

# Solving Sudoku with Dancing Links

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## Example: Combinatorial Enumeration

Create all permutations of the set  $\{0, 1, 2, 3\}$

- Simple example to demonstrate key ideas
- Creation, cardinality, existence?
- There are more efficient methods for this example

# Brute Force Backtracking

BLACK = Forward

BLUE = Solution

RED = Backtrack

root	0 1 2	0 1 3 3	0 2 1	0 2 3 3	0 3 1	0 3	1 0 2
0	0 1 2 2	0 1 3	0 2 1 3	0 2 3	0 3 1 3	0 3 3	1 0 2 1
0 0	0 1 2	0 1	0 2 1	0 2	0 3 1	0 3	1 0 2
0	0 1 2 3	0	0 2	0	0 3	0	1 0 2 2
0 1	0 1 2	0 2	0 2 2	0 3	0 3 2	root	1 0 2
0 1 0	0 1	0 2 0	0 2	0 3 0	0 3 2 0	1	1 0 2 3
0 1	0 1 3	0 2	0 2 3	0 3	0 3 2	1 0	1 0 2
0 1 1	0 1 3 0	0 2 1	0 2 3 0	0 3 1	0 3 2 1	1 0 0	:
0 1	0 1 3	0 2 1 0	0 2 3	0 3 1 0	0 3 2	1 0	:
0 1 2	0 1 3 1	0 2 1	0 2 3 1	0 3 1	0 3 2 2	1 0 1	:
0 1 2 0	0 1 3	0 2 1 1	0 2 3	0 3 1 1	0 3 2	1 0	:
0 1 2	0 1 3 2	0 2 1	0 2 3 2	0 3 1	0 3 2 3	1 0 2	:
0 1 2 1	0 1 3	0 2 1 2	0 2 3	0 3 1 2	0 3 2	1 0 2 0	:

## A Better Idea

- Avoid the really silly situations, such as: 1 0 1
- “Remember” that a symbol has been used already
- Additional data structure: track “available” symbols
- Critical: must maintain this extra data properly
- (Note recursive nature of backtracking)

