15-213 Recitation 15: Final Exam Preparation

25 April 2016

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Agenda

- Reminders
- Final Exam Review
- Fall 2012 exam

Reminders

- Proxy lab is due **tomorrow**!
- NO GRACE DAYS
- Penalty late days are allowed
- We will test your proxy manually
- We will read your code
- correctness: race conditions, robustness
- style: write clean, well-documented, modularized code – make it shine!
- Final exam is next week

Final Exam Details

- May 2 through 5
- sign-ups are open
- Eight to ten problems
- nominal time is 90-120 minutes, but you get five hours
- problems cover material from the entire semester
- Notes
- you are allowed two 8.5x11 double-sided sheets of notes
- no pre-worked problems allowed

Fall 2012 Final Exam – Multiple Choice (1)

- Each thread has its own ________.
- heap
- stack
- global values
- text data

- Simply decreasing the size of block headers used internally by malloc:
- decreases internal fragmentation
- increases internal fragmentation
- decreases external fragmentation
- increases external fragmentation

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- **stack**
- global values
- text data

- Simply decreasing the size of block headers used internally by malloc:
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- increases internal fragmentation
- decreases external fragmentation
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Fall 2012 Final Exam – Multiple Choice (2)

- Which of the following sentences about reader-writer locks is **not** true?
- Many readers can hold the same rwlock at the same time
- Two writers cannot hold the same rwlock at the same time
- Many readers and exactly one writer can hold the same rwlock at the same time
- An rwlock can be used as a mutex

- Which of the following is the correct ordering (left-to-right) of a file's compilation cycle?
- foo.c -> foo.o -> foo.s -> foo
- foo -> foo.s -> foo.o -> foo.c
- foo.c -> foo.s -> foo -> foo.o
- foo.c -> foo.s -> foo.o -> foo

Fall 2012 Final Exam – Multiple Choice (2)

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- foo.c -> foo.o -> foo.s -> foo
- foo -> foo.s -> foo.o -> foo.c
- foo.c -> foo.s -> foo -> foo.o
- **foo.c -> foo.s -> foo.o -> foo**

Fall 2012 Final Exam – Multiple Choice (3)

- Suppose an `int` A is stored at virtual address 0xff987cf0, while another `int` B is stored at virtual address 0xff987d98. If the size of a page is 0x1000 bytes, then A's physical address is numerically less than B's physical address.
- always true
- always false
- sometimes true, sometimes false
- not enough information

- Assuming no errors, which of the following functions returns exactly once?
- fork()
- execve()
- exit()
- longjmp()
- waitpid()

Fall 2012 Final Exam – Multiple Choice (3)

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- execve()
- exit()
- longjmp()
- **waitpid()**

returns twice

Fall 2012 Final Exam – Multiple Choice (4)

- On a 64-bit system, which of the following C expressions is equivalent to the C expression (x[2]+4)[3] ? Assume x is declared as int **x.
- *((*(x+16)) + 28)
- *((*(x + 2)) + 7)
- **(x * 28)
- *(((*x) + 2) + 7
- (**(x + 2) + 7)

- When can short counts occur?
- when an EOF is encountered during a read
- when a short int is used as a counter
- when writing to disk files
- when the kernel runs out of kernel memory

Fall 2012 Final Exam – Multiple Choice (4)

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- when the kernel runs out of kernel memory

Fall 2012 Final Exam – Multiple Choice (5)

- A program blocks SIGCHLD and SIGUSR1. It is then sent a SIGCHLD, a SIGUSR1, and another SIGCHLD, in that order. What signals does the program receive after it unblocks those signals (you may assume it is not sent any further signals afterward)?
- none, signals are discarded while blocked
- just a single SIGCHLD, since all subsequent signals are discarded

- Which of the following events does **not** generate a signal?
- division by zero
- a new connection arrives on a listening socket
- a write is attempted on a disconnected socket
- NULL is dereferenced
- a process whose parent has already terminated exits

