

Mathematical Logic

Part One

Question: How do we formalize the logic we've been using in our proofs?

Where We're Going

- **Propositional Logic** (Today)
 - Basic logical connectives.
 - Truth tables.
 - Logical equivalences.
- **First-Order Logic** (Friday/Monday)
 - Reasoning about properties of multiple objects.

Propositional Logic

A ***proposition*** is a statement that is,
by itself, either true or false.

Some Sample Propositions

- Puppies are cuter than kittens.
- Kittens are cuter than puppies.
- Usain Bolt can outrun everyone in this room.
- CS103 is useful for cocktail parties.
- This is the last entry on this list.

More Propositions

- I stay out too late.
- Got nothing in my brain.
- That's what people say.
- You should've known better than to mess with me.
- I'm gonna love ya like a black widow.

Things That Aren't Propositions

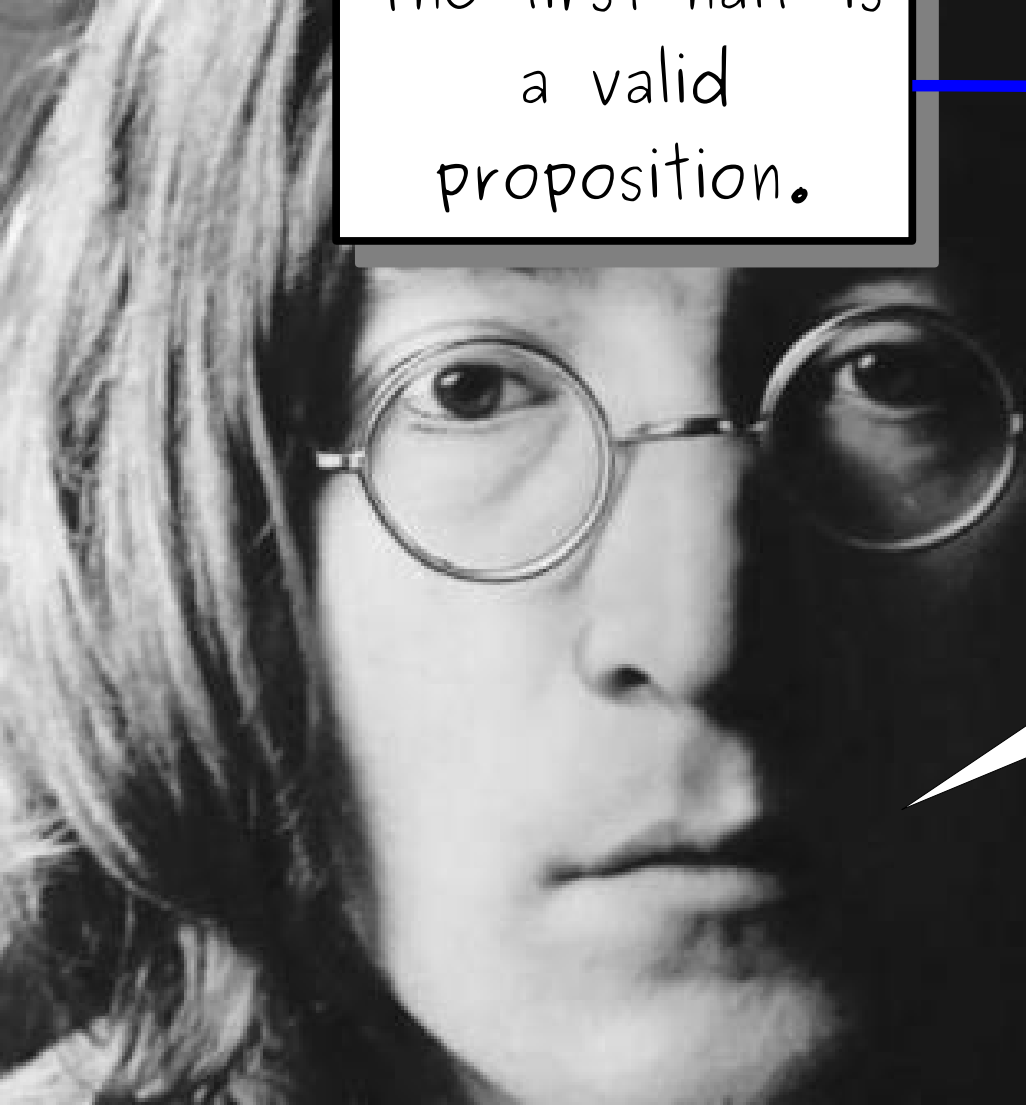


Commands
cannot be true
or false.

Things That Aren't Propositions



Things That Aren't Propositions



The first half is
a valid
proposition.

I am the walrus,
goo goo g'joob

Jibberish cannot
be true or
false.

Propositional Logic

- ***Propositional logic*** is a mathematical system for reasoning about propositions and how they relate to one another.
- Every statement in propositional logic consists of ***propositional variables*** combined via ***propositional connectives***.
 - Each variable represents some proposition, such as “You liked it” or “You should have put a ring on it.”
 - Connectives encode how propositions are related, such as “If you liked it, then you should have put a ring on it.”

