INDIAN INSTITUTE OF TECHNOLOGY, KHARAGPUR

Department of Computer Science & Engineering

Programming and Data Structures (CS11001)

Midsem (Autumn, 1st Year)

Date: Tue, Sep 27, 2011

Students: 660

Time: 09:00-11:00am

Marks: 55

-1 of 4 pages --

Answer ALL the questions.

Do all rough work on separate rough sheets which you should not submit.

Answer on the question paper itself in the spaces provided.

as: float x[3] code). Assume the	o-ordinates of X and \overrightarrow{i} , \overrightarrow{j} , \overrightarrow{k} are the unit vectors. It may be represented using an a $\{X_x, X_y, X_z\}$ (substituting X_x, X_y, X_z with their actual numerical values in the \overrightarrow{A} , \overrightarrow{B} , \overrightarrow{C} , \overrightarrow{P} , etc are similarly declared/defined (with initialisation, if necessary).	e 'C'
The dot product	\overrightarrow{B} may be computed and <i>returned</i> via the following 'C' function vecDP () defined	as:
ı		
		- [
Also, the cross preccp() as:	oduct $\overrightarrow{C}=\overrightarrow{A} imes \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fund	ction
-		ction
-	oduct $\overrightarrow{C} = \overrightarrow{A} imes \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fund	ection
-	oduct $\overrightarrow{C}=\overrightarrow{A} imes\overrightarrow{B}$ may be computed and stored in C[] via the following 'C' function of the following functi	ection
vecCP() as:	oduct $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{A} \times \overrightarrow{B}$ may be computed and stored in C[] via the following 'C' fundamental formula: $\overrightarrow{C} = \overrightarrow{C} = \overrightarrow{C} \times \overrightarrow$	ection

Roll:

The cross product 2	$\overrightarrow{Z}_{P,AB} = \overrightarrow{AB} \times \overrightarrow{AP}$ (fl	<pre>/ AP[1]=; AP[2]=; /* done similarly */ oat Z_P_AB[3]) may be computed using an above defined</pre>
unction as:		· · · · · · · · · · · · · · · · · · ·
Suppose that A, B, C may be computed as		which are co-planer to P . The vector \overrightarrow{AC} (float AC[3]
	; AC	C[1]=; AC[2]=; /* done similarly */
The cross product $\overline{2}$	$\overrightarrow{Z}_{C,AB} = \overrightarrow{AB} \times \overrightarrow{AC}$ (flo	C[1]=; AC[2]=; /* done similarly */ oat Z_C_AB[3]) may be computed using an above define
unction as:		· · · · · · · · · · · · · · · · · · ·
The dot product $d_{P,C}$	$Z_{C,AB} = \overrightarrow{Z}_{P,AB} \cdot \overrightarrow{Z}_{C_{\bullet}AB}$ (f	float d_Z_PC_AB) may be computed using an above defined
unction as:	**.	;
he condition that P	and C are on the same side	de of \overline{AB} is: () $\overrightarrow{Z}_{P,AC} \cdot \overrightarrow{Z}_{B,AC} \text{ (d_Z_PB_AC)}, \overrightarrow{Z}_{A,BC} = \overrightarrow{BC} \times \overrightarrow{A} \text{ and}$
imilarly, let $Z_{B,AC}$	$\overrightarrow{Z}_{A,BC}$ (d_Z_PA_BC) be:	$Z_{P,AC} \cdot Z_{B,AC}$ (d_z_PB_AC), $Z_{A,BC} = DC \times A$ and available.
, , .	determine whether P is in	
()
s: $p(x) = (((a_n x \text{ dvantage of this metarious powers of } x))$	of degree n , $p(x) = a_n x^n + a_{n-1} x + \dots + a_1 x + \dots + a_1 x + \dots$ ethod of evaluation is that	$a + a_{n-1}x^{n-1} + \ldots + a_1x^1 + a_0, a_n \neq 0$, it can be rewritten a_0 . This is the Horner's scheme for evaluating $p(x)$. The explicit exponentiation is avoided. Let the coefficients for the gray (say) P[], as float P[]={ $a_n, a_{n-1}, \ldots, a_1, a_0$ }
s: $p(x) = (((a_n x))^n x)$ dvantage of this mearious powers of x An iterative function	of degree n , $p(x) = a_n x^n + a_{n-1} x + + a_1 x +$ ethod of evaluation is that of $p(x)$ be stored in an area based on the above scheme	$a + a_{n-1}x^{n-1} + \ldots + a_1x^1 + a_0, a_n \neq 0$, it can be rewritten a_0 . This is the Horner's scheme for evaluating $p(x)$. The explicit exponentiation is avoided. Let the coefficients for the gray (say) P[], as float P[]={ $a_n, a_{n-1}, \ldots, a_1, a_0$ }
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s: $p(x) = (((a_n x \text{ dvantage of this mearious powers of } x \text{ an iterative function}$	of degree n , $p(x) = a_n x^n + a_{n-1} x + + a_1 x +$ ethod of evaluation is that of $p(x)$ be stored in an area based on the above scheme	$a_n + a_{n-1}x^{n-1} + \ldots + a_1x^1 + a_0, \ a_n \neq 0$, it can be rewrittent $a_n + a_0$. This is the Horner's scheme for evaluating $p(x)$. The explicit exponentiation is avoided. Let the coefficients for the tray (say) P[], as float P[]= $\{a_n, a_{n-1}, \ldots, a_1, a_0\}$ are is as follows:
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(CS11001)	Roll:	Sec:
<pre>int isPerfect (int n) { // ret</pre>	urn values: l if perfect. 0 (otherwise
int isreffect (int ii) (// icc	din varues. I il periodo, o	50
	// declarations with initial:	isations
repeat { // stepwise comput	ation of dvisors of n	
	·	
*		
	// extr	`a
} until ();	
} // final result is computed	and returned, in the last st	ep
(a) Representation of NaN in IEEE floating	g point 754 format is:	
(b) Representation of ∞ in IEEE floating j	point 754 format is:	
(c) Representation of $(1.4)_{10}$ in IEEE float	ting point 754 format is:	
(d) Decimal value of the IEEE floating poi	int 754 number 0, 10000101, 101111000	
(a) Decimal value of the tere hoaning hor	TUE 124 HOTHINGS A TOUGHTOT TATITION	22222222222222222

(CS11001)	Roll:	Sec:
A recursive function based on the above scheme is as i	follows:	
hornerPoly() {
		
}		
The digits of a given decimal integer n may be rotated significant digit of n , let n' represent the number after non-zero) (iii) required result is $10^k d + n'$. Complete comments and computing $10^k d$ in steps as you determ	d is removed from n (ii) let n' have the function rotRight () given time k (rather than compute $10^k d$ so	e k digits (k th digit is below, following the eparately).
	{ // function c	eclaration
int d, nDash;	// any more dec	larations
	// extra	act d
	// comp	nte nDash
// stepwise computation of k and 1	10^k * d	
		
} // final result is computed and ret	turned, in the last ste	p