Fall 2012 Final Exam – Multiple Choice (5)

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Fall 2012 Final Exam – Multiple Choice (6)

- In an x86-64 system, how many integers can be stored in a cache line if your cache is 4KB, is 4-way set associative, and contains 4 sets?
- 8
- 16
- 32
- 64
- 128

- What types of locality are leveraged by virtual memory?
- spatial locality
- temporal locality
- prime locality
- both spatial and temporal
- both temporal and prime locality

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- 32
- **64**
- 128

16 total cache lines
=
256 bytes per line

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- spatial locality
- temporal locality
- prime locality
- **both spatial and temporal**
- both temporal and prime locality

Fall 2012 Final Exam – Multiple Choice (7)

- Which of the following is **not** a section of an ELF file?
- .text
- .static
- .rodata
- .data
- .bss

- Choose the true statement.
- All thread-safe functions are reentrant.
- Some reentrant functions are not thread safe.
- It is never a good idea to use persistent state across multiple function calls
- It is impossible to have a race condition between two threads with no shared state.

Fall 2012 Final Exam – Multiple Choice (7)

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- It is never a good idea to use persistent state across multiple function calls
- **It is impossible to have a race condition between two threads with no shared state.**

not strictly true (*why?*) but appears in textbook and lecture notes

Fall 2012 Final Exam – Multiple Choice (8)

- We use dynamic memory because:
- the heap is significantly faster than the stack
- the stack is prone to corruption from buffer overflow
- storing data on the stack requires knowing the size of that data at compile time
- none of the above

Fall 2012 Final Exam – Multiple Choice (8)

- We use dynamic memory because:
- the heap is significantly faster than the stack
- the stack is prone to corruption from buffer overflow
- **storing data on the stack requires knowing the size of that data at compile time**
- none of the above

Fall 2012 Final Exam – Multiple Choice (9)

- In the following code, a parent opens a file twice, then the child reads a character:

```
char c;

int fd1 = open("foo.txt", O_RDONLY);
int fd2 = open("foo.txt", O_RDONLY);

if (!fork()) {
  read(fd1, &c, 1);
}
```

- Clearly, in the child, fd1 now points to the second character of foo.txt. Which of the following is now true in the parent?
- fd1 and fd2 both point to the first character
- fd1 and fd2 both point to the second character
- fd1 points to the first character, fd2 points to the second character
- fd2 points to the first character, fd1

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Fall 2012 Final Exam – Multiple Choice (10)

- Which of the following is true about races?
- A race occurs when correctness of the program depends on one thread reaching point A before another thread reaches point B.
- Exclusive access to all shared resources eliminates race conditions.
- Race conditions are the same as deadlocks.
- All race conditions occur inside

Fall 2012 Final Exam – Multiple Choice (10)

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Fall 2012 Final Exam – Multiple Choice (11)

- Consider the following two blocks of code, which are contained in *separate files*.

```
/* main.c */
int i = 0;
int main() {
  foo();
  return 0;
}
```

```
/* foo.c */
int i = 1;
void foo() {
  printf(“%d”, i);
}
```

- What will happen when you attempt to compile, link, and run this code?
- it will fail to compile
- it will fail to link
- it will raise a segmentation fault
- it will print “0”
- it will print “1”
- it will sometimes print “0” and sometimes print “1”

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- **it will fail to link**
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Fall 2012 Final Exam – Problem 2: Floating Point

- In this problem, you will work with floating point numbers based on the IEEE floating point format. We consider two different 6-bit formats:
- Format A:
- There is one sign bit S
- There are k=3 exponent bits. The bias is $2^{k-1} - 1 = 3$.
- there are n=2 fraction bits.
- Format B:
- There is one sign bit S
- Please write down the binary representation for the following (use round-to-even).
- Recall that for denormalized numbers, E = 1-bias. For normalized numbers, E = e-bias.

