

ALL-PAIR SHORTEST PATHS

$|V| = n$

$\{+ve/-ve\}$

INPUT: Graph $G = (V, E)$ with edge weights $\{w_e\}$

GOAL: Compute shortest paths/distances between all pairs.

(+ve)

Naive: Run Dijkstra from every node i

$\rightarrow n$ - Dijkstras

$\rightarrow |V| \cdot (|E| \log |V|) = |V| \cdot |E| \log |V|.$

Subproblem:

$D[i, j, k]$ = length of the shortest path from $i \rightsquigarrow j$
which only uses $v_1 \dots v_k$ as intermediate vertices



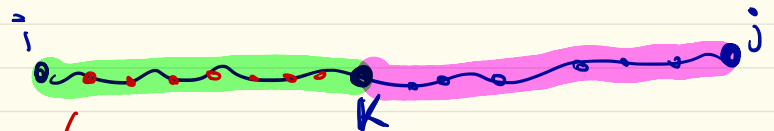
$D[i, j, 0]$ = length of shortest path with NO intermediate vertices

$D[i, j, n] = w_{ij}$
= length of shortest path $i \rightsquigarrow j$



$k \notin \text{Shortest path } i \sim j$
 every intermediate vertex $\in \{1 \sim k-1\}$
 $D[i, j, k-1]$

$$D[i, j, k] = \min$$



$k \in \text{Shortest path } i \sim j \text{ using } \{1..k\}.$

$$D[i, k, k-1] + D[k, j, k-1]$$

$$D[i, j, k] = \min \{ D[i, j, k-1], D[i, k, k-1] + D[k, j, k-1] \}$$

PSEUDO CODE

```
for i = 1 to n
  for j = 1 to n
     $D[i,j,0] = w_{ij}$ 
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```
for k = 0 to n
  for i = 1 to n
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```
    for j = 1 to n
       $D[i,j,k] = \min \left\{ \begin{array}{l} D[i,j,k-1], \\ D[i,k,k-1] \\ \quad + D[k,j,k-1] \end{array} \right\}$ 
```

LINEAR PROGRAMMING - family of optimisation problems

Variables: $x_1, \dots, x_n \in \mathbb{R}$

Max / Min linear function (~~x_1, \dots, x_n~~)

Subject to

linear constraints

COFFEE SHOP:

	COFFEE	MILK
$x = \text{unit of CAPPUCCINO}$	$3x$	$1x$
	$+$	$+$
$y = \text{unit of LATTE}$	$2y$	$4y$
	<hr/>	<hr/>
	≤ 6	≤ 8

