## Performance

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Amount and Gustafson's law are a particular case of of the General Law:

$$S = \left[ \frac{\alpha + \frac{\beta}{\alpha} (1 - \alpha)}{\alpha + \frac{\beta}{\alpha} \left( \frac{1 - \alpha}{n \rho} \right)} \right]$$

B = factor that multiplies the time that can be parallelized

x = factor that multiplies the time that can't be parallel

the scaling and proportion of the two times are equal.

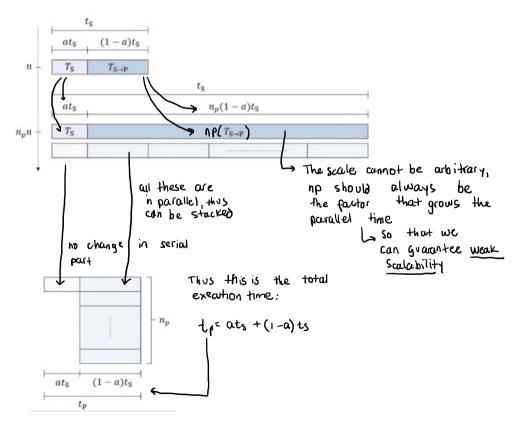
 $V = \frac{\beta}{\alpha} = n_p \rightarrow \text{The scaling factor is the amount}$ of processors

i.e. 
$$np=10$$
 $\gamma = \frac{6}{8} = \frac{np}{1}$ 

The parallel part is  $np$  times bigger than the serial part doesn't change

But to obtain 4 = np can involve othe values:

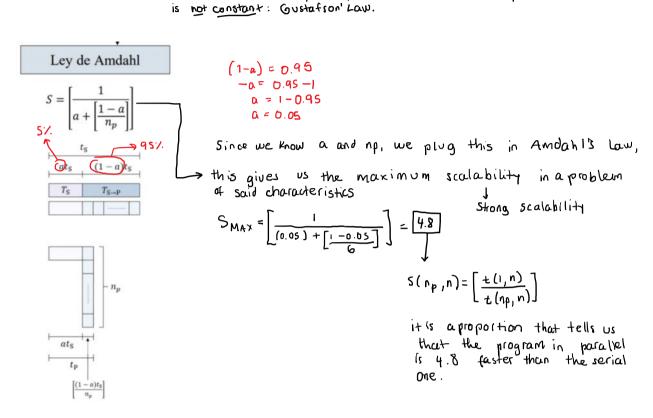
$$Y = \frac{12}{N} = \frac{20}{2}$$
 — But still the proportion is 1:10, thus parallel time is 10 times bigger than serial time



Exercises:

1) The 95% of execution time of a program, occurs inside a loop that we can parallelise. What is the maximum acceleration we can expect from the parallel program, executed in 6 processors?

95% to -> to = Total execution time of the program, with 1 processor = np = 1 = t, (serial time) 95% of to is parallelizable np= 6 n=constant -> We will use Amdahl's Law, since we only have 95%. of proportion, and they don't tell us anything about weak scalability. If they tell you that the size of the problem



10% of a program's execution time is spent within sequential code. What is the limit to the speedup achievable by a parallel version of the program. Suppose n is constant.

n = constant > so that we can use ambdahl's : سما  $S = \frac{1}{\alpha + (\frac{1-\alpha}{2n})}$ 

W=0.1 (10%)

np = since we don't have this, and we want the

Therefore, 
$$S_{MAx} = \lim_{np \to \infty} \left[ \frac{1}{a + \left(\frac{1-\alpha}{np}\right)} \right] = \frac{1}{a} = \frac{1}{0.1} = 10$$

if we got  $np \to \infty$  process
the fastest is 10 times
faster than the program in sequential mode.

if we got np ->00 processors liw yllaveu 2 = 012 xan2 smaller than 10.

(3) Example 1: Suppose we have a parallel program that is 15% serial and 85% linearly parallelizable. for a given problem size. Assume that the (absolute) serial time does not grow as the problem size (i) How much speedup can we achieve if we use 50 processors without scaling the problem? (ii) Suppose we scale up the problem size by a factor of 100. How much speedup could we achieve with 50 processors? this means < Ts does not that the slave

n ≠ constant of the problem n= noin) - as the mocessors alow, n change

$$n \neq constan+$$
 $\bar{n} = np(n) \rightarrow as the processors grow, n
grows by  $np(n)$$ 

of the problem will grow linearly.

Gustafson's Law used for problems were

The parallelizable time is scaled by np, while to does not change

(i) if n = constant and np = 50: Amdahl's Low

$$S = \frac{1}{\alpha + (\frac{1-\alpha}{np})}$$

 $S = \frac{1}{\alpha + (\frac{1-\alpha}{np})}$  We know that 15% is serial and 85% parallel,

We can now plug this into Amdahl's Law:

$$S = \frac{1}{0.15 + \left(\frac{1 - 0.15}{50}\right)} = 5.98$$

- (ii) Scale up the problem by a factor of 100. How much speedup (acceleration) can we have with 50 processors?
  - The scale does not follow n=hpxn, it grows as n = 100 × n

And therefore, since npxn is not ں ہوگی we cannot use Gustafson's Law, thus we need the general law:

$$S = \left[ \frac{\alpha + \frac{\beta_1}{\alpha} (1-\alpha)}{\alpha + \frac{\beta_2}{\alpha} (\frac{1-\alpha}{p_0})} \right]$$

v: multiplier for sequential part 15: multiplier for parallel part

d=1, since the problem says that to does not change  $\beta=100 \implies \beta T_{S\rightarrow p}=100 (T_{S\rightarrow p})$ a = 0.15 proportion that cont be parallelized

Thos,
$$S_{MAX} = \left[ \frac{0.15 + \frac{100}{1} \left( 1 - 0.15 \right)}{0.15 + \frac{100}{1} \left( \frac{1 - 0.15}{1 - 0.15} \right)} \right] = \frac{85.65}{1.85} = \frac{46.03}{1}$$

Example 2: Assume that you want to write a program that should achieve a speedup of 100 on 128 (4) processors.

(i) What is the maximum sequential fraction of the program when this speedup should be achieved under the assumption of strong scalability? -> Amdual's We start with Amdahl's law and then isolate f as follows:

Thus, only less than 1% of your program can be serial in the strong scaling scenario!

(ii) What is the maximum sequential fraction of the program when this speedup should be achieved under the assumption of weak scalability whereby the ratio  $\gamma$  scales linearly?

We now start with Gustafson's law and then isolate f as follows:

Thus, in this weak scaling scenario a significantly higher fraction can be serial!

Strong scalability means as up grows, n is constant, thus, Amdahl's Law.

001=2 no = 128 what is the maximum sequential fraction?

Is we look for a sequential fraction

$$100 = \left[ \frac{1}{\alpha + \left[ \frac{1-\alpha}{128} \right]} \right]$$

$$100 = \frac{1}{\alpha + \left(\frac{1-\alpha}{12\delta}\right)}$$

$$100 \left( a + \left( \frac{1-a}{12b} \right) \right) = 1$$

$$100a + 100\left(\frac{1-a}{128}\right) = 1$$

$$\frac{100 a + 100 - 100 a}{126} = 1$$

$$\frac{100a}{1} + \frac{100}{128} - \frac{100a}{128} = 1$$

$$\frac{126(100a)}{128} + \frac{100}{128} - \frac{100a}{128} = 1$$

$$Q = 0.0022$$
 $Q = 0.12^{\circ}/.$