MATH70132: Mathematical Logic

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Chapter 1

Propositional Logic

Propositional logic is the logic of reasoning and proof. Before we get started with anything formal, here's a motivating example.

Consider the following statement:

If Mr Jones is happy, then Mrs Jones is unhappy, and if Mrs Jones is unhappy, then Mr Jones is unhappy. Therefore, Mr Jones is unhappy.

One can ask ourselves whether it is logically valid to conclude that Mr Jones is unhappy based on the relationship between the happiness of Mr Jones and that of Mrs Jones expressed in the sentence preceding it.

Putting this into symbols, let P denote the statement that Mr Jones is happy, and let Q denote the statement that Mrs Jones is unhappy. We can express the statement as follows:

$$((P \implies Q) \land (Q \implies \neg P)) \implies (\neg P) \tag{1.0.1}$$

This disambiguation, by removing any question of marital harmony from what is otherwise a purely logical question, we can manually check whether (1.0.1) is a valid statement by constructing a truth table.

We will begin by developing some machinery to reason about these sorts of statements more formally.

1.1 Propositional Formulae

1.1.1 Propositions and Connectives

We begin by defining the notion of a proposition.

Definition 1.1.1 (Proposition). A **proposition** is a statement that is either true or false.

Convention. We will denote the state of being true by **T** and that of being false by **F**.

Propositions can be connected to each other using tools known as **connectives**. These can be thought of as **truth table rules**.

Convention. Before we define the actual connectives we shall use, we list them down, along with notation.

- 1. Conjunction (\land)
- 2. Disjunction (∨)
- 3. Negation (\neg)
- 4. Implication (\rightarrow)
- 5. The Biconditional (\leftrightarrow)

In particular, we will only use the \implies and \iff symbols when reasoning **informally**. For **formal** use, we will stick to the \rightarrow and \leftrightarrow symbols.

We define them exhaustively as follows.

Definition 1.1.2 (Connectives). Let p and q be true/false variables. We define each of the connectives listed above to take on truth values depending on those of p and q as follows.

p	q	$(\neg p)$	$(\neg q)$	$(p \wedge q)$	$(p \lor q)$	(p o q)	$(p \leftrightarrow q)$
Т	Т	F	F	Т	Т	Т	Т
						F	
						Т	
F	F	Т	Т	F	F	Т	Т

We are now ready to define the main object of study in this section: propositional formulae.

Definition 1.1.3 (Propositional Formula). A **propositional formula** is obtained from propositional variables and connectives via the following rules:

- (i) Any propositional variable is a propositional formula.
- (ii) If ϕ and ψ are formulae, then so are $(\neg \phi)$, $(\neg \psi)$, $(\phi \land \psi)$, $(\phi \lor \psi)$, $(\phi \to \psi)$, $(\psi \to \phi)$, and $(\phi \leftrightarrow \psi)$.
- (iii) Any formula arises in this manner after a finite number of steps.

What this means is that a propositional formula is a string of symbols consisting of

- 1. variables that take on true/false values,
- 2. connectors that express the relationship between these variables, and
- 3. parentheses/brackets that separate formulae within formulae and specify the order in which they must be evaluated when the constituent variables are assigned specific values.

In particular, every propositional formula is either a propositional variable or is built from 'shorter' formulae, where by 'shorter' we mean consisting of fewer symbols.

Convention. Throughout this module, we will adopt two important conventions when dealing with propositional formulae.

- 1. All propositional formulae, barring those consisting of a single variable, shall be enclosed in parentheses.
- 2. When we want to denote a propositional formula by a certain symbol, we will use the notation "symobol: formula".

As a concluding remark on the nature of propositional formulae, we will note that just as we use trees to evaluate expressions on the computer when performing arithmetic, we can use them to express and evaluate propositional formulae as well. We will not usually do this, however, as it takes up a lot of space. In any event, we would first need to make precise the notion of *evaluating* a propositional formula. For this, we will turn to the concept of a truth function.

1.1.2 Truth Functions

Any assignment of truth values to the propositional variables in a formula ϕ determines the truth value for ϕ in a **unique** manner, using the exhaustive definitions of the connectives given in Definition 1.1.2. We often express all possible values of a propositional formula in a **truth table**, much like we did in Definition 1.1.2 when defining the connectives.

Example 1.1.4. Consider the formula $\phi:((p\to (\neg q))\to p)$, where p and q are propositional variables. We construct a truth table as follows.

p	q	$(\neg q)$	$(p \to (\neg q))$	((p o (eg q)) o p)
Т	Т	F	F	Т
_	-	Т	Т	Т
F	Т	F	Т	F
F	F	F	Т	F

From this table, it is clear that the truth value of ϕ depends on the truth values of p and q in some manner (to be perfectly precise, it only depends on the truth value of p, and is, in fact, biconditionally equivalent to p). We would like to have a formal notion of navigating this dependence to 'compute a value for ϕ given values of p and q'.

Throughout this subsection, *n* will denote an arbitrary natural number.

Definition 1.1.5 (Truth Function). A truth function of *n* variables is a function

$$f: \{\mathsf{T}, \mathsf{F}\}^n \to \{\mathsf{T}, \mathsf{F}\}$$

These are very directly related to propositional formulae.

Definition 1.1.6 (Truth Function of a Propositional Formula). Let ϕ be a propositional formula whose variables are p_1, \ldots, p_n . We can associate to ϕ a truth function whose truth value at any $(x_1, \ldots, x_n) \in \{\mathbf{T}, \mathbf{F}\}^n$ corresponds to the truth value of ϕ that arises from setting p_i to x_i for all $1 \leq i \leq n$. We define this truth function to be the **truth function** of ϕ , denoted F_{ϕ} .

Bibliography

These lecture notes are based heavily on the following references:

- [1] Nicolas Bourbaki. *Groupes et Algèbres de Lie. Chapitre 1*. 1st ed. Éléments de Mathématique. Édition originale publiée par Hermann, Paris, 1972. Springer Berlin, Heidelberg, 2007. ISBN: 978-3-540-35337-9. DOI: https://doi.org/10.1007/978-3-540-35337-9.
- [2] Karin Erdmann and Mark J. Wildon. *Introduction to Lie Algebras*. Springer Undergraduate Mathematics Series. Springer, 2006.
- [3] Anna Lachowska et al. *MATH-335: Coxeter Groups*. Lecture Notes. École Polytechnique Fédérale de Lausanne, Spring 2020.
- [4] Martin Liebeck and Aleksander Horawa. *M4P46: Lie Algebras.* Lecture Notes. Imperial College London, Winter 2016.
- [5] Aluna Rizzoli and Sidharth Hariharan. *MATH-314: Representation Theory of Finite Groups*. Lecture Notes. École Polytechnique Fédérale de Lausanne, Spring 2024. URL: https://github.com/thefundamentaltheor3m/RepTheoryEPFL.

For the latest version of these notes, visit https://github.com/thefundamentaltheor3m/LieAlgebrasNotes. For any suggestions or corrections, please feel free to fork and make a pull request to my repository.