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**2.1 Order the following functions by growth rate:  $N$ ,  $\sqrt{N}$ ,  $N^{1.5}$ ,  $N^2$ ,  $N \log N$ ,  $N \log \log N$ ,  $N \log^2 N$ ,  $N \log(N^2)$ ,  $2/N$ ,  $2N$ ,  $2^{N/2}$ ,  $37$ ,  $N^2 \log N$ ,  $N^3$ . Indicate which functions grow at the same rate.**

Lowest to Highest Growth: None of the functions have the exact same growth, but they are tiered by their level of growth. Functions on the same line grow with similar trends.

$2/N$

$37$

$\text{sqrt}(N)$

$N$

$N \log \log N$     $N \log N$     $N \log(N^2)$

$N \log^2 N$

$N^{1.5}$

$N^2$     $N^2 \log N$

$N^3$

$2^{(N/2)}$     $2^N$

- . 2.6 In a recent court case, a judge cited a city for contempt and ordered a fine of \$2 for the first day. Each subsequent day, until the city followed the judge's order, the fine was squared (that is, the fine progressed as follows: \$2, \$4, \$16, \$256, \$65, 536, . . .).**

**a. What would be the fine on day  $N$ ?**

$$F(N) = 2^{2^{(N-1)}}$$

- b. How many days would it take the fine to reach D dollars? (A Big-Oh answer will do.)

$$N = \{(\log[\log D]) / (\log[2\log 2])\} + 1$$

(10 pts) Give an analysis of the Big-Oh running time for each of the following program fragments:

```
int sum = 0;
for ( int i = 0; i < 23; i ++ )
    for ( int j = 0; j < n ; j ++ )
        sum = sum + 1;
```

1 for initialization of sum : 1  
1 to set int i, 23+1=24 tests, 23 increments: 48  
1 to set int j, n+1 tests, n increments, executed 23 times:  
23(2n+2)  
1 addition, 1 initialization, executed n times: 2n  
TOTAL=1+48+23(2n+2)+2n=95+48n -> 48n-> **n**

```
int sum = 0;
for ( int i = 0; i < n ; i ++ )
    for ( int k = i ; k < n ; k ++ )
        sum = sum + 1;
```

1 for initialization of sum: 1  
1 to set int i, n+1 tests, n increments: 2n+2  
1 to set k, n+1 checks, n increments, n executions: n\*(2n+2)  
1 addition, 1 initialization, n executions: 2n  
TOTAL=1+(2n+2)+n\*(2n+2)+2n=2n<sup>2</sup> +6n +3 -> 2n<sup>2</sup>-> **n<sup>2</sup>**

```
public int foo(int n, int k) {
    if(n<=k)
        return 1;
    else
        return foo(n/k,k) + 1;
}
```

else loop is worse case:  
calls foo log(n)/log(k) times  
each call of foo has 1 check and 1 return: 2

Total =  $2 * [\log(n) / \log(k)] \rightarrow \log(n) / \log(k) \rightarrow \log(n)$

- . **2.11** An algorithm takes 0.5 ms for input size 100. How long will it take for input size 500 if the running time is the following (assume low-order terms are negligible):

**a. linear**

$$(500/100) = O(500)/O(100)$$

$$O(500) = 5 * O(100) = 5 * (.5e-3s) = 2.5ms = O(500)$$

**b.  $O(N \log N)$**

$$[500 \log(500)] / [100 \log(100)] = O(500) / O(100)$$

$$O(500) = O(100) * [500 \log(500)] / [100 \log(100)] = (.5ms) * 6.74 = 3.374ms = O(500)$$

**c. quadratic**

$$O(100) * (500^2) / (100^2) = O(500)$$

$$(.5ms) * (250000/10000) = (.5ms) (25) = 12.5ms = O(500)$$

**d. cubic**

$$O(100) * (500^3) / (100^3) = O(500)$$

$$(.5ms) * (125) = 62.5ms = O(500)$$

**2.15** Give an efficient algorithm to determine if there exists an integer  $i$  such that  $A_i = i$  in an array of integers  $A_1 < A_2 < A_3 < \dots < A_N$ . What is the running time of your algorithm?

```
Public boolean findIt(int[] arr){  
    int upperbound=arr.length-1;  
    int lowerbound=0;  
    while(upperbound>=lowerbound){  
        index=(upper_bound + lower_bound)/2;
```

```
    if(arr[index]==index+1){ return True;}  
    else if(arr[index]>index+1){upperbound=index-1;}  
    else if(arr[index]<index+1){lowerbound=index+1;}  
    return False;
```

Run Time:

$O(\text{findIt}) = 4[\log(N)/\log(2)] + 7 \rightarrow \log_2(N) \rightarrow \log(N)$