# Sophisticated Backtracking

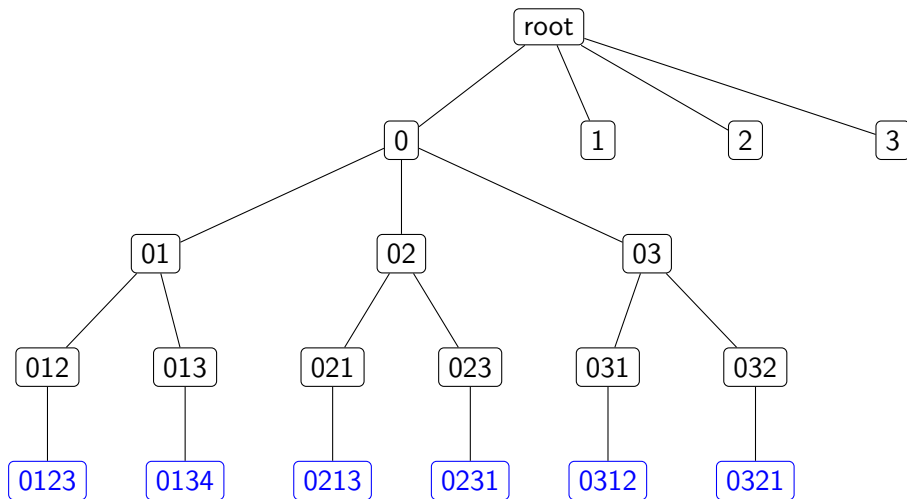
BLACK = Forward

BLUE = Solution

RED = Backtrack

root {0,1,2,3}	0 2 1 3 {}	0 3 2 1 {}	1 0 {2,3}	1 3 0 2 {}
0 {1,2,3}	0 2 1 {3}	0 3 2 {1}	1 {0,2,3}	1 3 0 {2}
0 1 {2,3}	0 2 {1,3}	0 3 {1,2}	1 2 {0,3}	1 3 {0,2}
0 1 2 {3}	0 2 3 {1}	0 {1,2,3}	1 2 0 {3}	1 3 2 {0}
0 1 2 3 {}	0 2 3 1 {}	root {0,1,2,3}	1 2 0 3 {}	1 3 2 0 {}
0 1 2 {3}	0 2 3 {1}	1 {0,2,3}	1 2 0 {3}	1 3 2 {0}
0 1 {2,3}	0 2 {1,3}	1 0 {2,3}	1 2 {0,3}	1 3 {0,2}
0 1 3 {2}	0 {1,2,3}	1 0 2 {3}	1 2 3 {0}	1 {0,2,3}
0 1 3 2 {}	0 3 {1,2}	1 0 2 3 {}	1 2 3 0 {}	root {0,1,2,3}
0 1 3 {2}	0 3 1 {2}	1 0 2 {3}	1 2 3 {0}	2 {0,1,3}
0 1 {2,3}	0 3 1 2 {}	1 0 {2,3}	1 2 {0,3}	:
0 {1,2,3}	0 3 1 {2}	1 0 3 {2}	1 {0,2,3}	:
0 2 {1,3}	0 3 {1,2}	1 0 3 2 {}	1 3 {0,2}	:
0 2 1 {3}	0 3 2 {1}	1 0 3 {2}	1 3 0 {2}	:

# Depth-First Search Tree



# Algorithm

```
n=4
available=[True]*n    # [True, True, True, True]
perm=[0]*n            # [0, 0, 0, 0]
```

```
def bt(level):
    for x in range(n):
        if available[x]:
            available[x]=False
            perm[level]=x
            if level+1 == n:
                print perm
            bt(level+1)
            available[x]=True
```

```
bt(0)
```

# Sudoku Basics

- $n^2$  symbols
- $n^2 \times n^2$  grid
- $n^2$  subgrids (“boxes”) each  $n \times n$
- Classic Sudoku is  $n = 3$
- Each symbol once and only once in each row
- Each symbol once and only once in each column
- Each symbol once and only once in each box
- The grid begins partially completed
- A Sudoku puzzle should have a unique completion



# Example

5		8	4	9
	6	7	5	3
1	5			
		2	8	
7			4	1
	3		2	5
4	9			3

⇒

5	1	3	6	8	7	2	4	9
8	4	9	5	2	1	6	3	7
2	6	7	3	4	9	5	8	1
1	5	8	4	6	3	9	7	2
9	7	4	2	1	8	3	6	5
3	2	6	7	9	5	4	1	8
7	8	2	9	3	4	1	5	6
6	3	5	1	7	2	8	9	4
4	9	1	8	5	6	7	2	3

# Sudoku via Backtracking

- Fill in first row, left to right, then second row, ...
- For each blank cell, maintain possible new entries
- As entries are attempted, update possibilities
- If a cell has just one possibility, it is forced
- Lots to keep track of, especially at backtrack step

# Sudoku via Backtracking

- Fill in first row, left to right, then second row, ...
- For each blank cell, maintain possible new entries
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- If a cell has just one possibility, it is forced
- Lots to keep track of, especially at backtrack step
- Alternate Title: “Why I Don’t Do Sudoku”

Top row, second column: possibilities?

5			8	4	9
	6	7	5	3	
1	5		3		1
			2	8	
				1	8
7			4	1	5
	3		2		
4	9		5		3

$\{1, 2, 3, 6, 7\}$

$\{1, 2, 4, 7, 8\} \longrightarrow \{1, 2, 4, 7, 8\} \cap \{1, 2, 3, 6, 7\} = \{1, 2, 7\}$

Suppose we try 2 first.

Seventh row, second column: possibilities?

5	2	8	4	9
	6	7	5	3
1	5	2	8	
7		4	1	5
	3	2		
4	9	5		3

$\{2, 3, 6, 8, 9\}$

$$\{1, 4, 7, 8\} \longrightarrow \{1, 4, 7, 8\} \cap \{2, 3, 6, 8, 9\} = \{8\}$$

One choice!

This may lead to other singletons in the affected row or column.

