UNIT-E

ABSTRACT DATA TYPES

Bruy or total

ಳವಾಗೂ ಗಾಗಿಕಟ್ಕ ಕೃತ್ತ ಗಾಚಿಕ್ಕ ಚಿ Abstract Data Types (ADTS): -

ADT is a high-level concept in data structules that defined a set of values and operations that can be performed on those values, without specifying how the values are stored or implemented.

In other word, an ADT is a way of describing what a data structures can do, without specifying how the values are stoned or implemented. Examples of ADES durable RUBE De Maria

1. Stack in down in some state of the

Last In, first out (LIFO) data stocectules with push and pop operations,

2. Quece in do to danders permissioner Fist In, First out (FIFO) data structules with enqueue and dequeue operations.

3. set ! Unordered collection of unique elements with add, remove, and search operations.

4. Map (or dictionary;

Key-value pairs with insert, delete and lookup operations. o jerry on Bras) Amel

5. Graph!

Non-linear dalon, eternetule with nodes and edges, supporting travelsal and sealch operations,

To post a nin prim BA Sunda with

b. List or Array;

ordered collection of elements with indexing Insertion and deletion operations. The work to be during a

ADIS and classes: 1 wat agent the MA

ADTS and classes are related concepts in programming but they serve different purposes; Abstract Data Type (ADT).

- * A high-level concept that defines a set of values and operations that can be performed transition those Walkes! said a dark
- * Specifies what a data structures can do, without worrying about how its done. A perfines the interface or contract of a data structures.) in terif of facility

class;

A programming construct that implements an ADT.

with push and trop operations.

- * A bile print or template that defines the properties (data) and Methods (operations) of an objection to messiles besshrond
 - * Encapsulate the data and behavior of an object, Making it a self-contained unit.

Examples soni as too a ving white gold

* ADT;

and hardens of the tions. Stack (Last In, First out data structules with push and pop operations)

Astrolassas beinger pillingger Stack Implementation (a class that implements the stack AST using an array or (inted list).

Introduction to opp;

Key concepts:

1. objects:

Reepresent real-world entitles or abstract concepts with properties (data) and Methods (Swortiens that operate on that data).

an pain.

the property of the

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1000 1000 1000

2. classes;

Blueprints or templates that defines the properties and Methods of an object.

3- Inheritance:

A class can inherit properties and methods from a parent class allowing for code rouse and a hierachical relationship.

4. polymorphism; en in the day.

objects of different classes can be treated as object of a common super class, enabling more flexibility in programming.

5- Incapsulation!

objects hide their Internal details and expose only necessary Information through public methods.

6. Abstraction!

Focus on essential features whiles hiding hon-essential details.

Widowi hamped was death railings

Benefits of oop!

1. Modul ariby:

Code is organized into Self-Contained units (objects) that are easy to maintain. and reuse.

in i milaukar s

2. Reusability!

Inheritance and polymorphism enables code reuse and reduce duplication.

3. Basies Maintenance:

Encapsulation and abstraction Make it easier to modify and extend code without affecting other parts of the program is the

4. Improved Readability:

oop concepts help make code motie intuitive and easier to understand. May day wag - sout

classes in python!

A class is a shieppint of template that defines the properties and behaviour of an objects. inaldward day

Key components of a class;

1. class name: heliated pristation prista A unique name for the class. d. Altributes (data):

Variables that are defined inside

the class, representing the object's properties.

3. Methods ! functions!

functions that are defined inside the class, representing the object's behavior.

syntax!

class class Name:

attributes

altributel = value 1

attributes = Values

Nethods

def method (Self):

code here

def prethod & (self):

code here

Example!

class Bank Account:

altributes

account-number = " "

account - name = " "

balance = 0.0

rue thous

det deposité self, amount):

self. balance t= amount

print ("Deposited & famount 3. New balance: \$

2 self · balance " 3.")

```
def withdraw (self, amount):
     if amount y solf balance:
       print ("Insufficient funds!")
     else
         Self = balance -= amount
        print ("withdraw & Eamount &. New balance:
               & & self balance &")
 def display details (self)!
   Print (" Account Number: & solf-account_number?
   print ("Account Name: & self. a crount_name f")
    Print ("Balance; & (self balance 3)
# Create an object of this class
   my account = BankAccount ()
   my account account Number = " 1234567890"
   my_ account · a count_name = " Kasthi"
   my-account balance
                             = 10,000
   my-account · display details ()
   my account deposit (5,000)
   my-account withdraw (7,000)
output!
   Account Number: 1234567890
    Account Name: Karthi
    Balance
                     10,000
    peposited
               : 5000
```

New Balance : 15,000

withd have

Maria Balanco

17,000

1 8,000

Inheritance!

