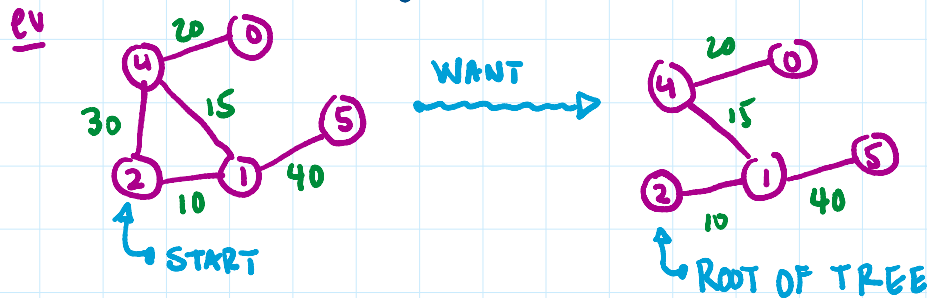


DIJKSTRA'S ALGORITHM

- ① Idea: we have a weighted, undirected simple ^{Connected} graph.
 we choose a starting vertex s .
 we'd like to know how to get from s to every other vertex using paths of min. total weight.



- ② Algorithm: Given a graph, and a starting vertex s .
 Create a set $S = \{s\}$ (will contain vertices)
 Assign each vertex a distance of ∞ except $d[s] = 0$.
 Create a list of predecessors p of size V all $NULL$.
 Then:

Choose a vertex x of min dist. which is not in S .

Put it in S

For all y adjacent to x :

if $d[x] + w(x,y) < d[y]$:

$d[y] = d[x] + w(x,y)$

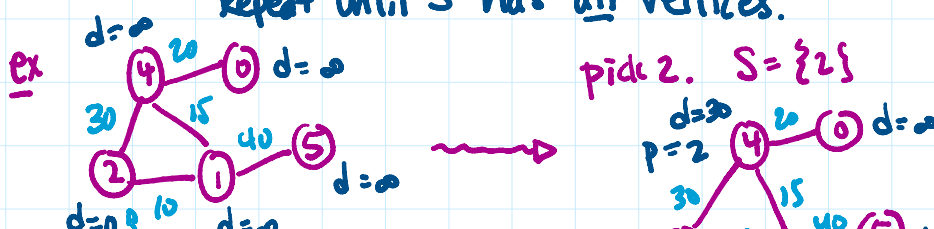
$p[y] = x$

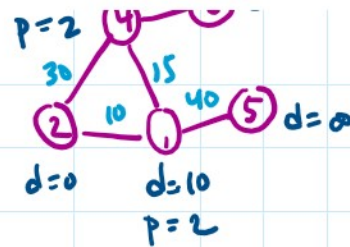
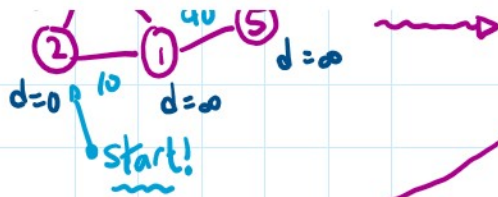
endif

end for

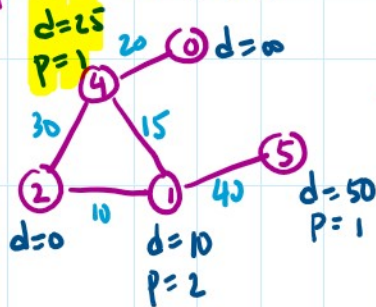
Repeat until S has all vertices.

$w(x,y)$ = edge weight

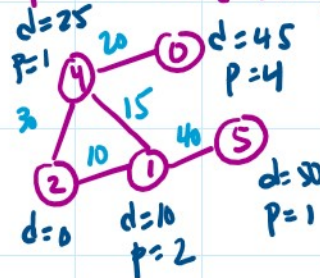




pick 1: $S = \{2, 1\}$



pick 4: $S = \{2, 1, 4\}$



pick 0: $S = \{2, 1, 4, 0\}$

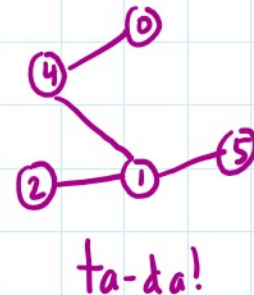
no change

pick 5: $S = \{2, 1, 4, 0, 5\}$

no change

DONE!

Note $P[0]=4$
 $P[1]=2$
 $P[2]=NULL$
 $P[4]=1$
 $P[5]=1$



③ Pseudocode.

\\ PRE: G is a graph with V vertices.

\\ PRE: s is the starting vertex.

def dijkstra(G, start):

 dist = [inf, ..., inf] of length V.

 pred = [NULL, ..., NULL] of length V.

 S = []

 dist[start] = 0

 while len(S) != V $\leftarrow V \leftarrow V$ times

 x = vertex in G-S with smallest distance

 for each edge in the edge set of x

 let y be the connected vertex

 if dist[x] + (Weight of Edge x,y) < dist[y]:

 dist[y] = dist[x] + (Weight of Edge x,y)

 pred[y] = x

 endif

 endfor

 append x to S

 endwhile

 return(pred)

end

} $\Theta(V)$

\leftarrow if S=list, this is $\Theta(V)$.

W/out for loop it's $\Theta(V^2)$;
 body of for loop takes
 $\Theta(2E)$ time (as before)

total $\Theta(V^2 + 2E)$
 $= \Theta(V^2 + E)$.

note: can be mod. using min heap (+ other)
 to make it $\Theta(E \log V)$.

note: can be mod. using min heap (1 time)
to make it $\theta(E \lg V)$.