# Exact Cover Problem

- Given: matrix of 0's and 1's
- Find: subset of rows
- Condition: rows sum to exactly the all-1's vector
- Amenable to backtracking (on columns, not rows!)
- Example: (Knuth)

0	0	1	0	1	1	0
1	0	0	1	0	0	1
0	1	1	0	0	1	0
1	0	0	1	0	0	0
0	1	0	0	0	0	1
0	0	0	1	1	0	1

# Solution

Select rows 1, 4 and 5:

⇒	0	0	1	0	1	1	0
	1	0	0	1	0	0	1
	0	1	1	0	0	1	0
⇒	1	0	0	1	0	0	0
⇒	0	1	0	0	0	0	1
	0	0	0	1	1	0	1
↓							
	1	1	1	1	1	1	1

# Sudoku as an Exact Cover Problem

- Matrix rows are per symbol, per grid location ( $n^2 \times (n^2 \times n^2) = n^6$ )
- Matrix columns are conditions: ( $3n^4$  total)
  - ▶ Per symbol, per grid row: symbol in row ( $n^2 \times n^2$ )
  - ▶ Per symbol, per grid column: symbol in column ( $n^2 \times n^2$ )
  - ▶ Per symbol, per grid box: symbol in box ( $n^2 \times n^2$ )

Place a 1 in entry of the matrix  
if and only if

matrix row describes symbol placement satisfying matrix column condition

- Example:  
Consider matrix row that places a 7 in grid at row 4, column 9
  - ▶ 1 in matrix column for “7 in grid row 4”
  - ▶ 1 in matrix column for “7 in grid column 9”
  - ▶ 1 in matrix column for “7 in grid box 6”
  - ▶ 0 elsewhere



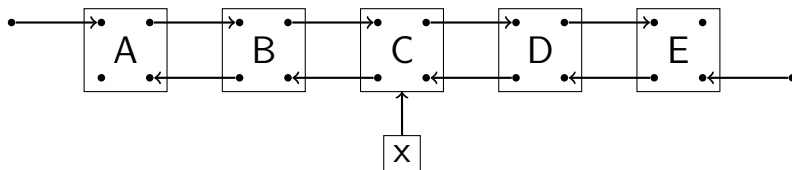
# Sudoku as an Exact Cover Problem

- Puzzle is “pre-selected” matrix rows
- Can delete these matrix rows, and their “covered matrix columns”
- $n = 3$ : 729 matrix rows, 243 matrix columns
- Previous example: Remove 26 rows, remove  $3 \times 26 = 78$  columns
- Select  $81 - 26 = 55$  rows, from 703, for exact cover (uniquely)
- Selected rows describe placement of symbols into locations for Sudoku solution

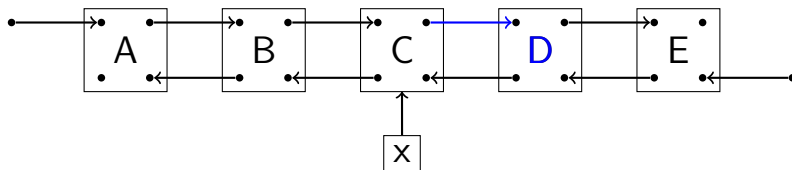
# Dancing Links

- Manage lists with frequent deletions and restorations
- Perfect for descending, backtracking in a search tree
- Hitotumatu, Noshita (1978, Information Processing Letters)
  - ▶ “pointers of each already-used element are still active while. . . removed”
  - ▶ Two pages,  $N$  queens problem
  - ▶ Donald Knuth listed in the Acknowledgement
- Popularized by Knuth, “Dancing Links” (2000, arXiv)
  - ▶ Algorithm X = “traditional” backtracking
  - ▶ Algorithm **DLX** = **D**ancing **L**inks + Algorithm **X**
  - ▶ 26 pages, applications to packing pentominoes in a square