Value	Format A Bits	Format B Bits
One	0 011 00	0 01 000
Three		
7/8		
15/8		

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Value	Format A Bits	Format B Bits
One	0 011 00	0 01 000
Three	**0 100 10**	**0 10 100**
7/8	**0 010 11**	**0 00 111**
15/8	**0 100 00**	**0 01 111**

both exact

both exact
A norm
B denorm

A round-to-even, B exact

Problem 3. (6 points):

Arrays. Consider the C code below, where H and J are constants declared with #define .

```
int array1[H][J];
int array2[J][H];

void copy_array(int x, int y) {
        array2[x][y] = array1[y][x];
}
```

Suppose the above C code generates the following x86-64 assembly code:

```
# On entry:
#       %edi = x
#       %esi = y
#
copy_array:
        movslq    %esi,%rsi
        movslq    %edi,%rdi
        movq      %rdi, %rax
        salq      $4, %rax
        subq      %rdi, %rax
        addq      %rsi, %rax
        leaq      (%rsi,%rsi,4), %rsi
        leaq      (%rdi,%rsi,2), %rsi
        movl      array1(,%rsi,4), %edx
        movl      %edx, array2(,%rax,4)
        ret
```

What are the values of H and J ?

H =

J =

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        addq      %rsi, %rax
        leaq      (%rsi,%rsi,4), %rsi
        leaq      (%rdi,%rsi,2), %rsi
        movl      array1(,%rsi,4), %edx
        movl      %edx, array2(,%rax,4)
        ret
```

rax = 16*x – x + y
rax = 15*x + y

rsi = 5 * y

rsi = x + 2*(5*y)

What are the values of H and J ?

H = **15**

J = **10**

Problem 4. (8 points):

Loops. Consider the following x86-64 assembly function:

```
loop:
        # on entry: a in %rdi, n in %esi
        movl        $0, %r8d
        movl        $0, %ecx
        testl       %esi, %esi
        jle .L3
.L6:
        movl        (%rdi,%rcx,4), %edx
        leal        3(%rdx), %eax
        testl       %edx, %edx
        cmovns      %edx, %eax
        sarl        $2, %eax
        addl        %eax, %r8d
        addq        $1, %rcx
        cmpl        %ecx, %esi
        jg      .L6
.L3:
        movl        %r8d, %eax
        ret
```

```
int loop(int a[], int n)
{
        int i, sum;

        sum = _____;

        for (i = ____________; ____________; ____________) {

                sum += ______________;

        }

        return ____________;
}
```

Fill in the blanks of the corresponding C code.

- You may only use the C variable names `n`, `a`, `i` and `sum`, not register names.
- Use array notation in showing accesses or updates to elements of `a`.

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loop:
        # on entry: a in %rdi, n in %esi
        movl        $0, %r8d
        movl        $0, %ecx
        testl       %esi, %esi
        jle .L3
.L6:
        movl        (%rdi,%rcx,4), %edx
        leal        3(%rdx), %eax
        testl       %edx, %edx
        cmovns      %edx, %eax
        sarl        $2, %eax
        addl        %eax, %r8d
        addq        $1, %rcx
        cmpl        %ecx, %esi
        jg     .L6
.L3:
        movl        %r8d, %eax
        ret
```

eax = a[i]+3

edx = a[i]

edx zero or negative?

move if non-negative

```
int loop(int a[], int n)
{
        int i, sum;

        sum = __0__;

        for (i = ____0____; ____i < n____; ____i++____) {

                sum += ____(a[i] < 0 ? a[i]+3 : a[i]) >> 2____;

        }

        return ____sum____;
}
```

Fill in the blanks of the corresponding C code.

