

UNIT-3

Linked List

Linked List can be defined as collection of data elements called nodes that are randomly stored in the memory.

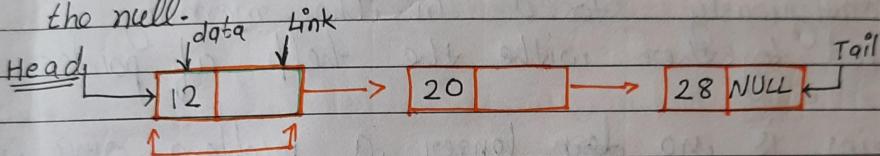
A node contain two fields



→ data stored at that particular address

→ the pointer which contain the address of the next node in the memory.

- The last node of the list contains ~~address~~ pointer to the null.

Uses of Linked List :-

- 1) The list is not required to be contiguously present in the memory. The node can be present/reside anywhere in the memory and linked together to make a list. This achieves optimized utilization of space.
- 2) List size is limited to memory size and doesn't need to be declare in advance.
- 3) Empty node can not be present in linked list.
- 4) We can store values of primitive types in the singly linked list.

Why use Linked List over array

An array contain the following limitations:-

- Size of array must be known in advance before using it in the program.

- 2) Increasing size of the array is time taking process. It is ~~also~~ almost impossible to expand the size of array at run time.
- (3) All the elements in array need to be contiguously stored in the memory, inserting any element in the array need shifting of all the predecessors.
- Linked List overcome all the limitations of an array. Using linked list is useful bcz
- 1) It allocates memory dynamically.
 - 2) All the nodes of linked list are non-contiguously stored in the memory & linked together with the help of pointer.
 - 3) Sizing is no longer a problem since we do not need to define its size at the time of declaration. List grows as per the program's demand and limited to the available memory space.

Types

- 1) Singly / one way Linked List
- 2) Doubly or two way Linked List
- 3) Circular Linked List.
- 4) Doubly Circular Linked List.

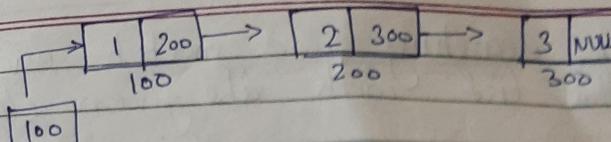
(1) Singly Linked List or one way Linked List \Rightarrow This is the simple type of linked list. The singly linked list is a DS that contains two parts i.e., data part & other one is address part also known as a pointer.

data/info



Link.

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head is first or start

- o It uses one external pointer ~~after~~ FIRST / HEAD for storing address of first node
- o The pointer of last node is set to NULL while other nodes contain the address of next nodes in their link.

Features \Rightarrow contain only single link.

- \rightarrow Only forward traversal is possible
- \rightarrow We can't traverse in backward direction.

Representation of the node in a singly linked list

```
struct node  
{ int data;  
  struct node *next;  
};
```

Operations on Singly Linked List :=>

- 1) Node Creation \Rightarrow A linked list is formed when many nodes are linked together to form a chain.
- 2) Insertion \Rightarrow Adding an element to the linked list.
- 3) Deletion \Rightarrow It involves deleting a node from the linked list.
- 4) Traversing \Rightarrow Traversing means we simply visit each node of the list at least once in order to perform some operation.

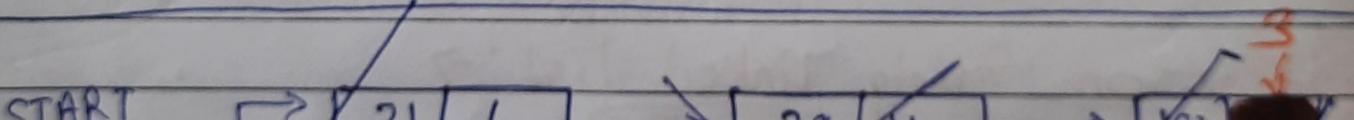
- 5) Searching \Rightarrow In searching, we match each element of the list with the given element.
- 6) Sorting \Rightarrow To arrange nodes in a specific order.
- 7) Updating \Rightarrow To update a node.