AVAILABLE COFFEE = 6

AVAILABLE MILK = 8

Maximum # of drinks = $x + y$

$x = \# \text{ of units of Cappuccino}$

$y = \# \text{ of units of Latte}$

$\in \text{positive real numbers.}$

Max $x + y$

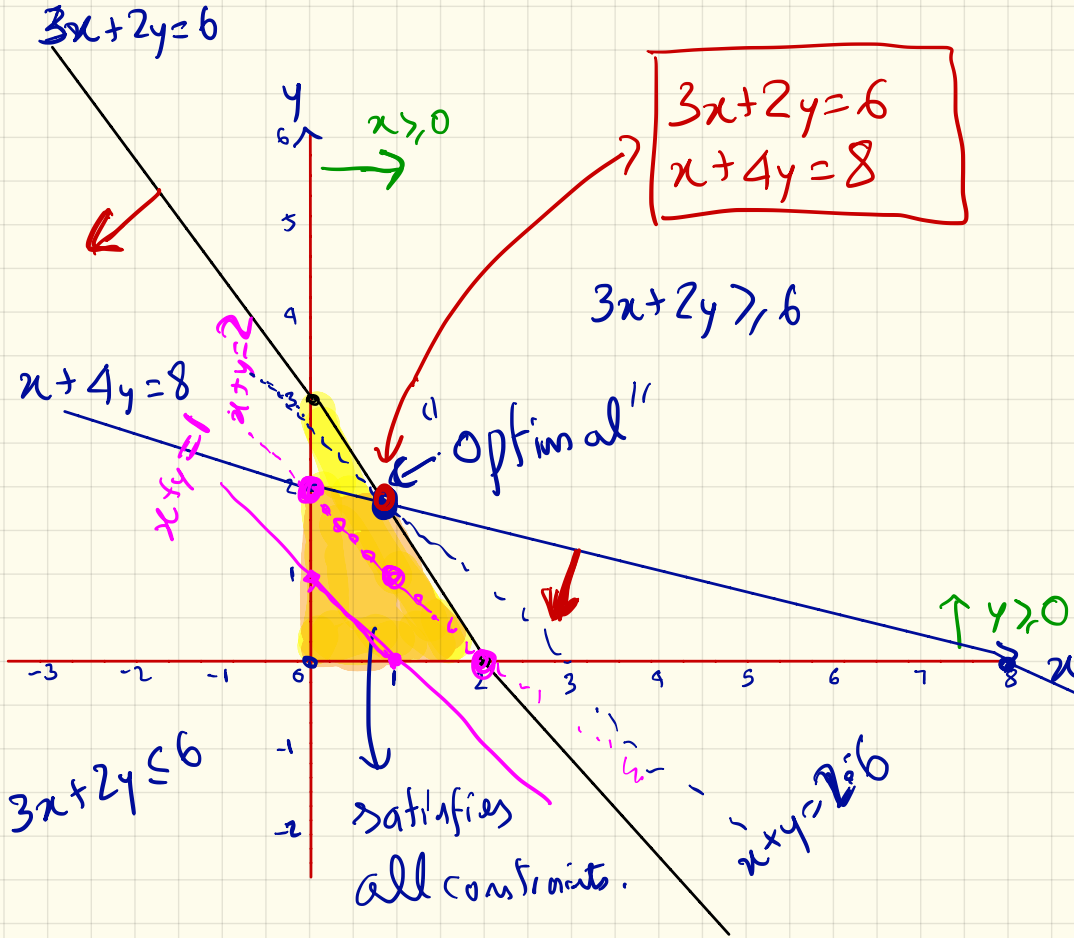
$3x + 2y \leq 6$

$x + 4y \leq 8$

$x \geq 0$

$y \geq 0$

$3x + 2y = 6$



Feasible = satisfy all constraints

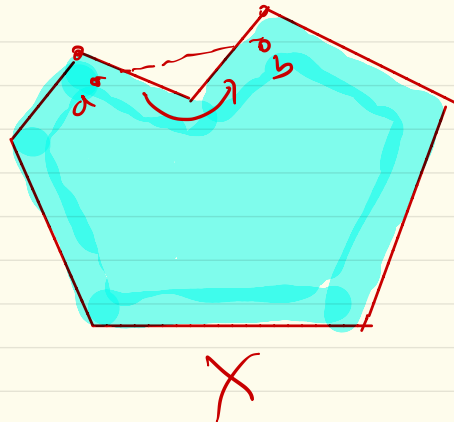
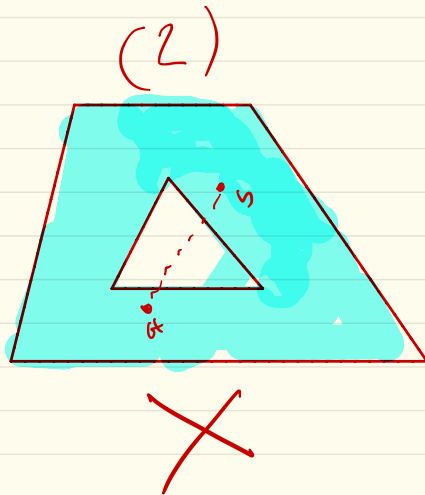
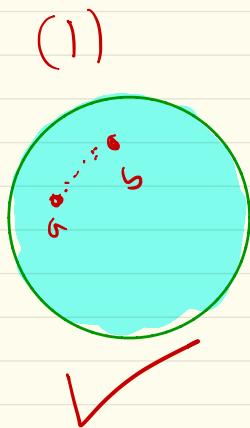
Feasible region = range of variables
that satisfy all constraints.

FACT 1: Feasible region of a linear program
is a **CONVEX POLYGON / Polyhedron**.

FACT 2: The optima of a linear program
occurs at vertex of a polygon /
polyhedron.

FACT 3: Corner / Vertex of polygon
is given by the intersection of constraints

CONVEX: A subset of points $S \subseteq \mathbb{R}^d$ is convex
if $\forall a, b \in S$
(Line joining a, b) $\in S$.



Maximije

$$x + y$$

$$\text{OPT} = "2.6"$$

$$(0.3) \cdot (3x + 2y \leq 6)$$

$$(0.1) \cdot (x + 4y \leq 8)$$

$$0 \cdot (-x \leq 0)$$

$$0 \cdot (-y \leq 0)$$

$$\boxed{x + y \leq 2.6}$$