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01-13 4192 030

BSCS-4B

Q3

Part a)

$$L = \{0^m 1^n \mid n = 3m\}$$

Proof

- - By contradiction, let L be regular
- - P/L constant should exist

→ Let $N = P/L$ constant

- - consider $w = 0^{3N} 1^N$ s.t. $|w| \geq N$
- - By pumping lemma, we should be able to break $w = xyz$ such that

1) $y \neq \epsilon$

2) $|xy| \leq N$

3) for all $k \geq 0$, the string xy^kz is also in L .

- - Now divide the w into 3 parts by the rule

$$x = 0^{2N}, y = 0^N, z = 1^N$$

- - By (3), any string of the form $xy^kz \in L$ for all $k \geq 0$

$$w = xy^kz$$

$$= xy^0z$$

$$= 0^{2N} \cdot (0^N)^0 \cdot 1^N$$

$$= 0^{2N} 1^N \notin L$$

L is not regular so is a contradiction ~~so~~ then

$L \{ 0^m 1^n \mid m = 3n \}$ is not regular language

Q3

Part b

$$w = 0^{3N} 1^N$$

$$|w| = 3N + 1 = \text{~~4N + 1~~}$$

then $|w| > N \Rightarrow 4N > N$

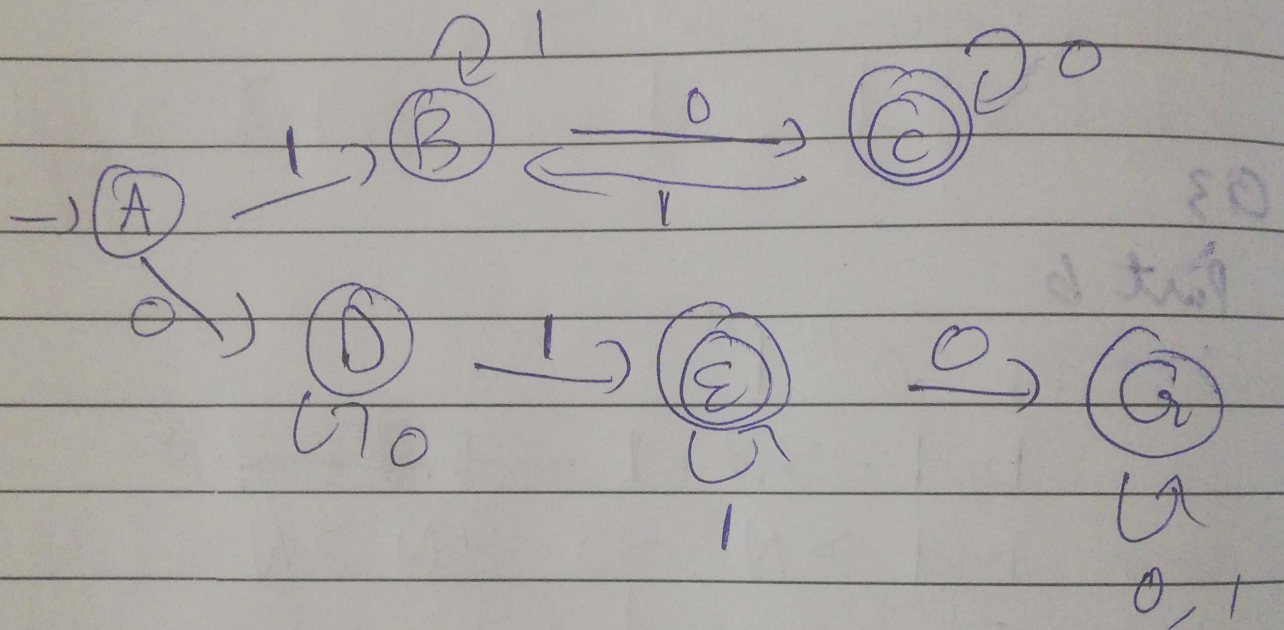
Proved because it is true i.e. $|w| > N$ is true