Inheritance in python is a mechanism that allow one class to inherit the altributes and Methods of another class.

The class that is being inherited from is called the parent or superclass and the class that is doing the inheriting is called child or subclass.

Example:

#Palent class

class shape:

def_int - (Self /color):

Self. color = color

def. get_color (Self);

return soft, color

Child class

class circle (Shape):

det lint- (self, color, radius):

Super (). Int. (color)

self · radius = radius

def gel area (salf):

return 3.14 * (Self radius * * 2)

child class

class Rectargle (Shape):

det -int- (self, cdor, width, height)!

Super (). - init- (color)

self . wildth = width

Self.height = height

def get_aren (self)!

return self.width * self.height

Circle = Circle ("red",5)

Print (circle get-Lolor())

Print (circle get-alea())

rectangle - Rectangle ("blue", 416)

Print (rectangle get-Color())

Print (rectangle get- area L)

output; red 78.5 blue 24

Types of inheritance:

1. gingle Inheritance:
one child class inherits from the palent
class.

2. <u>Multiple Inheritance</u>?

one child class inherits from multiple

parent classes.

3. Multilevel Inhelitance:

A child class inhelits from a palent class that itself inhelits from another parent class.

4. Hierarchical Inheritance:

Multiple subclasses inherit from a single Palent class.

5. Hybrid inheritance:

Combination of kultiple and kultileval

Namespace:

In python, a namespace is a Mapping of names to objects. It's a way to organized and scope names to avoid name clashes.

Types of Namespaces;

1. Cilobal Namespace!

The global namespace contains built-in names such as len, print and range.

2. Local Namespace:

Each function or Heethod has its own.

Local namespace, which contains names defined within that function

3. Module Namespace:

Each Module has its own namespace which contains names defined in that Module.

4. class Nancepace!

Each class has its own namespace which contains names defined in that class.

5. Instance warnespace!

Bach instance of a class has its own namespace, which contains names defined in that instance

```
Example,
 # Global namespace
  XZID
 def my-function ():
    # local namespace
   Print ("Local:", locale ())
   Print ("Global:", globals ())
  class myclass:
     # class namespace
     Z=30
    def _init_ (self):
      #Instance namespace
       Self. 10=40
   Obj = myches ()
    Print ("class:", myclass-dict_)
    Print ("Instance:", obj. did-)
   Print (" Global x: ", x)
    Paint (" Local y (not accessible): ", y)
   Print (" (lass z:", Myclass. Z)
    Print ("Instance w:", obj. w)
  output:
     Crlobal: g..., X': 10, ...}
      Local: f'y': 203
class f... z'; 30, ... }
      Instance: f..., w: 40, ... }
       Golobal X:10
```

Local y (not accessible);

cass 2:30 Instance w:40

shallow and deep copying:

In Python, when you assign a new variable to an Existing variable, it doesn't create a new copy of the original variable.

Instead, it creates a new reference to the same object. This can lead to unexpected behavior when Modifying the new variable.

To create a true copy of an object, you can use the copy module, which provides two types of copying;

I shallow copy:

creates a new object and then (to the extend possible) inserts reference into it to the objects found in the original object.

2. Deep copy:

creates a new object and then, recursively, inserts copies into it of the objects found in the original object.

Example:

import copy original -list = [C1/2], C3/4]

A shallow copy

shallow_copied_list = copy · copy (original_list) # Deep copy doep_copied_list = copy deepcopy (original_list) original_list append ([5,6])

original_list (o] (o] = 'x'

print ("original list:", original_list)

print ("shallow Copied list:", ehallow-copied_list)

print ("Deep copied list:", deep_copied_list)

output:

Original. List: [[x', 2], [3,4], [5,6]]

Shallow Copied List: [['x',2], [3,4]]

Deep copied list: [[1,2], [34]]

Introduction to analysis of algorithms!