- You may only use the C variable names n, a, i and sum, not register names.
- Use array notation in showing accesses or updates to elements of a.

negative integers must be biased before right shift to divide by a power of two with correct rounding

Problem 5. (10 points):

Stack discipline. Consider the following C code and its corresponding 32-bit x86 machine code. Please complete the stack diagram on the following page.

```
int fact(int n) {
        if (n == 1)
                return n;
        else
                return n      *  fact(n-1);
}
```

```
080483a4 <fact>:
  80483a4:      55                      push    %ebp
  80483a5:      89 e5                   mov     %esp,%ebp
  80483a7:      53                      push    %ebx
  80483a8:      83 ec 04                sub     $0x4,%esp
  80483ab:      8b 5d 08                mov     0x8(%ebp),%ebx
  80483ae:      83 fb 01                cmp     $0x1,%ebx
  80483b1:      74 0e                   je      80483c1 <fact+0x1d>
  80483b3:      8d 43 ff                lea     0xffffffff(%ebx),%eax
  80483b6:      89 04 24                mov     %eax,(%esp)
  80483b9:      e8 e6 ff ff ff          call    80483a4 <fact>
  80483be:      0f af d8                imul    %eax,%ebx
  80483c1:      89 d8                   mov     %ebx,%eax
  80483c3:      83 c4 04                add     $0x4,%esp
  80483c6:      5b                      pop     %ebx
  80483c7:      5d                      pop     %ebp
  80483c8:      c3                      ret
```

32-bit
stack
convention

A. Draw a detailed picture of the stack, starting with the caller invoking fact(4), and ending immediately **before** the call instruction that invokes fact(2).

- The stack diagram should begin with the argument for fact that the caller has placed on the stack. To help you get started, we have given you the first one.
- Use the actual values for function arguments, rather than variable names. For example, use 3 or 2 instead of n.
- For callee-saved registers that are pushed to the stack, simply note the register name (e.g, %ebx).
- Always label %ebp and give its value when it is pushed to the stack, e.g., old %ebp: 0xffff1400 .

Value of %ebp when fact(4) is called: 0xffffd848
Return address in function that called fact(4): 0x080483e6

this problem
is obsolete

```
Stack          The diagram starts with the
addresss       argument for fact(4)
               +-----------------------------------+
0xffffd830 |                    4                          |
               +-----------------------------------+
0xffffd82c |                                               |
               +-----------------------------------+
0xffffd828 |                                               |
               +-----------------------------------+
0xffffd824 |                                               |
               +-----------------------------------+
0xffffd820 |                                               |
               +-----------------------------------+
0xffffd81c |                                               |
               +-----------------------------------+
0xffffd818 |                                               |
               +-----------------------------------+
0xffffd814 |                                               |
               +-----------------------------------+
0xffffd810 |                                               |
               +-----------------------------------+
```

the answer will differ from what you learned for 64-bit

B. What is the final value of %ebp, immediately **before** execution of the instruction that calls fact(2) ?

%ebp=0x______________________

C. What is the final value of %esp, immediately **before** execution of the instruction that calls fact(2) ?

%esp=0x_____________________

```
Stack          The diagram starts with the
addresss       argument for fact(4)
               +-----------------------------------+
0xffffd830 |                 4                              |
               +-----------------------------------+
0xffffd82c |      caller ra: 0x080483e6                     |
               +-----------------------------------+
0xffffd828 |      old ebp: 0xffffd848                       |
               +-----------------------------------+
0xffffd824 |      ebx   (4)                                 |
               +-----------------------------------+
0xffffd820 |      3                                         |
               +-----------------------------------+
0xffffd81c |      caller ra: 0x8048ebe                      |
               +-----------------------------------+
0xffffd818 |      old ebp: 0xffffd828                       |
               +-----------------------------------+
0xffffd814 |      ebx   (3)                                 |
               +-----------------------------------+
0xffffd810 |      2                                         |
               +-----------------------------------+
```

