

0-1 Knapsack Problem

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Difficulty: Medium Accuracy: 31.76% Submissions: 474K+ Points: 4

Given n items, each with a specific **weight** and **value**, and a knapsack with a capacity of W , the task is to put the items in the knapsack such that the **sum of weights of the items** $\leq W$ and the **sum of values** associated with them is **maximized**.

Note: You can either place an item entirely in the bag or leave it out entirely. Also, each item is available in **single** quantity.

Examples:

Input: $W = 4$, $val[] = [1, 2, 3]$, $wt[] = [4, 5, 1]$
Output: 3
Explanation: Choose the last item, which weighs 1 unit and has a value of 3.

Input: $W = 3$, $val[] = [1, 2, 3]$, $wt[] = [4, 5, 6]$
Output: 0
Explanation: Every item has a weight exceeding the knapsack's capacity (3).

Input: $W = 5$, $val[] = [10, 40, 30, 50]$, $wt[] = [5, 4, 2, 3]$
Output: 80
Explanation: Choose the third item (value 30, weight 2) and the last item (value 50, weight 3) for a total value of 80.

Constraints:

$2 \leq val.size() = wt.size() \leq 10^3$
 $1 \leq W \leq 10^3$
 $1 \leq val[i] \leq 10^3$
 $1 \leq wt[i] \leq 10^3$

[Try more examples](#)

eg: $W = 4$
 $val[] = [1, 2, 3]$
 $wt[] = [4, 5, 1]$ } \rightarrow O/p: 3

we can either include an element or entirely exclude it.

parameters to be maintained.

i, W
 include: $help(i-1, W - wt[i]) + val[i]$
 exclude: $help(i-1, W)$

Base Condition:-

if ($i < 0$) return 0;

Brute Force Solution: T.C. $O(2^n)$

```
class Solution {
public:
    int knapsack(int W, vector<int> &val, vector<int> &wt) {
        // code here
        return help(val, wt, n-1, W);
    }
    int help(vector<int> &val, vector<int> &wt, int i, int W) {
        if (i < 0) return 0;
        int include = 0;
        if (W - wt[i] >= 0) include = val[i] + help(val, wt, i-1, W - wt[i]);
        int exclude = help(val, wt, i-1, W);
        return max(include, exclude);
    }
};
```

Optimized Solution:- Using DP + memoization:-

```

class Solution {
public:
    int knapsack(int W, vector<int> &val, vector<int> &wt) {
        // code here
        int n = val.size();
        vector<vector<int>> dp(n, vector<int>(W+1, -1));
        return help(val, wt, n-1, W, dp);
    }
    int help(vector<int> &val, vector<int> &wt, int i, int W, vector<vector<int>> &dp){
        if(i < 0) return 0;
        if(dp[i][W] != -1) return dp[i][W];
        int include = 0;
        if(W - wt[i] >= 0) include = val[i] + help(val, wt, i-1, W - wt[i], dp);
        int exclude = help(val, wt, i-1, W, dp);
        return dp[i][W] = max(include, exclude);
    }
};

```

T.C : $O(n \times W)$

S.C : $O(n \times W)$

Bottom-Up DP:-

[1, 2, 3]
[4, 5, 1]

W = 4

4, 5

	0	1	2	3	4
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	0	0	0
3	0	3	3	3	3

return dp[n][W];

include:-

for(int w = wt[i]; w <= W; w++) {
 $dp[i][w] = \max(dp[i-1][w],$
 $val[i] + dp[i-1][w - wt[i]])$

exclude:-
 $dp[i][w] = dp[i-1][w];$

```

class Solution {
public:
    int knapsack(int W, vector<int> &val, vector<int> &wt) {
        // code here
        int n = val.size();
        vector<vector<int>> dp(n+1, vector<int>(W+1, 0));
        for(int i = 1; i <= n; i++) {
            for(int j = 1; j <= W; j++) {
                if(j < wt[i]) {
                    dp[i][j] = dp[i-1][j];
                } else {
                    dp[i][j] = max(dp[i-1][j], val[i] + dp[i-1][j - wt[i]]);
                }
            }
        }
        return dp[n][W];
    }
};

```

Bottom-Up Soln \uparrow
 use val[i-1] and wt[i-1]