

Matrix Chain Multiplication :-

Given an array $arr[]$ which represents the dimensions of a sequence of matrices where the i^{th} matrix has the dimensions $(arr[i-1] \times arr[i])$ for $i=1$, find the most efficient way to multiply these matrices together. The efficient way is the one that involves the least number of multiplications.

Examples:

Input: $arr[] = [2, 1, 3, 4]$

Output: 20

Explanation: There are 3 matrices of dimensions 2×1 , 1×3 , and 3×4 . Let this 3 input matrices be M_1 , M_2 , and M_3 . There are two ways to multiply: $((M_1 \times M_2) \times M_3)$ and $(M_1 \times (M_2 \times M_3))$, note that the result of $(M_1 \times M_2)$ is a 2×3 matrix and result of $(M_2 \times M_3)$ is a 1×4 matrix.

$((M_1 \times M_2) \times M_3)$ requires $(2 \times 1 \times 3) + (2 \times 3 \times 4) = 30$

$(M_1 \times (M_2 \times M_3))$ requires $(1 \times 3 \times 4) + (2 \times 1 \times 4) = 20$.

The minimum of these two is 20.

Input: $arr[] = [1, 2, 3, 4, 3]$

Output: 30

Explanation: There are 4 matrices of dimensions 1×2 , 2×3 , 3×4 , and 4×3 . Let this 4 input matrices be M_1 , M_2 , M_3 and M_4 . The minimum number of multiplications are obtained by $((M_1 \times M_2) \times M_3) \times M_4$. The minimum number is $(1 \times 2 \times 3) + (1 \times 3 \times 4) + (1 \times 4 \times 3) = 30$.

Input: $arr[] = [3, 4]$

Output: 0

Explanation: As there is only one matrix so, there is no cost of multiplication.

Constraints:

$2 \leq arr.size() \leq 100$

$1 \leq arr[i] \leq 200$

[Try more examples](#)

$$arr[] = [2, 1, 3, 4] \rightarrow \text{size} = n$$

$$\text{mat } A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix} \begin{matrix} m \text{ rows} \\ n \text{ columns} \end{matrix} \rightarrow (n-1) \text{ matrices}$$

$$\text{dimension}(A) = m \times n$$

$$\text{dimension}(\text{mat}(i)) = arr[i-1] \times arr[i]$$

$$i = 1$$

$$\text{dim}(\text{mat}(1)) = arr[0] \times arr[1]$$

$$= 2 \times 1$$

$$\text{mat}(2) = 1 \times 3$$

$$\text{mat}(3) = 3 \times 4$$

$$\text{mat}(1) \times \text{mat}(2) \times \text{mat}(3) = ((\text{mat}(1) \times \text{mat}(2)) \times \text{mat}(3))$$

$$((\text{mat}(1) \times (\text{mat}(2) \times \text{mat}(3)))) \quad \text{arr}[i-1] \times arr[i]$$

The Brute Force Approach:-

$$(((\text{mat}(1) \times \text{mat}(2)) \times (\text{mat}(3) \times (\text{mat}(4) \dots \text{mat}(n)))))) \quad a \times b \times c$$

$$\rightarrow i \quad \text{for(int } j = i+1; j < n; j++)$$

$$ans = \min(\text{help}(i, j) + \text{help}(j+1, n) + i \times j \times n)$$

$$i \times j \times n$$

$$\begin{bmatrix} a & b \\ a & b \end{bmatrix} \begin{bmatrix} b & c \\ b & c \end{bmatrix} \rightarrow a \times b \times c$$

$$\begin{bmatrix} b & b & b \\ b & b & b \end{bmatrix} \rightarrow a \times c$$

Base Case:-

if $(i+1 == j)$ return 0;

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C++ (g++ 5.4)
1 // Driver Code Ends
2
3 class Solution {
4 public:
5     int matrixMultiplication(vector<int> &arr) {
6         // code here
7         int n = arr.size();
8         vector<vector<int>> dp(n, vector<int>(n, INT_MAX));
9         return help(0, n-1, arr, dp);
10    }
11
12    int help(int i, int j, vector<int> &arr, vector<vector<int>> &dp) {
13        if (dp[i][j] != -1) return dp[i][j];
14        if (i == j) return dp[i][j] = 0;
15        for (int k = i+1; k <= j; k++) {
16            dp[i][j] = min(dp[i][j], help(i, k, arr, dp) + help(k, j, arr, dp) + arr[i]*arr[k]*arr[j]);
17        }
18        return dp[i][j];
19    }
20 };
21
22
23
24
25
26
```