B. What is the final value of %ebp, immediately **before** execution of the instruction that calls fact(2) ?

%ebp=0x___ffffd818___

C. What is the final value of %esp, immediately **before** execution of the instruction that calls fact(2) ?

%esp=0x___ffffd810___

Problem 6. (12 points):

Cache memories. Consider the following matrix transpose function

```
typedef int array[2][2];

void transpose(array dst, array src) {
    int i, j;

    for (j = 0; j < 2; j++) {
        for (i = 0; i < 2; i++) {
            dst[i][j] = src[j][i];
        }
    }
}
```

running on a hypothetical machine with the following properties:

- sizeof(int) == 4 .
- The src array starts at address 0 and the dst array starts at address 16 (decimal).
- There is a single L1 data cache that is direct mapped and write-allocate, with a block size of 8 bytes.
- Accesses to the src and dst arrays are the only sources of read and write accesses to the cache, respectively.

A. Suppose the cache has a total size of 16 data bytes (i.e., the block size times the number of sets is 16 bytes) and that the cache is initially empty. Then for each row and col, indicate whether each access to src[row][col] and dst[row][col] is a hit (h) or a miss (m). For example, reading src[0][0] is a miss and writing dst[0][0] is also a miss.

src array		
	col 0	col 1
row 0	m	
row 1		

dst array		
	col 0	col 1
row 0	m	
row 1		

B. Repeat part A for a cache with a total size of 32 data bytes.

src array		
	col 0	col 1
row 0	m	
row 1		

dst array		
	col 0	col 1
row 0	m	
row 1		

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src array		
	col 0	col 1
row 0	m	m
row 1	m	h

dst array		
	col 0	col 1
row 0	m	m
row 1	m	m

B. Repeat part A for a cache with a total size of 32 data bytes.

src array		
	col 0	col 1
row 0	m	h
row 1	m	h

dst array		
	col 0	col 1
row 0	m	h
row 1	m	h

Problem 7. (6 points):

Linking. Consider the executable object file a.out , which is compiled and linked using the command

```
unix> gcc -o a.out main.c foo.c
```

and where the files main.c and foo.c consist of the following code:

```
/* main.c */
#include <stdio.h>

int a = 1;
static int b = 2;
int c = 3;

int main()
{
        int c = 4;

        foo();
        printf("a=%d b=%d c=%d\n", a, b, c);
        return 0;
}
```

```
/* foo.c */
int a, b, c;

void foo()
{
        a = 5;
        b = 6;
        c = 7;
}
```

What is the output of a.out ?

Answer: a=____, b=_____, c=_____

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What is the output of a.out ?

Answer: a= 5 , b= 2 , c= 4

Problem 8. (10 points):

Exceptional control flow. Consider the following C program. (For space reasons, we are not checking error return codes, so assume that all functions return normally.)

```
int main()
{
        int val = 2;

        printf("%d", 0);
        fflush(stdout);

        if (fork() == 0) {
                val++;
                printf("%d", val);
                fflush(stdout);
        }
        else {
                val--;
                printf("%d", val);
                fflush(stdout);
                wait(NULL);
        }
        val++;
        printf("%d", val);
        fflush(stdout);
        exit(0);
}
```

For each of the following strings, circle whether (Y) or not (N) this string is a possible output of the program. You will be graded on each sub-problem as follows:

- If you circle no answer, you get 0 points.
- If you circle the right answer, you get 2 points.
- If you circle the wrong answer, you get −1 points (so don't just guess wildly).

