

Expectation Entropy as a Password Strength Metric

Khan Reaz, Gerhard Wunder
 Cybersecurity and AI Research Group
 Freie Universität Berlin, Germany
 Email: khanreaz@ieee.org, g.wunder@fu-berlin.de

Abstract—The classical combinatorics-based password strength formula provides a result in tens of bits, whereas the NIST Entropy Estimation Suite give a result between 0 and 1 for Min-entropy. In this work, we present a newly developed metric – *Expectation entropy* that can be applied to estimate the strength of any random or random-like password. *Expectation entropy* provides the strength of a password on the same scale as an entropy estimation tool. Having an *Expectation entropy* of a certain value, for example, 0.4 means that an attacker has to exhaustively search at least 40% of the total number of guesses to find the password.

Index Terms—Password, Entropy, Randomness, NIST Entropy Estimation Tool

I. MOTIVATION

The *recipe* for constructing a strong password has two ingredients: randomness, and length, i.e., the characters of the password must have high randomness (ideally truly random) and the number of characters should be large. A good combination of these properties can make brute-force attacks infeasible for that password.

State-of-the-art password strength estimation methods can be grouped into three main categories [2]: (1) Attack-based methods estimate the time it takes to break it with brute force. (2) Heuristic-based methods are based on Shannon’s notion of entropy bits. (3) Probabilistic-based methods consider human intrinsic nature, context, and imposed password composition rules. NIST’s Entropy Estimation Suite [3] is the industry standard solution to estimate min-entropy using 10 tests as described in [4], [5].

In our previous work [6], the device provisioning problem of Wi-Fi personal mode has been addressed and provided with *ComPass* protocol to supplement WPA2/WPA3. *ComPass* foregoes the pre-printed or user-generated password with an automatically generated strong symmetric password for the participating devices. The generated password is between 16 and 32 characters long and is generated by extracting signal parameters from typical Wi-Fi OFDM signals using Physical Layer Security methods. As we wanted to analyse the randomness of the quantise bit-string, and the strength of the generated password, we noticed that the classical combinatorics-based password strength formula: $\log_2(\text{character space}^{\text{length of the password}})$ provides the result in tens of bits, whereas entropy estimation

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formulae (and NIST entropy estimation tool) give a result between 0 and 1.

In this work, we present a newly developed metric – *Expectation entropy* that can be applied to estimate the strength of any random or random-like password. It captures the composition of the characters and estimate the strength from a single password. *Expectation entropy* provides the strength of the password on the same scale as entropy estimation formulae and NIST Entropy Estimation Suite.

II. VARIOUS DEFINITIONS OF ENTROPY

Entropy is a measure of the disorder, randomness or variability in a closed system. The larger the amount of entropy, the greater the uncertainty in predicting the value of an observation. It is usually denoted by H . If p_i is the probability of an element of a random discreet variable, then the min-entropy, H_∞ is the largest value m having the property that each observation of the variable guarantees at least m bits of information: $H_\infty = -\log_2(\max(p_i))$.

Shannon entropy, H_1 (or just H) was introduced by Claude Shannon in [7] as : $H_1 = -\sum_{i=1}^N p_i \log_2 p_i$, where $N = 2$ for binary digits, and $N = 26$ for English letters. Shannon entropy gives an average estimate and ignores the length of the password. Ralph Hartley proposed a quantitative measure of “information” two decades prior to Shannon [8]. Hartley entropy measures only the size of the distribution and ignores the probabilities: $H_0 = \log_2 N$.

However, in a password cracking case, an attacker must guess each value one at a time which makes all previously mentioned metrics unrelated to guessing difficulty [9]. Massey [10], Cachin [11], and Pliam [12] individually developed the *Guessing entropy* concept which states the expected number of guesses required if the attacker optimally tries: $G = \sum_{i=1}^N p_i \cdot i$.

III. EXPECTATION ENTROPY

Let us define four disjoint element sets such that \mathcal{L} be the set of lower-case letters, \mathcal{U} be the set of upper-case letters, \mathcal{D} be the set of digits, and \mathcal{S} be the set of symbols. Which means, $|\mathcal{L}| = 26$, $|\mathcal{U}| = 26$, $|\mathcal{D}| = 10$, $|\mathcal{S}| = 32$, and their disjoint union resulting in the total character space \mathcal{K} of cardinality $|\mathcal{K}| = |\mathcal{L}| + |\mathcal{U}| + |\mathcal{D}| + |\mathcal{S}| = 94$ for English language. Naturally, one can choose the character space for a different language.

