

Probabilistic program induction of symbols

Or: More complex models!

- The Language of Thought: computational cognitive science approaches to category learning
- Who: Fausto Carcassi
- When: Sommer semester 2022



Where are we?

- Last week, we have seen two applications of the idea of pLoT to new conceptual domains: handwritten digits and kinship systems.
- This week, we'll look at a few more applications.
- It'll be a bit of a whirlwind!
- We will look at:
 - Learning numerals in an LoT
 - Learning abstract visual concepts in an LoT
 - Learning sequences in an LoT



- Piantadosi et al (2012), Bootstrapping in a language of thought: A formal model of numerical concept learning.
- Children exhibit very regular patterns in the way they learn number systems.
- The first learn to recognize small sets of size 1, then 2, then 3, etc.
 - In Carey's formulation, early number-word meanings are represented using mental models of small sets. For instance two-knowers might have a mental model of "one" as {X} and "two" as {X, X}. These representations rely on children's ability for enriched parallel individuation, a representational capacity that Le Corre and Carey (2007) argue can individuate objects, manipulate sets, and compare sets using one-to-one correspondence. Subset-knowers can, for instance, check if "two" applies to a set S by seeing if S can be put in one-to-one correspondence with their mental model of two, {X, X}
- Then at about 3;6 they learn the full recursive system (Cardinal Principal learners)
- This is a qualitative jump rather than continuous smooth progress.



- Let's see if we can reproduce this qualitative conceptual jump with an LoT!
- We have at least three choices for how to set up the LoT. Each sentence in the LoT could be:
 - A function from a number word to a predicate of sets
 - A function from a set to a number word
 - A function that constructs a set from the objects in the situation
- Children can do all these three things, and there is no clear empirical evidence one way or the other.
- In the paper, they go the second way: from a set to a number word.



• Rules in the LoT:

Functions mapping sets to truth values

(singleton? X)Returns true iff the set X has exactly one element(doubleton? X)Returns true iff the set X has exactly two elements(tripleton? X)Returns true iff the set X has exactly three elements

Functions on sets

(set-difference X Y)Returns the set that results from removing Y from X(union X Y)Returns the union of sets X and Y(intersection X Y)Returns the intersect of sets X and Y(select X)Returns a set containing a single element from X

Logical functions

(and P Q)Returns TRUE if P and Q are both true(or P Q)Returns TRUE if either P or Q is true(not P)Returns TRUE iff P is false(if P X Y)Returns X iff P is true, Y otherwise

Functions on the counting routine

(next W)Returns the word after W in the counting routine(prev W)Returns the word before W in the counting routine(equal-word? W V)Returns TRUE if W and V are the same word

Recursion

(L S) Returns the result of evaluating the entire current lambda expression on set S



• Recursion is the only one that's a bit complicated. What do you think the following does?

```
\lambda \ S \cdot (if(singleton? \ S)
"one"

(next \ (L \ (select \ S)))).
```

```
One-knower
                                       Two-knower
                                                                                                           Singular-Plural
                                                                                                                                                  Mod-5
      \lambda S. (if (singleton? S)
                                                       \lambda S. (if (singleton? S)
                                                                                                                                                        \lambda S. (if (or (singleton? S)
                                                                                                                  \lambda S. (if (singleton? S)
                 "one"
                                                                 "one"
                                                                                                                                                                    (equal-word? (L (set-difference S)
                                                                                                                            "one"
                                                                (if (doubleton? S)
                undef)
                                                                                                                                                                                      (select S)
                                                                                                                            "two")
                                                                     "two"
                                                                                                                                                                                 "five"))
                                                                    undef))
                                                                                                                                                                   "one"
                                                                                                                                                                   (next (L (set-difference S
Three-knower
                                       CP-knower
                                                                                                                                                                                           (select S)))))
                                             \lambda S . (if (singleton? S)
    \lambda S . (if (singleton? S)
                                                                                                           2-not-1-knower
                                                                                                                                                  2N-knower
              "one"
                                                       "one"
                                                                                                                                                    \lambda S . (if (singleton? S)
              (if (doubleton? S)
                                                       (next (L (set-difference S
                                                                                                                  \lambda S . (if (doubleton? S)
                  "two"
                                                                                                                                                              "one"
                                                                               (select S)))))
                                                                                                                            "two"
                 (if (tripleton? S)
                                                                                                                           undef)
                                                                                                                                                              (next (next (L (set-difference S (select S))))))
                    "three"
                    undef))
```