Analysis of Algorithms is the process of evaluating the performance of an algorithm, typically in terms of its new time and space complexity. This involves undesstanding how the algorithms running time and remory usage scale as the input size increases.

Concepts of Analysis of Algorithms;

1. Time complexity:

The amount of time an algorithm takes to complete usually expressed as a handion of the input size.

2. Space Complexity:

The amount of memory an algorithm 11908.

Usually expressed as a function of the input size.

3. Big o Notation:

the upper bound of an algorithm's time or

space complexity.

4. Bost Case:

The Minimum amount of time or space an algorithm requires, usually occurring when the input is in a specific order or has a specific structure.

5. Average Lase;

The expected amount of time or space an algorithm nequires, usually occurry calculated by averaging the time or space required for a large number of inputs.

6. Worst case?

The maximum amount of time or space an algorithm requires usually occurring when the input is in a specific order or has a specific structule that causes the algorithm to perform poorly.

Some Common time Complosoities include;

* O(1) - constant time

* Ollogn) - Logarithmic Line

* O(n) - Linear line

* o(nlogn) - Linear/Hmic time

* O(n^2) - quadratic time

* O(211) - exponential time

By analyzing on algorithm's time and space complexity, we can;

* predict its performance on large inputs *Compare the officiency of different algority 4 Identify bottlenecks and areas for optimization

Asymptotic notations!

Asymptotic notations are used to describe the growth rate of an algorithm! time or space complexity as the input size increases.

Hele are the most common asymptotic notations;

1. Big to (o(n)):

Upper bound, worst case scenario. It gives the maximum time or space an algorithm

can take. 2. Big 2 (I(n)):

Lower bound, best case scenario- It gives the minimum time or space an algorithms can take.

3. Big b (o(n)):

Exact bound, avelage-case scenario. It gives the exact time or space an algorithm takes.

4. Little 6 (oln):

Loose upper bound. It gives a upper bound that is not tight. 5. Little w (w(r));

Loose lower bound. It gives a

lower bound that is not tight.

Important points;

- 1. Big o 1s the most commonly used notation.
- d. Big o and Big 51 can be used to describe both time and space complexity.
- 3. Rig or is used when the upper and lower bounds are the same.
- 4. Little o and little wale used when the bounds are not light.

Example;

def best case (n);

Best-Case Scanario: I (n)

for i in range (n): print Li)

def overage - case (n):

Anelage - case scenario: @ (n logn)
for i in range (n):

for j in range (int (n/2)):

Print (i,i)

def worst_case (n):

webst-case scenario: 0 (n^2)
for i in range (n):

for j'in range (n)',
print('i,j)

#Test the functions

best_case (5)

average_case (5)

worst_case (5)

Benefits:

1. Simplifies analysis:

Asymptotic notations simplify the analysis of algorithms by tousing on the gravity rate rather than the exact running time.

2. Hides constants Focus on scalability:

Asymptotic notations help us understand how an algorithm's running time or space usage scales as the input size increases.

3. Hides Constant!

Asymptotic notations ignore constants
factors, which can vary depending on the
implementation and hardwards

4. Comparbility:

Asymptotic notations provide a Common language for comparing the efficiency of different algorithms.

5. Predicts performance:

Asymptotic notations help predict an algorithm's performance on large inputs even if we coun't measure it directly.

6. Helps in optimization!

Asymptotic notations identify performance bottlenecks and guide optimization effects, T. Theosetical foundations:

Asymptotic notations provide a the workical foundation for understanding algorithmie complexity.

8. platform independence:

Asymptotic notations are platform independent, making than useded for analyzing algorithms across different harduale and softwale environment.

Divide and Longuer!

It is a popular algorithmic paradigm used to solve complex problems by breaking them down into smaller sub-problems, solving each sub program, and then combining the Solutions to the original problem. Three Main steps:

Divide the problem into smaller Sub problème that are more, manageable and Similar in etructure to the original problems. 2. conques!

solve each sub-problem recursively or itelatively using the same approach.

