

④ Subsequences of strings

You are given a string 'STR' containing lowercase English letters from a to z inclusive. Your task is to find all non-empty possible subsequences of 'STR'.

A Subsequence of a string is the one which is generated by deleting 0 or more letters from the string and keeping the rest of the letters in the same order.

Detailed explanation (Input/output format, Notes, Images)

Constraints:

1 <= T <= 10
1 <= |STR| <= 16

Where |STR| represents the length of the string 'STR'.

Time Limit: 1 sec

Sample Input 1:

1
abc

Sample Output 1:

a ab abc ac b bc c

Explanation of sample input 1:

All possible subsequences of abc are :
"a", "b", "c", "ab", "bc", "ac", "abc"

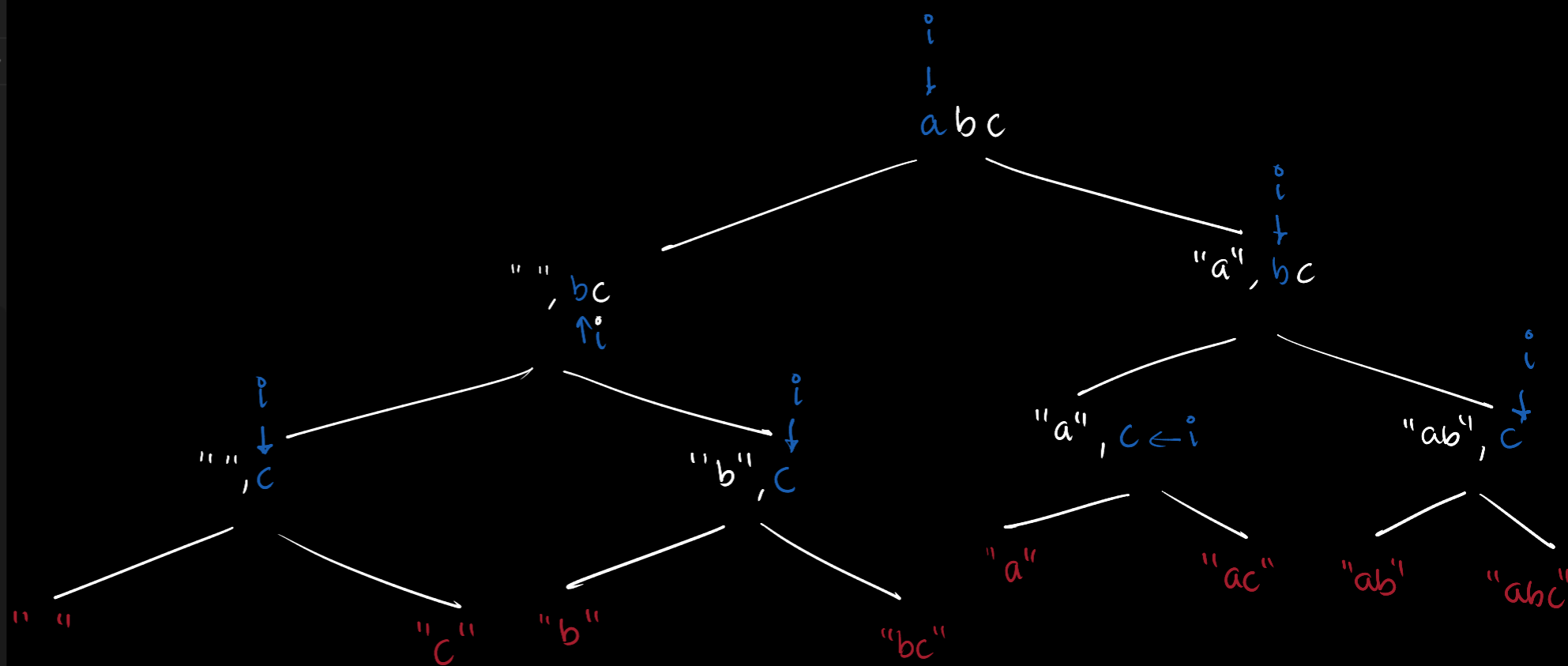
Sample Input 2:

1
bbb

Sample Output 2:

b b b bb bb bbb

→ Using recursion:-



→ The base case will be as soon as $i \geq n$, we will push the string to our vector of string.

$i \geq n$, we will push
↓
size of str.

```
#include <bits/stdc++.h>

void getTheSubsequence(vector<string> &ans, string str, int i, int n, string s) {
    if(i >= n) {
        if(s != "")
            ans.push_back(s);
        return;
    }

    getTheSubsequence(ans, str, i + 1, n, s);
    s = s + str[i];
    getTheSubsequence(ans, str, i + 1, n, s);
}

vector<string> subsequences(string str){
    vector<string> ans;
    string s = "";
    int i = 0, n = str.length();
    getTheSubsequence(ans, str, i, n, s);

    return ans;
}
```

Time Complexity	$O(2^N)$
Space Complexity	$O(2^N)$

→ Bitwise

0 0 0 → ""
0 0 1 → "c"
0 1 0 → "b"
0 1 1 → "bc"
1 0 0 → "a"
1 0 1 → "ac"
1 1 0 → "ab"
1 1 1 → "abc"

```
#include <bits/stdc++.h>
vector<string> subsequences(string str) {

    vector<string> ans;

    int length = str.length();

    int endIndex = 1 << length;

    for (int i = 1; i < endIndex; i++) {
        int j = length - 1, no = i;

        string output = "";

        while (no) {
            if ((no & 1)) {
                output = str[j] + output;
            }
            j--;
            no >>= 1;
        }
        if (output != "") {
            ans.push_back(output);
        }
    }

    return ans;
}
```

→ 2^N

→ $\log 2^N = N \cdot \log 2$

Time Complexity	$O(N \cdot 2^N)$
Space Complexity	$O(N)$