Q2

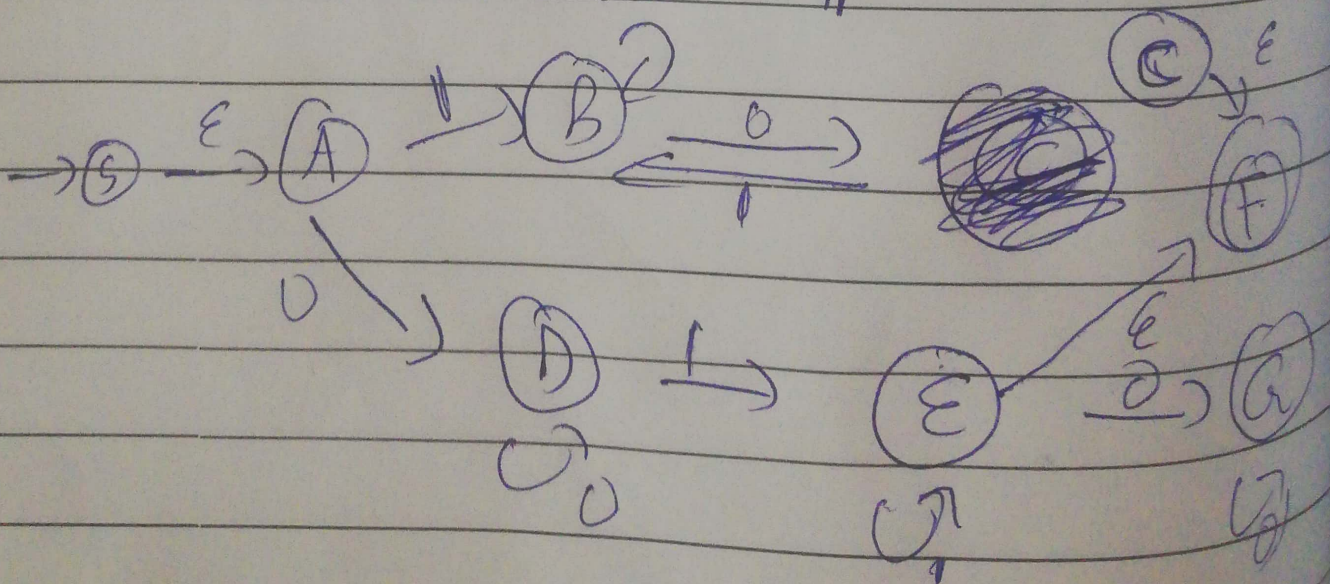
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even

So states will be removed according to the order ϵ, C, B, D, G, A



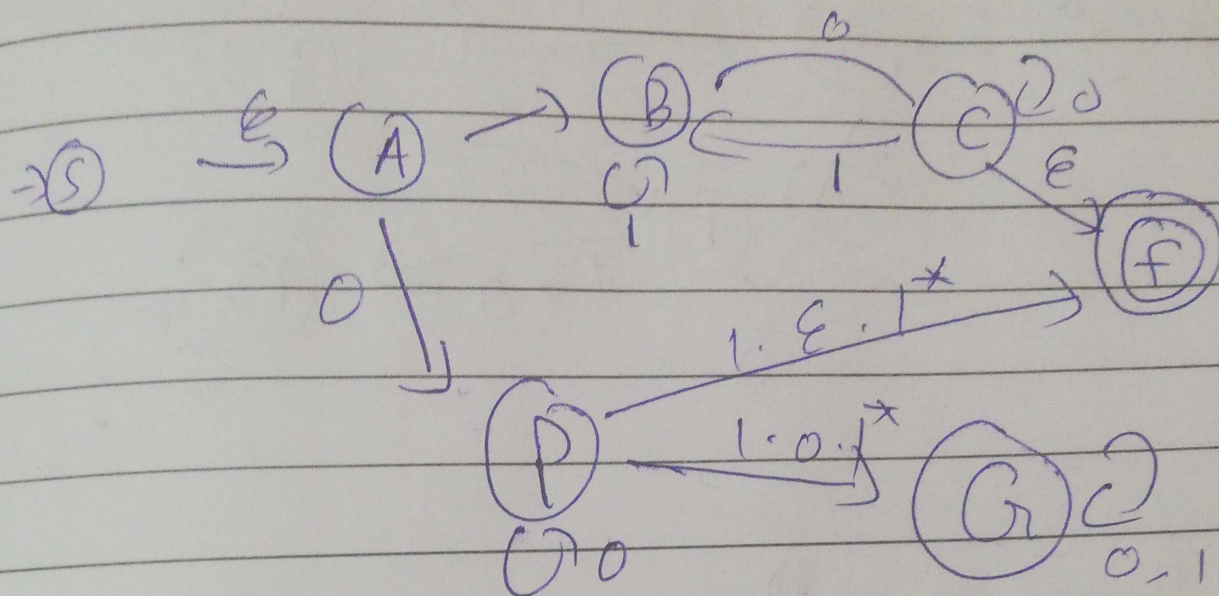
1) First we insert start node s and end node 'F'



