

Minimax Algorithm

① Intro: The minimax algorithm is an algorithm which can help make a best choice in a situation involving worst-case scenarios.

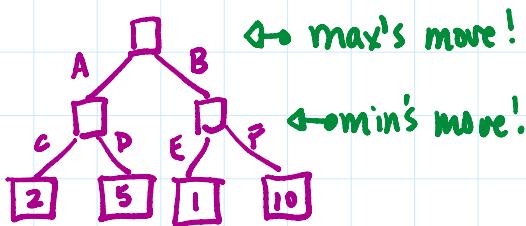
Sps Max and Min are playing a game.

Sps the state of the game can be rep. by a number
Sps it's max's turn and max has 2 choices:



Best choice for Max is move A to get 8.

Now let's complicate it:

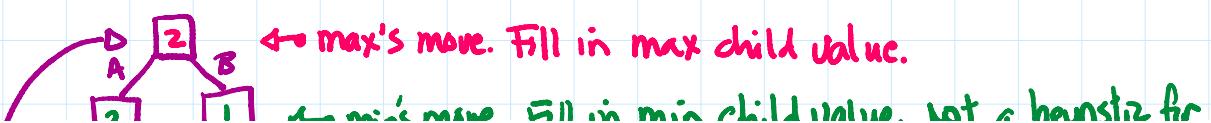


If Max sees 10 and picks B then Min will choose E, yielding 1.
So Max should choose A then min will choose C, yielding 2.
That's the best poss. for max!

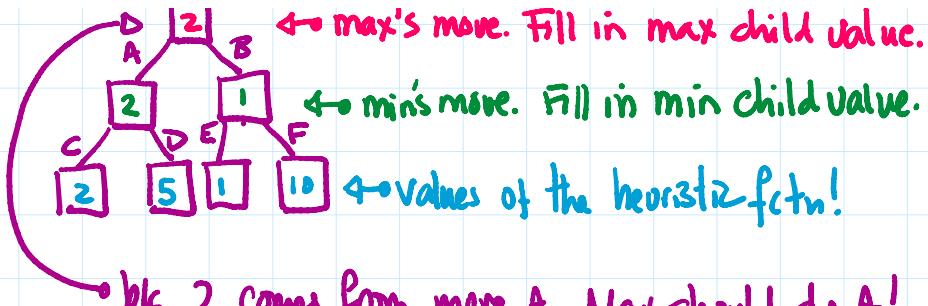
② Algorithm: The situation is:

- Max's move
- Max can look some # moves ahead, rep. by a tree.
- There is a heuristic fact which assigns a number to each leaf situation (game position).
- if we fill in those leaves, how do we then get back to the root to determine Max's best move?!

By



by



• blk 2 comes from move A, Max should do A!

③ Pseudocode: For filling in values:

```
function minimax(node nd, depth):
    if nd.value is not undefined: ← it's a leaf
        return(nd.value)
    else:
        if depth is even: ← max's turn
            return( max(minimax(ndc, depth+1) for all children ndc of nd) )
        else: ← min's turn.
            return( min(minimax(ndc, depth+1) for all children ndc of nd) )
        end if
    end if
end function
```

Have a tree!
Heuristic has been applied
to the leaves!

Call: minimax(root, 0)

Time Complexity: Function executes once for each node,
hence if n nodes its $\Theta(n)$.

Another way: if tree depth = d
if max branch factor (# children) = b
then

$$n \leq \underbrace{1 + b + b^2 + \dots + b^d}_{\frac{b^{d+1} - 1}{b - 1}}$$

about the best we can do!

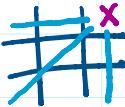
④ Tic-Tac-Toe

Max plays X, Min plays O. Max goes first.

Say: for a board, an open winning line (OWL) is a row/col/diag which contains at least one of the player's marks and none of the opponent's marks.