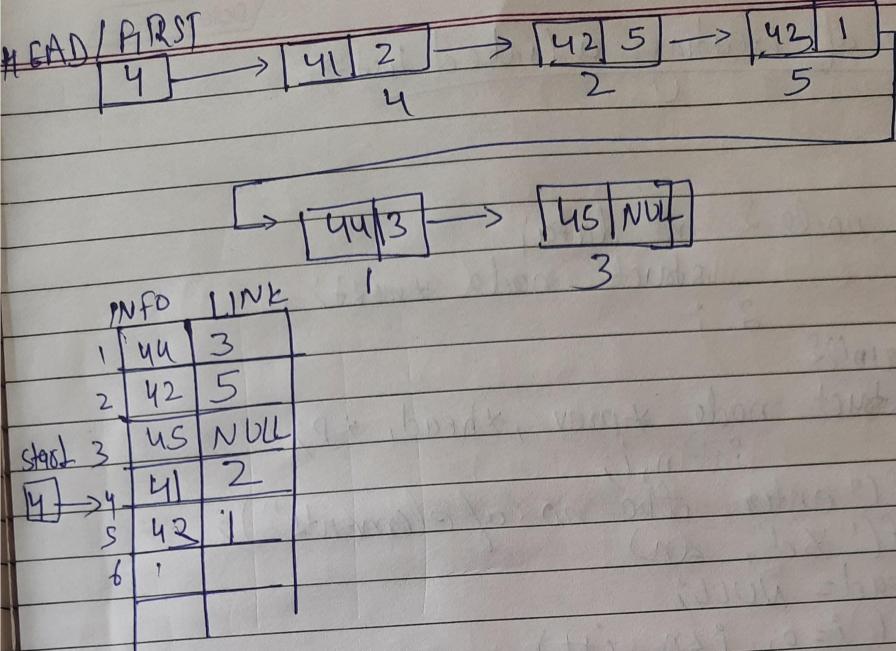
QUESTION Representation of Link List

Linked list can be represented in memory by using two array \Rightarrow data/info and link.

The linked list can be represented in memory by

- In linked list it is not necessary that the list must start from node [0]. Rather it can start from anywhere.





Let's create a node 'P'

```
# include < stdio.h >
# include < stdlib.h >
struct node {
    int data;
    struct node *next;
};

int main() {
    struct node *prev, *head, *P;
    P = (struct node*) malloc (sizeof (struct node));
    (struct node*)
    pointf ("Enter the data");
    scanf ("%d", &P->data);
    P->next = NULL;
    return 0;
}
```

3.

Example of creating a linked list

#

#

```
struct node { int data;  
    struct node *next};
```

{ };

void main()

```
struct node *prev, *head, *p;
```

```
int n, i;
```

```
printf("Enter the no. of elements");
```

```
scanf("%d", &n)
```

```
head = NULL;
```

```
for (i=0; i<n; i++)
```

{ };

```
p = (struct node*) malloc(sizeof(struct node))
```

```
scanf("%d", &p->data);
```

```
p->next = NULL
```

```
if (head == NULL)
```

```
{ head = p; }
```

{ };

```
else
```

```
prev->next = p;
```

```
prev = p;
```

{ };

{ };

Dynamic memory allocation → malloc()

```
pointer = (cast type *) malloc( byte-size )
```

```
int *ptr = (int *) malloc( sizeof(int) );
```

Traversal \Rightarrow

Algorithm \Rightarrow

- 1) set $\text{ptr} = \text{head}$;
- 2) if $\text{ptr} = \text{NULL}$
 write "empty List"
 goto step 6;
end of if

3. Repeat step 4 & 5 until $\text{ptr} = \text{NULL}$

step 4: print $\text{ptr} \rightarrow \text{data}$

step 5: $\text{ptr} = \text{ptr} \rightarrow \text{next}$
(end of loop)

step 6: exit.

