may contain **negative edge weights**, but it does not contain any **negative weight cycles**.

Your task is to find the **shortest distance** between every pair of nodes **i** and **j** in the graph.

Note: Modify the distances for every pair **in place**.

Examples :

**Input:** **dist[][]** =  $[[0, 4, 10^8, 5, 10^8], [10^8, 0, 1, 10^8, 6], [2, 10^8, 0, 3, 10^8], [10^8, 10^8, 1, 0, 2], [1, 10^8, 10^8, 4, 0]]$



**Output:**  $[[0, 4, 5, 5, 7], [3, 0, 1, 4, 6], [2, 6, 0, 3, 5], [3, 7, 1, 0, 2], [1, 5, 5, 4, 0]]$

	0	1	2	3	4
0	0	4	5	5	7
1	3	0	1	4	6
2	2	6	0	3	5
3	3	7	1	0	2
4	1	5	5	4	0

**Explanation:** Each cell **dist[i][j]** in the output shows the shortest distance from node **i** to node **j**, computed by considering all possible intermediate nodes.

Input: dist[][] = [[0, -1, 2], [1, 0, 10<sup>8</sup>], [3, 1, 0]]

	0	1	2
0	0	-1	2
1	1	0	10 <sup>8</sup>
2	3	1	0

Output: [[0, -1, 2], [1, 0, 3], [2, 1, 0]]

	0	1	2
0	0	-1	2
1	1	0	3
2	2	1	0

**Explanation:** Each cell dist[i][j] in the output shows the shortest distance from node i to node j, computed by considering all possible intermediate nodes.

From 2 to 0 shortest distance should be 2 by following path 2 -> 1 -> 0

From 1 to 2 shortest distance should be 3 by following path 1 -> 0 -> 2

**Constraints:**

1 ≤ dist.size() ≤ 100

-1000 ≤ dist[i][j] ≤ 1000

[Try more examples](#)

class Solution {

public:

void floydWarshall (vector<vector<int>> &dist) {

int n = dist.size();

for (int k = 0; k < n; k++) {

for (int i = 0; i < n; i++) {

for (int j = 0; j < n; j++) {

if (dist[i][k] != 1e8 && dist[k][j] != 1e8 &&

dist[i][k] + dist[k][j] < dist[i][j]) {

$$\text{dist}[i][j] = \text{dist}[i][k] + \text{dist}[k][j];$$

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