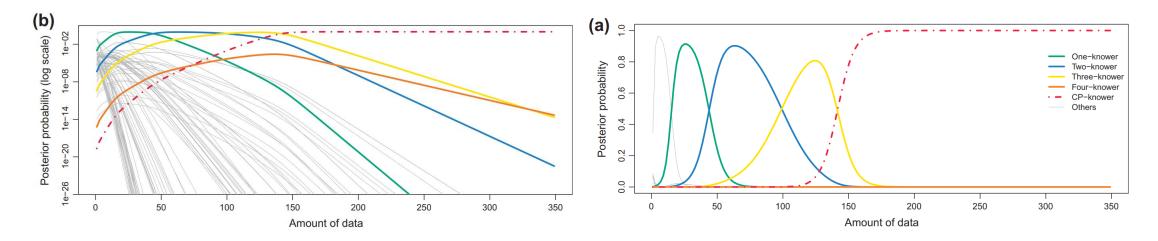


- The likelihood function is pretty typical:
 - First, a set of objects is chosen from the universe of object, e.g. 'cats'
 - Second, the hypothesis is evaluated on the set.
 - This can result either in a number word or 'undef'
 - If the result is 'undef', a random number word is produced
 - If the result is a number word, the word is produced with probability α and with 1α a random word is picked.
- Resulting likelihood function:

$$P(w_i|t_i,c_i,L) = \begin{cases} \frac{1}{N} & \text{if } L \text{ yields } undef \\ \alpha + (1-\alpha)\frac{1}{N} & \text{if } L \text{ yields } w_i \\ (1-\alpha)\frac{1}{N} & \text{if } L \text{ does not yield } w_i \end{cases}$$



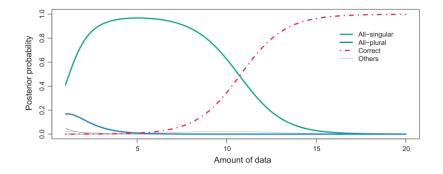
- The model also penalizes recursive functions with a parameter γ
- Results look strikingly like human learning patterns:



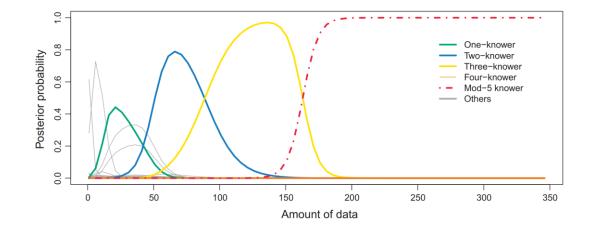
• In particular, note all the hypotheses that are considered and then disregarded!



• The same LoT can also learn things like singular/plural morphology:

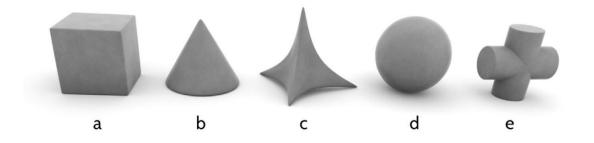


• And mod-n systems (e.g. days of the week)



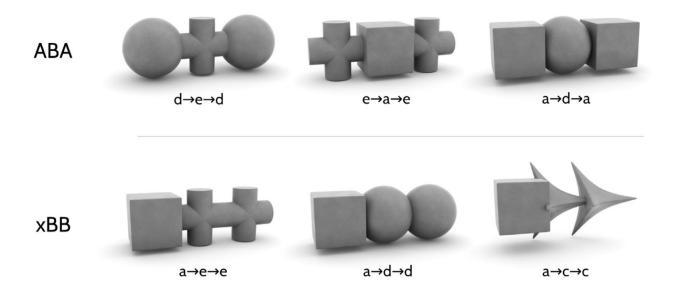


- Overlan et al (2017), Learning abstract visual concepts via probabilistic program induction in a Language of Thought
- Last week, we have seen a model that can learn and do various other things with handwritten characters.
- Now let's see if we can work with something else in the visual modality, namely the structure of 3-d objects.
- Suppose we have the following primitive objects, and we can combine them in various ways:



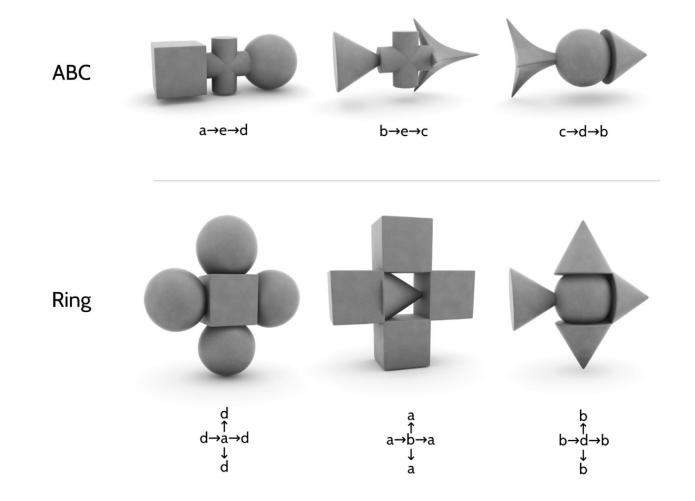


• For instance, we can produce the following categories:



• Note that each category identifies a *class* of objects!