Propositional Variables

- Each proposition will be represented by a ***propositional variable***.
- Propositional variables are usually represented as lower-case letters, such as p , q , r , s , etc.
- Each variable can take one of two values: true or false.

Propositional Connectives

- **Logical NOT: $\neg p$**
 - Read “**not** p ”
 - $\neg p$ is true if and only if p is false.
 - Also called **logical negation**.
- **Logical AND: $p \wedge q$**
 - Read “ p **and** q .”
 - $p \wedge q$ is true if both p and q are true.
 - Also called **logical conjunction**.
- **Logical OR: $p \vee q$**
 - Read “ p **or** q .”
 - $p \vee q$ is true if at least one of p or q are true (inclusive OR)
 - Also called **logical disjunction**.

Truth Tables

- A ***truth table*** is a table showing the truth value of a propositional logic formula as a function of its inputs.
- Useful for several reasons:
 - Formally defining what a connective “means.”
 - Deciphering what a complex propositional formula means.

The Truth Table Tool

Summary of Important Points

- The `v` operator is an *inclusive* “or.” It's true if at least one of the operands is true.
 - Similar to the `||` operator in C, C++, Java and the `or` operator in Python.
- If we need an exclusive “or” operator, we can build it out of what we already have.

Mathematical Implication

Implication

- The \rightarrow connective is used to represent implications.
 - Its technical name is the *material conditional* operator.
- What is its truth table?

Why This Truth Table?

- The truth values of the \rightarrow are the way they are because they're *defined* that way.
- We want $p \rightarrow q$ to mean “whenever p is true, q is true as well.”
- The only way this *doesn't* happen is if p is true and q is false.
- In other words, $p \rightarrow q$ should be true whenever $\neg(p \wedge \neg q)$ is true.
- What's the truth table for $\neg(p \wedge \neg q)$?

Truth Table for Implication

| p | q | $p \rightarrow q$ |
|-----|-----|-------------------|
| F | F | T |
| F | T | T |
| T | F | F |
| T | T | T |

The only way for $p \rightarrow q$ to be false is for p to be true and q to be false. Otherwise, $p \rightarrow q$ is by definition true.

The Biconditional Operator

The Biconditional Operator

- The biconditional operator \leftrightarrow is used to represent a two-directional implication.
- Specifically, $p \leftrightarrow q$ means that p implies q and q implies p .
- What should its truth table look like?

The Biconditional

- The ***biconditional*** connective $p \leftrightarrow q$ is read “ p if and only if q . ”
- Here's its truth table:

| p | q | $p \leftrightarrow q$ |
|-----|-----|-----------------------|
| F | F | T |
| F | T | F |
| T | F | F |
| T | T | T |

The Biconditional

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- Here's its truth table:

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|-----|-----|-----------------------|
| F | F | T |
| F | T | F |
| T | F | F |
| T | T | T |

One interpretation of \leftrightarrow is to think of it as equality: the two propositions must have equal truth values.

True and False

- There are two more “connectives” to speak of: true and false.
 - The symbol \top is a value that is always true.
 - The symbol \perp is value that is always false.
- These are often called connectives, though they don't connect anything.
 - (Or rather, they connect zero things.)

Operator Precedence

- How do we parse this statement?

$$\neg x \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

- Operator precedence for propositional logic:

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow y \vee z \rightarrow x \vee (y \wedge z)$$

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\rightarrow

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Operator Precedence

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$$(\neg x) \rightarrow y \vee z \rightarrow x \vee (y \wedge z)$$

- Operator precedence for propositional logic:

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\vee

\rightarrow

\leftrightarrow

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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow (y \vee z) \rightarrow (x \vee (y \wedge z))$$

- Operator precedence for propositional logic:

\neg

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\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow (y \vee z) \rightarrow (x \vee (y \wedge z))$$

- Operator precedence for propositional logic:

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\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow ((y \vee z) \rightarrow (x \vee (y \wedge z)))$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
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Operator Precedence

- How do we parse this statement?

$$(\neg x) \rightarrow ((y \vee z) \rightarrow (x \vee (y \wedge z)))$$

- Operator precedence for propositional logic:

\neg

\wedge

\vee

\rightarrow

\leftrightarrow

- All operators are right-associative.
- We can use parentheses to disambiguate.

Operator Precedence

- The main points to remember:
 - \neg binds to whatever immediately follows it.
 - \wedge and \vee bind more tightly than \rightarrow .
- We will commonly write expressions like $p \wedge q \rightarrow r$ without adding parentheses.
- For more complex expressions, we'll try to add parentheses.
- Confused? Just ask!

Time-Out for Announcements!

Graded Assignments

- Problem Set 1 and Problem Set 2
Checkpoints are graded and available online through Scoryst.
- Be sure to review the feedback and not just the grade – you'll get a lot more out of this course if you do!