2) Remove E

$$D \cdot E \cdot G = D \cdot G = 1 + 1 \cdot 0 \cdot 1^*$$

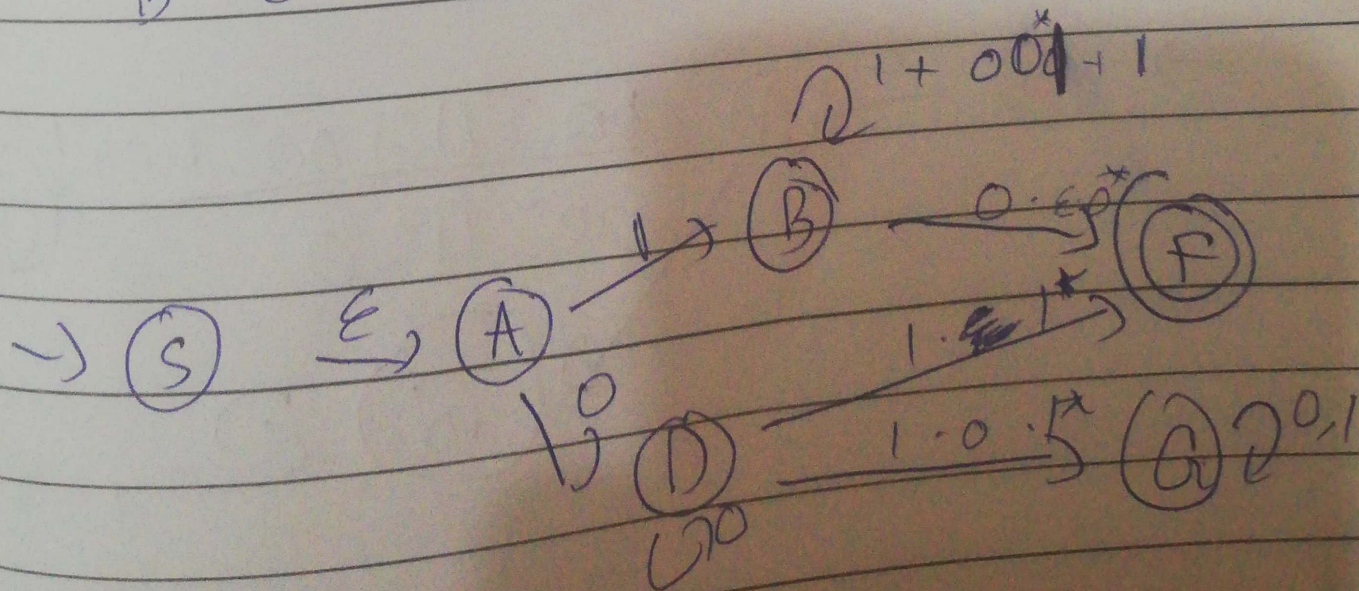
$$D \cdot E \cdot F = D \cdot F = \emptyset + 0 = \emptyset + 1 \cdot E \cdot 1^*$$



3) Remove C

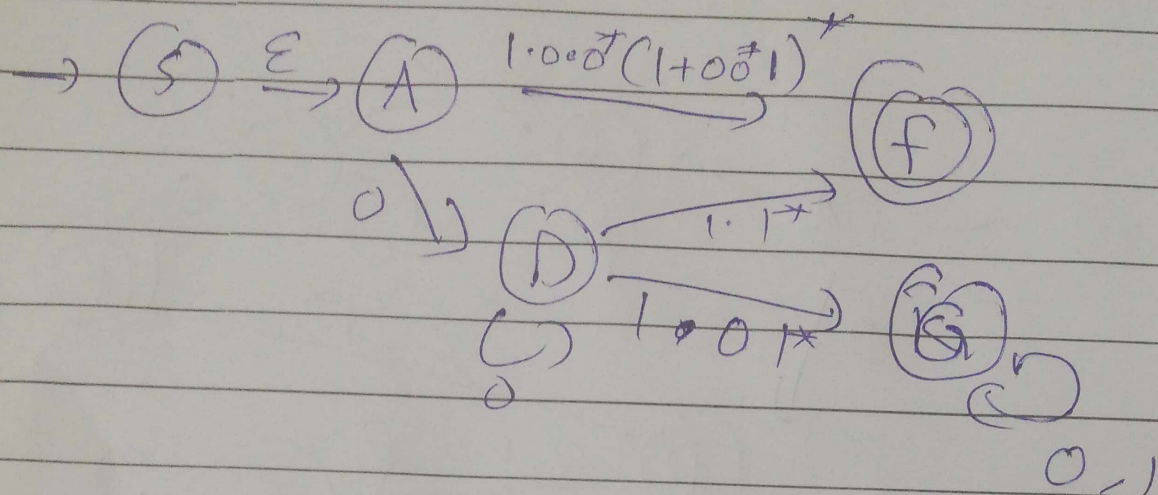
$$B \cdot C \cdot F \rightarrow B \cdot F = \emptyset + 0 \cdot E \cdot 0^*$$