A.	01432	Y	N
B.	01342	Y	N
C.	03142	Y	N
D.	01234	Y	N
E.	03412	Y	N

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        val++;
        printf("%d", val);
        fflush(stdout);
    }
    else {
        val--;
        printf("%d", val);
        fflush(stdout);
        wait(NULL);
    }
    val++;
    printf("%d", val);
    fflush(stdout);
    exit(0);
}
```

For each of the following strings, circle whether (Y) or not (N) this string is a possible output of the program. You will be graded on each sub-problem as follows:

- If you circle no answer, you get 0 points.
- If you circle the right answer, you get 2 points.
- If you circle the wrong answer, you get −1 points (so don't just guess wildly).

A. 01432 Y [✓ N]

B. 01342 [✓ Y] N

C. 03142 [✓ Y] N

D. 01234 Y [✓ N]

E. 03412 [✓ Y] N

Problem 9. (12 points):

Address translation. This problem concerns the way virtual addresses are translated into physical addresses.
Imagine a system has the following parameters:

- Virtual addresses are 20 bits wide.
- Physical addresses are 18 bits wide.
- The page size is 1024 bytes.
- The TLB is 2-way set associative with 16 total entries.

The contents of the TLB and the first 32 entries of the page table are shown as follows. **All numbers are given in hexadecimal**.

TLB			
Index	Tag	PPN	Valid
0	03	C3	1
	01	71	0
1	00	28	1
	01	35	1
2	02	68	1
	3A	F1	0
3	03	12	1
	02	30	1
4	7F	05	0
	01	A1	0
5	00	53	1
	03	4E	1
6	1B	34	0
	00	1F	1
7	03	38	1
	32	09	0

Page Table					
VPN	PPN	Valid	VPN	PPN	Valid
000	71	1	010	60	0
001	28	1	011	57	0
002	93	1	012	68	1
003	AB	0	013	30	1
004	D6	0	014	0D	0
005	53	1	015	2B	0
006	1F	1	016	9F	0
007	80	1	017	62	0
008	02	0	018	C3	1
009	35	1	019	04	0
00A	41	0	01A	F1	1
00B	86	1	01B	12	1
00C	A1	1	01C	30	0
00D	D5	1	01D	4E	1
00E	8E	0	01E	57	1
00F	D4	0	01F	38	1

Part 1

1. The diagram below shows the format of a virtual address. Please indicate the following fields by labeling the diagram:

VPO	The virtual page offset
VPN	The virtual page number
TLBI	The TLB index
TLBT	The TLB tag

19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

2. The diagram below shows the format of a physical address. Please indicate the following fields by labeling the diagram:

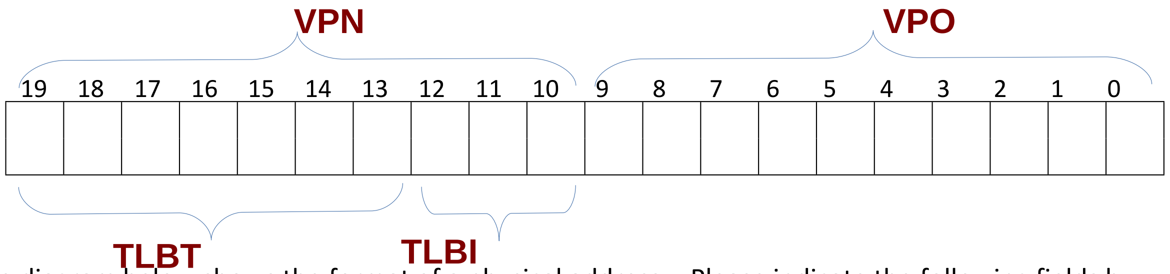
PPO	The physical page offset
PPN	The physical page number

17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

Part 1

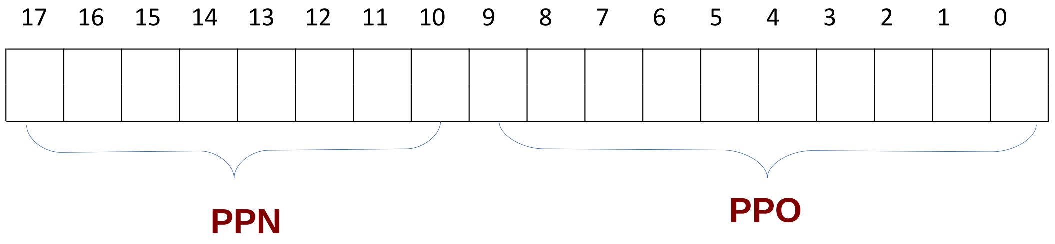
1. The diagram below shows the format of a virtual address. Please indicate the following fields by labeling the diagram:

VPO	The virtual page offset
VPN	The virtual page number
TLBI	The TLB index
TLBT	The TLB tag



2. The diagram below shows the format of a physical address. Please indicate the following fields by labeling the diagram:

PPO	The physical page offset
PPN	The physical page number



Part 2

For the given virtual addresses, please indicate the TLB entry accessed and the physical address. Indicate whether the TLB misses and whether a page fault occurs. If there is a page fault, enter "-" for "PPN" and leave the physical address blank.

