

EDIT DISTANCE

INPUT: Two strings $x[1..n]$ and $y[1..m]$

GOAL: Minimum # of operations/keystrokes to edit

↑ x into y

- a) Delete Characters
- b) Insert Character
- c) Substitute / Replace

$x = \text{FAST}$

F A S T _ _ C A T S
↑ ↑
4 deletions + 4 insertions
8 operations

$y = \text{CATS}$

F _ A _ S T
_ C A T S _
↑ ↑ ↑ ↑
deletions insertions deletions
2 insertions + 2 deletions
= 4.

DEFINE SUBPROBLEMS:

{ F _ A _ S T
_ C A T S _ }

Insertion

.
A

Deletion

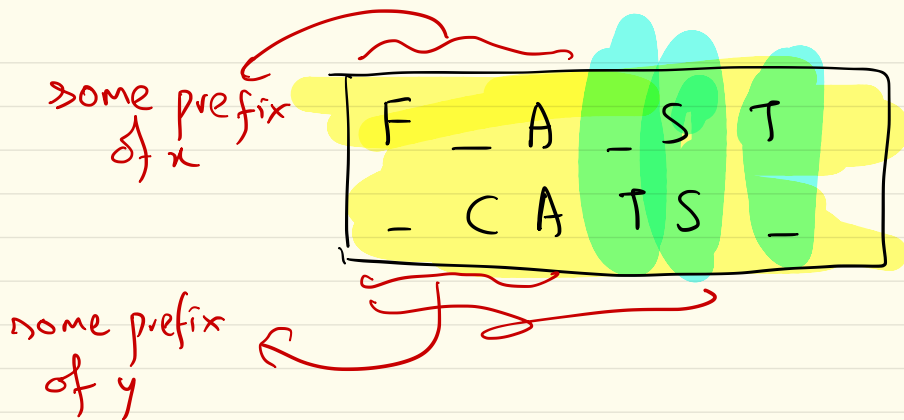
A
-

Substitution

A
B

Carry-Over

A
A



Imagine the optimal solution

SUBPROBLEM

(edit distance)

$E[i, j] =$ # of keystrokes to go from $x[1 \dots i]$ to $y[1 \dots j]$

ANSWER: $E[m, n]$

RECURRENCE RELATION:

$$E[i, j] = \min$$

$x[1..i]$

$y[1..j]$

$?$

$-$
 A

$$1 + E[i, j-1]$$

A
 $-$

$$1 + E[i-1, j]$$

A
 B

if $(x[i] \neq y[j])$

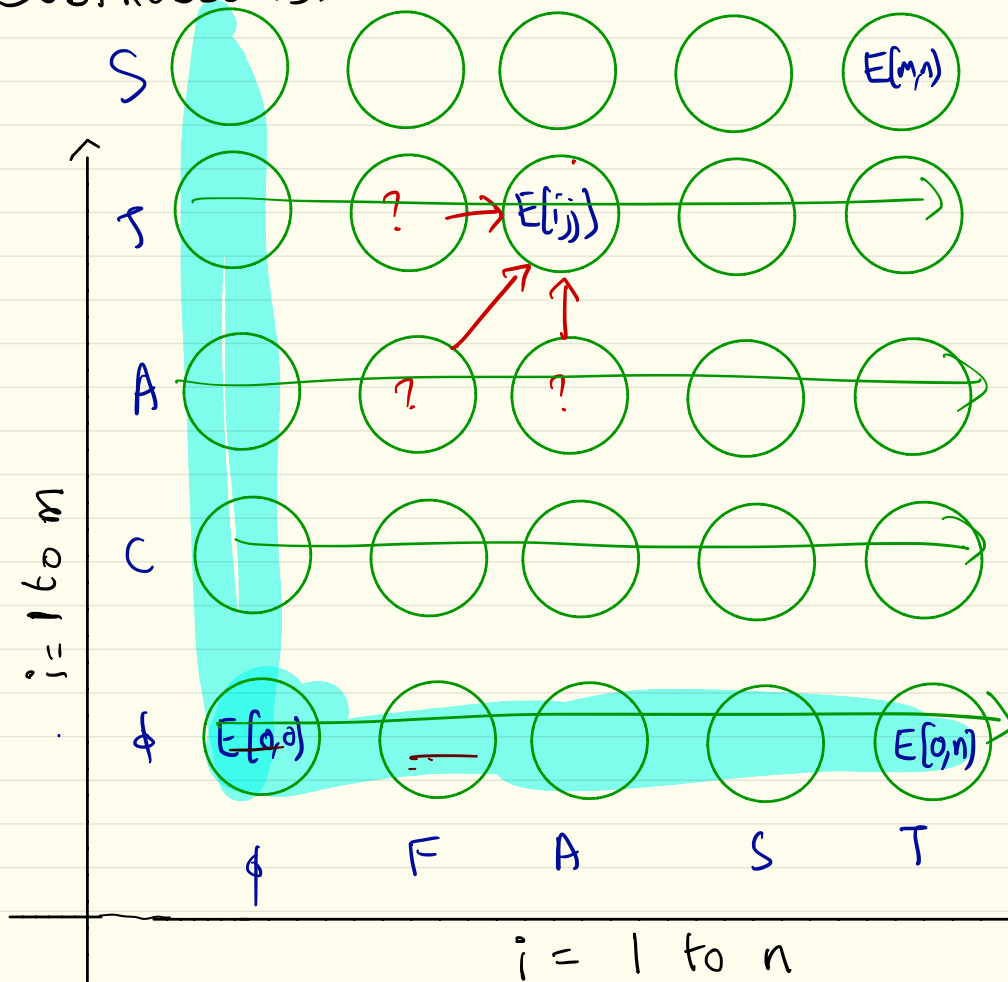
$$1 + E[i-1, j-1]$$

A
 A

if $(x[i] = y[j])$

$$E[i-1, j-1]$$

SUBPROBLEMS:



$$E(0,i) = i$$

$$E(j,0) = j$$

$$E[i,0] = i \quad \forall \quad i = 1 \dots m$$

$$E[0,j] = j \quad \forall \quad j = 1 \dots n$$

for $i = 1$ to m

for $j = 1$ to n

$$E[i,j] =$$

recurrence
relation

return $E[m,n]$

GAMBLING STRATEGY:

Play $n=500$ games in a Casino

Game A: w.p $\frac{1}{2}$ earn 2\$
 w.p $\frac{1}{2}$ lose 2\$

Game B:

w.p	$\frac{2}{3}$	earn	5\$
w.p	$\frac{1}{3}$	lose	5\$

GOAL: SUCCEED if you win exactly 170\$ after n games.

· COMPUTE OPTIMAL STRATEGY.

STRATEGY = ??

$$A \left[\underset{\text{money}}{m \$}, \underset{\text{games left}}{l} \right] = \text{Game A} / \text{Game B}$$

OPTIMAL STRATEGY = Maximizes probability of "SUCCESS"
(having exactly 170\$)
at end of n games

$P[\underset{\substack{\uparrow \\ \text{money} \\ \text{earned}}}{m\$}, \underset{\substack{\uparrow \\ \text{games} \\ \text{left}}}{l}] =$ Probability of Success
for optimal strategy
starting with $m\$$ and
 l games left.

$P[0, n] \leftarrow \text{ANSWER}$

$$P[m \$, l_{\text{games}}^{\text{left}}] = \max \left\{ \begin{array}{l} \text{Game A} \quad \frac{1}{2} P[m+2, l-1] \\ \quad \quad \quad + \frac{1}{2} P[m-2, l-1] \\ \\ \text{Game B} \quad \frac{2}{3} P[m+5, l-1] \\ \quad \quad \quad + \frac{1}{3} P[m-5, l-1] \end{array} \right.$$

"Solve subproblems in increasing order of l ."

$$P(m, 0) = \begin{cases} 1 & \text{if } m = 170 \\ 0 & \text{otherwise} \end{cases}$$