$$B \cdot C \cdot B \rightarrow B \cdot B = 1 + 0 \cdot 1 \cdot 0^*$$



4) Remove B

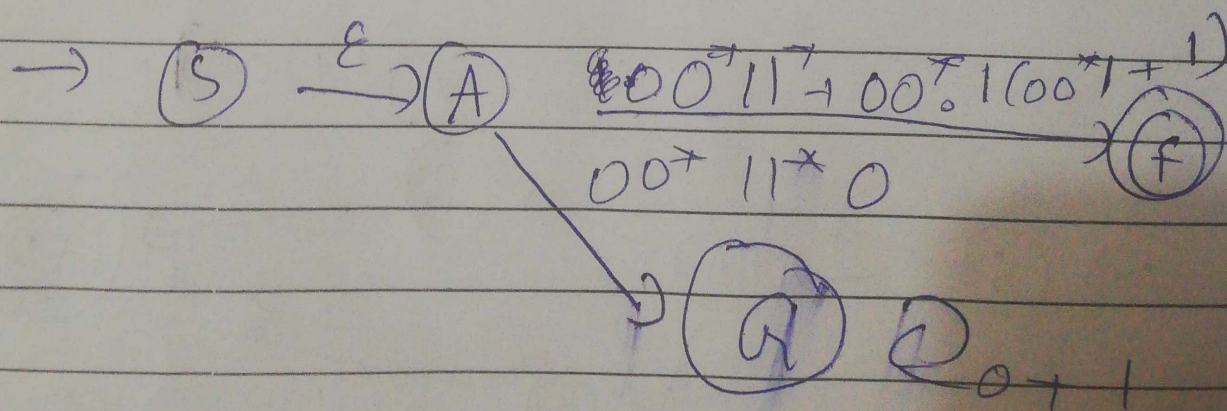
$$A \cdot B \cdot f \Rightarrow A f = 1 + 1 \cdot 0 \cdot 0^* (1 + 00^*1)^*$$



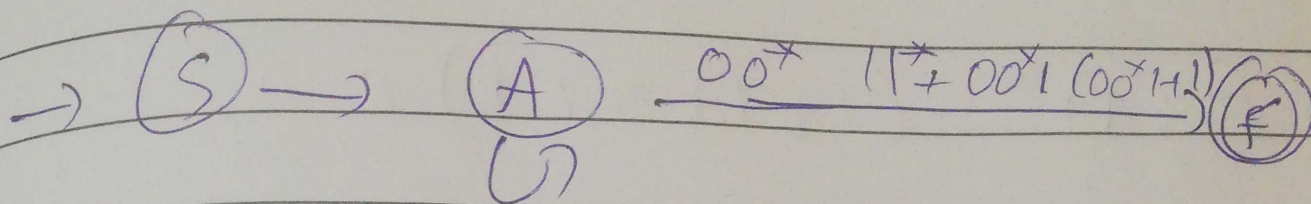
5) Remove D

$$A D G = A G = \emptyset + 0 \cdot 0 \cdot 0^* (1; 01^*)$$

$$A D f = A f = 1 \cdot 00^* (1 + 00^*1)^* + 00^*1 \cdot 1$$

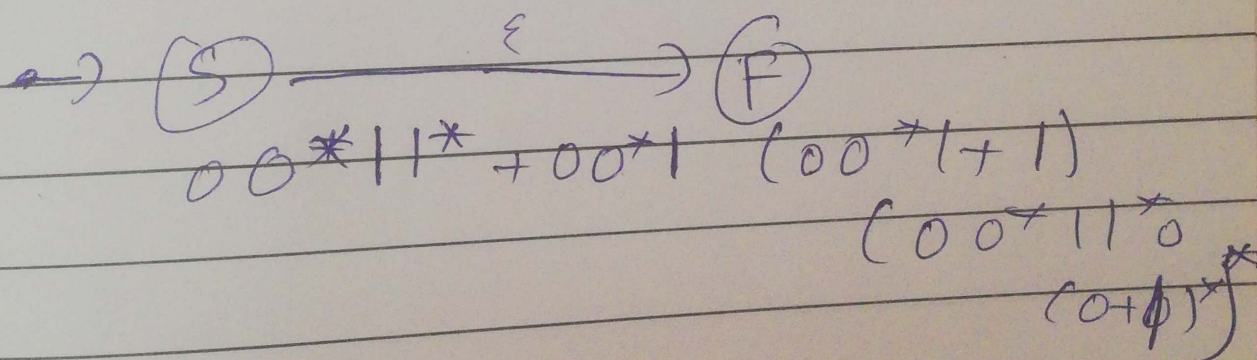


6) Remove A



$$0 \cdot 0^* (1 \cdot \phi^* \cdot 0^*) (0 + 1)^*$$

7)



a1 Part b

030

11110

$$(0+1)^* (11110) (0+1)^*$$

Part 2