Virtual address: 078E6

1. Virtual address (one bit per box)

19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

2. Address translation

Parameter	Value	Parameter	Value
VPN	0x	TLB Hit? (Y/N)	
TLB Index	0x	Page Fault? (Y/N)	
TLB Tag	0x	PPN	0x

3. Physical address(one bit per box)

17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

Part 2

For the given virtual addresses, please indicate the TLB entry accessed and the physical address. Indicate whether the TLB misses and whether a page fault occurs. If there is a page fault, enter "-" for "PPN" and leave the physical address blank.

Virtual address: 078E6

1. Virtual address (one bit per box) **0000 0111 1000 1110 0110**

19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

2. Address translation

Parameter	Value	Parameter	Value
VPN	0x **01E**	TLB Hit? (Y/N)	**N**
TLB Index	0x **6**	Page Fault? (Y/N)	**N**
TLB Tag	0x **03**	PPN	0x **57**

3. Physical address(one bit per box)

01 0101 1100 1110 0110

17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

Virtual address: 04AA4

1. Virtual address (one bit per box)

19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

2. Address translation

Parameter	Value	Parameter	Value
VPN	0x	TLB Hit? (Y/N)	
TLB Index	0x	Page Fault? (Y/N)	
TLB Tag	0x	PPN	0x

3. Physical address(one bit per box)

17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

Virtual address: 04AA4

1. Virtual address (one bit per box)

0000 0100 1010 1010 0100

19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

2. Address translation

Parameter	Value	Parameter	Value
VPN	0x **012**	TLB Hit? (Y/N)	**Y**
TLB Index	0x **2**	Page Fault? (Y/N)	**N**
TLB Tag	0x **02**	PPN	0x **68**

3. Physical address(one bit per box)

01 1010 0010 1010 0100

17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

Problem 10. (10 points):

Concurrency, races, and synchronization. Consider a simple concurrent program with the following specification: The main thread creates two peer threads, passing each peer thread a unique integer *thread ID* (either 0 or 1), and then waits for each thread to terminate. Each peer thread prints its thread ID and then terminates.

Each of the following programs attempts to implement this specification. However, some are incorrect because they contain a race on the value of myid that makes it possible for one or more peer threads to print an incorrect thread ID. Except for the race, each program is otherwise correct.

You are to indicate whether or not each of the following programs contains such a race on the value of myid . You will be graded on each subproblem as follows:

- If you circle no answer, you get 0 points.
- If you circle the right answer, you get 2 points.
- If you circle the wrong answer, you get −1 points (so don't just guess wildly).

A. **Does the following program contain a race on the value of myid ?**

Yes

No

