

Maximum Sum of Non-adjacent Nodes :-

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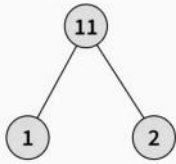


Difficulty: Medium Accuracy: 55.35% Submissions: 83K+ Points: 4 Average Time: 45m

Given a **binary tree** with a value associated with each node. Your task is to select a **subset of nodes** such that the sum of their values is **maximized**, with the condition that no two selected nodes are **directly connected** that is, if a node is included in the subset, neither its **parent** nor its **children** can be included.

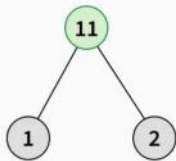
Examples:

Input: root[] = [11, 1, 2]

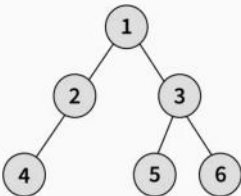


Output: 11

Explanation: The maximum sum is obtained by selecting the node 11.

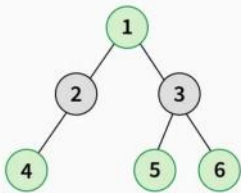


Input: root[] = [1, 2, 3, 4, N, 5, 6]



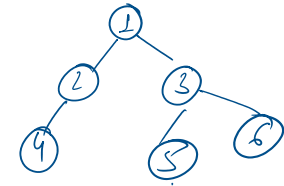
Output: 16

Explanation: The maximum sum is obtained by selecting the nodes 1, 4, 5, and 6, which are not directly connected to each other. Their total sum is 16.

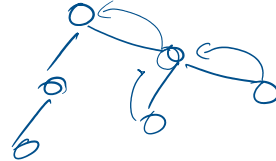
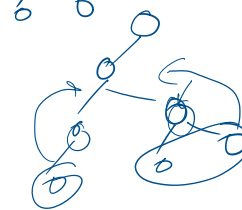


Constraints:

$1 \leq \text{no. of nodes in the tree} \leq 10^4$
 $1 \leq \text{Node.val} \leq 10^5$



def node: return root->data
 max(l+r, root->data)



[Naive Approach] Using Recursion - $O(2^n)$ Time and $O(n)$ Space

We can solve this problem by considering the fact that **both node and its immediate children** can't be in **sum** at the same time.

- **Include the current node's value in the sum:** In this case, we cannot include the values of its immediate children in the sum. Therefore, we recursively call the function on the **grandchildren** of the current node.
- **Exclude the current node's value in the sum:** In this case, we are allowed to include the values of its **immediate children** in the sum. So, we recursively call the function on the immediate children of the current node.
- Finally we will choose **maximum** from both of the results.

```
class Solution {
public:
    // Function to return the maximum sum of non-adjacent nodes.
    int getMaxSumUtil(Node* node) {
        if (node == nullptr) {
            // If the node is null, the sum is 0
            return 0;
        }

        // Calculate the maximum sum including the
        // current node
        int incl = node->data;

        // If the left child exists, include its contribution
        if (node->left) {
            incl += getMaxSumUtil(node->left->left) +
                getMaxSumUtil(node->left->right);
        }

        // If the right child exists, include its contribution
        if (node->right) {
            incl += getMaxSumUtil(node->right->left) +
                getMaxSumUtil(node->right->right);
        }

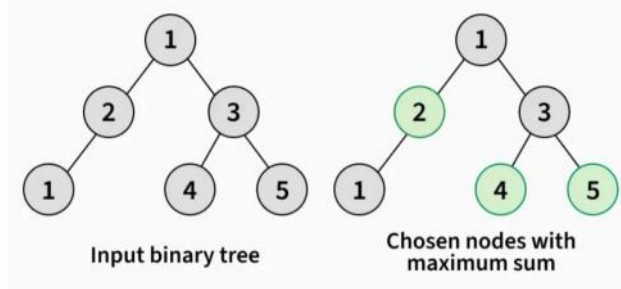
        // Calculate the maximum sum excluding
        // the current node
        int excl = 0;
        if (node->left) {
            excl += getMaxSumUtil(node->left);
        }
        if (node->right) {
            excl += getMaxSumUtil(node->right);
        }

        // The result for the current node is the
        // maximum of including or excluding it
        return max(incl, excl);
    }

    int getMaxSum(Node* root) {
        // If the tree is empty, the maximum sum is 0
        if (root == nullptr) {
            return 0;
        }

        // Call the utility function to compute the
        // maximum sum for the entire tree
        return getMaxSumUtil(root);
    }
};
```