What could a grammar for this look like?

```
START \rightarrow let < BV\_PART > :x_1 = FIRST\_PART; EXPR
           \mathbf{EXPR} \rightarrow \mathbf{let} < \mathbf{BV\_PART} > : x_n = \mathbf{PART}; \mathbf{EXPR}
                     \rightarrow STRING
       STRING \rightarrow BV\_PART
                     \rightarrow STRING CONNECT STRING
                     \rightarrow \{ \mathbf{STRING} \}
FIRST\_PART \rightarrow \texttt{sample}(FIRST\_SET)
                                                                                                      1 - p_{single}

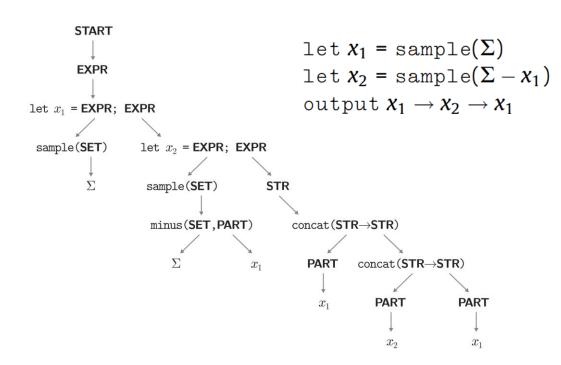
ightarrow SINGLE
                                                                                                      p_{single}
           \mathbf{PART} \to \mathbf{BV\_PART}
                                                                                                      (1-p_{single})/2

ightarrow sample(SET)
                                                                                                      (1-p_{single})/2

ightarrow SINGLE
                                                                                                      p_{single}
  \mathbf{FIRST\_SET} \to \Sigma
                                                                                                      1-p_{minus}
                     \rightarrow minus(FIRST_SET, FIRST_PART)
                                                                                                      p_{minus}
              \mathbf{SET} \to \Sigma
                                                                                                      1-p_{minus}
                     \rightarrow minus (SET, BV_PART)
                                                                                                      p_{minus}
   \mathbf{CONNECT} \to '\uparrow' \ | \ '\downarrow' \ | \ '\leftarrow' \ | \ '\to'
       SINGLE \rightarrow 'a' | ... | 'e'
```



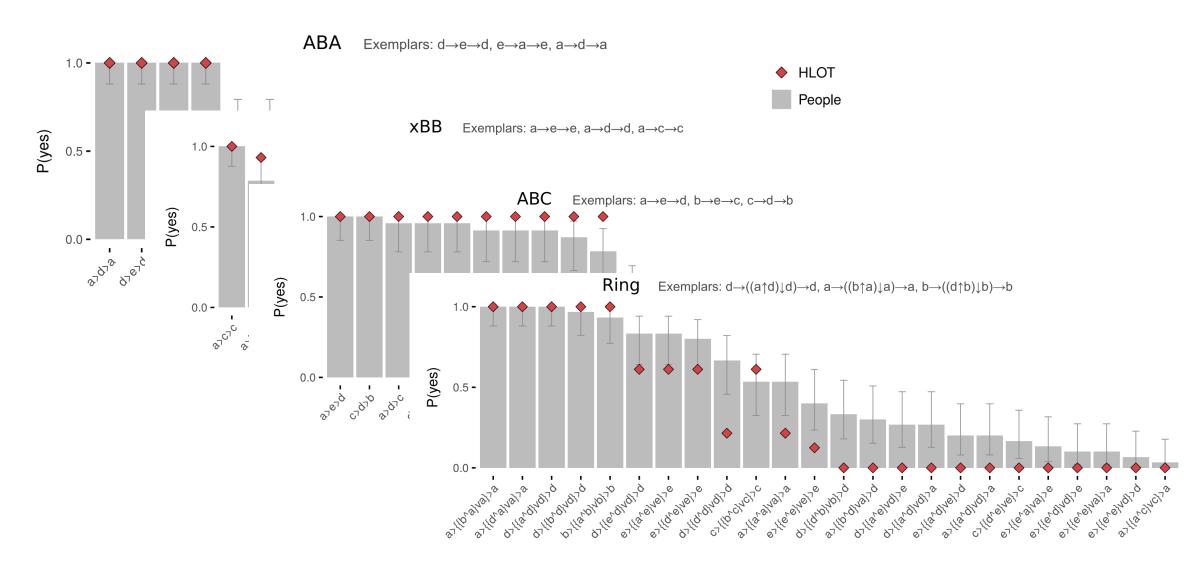
• Example of a derivation and some categories:



ABAxBBlet $x_1 = \text{sample}(\Sigma_R)$ let $x_1 = \text{sample}(\Sigma)$ let x_2 = sample($\Sigma_R - x_1$) let $x_2 = `a'$ output $x_1 \rightarrow x_2 \rightarrow x_1$ output $x_2 \rightarrow x_1 \rightarrow x_1$ ABCRinglet $x_1 = \text{sample}(\Sigma)$ let $x_1 = \text{sample}(\Sigma_R)$ let x_2 = sample($\Sigma - x_1$) let x_2 = sample($\Sigma_R - x_1$) let x_3 = sample($\Sigma - x_2 - x_1$) output $x_1 \to ((x_2 \uparrow x_1) \downarrow x_1) \to x_1$

output $x_2 \rightarrow x_3 \rightarrow x_1$



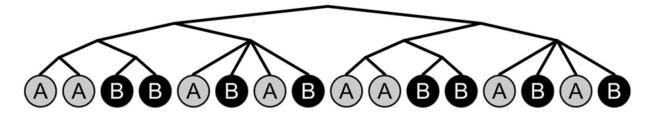




- Planton et al (2021), A theory of memory for binary sequences: Evidence for a mental compression algorithm in humans
- The scope here is to understand how humans deal with *binary sequences*, i.e. sequences composed of only two elements.
 - For instance, {0,1}*
 - But not that it doesn't need to be symbols. E.g. it can be two pitches.
- Much literature has been devoted to understand how humans learn these.
- Usually, the general strategy is to find a mechanism where strings can be encoded, which explains how long strings can be learned, which would be impossible with pure memorization.



• In this case, we'll use an LoT where each sentence can encode a binary series. For instance:



LoT program expression:

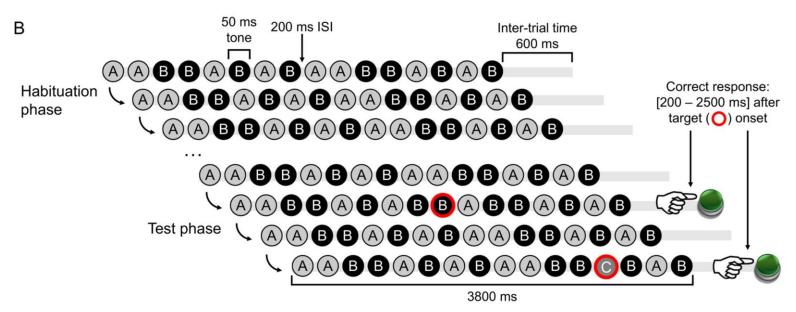
[[[+0]^2]^2,[b]^4]^2<+0>

LoT complexity = 12

- LoT:
 - Staying ("+o")
 - Moving to the other item (here denoted 'b')
 - Repetition ("^n", where n is any number),
 - Possibly with a variation in the starting point
 - Denoted by <x> where x is an elementary instruction, either +o or b
 - Embedding of expressions is represented by brackets ("[...]") and concatenation by commas (",").



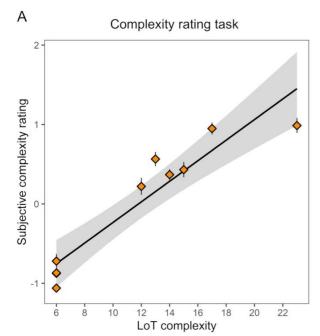
• 'Sequence violation' experimental paradigm:

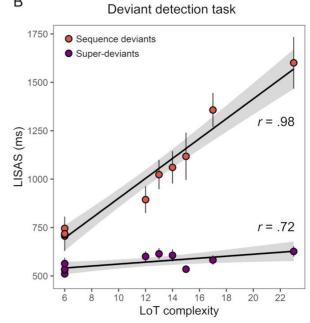


• The basic hypothesis was that, for equal sequence length, error rate and response time in violation detection would increase with sequence complexity.



- This is the results of the first experiment
- Participants were asked subjective complexity (left)
- And to identify deviations in:
- Sequence deviants
 - A note is replaces with the other note
- Superdeviant
 - A new note is introduced
- LISAS: Linear Integrated Speed-Accuracy Score







Summary

- This week, we saw some other applications of the pLoT idea
- First, we saw a model of numeral acquisition
- Then, a model of learning visual concepts (categories of 3-d objects)
 - This could easily be extended, if you're into 3-d rendering
 - Possible final project?
- Finally, we saw a model of learning binary sequences of tones
- In the lab this week, we'll try to implement a model from scratch
 - I am not going to prepare beforehand so we can all write it together
- Next week is the last week, so the lecture is going to be a review of what we've done, with some concluding remarks on the general project.
- In the lab next week we'll implement another model.