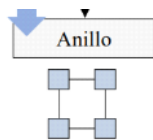


(b) Commuted - Direct - Ring  
Solution



Considering computation time complexity as:

$$t_c(n_p, n) = \left\lceil \frac{n}{n_p} + 1 \right\rceil + k_p = \left\lceil \frac{n}{n_p} + 1 \right\rceil + \log_2 n_p$$

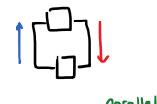
Case  
(serial)  $k_p = 0$   $2^{k_p} = 2^0 = 1$

Data/Diagram

time  
 $t = 0$

Tree (sums)

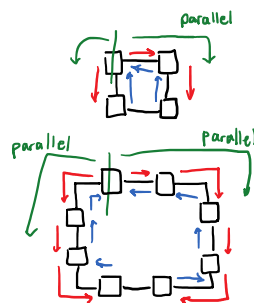
(parallel)  $k_p = 1$   $2^{k_p} = 2^1 = 2$



$$t = 6 = 1(2)(3) + \left( \frac{2}{2} - 1 \right) 1(2)(3)$$

1]

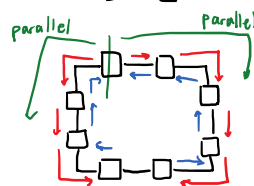
$k_p = 2$   $2^{k_p} = 2^2 = 4$



$$t = 1(2)(3) + 1(2)(3) = 12$$

2]

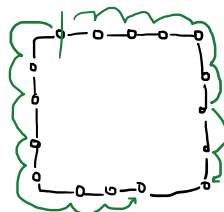
$k_p = 3$   $2^{k_p} = 2^3 = 8$



$$t = 1(2)(3) + 3(2)(3) = 24$$

3]

$k_p = 4$   $2^{k_p} = 2^4 = 16$



$$t = 1(16)(3) + \left( \frac{16}{2} - 1 \right) 1(2)(3)$$

4]

Thus, in a general way, we see that

$$t_c = (2)(3) + \left[ \frac{n_p}{2} - 1 \right] (2)(3)$$

$$t_c = 6 + \left[ \frac{n_p}{2} - 1 \right] 6$$

$$t_c = 6 + 6 \frac{n_p}{2} - 6$$

$$t_c(n_p, n) = \boxed{t_c = 3n_p}$$

for communication time complexity  
in a commuted - direct - ring network.