Code \Rightarrow

```
#include <iostream.h>
struct node {
    int data;
    struct node* next;
};
```

```
void display (struct node *head)
{
    struct node *tmp;
    tmp = head;
    while (tmp != NULL)
    {
        printf ("%d", tmp->data);
        tmp = tmp->next;
    }
}
```

```
void main () {
    struct node *prev;
    int n;
}
```

printf(" enter the no. of elements");

scanf(" %d", &n);

head = NULL;

for (i=0; i<n; i++)

{

p=(struct node*)malloc (sizeof(struct node));

scanf(" %d", &p->data);

p->next = null;

if (head == NULL)

{ head = p; }

else

{ prev->next = p; }

prev = p;

{ display(head); }

}

Inserion :- There are 3 different possibilities for inserting a node into a linked list

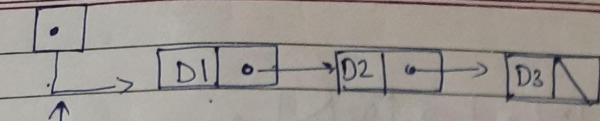
→ 1. Inserting at the beginning of the list.

→ 2. Inserting at the end of the list

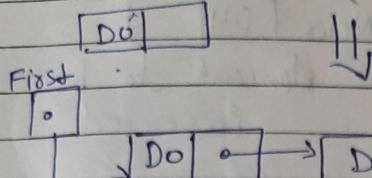
→ 3. Inserting a new node except the above mentioned position or we can say inbetween somewhere.

Case 1 Inserction at the beginning.

First



To be inserted here



New

Algorithm =) 1) Make a new node & set the fields

NEW ← NODE

INFO(NEW) ← X

2) Add elements and reset first.

LINK(NEW) ← FIRST

FIRST ← NEW

3. RETURN .

Note

* FIRST => is the external pointer containing address of first node.

* NEW => A pointer containing address of NEW node.

* X => Data to be stored in new node.

Code =>

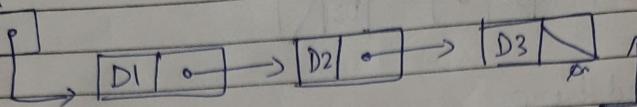
```

struct node *front( struct node *head, int value )
{
    struct node *P;
    P = (struct node*) malloc(sizeof(struct node));
    P->data = value;
    P->next = head;
    front(P); head = P;
}
  
```

3.

Case 2: Inserting at the end.

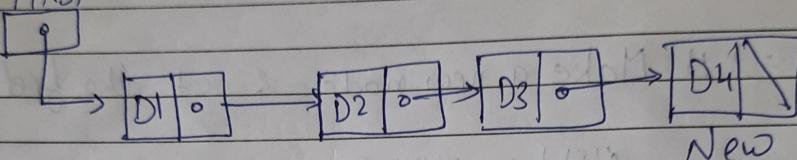
first



New [D4]

element to be inserted

FIRST



Algorithm: → 1. Get a new node & set the fields

New ← NODE

info(New) ← X

LINK(New) ← NULL

2. Loc ← FIRST.

3. //Pass all nodes & search the last node.

while (LINK(LOC) != NULL) etc

LOC ← LINK(LOC)

end while.

4. LINK(LOC) ← NEW //element inserted.

5. Return.

(*) Loc ⇒ Location of each node one by one.

Code :

void end(struct node *head, int value)

{

 struct node *p, *q;

 p = (struct node *) malloc (sizeof (struct node));

 p->data = value;

 p->next = NULL;

 q = head;

 while (q->next != NULL)

{

 q = q->next;

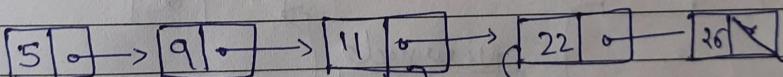
}

 q->next = p;

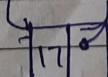
}

* Case 3: Insertion at place.