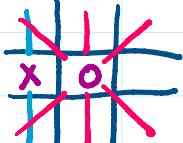
Ex  max has 3 OWLs.

ex



max has 3 owls.

ex



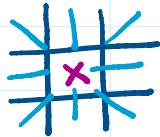
max has 1 owl

min has 3 owls.

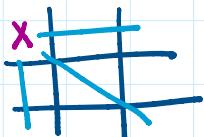
define heuristic: $h = (\# \text{owls X has}) - (\# \text{owls O has})$

- First, max can look 1 move ahead!

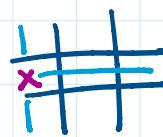
Max has 3 poss moves: Center, corner, edge (up to symmetry)



$$h = 4 - 0$$



$$h = 3 - 0$$



$$h = 2 - 0$$

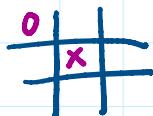
in a tree:



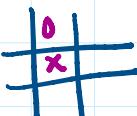
→ play center!

- Now two moves ahead!

If max plays center, min has 2 moves:

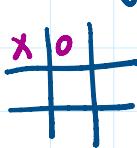


$$h = 3 - 2 = 1$$

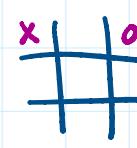


$$h = 3 - 1 = 2$$

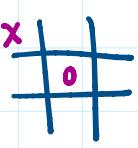
If max plays corner, min has 5 moves



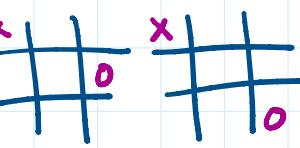
$$h = 2 - 1 = 1$$



$$h = 2 - 2 = 0$$



$$h = 2 - 3 = -1$$

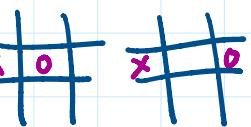
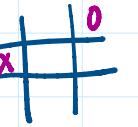
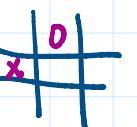
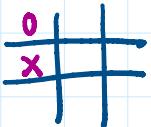


$$h = 3 - 2 = 1$$



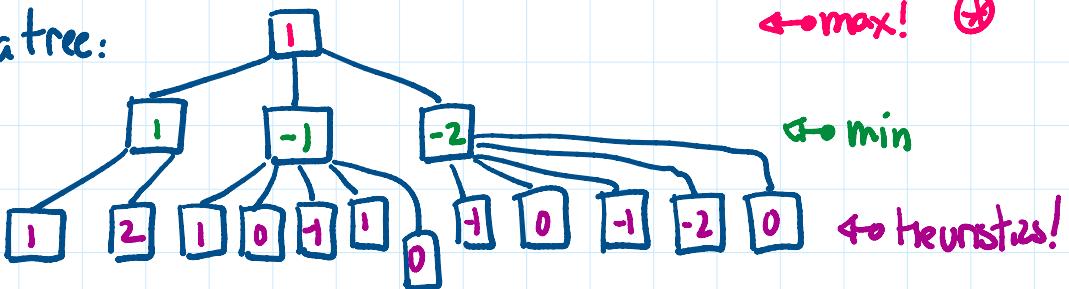
$$h = 2 - 2 = 0$$

If max plays edge, min has 5 moves



$$\begin{array}{c} \times \quad + \\ \cancel{\times} \quad \cancel{+} \\ h = 1 - 2 = -1 \end{array}
 \quad
 \begin{array}{c} \times \quad + \\ \cancel{\times} \quad \cancel{+} \\ h = 2 - 2 = 0 \end{array}
 \quad
 \begin{array}{c} \times \quad + \\ \cancel{\times} \quad \cancel{+} \\ h = 2 - 3 = -1 \end{array}
 \quad
 \begin{array}{c} \times \quad 0 \\ \cancel{\times} \quad + \\ h = 1 - 3 = 2 \end{array}
 \quad
 \begin{array}{c} \times \quad + \\ \cancel{\times} \quad \cancel{+} \\ h = 1 - 1 = 0 \end{array}$$

As a tree:



⊗ Says max should play the center!

Note: Not always one best move!