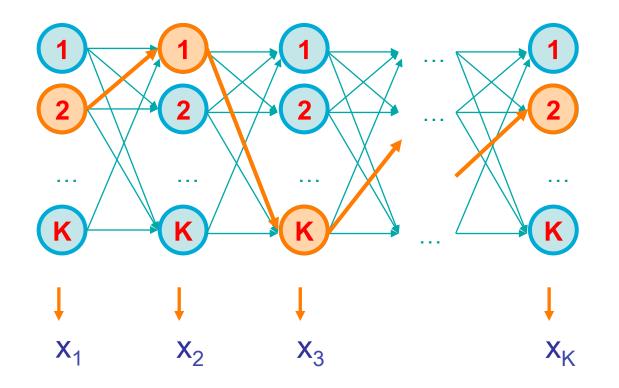


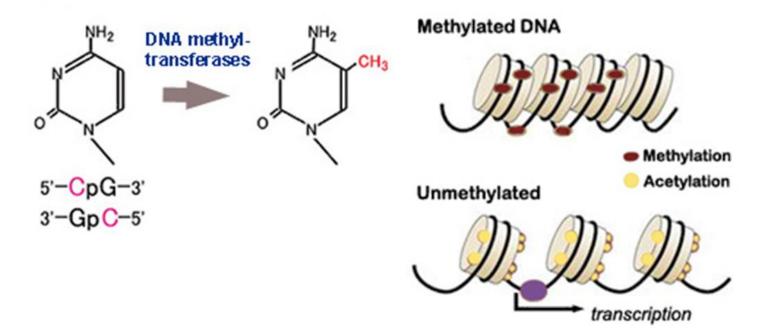
### **Hidden Markov Models**



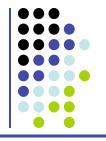
### **Motivating Example**

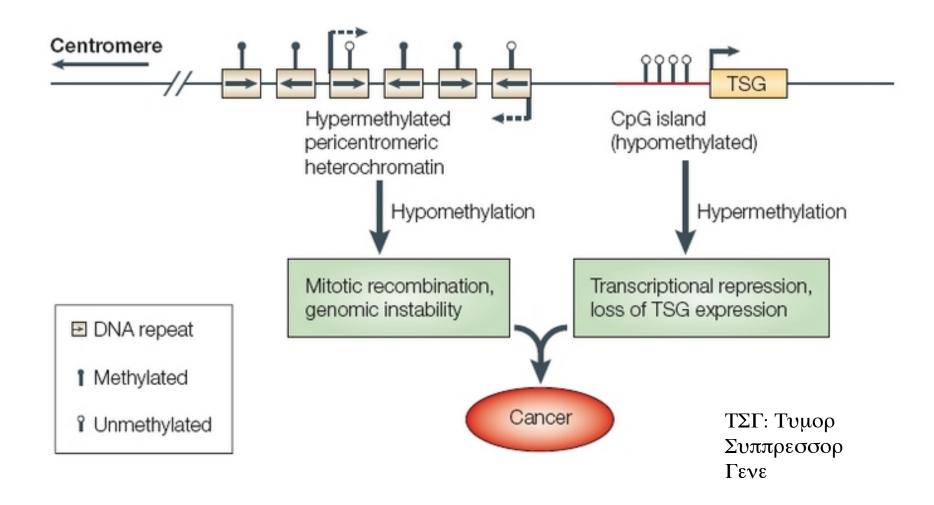


- CpG sites (-cytosine-phosphate-guanine-)
  - C in CpG can be methylated -> 5-methylcytosine
  - 70%-80% of CpG cytonsines are methylated