A password P , where each of its character c is chosen uniformly at random, is called valid if it contains characters from at least two disjoint element sets, and the password length

$|P|$ satisfies $\min(|\mathcal{L}|, |\mathcal{U}|, |\mathcal{D}|, |\mathcal{S}|) \leq |P|$. Then the probability of a character from each element sets would be: $p_{\mathcal{L}} = \mathbb{P}(c \in \mathcal{L}) = \frac{26}{94}$, $p_{\mathcal{U}} = \mathbb{P}(c \in \mathcal{U}) = \frac{26}{94}$, $p_{\mathcal{D}} = \mathbb{P}(c \in \mathcal{D}) = \frac{10}{94}$, and $p_{\mathcal{S}} = \mathbb{P}(c \in \mathcal{S}) = \frac{32}{94}$.

If l, u, d, s are the number of characters chosen at random from $\mathcal{L}, \mathcal{U}, \mathcal{D}, \mathcal{S}$ sets respectively such that $l, u, d, s \leq |\mathcal{K}|$. We find the expectation of a character c appearing in a password as:

$$E(c(P)) = p_{\mathcal{L}} \cdot l + p_{\mathcal{U}} \cdot u + p_{\mathcal{D}} \cdot d + p_{\mathcal{S}} \cdot s \quad (1)$$

The maximum entropy (which is traditionally defined as Hartley entropy) of the total character space, \mathcal{K} is $H_0(\mathcal{K}) = \log_2 |\mathcal{K}|$. Then we express the *Expectation entropy* H_E of the password P as:

$$H_E(P) = \frac{\log_2 E(c(P))}{H_0(\mathcal{K})} \quad (2)$$

It is not difficult to show that the upper bound is achieved for the random variable $E(c(P))$ if and only if l, u, d, s are equal to $4 \cdot |\mathcal{K}|$. When the password length is larger than the restriction, H_E gives a value more than 1, and for the smaller length it gives a negative value.

IV. RESULTS

Empirical evaluation has been done using two different types of datasets. In the first case, 100,000 random passwords are generated using a computer with a dedicated hardware random number generator. These passwords are labelled with RandomMin, Random16ch, Random32ch, Random128ch, and RandomMax based on the length of the password. In the second case, three publicly leaked password databases (LinkedIn, 10Million, and WPA2) were used for evaluation.

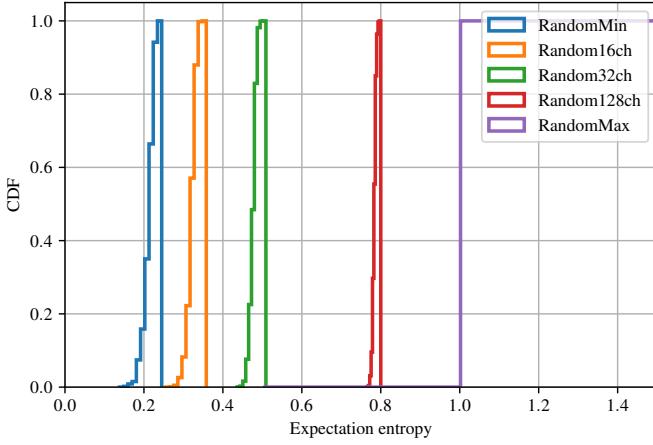


Fig. 1. Comparison of Expectation entropy for randomly generated passwords with different length.

Fig. 1, and Fig. 2 summarise the result in terms of Cumulative Distributive Function (CDF) and *Expectation entropy*. It can be observed from Fig. 1 that the value of H_E increases or decreases according to the length of the password and satisfies the bounds. Fig. 2 shows that the publicly leaked databases

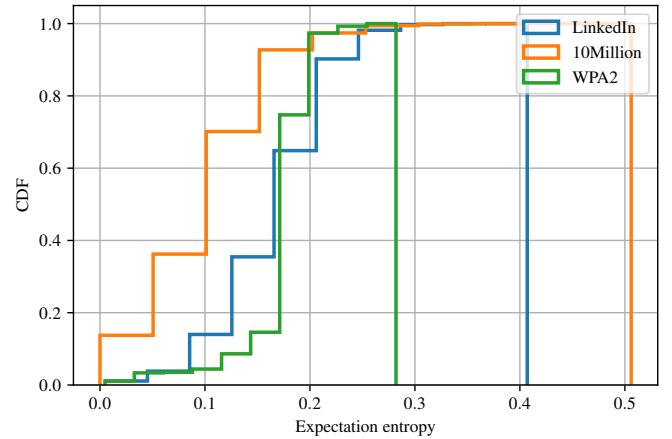


Fig. 2. Comparison of Expectation entropy for leaked passwords.

mostly contain passwords with a short length and the characters of the password are not chosen from all element sets, hence low *Expectation entropy*. Having an *Expectation entropy* of a certain value, for example, 0.4 means that an attacker has to exhaustively search at least 40% of the total number of guesses to find the password using brute-force.

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