First

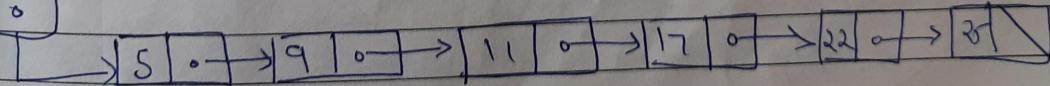


↓



New node to be inserted.

First



Algorithm *

Step1: Get new node & set fields

New \leftarrow ~~ptr~~ ptr.

Info(New) \leftarrow X

2. Set temp = head

3) set i = 0;

Set 4) Repeat step 5 & 6 until loc. ~~ptr~~ ==

5. ~~set~~ temp = temp \rightarrow next.

6) if temp = NULL

write desired node is not present
goto step 9.

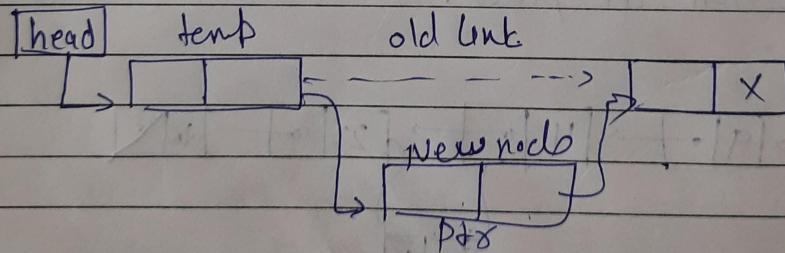
end of if

end of loop.

step7: ptr \leftarrow next = temp \rightarrow next
temp \rightarrow next = ptr.

step8 \Rightarrow set ptr = new

step9: exit



$ptr \rightarrow next = temp \rightarrow next$

$temp \rightarrow next = ptr$.

Algorithm: Inserting at nth node of lsd.

Code:

Step 1: \Rightarrow If $\text{ptr} = \text{NULL}$
write overflow
goto step
end of if

Step 2: set new_node = ptr.

Step 3: new_node \rightarrow data = data.

Step 4: set temp = head

Step 5: $i = 0$ set

Step 6: repeat step 5 & 6 until i

Step 7: temp = temp \rightarrow next.

Step 8: If temp = NULL

write desired node
is not found

goto step 12.

end of if

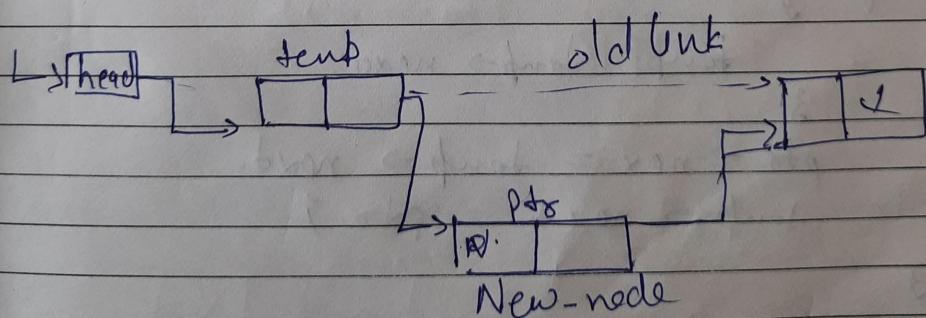
end of loop.

Step 9: ptr \rightarrow next = temp \rightarrow next.

Step 10: temp \rightarrow next = ptr

Step 11: set ptr = new_node.

Step 12: exit.



$\text{ptr} \rightarrow \text{next} = \text{temp} \rightarrow \text{next}$

$\text{temp} \rightarrow \text{next} = \text{ptr}$.