### **Motivating Example**





### **Motivating Example**



#### CpG islands: regions with a high frequency of CpG sites

#### near gene-promoter

CATTCCGCCTTCTCTCCCGAGGTGGCGCGTGGGA GGTGTTTTGCTCGGGTTCTGTAAGAATAGGCCAGG CAGCTTCCCGCGGGATGCGCTCATCCCCTCTCGG GGTTCCGCTCCCACCGCGCCGCGTTCGGCCGGTT CCGCCTGCGAGATGTTTTCCGACGGACAATGATTC CACTCTCGGCGCCTCCCATGTTGATCCCAGCTCCT CTGCGGGCGTCAGGACCCCTGGGCCCCGCCCCG CTCCACTCAGTCAATCTTTTGTCCCCGTATAAGGCG GATTATCGGGGTGGCTGGCTGATTCCGA CGAATGCCCTTGGGGGTCACCCGGGAGGGAACTC CGGGCTCCGGCTTTGGCCAGCCCGCACCCCTGGT TGAGCCGGCCCGAGGGCCCCCAGGGGGGCGCTCG ATGTTCCTGCAGCCCCCCGCAGCAGCCCCACTCC CCGGCTCACCCTACGATTGGCTGGCCGCCCCGAG CTCTGTGCTGTGATTGGTCACAGCCCGTGTCCGTC GCGGGCGCCGGGCGGATACGAGGTGACGCGCA GAGGCCCAGCTCGGGGCGGTGTCCCGCGCCGGC GACTGCGGGCGGAGTTTCGCGAGGGCCGAAG GGGCAGTGTGACGGCAGCGGTCCTGGGAGG CCGCGCGCGTCGGAGCAGCTCCCCGTCCTCCGCA GCCCACCTCCACCTCGATGCGGTGCCGGGCTGC TGCGTGATGGGGCTGCGGAGCGCCCCTGCGG CTCGCGGCGGCCGCTGCTCGCGCTGAGGTGCGT GCCCCCGCGCGCGCGCGCGC CGGTGCCCG GGCTCCTGTTGACCCGGTCGCCCGTCGGTCTGC AGCGCGGCTGAGGTAAGGCGGCGGGGCTGGCCG CGGTTGGCGCCGCGGTCGCGGGGTTGGGGAGGG GGTCCGGGCGGGGTCTGAGGGGA

#### normal example of the genome

CTCTTAGTTTTGGGTGCATTTGTCTGGTCTTCCAAA CTAGATTGAAAGCTCTGAAAAAAAAAAACTATCTTGT GTTTCTATCTGTTGAGCTCATAGTAGGTATCCAGGA AGTAGTAGGGTTGACTGCATTGATTTGGGACTACAC TGGGAGTTTTCTTCGCCATCTCCCTTTAGTTTTCCT TTTTTCTTTCTTTCTTTTCTTTTTTTCTTTTTTT TTGAGATGTCGTCTTGCTCAGTCCCCCAGGCTGGA GTGCAGTGGTGCGATCTTGGCTCACTGTAGCCTCC ACCTCCCAGGTTCAAGCAATTCTACTGCCTTAGCCT CCCGAGTAGCTGGGATTACAAGCACCCGCCACCAT TCCTGGCTAATTTTTTTTTTTTGTATTTTTAGTTGAGA CAGGGTTTCACCATGTTGGTGATGCTGGTCTCAGA CTCCTGGGGCCTAGCGATCCCCCTGCCTCAGCCT CCCAGAGTGTTAGGATTACAGGCATGAGCCACTGT ACCCGGCCTCTCTCCAGTTTCCAGTTGGAATCCAA GGGAAGTAAGTTTAAGATAAAGTTACGATTTTGAAAT CTTTGGATTCAGAAGAATTTGTCACCTTTAACACCT AGAGTTGAACGTTCATACCTGGAGAGCCTTAACATT AAGCCCTAGCCAGCCTCCAGCAAGTGGACATTGGT CAGGTTTGGCAGGATTCGTCCCCTGAAGTGGACT GAGAGCCACACCCTGGCCTGTCACCATACCCATCC CCTATCCTTAGTGAAGCAAAACTCCTTTGTTCCCTT CTCCTTCTCCTAGTGACAGGAAATATTGTGATCCTA AAGAATGAAAATAGCTTGTCACCTCGTGGCCTCAG GCCTCTTGACTTCAGGCGGTTCTGTTTAATCAAGT GACATCTTCCCGAGGCTCCCTGAATGTGGCAGATG AAAGAGACTAGTTCAACCCTGACCTGAGGGGAAAG CCTTTGTGAAGGGTCAGGAG

https://commons.wikimedia.org/wiki/File:Cpg\_islands.svg





#### 1. Decoding

GIVEN a HMM M, and a sequence x,

FIND the sequence  $\pi$  of states that maximizes P[ x,  $\pi$  | M ]