[Expected Approach 1] Using Top-Down DP (Memoization) - $O(n)$ Time and $O(n)$ Space



The naive approach leads to **recalculating** results for the same nodes multiple times. For example, if we **include** the root node, we **recursively** compute the sum for its **grandchildren (nodes 4 and 5)**. But if we **exclude** the root, we compute the sum for its children, and **node 3** also computes the sum for its children (**4 and 5 again**).

To avoid this redundancy, we use **memoization**:

- We store the result of each node in a **hashmap**.
- When a node's value is needed again, we directly **return** it from the map instead of recalculating.

```

100 ~/
101
102 class Solution {
103 public:
104     // Function to return the maximum sum of non-adjacent nodes.
105     int getMaxSumUtil(Node* node, unordered_map<Node*, int>& memo) {
106         if (node == nullptr) {
107             // If the node is null, the sum is 0
108             return 0;
109         }
110
111         // If the result is already computed, return it from memo
112         if (memo.find(node) != memo.end()) {
113             return memo[node];
114         }
115
116         // Calculate the maximum sum including the current node
117         int incl = node->data;
118
119         // If the left child exists, include its grandchildren
120         if (node->left) {
121             incl += getMaxSumUtil(node->left->left, memo) +
122                 getMaxSumUtil(node->left->right, memo);
123         }
124
125         // If the right child exists, include its grandchildren
126         if (node->right) {
127             incl += getMaxSumUtil(node->right->left, memo) +
128                 getMaxSumUtil(node->right->right, memo);
129         }
130
131         // Calculate the maximum sum excluding the current node
132         int excl = getMaxSumUtil(node->left, memo) +
133             getMaxSumUtil(node->right, memo);
134
135         // Store the result in memo and return
136         // the maximum of incl and excl
137         memo[node] = max(incl, excl);
138         return memo[node];
139     }
140 }
141
142 // Function to compute the maximum
143 // sum of non-adjacent nodes
144 int getMaxSum(Node* root) {
145     unordered_map<Node*, int> memo;
146     return getMaxSumUtil(root, memo);
147 }
148
149 };
150

```

[Expected Approach 2] Using Pairs - O(n) Time and O(h) Space

Return a **pair** for each node in the binary tree such that the **first** of the pair indicates the **maximum sum** when the data of a node is **Included** and the second indicates the maximum sum when the data of a particular node is not included.

```

1 //
2 class Solution {
3 public:
4     pair<int, int> maxSumHelper(Node* root){
5         if(root==nullptr){
6             pair<int, int> sum(0, 0);
7             return sum;
8         }
9         pair<int, int> sum1=maxSumHelper(root->left);
10        pair<int, int> sum2=maxSumHelper(root->right);
11        pair<int, int> sum;
12        // This node is included(left and right children are not included)
13        sum.first=sum1.second+sum2.second+root->data;
14        // This node is excluded(either left or right child is included)
15        sum.second=max(sum1.first, sum1.second)+max(sum2.first, sum2.second);
16        return sum;
17    }
18    // Function to return the maximum sum of non-adjacent nodes.
19    int getMaxSum(Node *root) {
20        // code here
21        pair<int, int> res=maxSumHelper(root);
22        return max(res.first, res.second);
23    }
24 };
25

```