Code

```
# include <stdio.h>
# include <conio.h>
struct node *head = NULL;
struct node {
    int data;
    struct node *next;
};
```

void insert_pos_n(int data, int position)

{
 struct node *ptr = (struct node*) malloc
(sizeof(struct node));

ptr->data = data;

int i;

struct node *temp = head;

if (position == 1)

{
 ptr->next = temp;

head = ptr;

return;

for (i=1; i< position; i++)

{

temp = temp->next;

}

3

ptr->next = temp->next.

temp->next = ptr;

3

void display()

{
 struct node *temp = head;
 printf("\nList : ");

```
while (temp != NULL)
    { printf("In Y.d", temp->data);
    }
```

temp = temp->next;

```
int main()
```

```
    { int i, n, pos, data;
    struct node *temp, *prev, *head;
    printf("Enter the no of nodes");
    scanf("%d", &n);
    head = NULL;
    printf("Enter the elements");
    for (i=0; i<n; i++)
        {
```

```
        temp = (struct node *) malloc(sizeof(struct node));
        temp->next = NULL;
```

```
        if (head == NULL)
            head = temp;
        else
            { prev->next = temp;
            temp->prev = prev;
            prev = temp;
            }
```

```
    } // end of for
```

```
    printf("Enter the data you want to
    insert at position");
    scanf("%d %d", &data, &pos);
    ins_at_pos(data, pos);
```

```
    display();
    getch();
```

Deletion operations.

- 1) Delete at beginning.
- 2) Deletion at the end of the list.
- 3) Deletion after specified node.

(1) Deletion at beginning : \Rightarrow It involves deletion of a node from the beginning of the list.

Logic \Rightarrow $\left[\begin{array}{l} \text{ptr} = \text{head} \\ \text{head} = \text{ptr} \rightarrow \text{next} \\ \text{free}(\text{ptr}) \end{array} \right]$

free the pointer ptr which were pointing to the head node of the list.

Algorithm

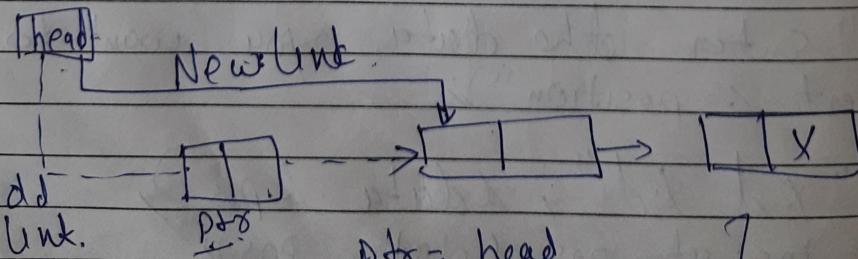
Step 1: if HEAD = NULL
 write underflow
 go to step 5
 (end of if)

Step 2: Set ptr = Head.

Step 3: set head = ptr \rightarrow next

Step 4: Free ptr;

Step 5: EXIT.



$\left[\begin{array}{l} \text{ptr} = \text{head} \\ \text{head} = \text{ptr} \rightarrow \text{next} \\ \text{free}(\text{ptr}) \end{array} \right]$

Code void begDelete()

{

struct node *ptr;
if (head == NULL)

{

printf("List is empty");

}

else { ptr = head;
head = ptr->next;
free(ptr);

}

3

2 Deletion at the end of the List => It involves deletion of a node from the end of the list.

Logic

ptr = head

while (ptr->next != NULL)

{ ptr1 = ptr

ptr = ptr->next;

3

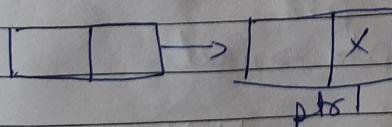
ptr1->next = NULL;

free(ptr);

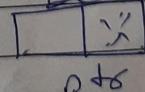
DeletedNode.

| Head |

L ->



ptr1->next = NULL
free(ptr)



Algorithm

Step 1: if head = NULL
 write underflow
 goto step 8
 (end of if)

Step 2: set ptr = head

Step 3: ~~repeat~~ repeat step 4 & 5 while
 $\text{ptr} \rightarrow \text{next} \neq \text{NULL}$

Step 4: set ptr1 = ptr

Step 5: set ptr = ptr \rightarrow next
 (end of loop)

Step 6: set ptr1 \rightarrow next = NULL

Step 7: free PTR.