# Doubly-Linked List

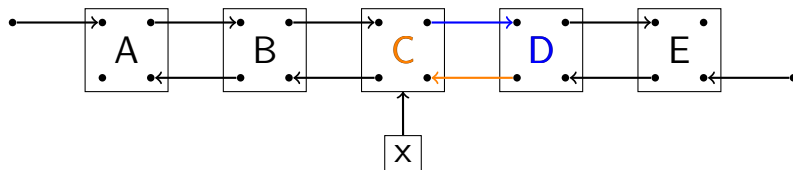


## Remove Node "C" From List



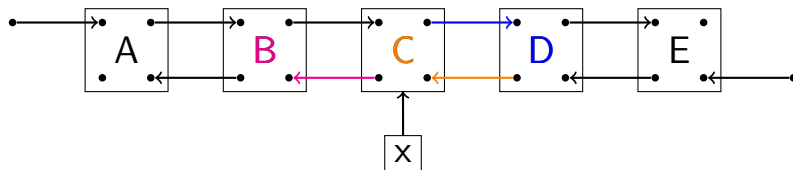
$R[x]$

## Remove Node "C" From List



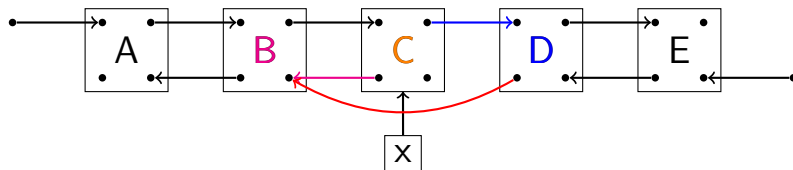
$L[R[x]]$

## Remove Node "C" From List



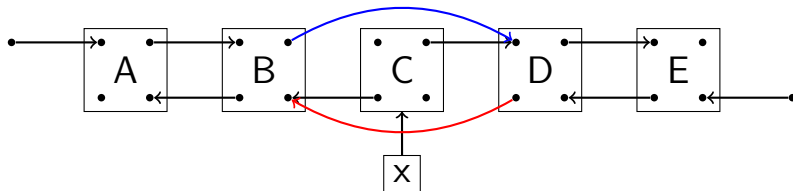
$L[R[x]]$        $L[x]$

## Remove Node "C" From List



$$L[R[x]] \leftarrow L[x]$$

## Two Assignments to Totally Remove "C"

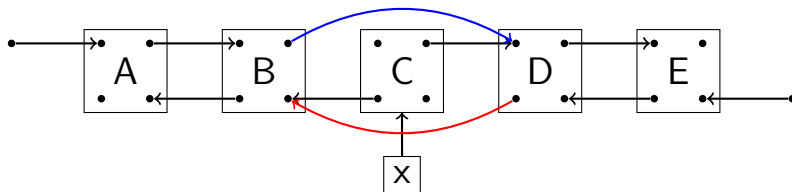


$$L[R[x]] \leftarrow L[x]$$

$$R[L[x]] \leftarrow R[x]$$



## Two Assignments to Totally Remove "C"

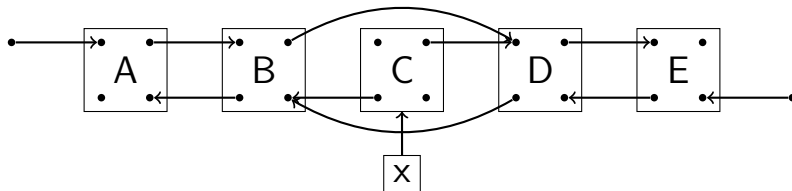


$$L[R[x]] \leftarrow L[x]$$

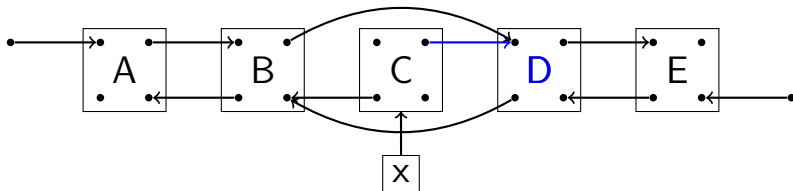
$$R[L[x]] \leftarrow R[x]$$

DO NOT CLEAN UP THE MESS

## List Without “C”, Includes Our Mess

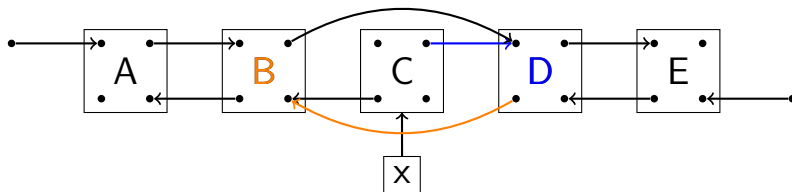


## Restore Node "C" to the List



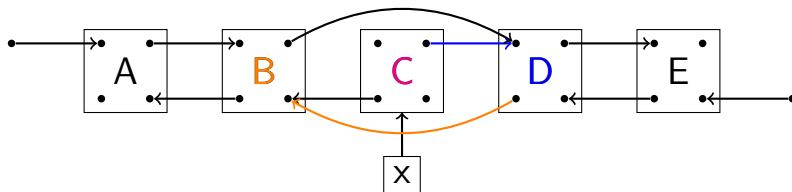
$R[x]$

## Restore Node "C" to the List



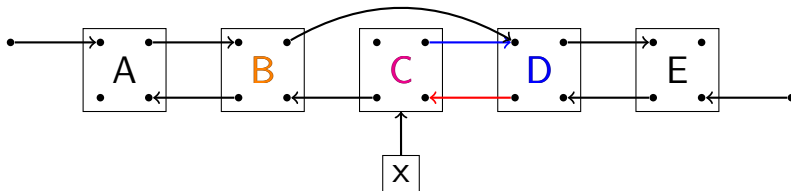