```
void *foo(void *vargp) {
    int myid;
    myid = *((int *)vargp);
    Free(vargp);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i, *ptr;

    for (i = 0; i < 2; i++) {
        ptr = Malloc(sizeof(int));
        *ptr = i;
        Pthread_create(&tid[i], 0, foo, ptr);
    }
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

B. **Does the following program contain a race on the value of myid ?**

Yes

No

```
void *foo(void *vargp) {
    int myid;
    myid = *((int *)vargp);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    for (i = 0; i < 2; i++)
        Pthread_create(&tid[i], NULL, foo, &i);
    Pthread_join(tid[0], NULL);
    Pthread_join(tid[1], NULL);
}
```

A. **Does the following program contain a race on the value of myid ?** Yes [✓ No]

```
void *foo(void *vargp) {
    int myid;
    myid = *((int *)vargp);
    Free(vargp);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i, *ptr;

    for (i = 0; i < 2; i++) {
        ptr = Malloc(sizeof(int));
        *ptr = i;
        Pthread_create(&tid[i], 0, foo, ptr);
    }
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

separate heap variable for each thread

B. **Does the following program contain a race on the value of myid ?** [✓ Yes] No

```
void *foo(void *vargp) {
    int myid;
    myid = *((int *)vargp);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    for (i = 0; i < 2; i++)
        Pthread_create(&tid[i], NULL, foo, &i);
    Pthread_join(tid[0], NULL);
    Pthread_join(tid[1], NULL);
}
```

C. **Does the following program contain a race on the value of myid ?**

Yes

No

```
void *foo(void *vargp) {
    int myid;
    myid = (int)vargp;
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    for (i = 0; i < 2; i++)
        Pthread_create(&tid[i], 0, foo, i);
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

D. **Does the following program contain a race on the value of myid ?**

Yes

No

```
sem_t s; /* semaphore s */

void *foo(void *vargp) {
    int myid;
    P(&s);
    myid = *((int *)vargp);
    V(&s);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    sem_init(&s, 0, 1); /* S=1 INITIALLY */

    for (i = 0; i < 2; i++) {
        Pthread_create(&tid[i], 0, foo, &i);
    }
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

C. **Does the following program contain a race on the value of myid ?** Yes [✓ No]

```
void *foo(void *vargp) {
    int myid;
    myid = (int)vargp;
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    for (i = 0; i < 2; i++)
        Pthread_create(&tid[i], 0, foo, i);
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

myid is passed directly on the stack

D. **Does the following program contain a race on the value of myid ?** [✓ Yes] No

```
sem_t s; /* semaphore s */

void *foo(void *vargp) {
    int myid;
    P(&s);
    myid = *((int *)vargp);
    V(&s);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    sem_init(&s, 0, 1); /* S=1 INITIALLY */

    for (i = 0; i < 2; i++) {
        Pthread_create(&tid[i], 0, foo, &i);
    }
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

the mutex doesn't actually protect myid (*why?*)

E. **Does the following program contain a race on the value of myid ?** Yes No

```
sem_t s; /  *  semaphore s  */

void  *foo(void  *vargp) {
        int myid;
        myid =  *((int  *)vargp);
        V(&s);
        printf("Thread %d\n", myid);
}

int main() {
        pthread_t tid[2];
        int i;

        sem_init(&s, 0, 0); /  *  S=0 INITIALLY  */

        for (i = 0; i < 2; i++) {
                Pthread_create(&tid[i], 0, foo, &i);
                P(&s);
        }
        Pthread_join(tid[0], 0);
        Pthread_join(tid[1], 0);
}
```

E. **Does the following program contain a race on the value of myid ?** Yes [✓ No]

```
sem_t s; /    *  semaphore s    */

void   *foo(void     *vargp) {
    int myid;
    myid =    *((int     *)vargp);
    V(&s);
    printf("Thread %d\n", myid);
}

int main() {
    pthread_t tid[2];
    int i;

    sem_init(&s, 0, 0); /         *  S=0 INITIALLY       */

    for (i = 0; i < 2; i++) {
        Pthread_create(&tid[i], 0, foo, &i);
        P(&s);
    }
    Pthread_join(tid[0], 0);
    Pthread_join(tid[1], 0);
}
```

this mutex **does** protect myid (*why?*)