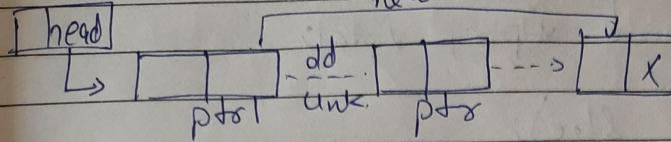
Step 8: exit.

Code

```

void end_delete() {
    struct node *ptr, *ptr1;
    if (head == NULL)
        printf ("List is empty");
    else if (head  $\rightarrow$  next == NULL)
        else
            ptr = head
            while (ptr  $\rightarrow$  next != NULL)
                ptr1 = ptr
                ptr = ptr  $\rightarrow$  next
                ptr1  $\rightarrow$  next = NULL
                free (ptr);
                printf ("Deleted node from list");
}
  
```

Deletion a node from specified position



$$\underline{ptr1 \rightarrow next} = \underline{ptr2 \rightarrow next}$$

`free(ptr2)`

Algorithm :

- step1 : if `head = NULL`
write underflow
goto step 10
end of if.

step2 : set `temp = head`

step3 : set `i = 0`

step4 : repeat step 5 to 8 until `i`

step5 : `temp1 = temp`

step6 : `temp = temp \rightarrow next`

step7 : if `temp = NULL`
write "desired node not
present"

goto step 12

end of if

step8 : `i = i + 1`

end of loop

step9 : `temp1 \rightarrow next = temp \rightarrow next`

step10 : `free temp`

step11 : exit

Code

```

void delete_specified()
{
    struct node *ptr, *ptr1;
    int loc, i;
    scanf("y.d", &loc);
    ptr = head;
    for (i = 0; i < loc; i++)
    {
        if (ptr == NULL)
            printf("There are less than %d elements in the list");
        return;
    }
    ptr1->next = ptr->next;
    free(ptr);
    printf("Deleted node is %d", loc);
}

```

Applications of Linked List:-

Linked lists are used in wide variety of applications because of their characteristic. They save wastage of memory. Some common application of linked list are:-

a) in computer Science:-

- 1) Implementing Stack.
- 2) Implementing Queue.
- 3) Implementing non-Linear structures
- 4) Maintaining directories
- 5) Performing arithmetic operations on

long integers.

6) Manipulating polynomials.

7) Representing sparse matrices etc.

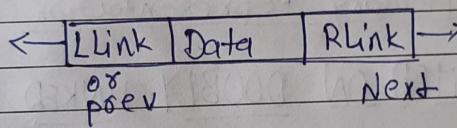
b) in real world :-

1) Image viewer.

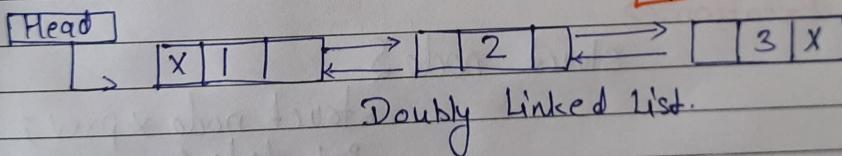
2) previous & next page in web browser.

3) Music player, etc.

Doubly Linked List :- Doubly linked list is a complex type of linked list in which a node contains 2 pointers to the previous as well as the next node in the sequence. A node consists of 3 parts.



The **Prev** part of first node & the **Next** part of the last node will always contain **NULL** indicates end of each direction.



In C, structure doubly linked list can be given as :-

struct node {

 struct node *prev;

 int data;

 struct node *next;

}

Memory Representation of a doubly Linked List =>

Generally, doubly linked list contain more space after every node & therefore cause more expensive basic operations.

Head



	Data	Prev	Next
1	13	-1	4
2			
3			
4	15	1	6
5			
6	19	4	8
7			
8	57	6	-1

Memory Representation

OPERATIONS PERFORMED ON DOUBLY LINKED LIST =>

1) Node Creation =>

struct node S

struct node *prev;

int data;

struct node *next;

3)

struct node *head;