$L[R[x]]$

## Restore Node "C" to the List



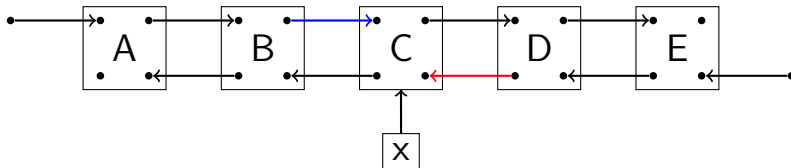
$L[R[x]]$  x

## Restore Node "C" to the List



$$L[R[x]] \leftarrow x$$

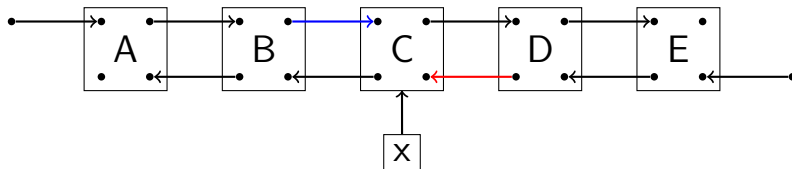
## Restore Node "C" to the List



$$L[R[x]] \leftarrow x$$

$$R[L[x]] \leftarrow x$$

## Restore Node "C" to the List



$$L[R[x]] \leftarrow x$$

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WE NEED OUR MESS, IT CLEANS UP ITSELF



# DLX for the Exact Cover Problem

- Backtrack on the columns
- Choose a column to cover, this will dictate a selection of rows

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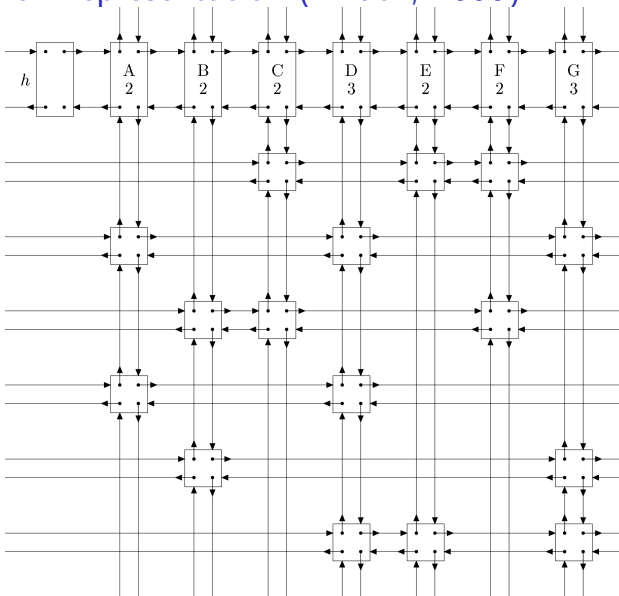
# DLX for the Exact Cover Problem

- Backtrack on the columns
- Choose a column to cover, this will dictate a selection of rows
- Loop over rows, for each row choice remove covered columns
- Recursively analyze new, smaller matrix
- Restore rows and columns on backtrack step