Solution Sets

- Problem Set 1 solutions, Problem Set 2 Checkpoint solutions, and Discussion Problems 2 solutions are all available in hardcopy outside.
- Missed lecture? Pick them up in the filing cabinet in the Gates B wing.
- SCPD students – you should get copies of these handouts soon.

Stanford HOPES



HUNTINGTON'S OUTREACH PROJECT
FOR EDUCATION, AT STANFORD

- Stanford HOPES (Huntington's Outreach Program for Education, at Stanford) is looking for web designers, graphic designers, and researchers for this academic year.
- Check out their website at <http://hopes.stanford.edu>.
- Interested? Contact Kristen Powers at kapowers@stanford.edu.

Some Logistics

- Maesen and I will be at the Grace Hopper Conference for the rest of this week.
 - Stephen Macke (smacke@cs.stanford.edu) is acting as Head TA while Maesen is out.
 - Kevin Crain will be giving Friday's lecture on logic.
- I will be a lot slower at responding to emails – sorry for the inconvenience!

Your Questions

A Quick Note on Questions

“How important is discrete math compared to calculus in the math world? In your opinion, what should be taught first, set theory or limits?”

“How important is it to learn LaTeX?
I know some people that use it and
wondered if it would be worth it to
learn it.”

“I really want to have a portfolio before my senior year and am seriously thinking of taking a gap year to really 'find myself' with computer science. Do you recommend this? How can I best utilize a gap year?”

“I find it very hard to engage with the material. Can you give some words of inspiration to make me excited about sets, induction, graphs, etc?”

Back to CS103!

Recap So Far

- A ***propositional variable*** is a variable that is either true or false.
- The ***propositional connectives*** are
 - Negation: $\neg p$
 - Conjunction: $p \wedge q$
 - Disjunction: $p \vee q$
 - Implication: $p \rightarrow q$
 - Biconditional: $p \leftrightarrow q$
 - True: \top
 - False: \perp

Translating into Propositional Logic

Some Sample Propositions

a: I will get up early this morning

b: There is a lunar eclipse this morning

c: There are no clouds in the sky this morning

d: I will see the lunar eclipse

"I won't see the lunar
eclipse if I don't get up
early this morning."

$$\neg a \rightarrow \neg d$$

“ p if q ”

translates to

$$q \rightarrow p$$

It does *not* translate to

$$p \rightarrow q$$

Some Sample Propositions

a: I will get up early this morning

b: There is a lunar eclipse this morning

c: There are no clouds in the sky this morning

d: I will see the lunar eclipse

"If I get up early this morning, but it's cloudy outside, I won't see the lunar eclipse."

$$a \wedge \neg c \rightarrow \neg d$$

“ p , but q ”

translates to

$$p \wedge q$$

The Takeaway Point

- When translating into or out of propositional logic, be very careful not to get tripped up by nuances of the English language.
 - In fact, this is one of the reasons we have a symbolic notation in the first place!
- Many prepositional phrases lead to counterintuitive translations; make sure to double-check yourself!

Propositional Equivalences

Quick Question

What would I have to show you to convince you that the statement $p \wedge q$ is false?

Quick Question

What would I have to show you to convince you that the statement $p \vee q$ is false?

De Morgan's Laws

- Using truth tables, we concluded that

$$\neg(p \wedge q)$$

is equivalent to

$$\neg p \vee \neg q$$

- We also saw that

$$\neg(p \vee q)$$

is equivalent to

$$\neg p \wedge \neg q$$

- These two equivalences are called ***De Morgan's Laws***.

Logical Equivalence

- Because $\neg(p \wedge q)$ and $\neg p \vee \neg q$ have the same truth tables, we say that they're **equivalent** to one another.
- We denote this by writing

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

- The \equiv symbol is not a connective. It's related to \leftrightarrow , but it's not the same:
 - The statement $\neg(p \wedge q) \equiv \neg p \vee \neg q$ means “these formulas are equivalent.”
 - The statement $\neg(p \wedge q) \leftrightarrow (\neg p \vee \neg q)$ is a propositional formula. If you plug in different values of p and q , it will evaluate to a truth value. It just happens to evaluate to true every time.

An Important Equivalence

- Earlier, we talked about the truth table for $p \rightarrow q$. We chose it so that

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- Later on, this equivalence will be incredibly useful:

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

Another Important Equivalence

- Here's a useful equivalence. Start with

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- By De Morgan's laws:

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

$$\equiv \neg p \vee \neg\neg q$$

$$\equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

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$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

- By De Morgan's laws:

$$p \rightarrow q \equiv \neg(p \wedge \neg q)$$

$$\equiv \neg p \vee \neg \neg q$$

$$\equiv \neg p \vee q$$

- Thus $p \rightarrow q \equiv \neg p \vee q$

If p is false, then $\neg p \vee q$ is true. If p is true, then q has to be true for the whole expression to be true.

Next Time

- **First-Order Logic**
- **Translating into Logic**
- **Manipulating Logical Statements**