3. Combine!

Compline the solutions to the sal problems to solve the original problem.

Example!

def merge_sort (asa)! # Divide If len (arr) z=1; return our

mld = len (ans) //2 left-half = arr [:Mid] -right_half=arrcmid:]

A Conques

left_half=mesge_sort (left_half) right-half= meage_sort (right-half)

ombine

rebusn mesge (left-half, right-half)

def merge (left, right)!

merged = CJ:

left_index = sight_index =0 while loft Index clan (left) and night index clan(mit) if lost Cleft_index] z=night [right_index]:

merged append (left [left_index])

left index = 1

else!

merged append (right [right-inless]) · Man might Indep 4= 1+, in plant of merged extend (left [left_indexi]) merged expland (Cologht Cologht_index: J) ratusi mesged

Test the program arr = [5/2/8/3/1,6/4] point (" original array:", arr) print ("sorted curray;", merge_sort (anr))

output?

original Array: [5,2,8,3,1,6,4] Sorted Arrang: [1,2,3,4,5,6,8] Benefits!

1. Edficient: often have a time complexible of o (n logn) or better.

can be applied to large problems by d. scalable! breaking them down into emaller sub-problems. Jahos Kirkhark .

3. Easy to implement:

Recursive solutions can be easier to conderstand and implements: Manjoros bee shirin &

drawbacks!

1. Recursive, overhead!

Peausive solutions can have Over head due to functions calle and returns.

a. Difficulty on solving emall problems!

some algorithms may have difficulty Solving small sub-problems efficiently.

Recursion!

Recursion le de programming technique where a function calls itself repeatedly un it reaches a base case that stops the recursion. Here's a breakdown of reculsion.

1-Base Case!

A trivial case that can be solved without recussion.

2 Receive rue!

A case that can be broken down into smaller scib- problems of the same type which are then solved nearstvely.

3. Recorreive call:

The function calls Itself with a Smaller input or a Modified reserion of the Benefits! collins du land many many

1 simplified code:

Recursion can lead to more elegant and concise code moderni la : hand to have

2. Divide ant conquer!

· Recussion naturally implements. the divide and conquer approach. into so sul, know esker

3. Tree traversals:

Recursion is well-suited for

traversing tree- like data struckere. , situation a walding the Mone palia drawbades!

1. stark overflows

Deep recussion can lead to a stack overflow error. Hadigan on Silver

Transfer stigers of group in

Chalexus if Line

· (Catera . tala

26 H 601 . 3 . 21

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Himmy in silica

2. Inefficiency:

Leaussian can be solved than i berative solutions due to function call overhead.

3. Difficulty in debugging!

Reaussion can make debugging more Challenging due to the repeated function calls.

Example!

calculates the factorial of a humber.

def factorial (n): # Base case

If U==0 or U==1 return 1

Reaveive case

neturn n * factorial (n-1)

print (factored (5))

output:

The factorial function calls Itself with decreasing values of a until 11- reaches the base case (n==0 or n==1) at which point it starts returning the results back up the all stack.

Analysing recessive Algorithms!

Identify the Envial case that stops the recusion

. 23/2

Plantin.

is were

Starp & Lynn .

2. Recursive case!

. Land we Understand how the problem is broken down into smaller sub-problems. this is at the fitting .

3. Recursive call:

'Arialy ze the function call with a smaller input or modified original input.

4. Telmination! Ensure the recursion terminates (base M. Terridon For case is readed).

5. Time Lamplexity:

Calculate the time complexity (5.9. O(n), O(21n)). sin whomas is

6- space complexible in delination in

Determine the space complexity (B.g. 0(1),0(1)).

7. optimization:

Look for apportunities to optimize the necusive algorithm. CEG. memoization, diffrance programming.)

and it and it is the transfer in a series of

. Dirato Dir. sat

Some common bechniques for analyzing reculsive algorithms include:

1- Recursion tree!

Visualize the recursive calls as a tree to condenstand the algorithm's behavior,

d. Reculsion depth:

Aralyze the maximum depth of the recursive calls.

3. Function call strehead:

of repeated function calls.

4. Cache performance:

Evaluate the impact of reccelsion