## Exact Cover Example (Knuth, 2000)

	A	B	C	D	E	F	G
1	0	0	1	0	1	1	0
2	1	0	0	1	0	0	1
3	0	1	1	0	0	1	0
4	1	0	0	1	0	0	0
5	0	1	0	0	0	0	1
6	0	0	0	1	1	0	1

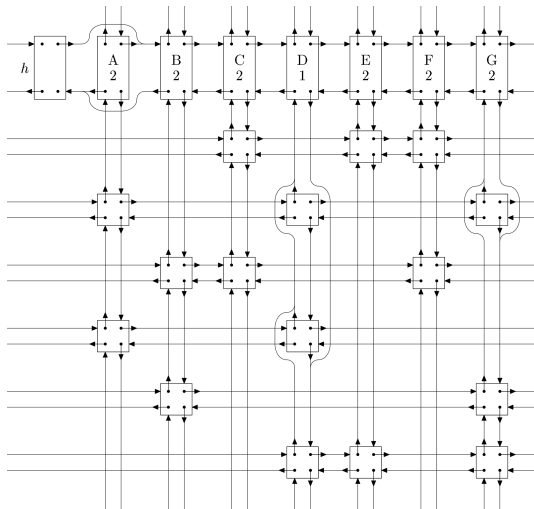
# Exact Cover Representation (Knuth, 2000)



# Exact Cover Representation (Knuth, 2000)

- Cover column **A**
- Remove rows **2, 4**

	<b>A</b>	B	C	D	E	F	G
1	0	0	1	0	1	1	0
<b>2</b>	1	0	0	1	0	0	1
3	0	1	1	0	0	1	0
<b>4</b>	1	0	0	1	0	0	0
5	0	1	0	0	0	0	1
6	0	0	0	1	1	0	1

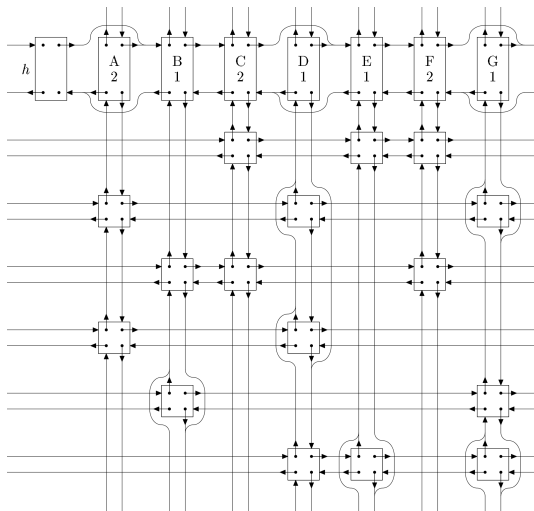


# Exact Cover Representation (Knuth, 2000)

- Loop through rows
- Row 2 covers **D**, **G**
- **D** removes row 4, 6
- **G** removes row 5, 6

	A	B	C	<b>D</b>	E	F	<b>G</b>
1	0	0	1	0	1	1	0
2	1	0	0	1	0	0	1
3	0	1	1	0	0	1	0
4	1	0	0	1	0	0	0
<b>5</b>	0	1	0	0	0	0	1
<b>6</b>	0	0	0	1	1	0	1

Recurse on  $2 \times 4$  matrix  
It has no solution,  
so will soon backtrack



# Implementation in Sage

The games module only contains code for solving Sudoku puzzles, which I wrote in two hours on Alaska Airlines, in order to solve the puzzle in the inflight magazine. — William Stein, Sage Founder

- Sage, open source mathematics software, [sagemath.org](http://sagemath.org)



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- Sage, open source mathematics software, [sagemath.org](http://sagemath.org)
- Stein (UW): naive recursive backtracking, run times of 30 minutes
- Carlo Hamalainen (Turkey/Oz): DLX for exact cover problems
- Tom Boothby (UW): Preliminary representation as an exact cover
- RAB: Optimized backtracking
  - ▶ lots of look-ahead
  - ▶ automatic Cython conversion of Python to C
- RAB: new class, conveniences for printing, finished DLX approach

# Timings in Sage

## Test Examples:

- Original doctest, provenance is Alaska Airlines in-flight magazine?
- 17-hint “random” puzzle (no 16-hint puzzle known)
- Worst-case: top-row empty, top-row solution 9 8 7 6 5 4 3 2 1
- All ~48,000 known 17-hint puzzles (Gordon Royle, UWA)

Equipment: R 3500 machine, 3 GHz Intel Core Duo

Puzzle	Time (milliseconds)		
	Naive	Custom	DLX
Alaska	34	0.187	1.11
17	1,494,000	441.0	1.20
Worst	4,798,000	944.0	1.21
48K 17			~60,000

Talk available at:

[buzzard.pugetsound.edu/talks.html](http://buzzard.pugetsound.edu/talks.html)