#### 2. Evaluation

GIVEN a HMM M, and a sequence x,

FIND Prob[x | M]

#### 3. Learning

GIVEN a HMM M, with unspecified transition/emission probs.,

and a sequence x,

FIND parameters  $\theta = (e_i(.), a_{ij})$  that maximize P[ x |  $\theta$  ]



### **Problem 1: Decoding**

Find the most likely parse of a sequence

### **Decoding - Review**

GIVEN 
$$x = x_1 x_2 \dots x_N$$

Find 
$$\pi = \pi_1, \dots, \pi_N$$
,  
to maximize P[x,  $\pi$ ]

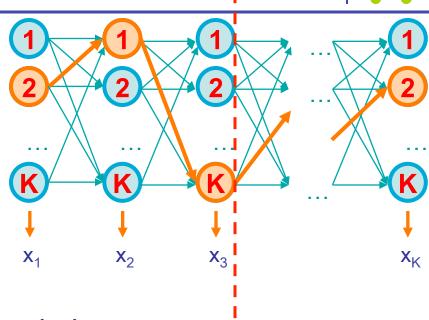
$$\pi^* = \operatorname{argmax}_{\pi} P[x, \pi]$$

Maximizes  $a_{0\pi 1} e_{\pi 1}(x_1) a_{\pi 1\pi 2} \dots a_{\pi N-1\pi N} e_{\pi N}(x_N)$ 

### Dynamic Programming!

$$V_k(i) = \max_{\{\pi_1...\pi_{i-1}\}} P[x_1...x_{i-1}, \pi_1, ..., \pi_{i-1}, x_i, \pi_i = k]$$

= Prob. of most likely sequence of states ending at state  $\pi_i$  = k



Given that we end up in state k at step i, maximize product to the left and right

### Decoding – Review



Inductive assumption: Given that for all states k, and for a fixed position i,

$$V_k(i) = \max_{\{\pi_1...\pi_{i-1}\}} P[x_1...x_{i-1}, \pi_1, ..., \pi_{i-1}, x_i, \pi_i = k]$$

What is  $V_i(i+1)$ ?

From definition,

$$\begin{split} V_{l}(i+1) &= \text{max}_{\{\pi 1 \dots \pi i\}} P[\; x_{1} \dots x_{i}, \; \pi_{1}, \; \dots, \; \pi_{i}, \; x_{i+1}, \; \pi_{i+1} = l \;] \\ &= \text{max}_{\{\pi 1 \dots \pi i\}} P(x_{i+1}, \; \pi_{i+1} = l \;|\; x_{1} \dots x_{i}, \; \pi_{1}, \dots, \; \pi_{i}) \; P[x_{1} \dots x_{i}, \; \pi_{1}, \dots, \; \pi_{i}] \\ &= \text{max}_{\{\pi 1 \dots \pi i\}} P(x_{i+1}, \; \pi_{i+1} = l \;|\; \pi_{i} \;) \; P[x_{1} \dots x_{i-1}, \; \pi_{1}, \; \dots, \; \pi_{i-1}, \; x_{i}, \; \pi_{i}] \\ &= \text{max}_{k} \left[ P(x_{i+1}, \; \pi_{i+1} = l \;|\; \pi_{i} = k) \; \textbf{max}_{\{\pi 1 \dots \pi i - 1\}} \textbf{P[x_{1} \dots x_{i-1}, \pi_{1}, \dots, \pi_{i-1}, \; x_{i}, \pi_{i} = k]] \\ &= \text{max}_{k} \left[ \; P(x_{i+1} \;|\; \pi_{i+1} = l \;) \; P(\pi_{i+1} = l \;|\; \pi_{i} = k) \; \textbf{V}_{k}(\textbf{i}) \; \right] \\ &= e_{l}(x_{i+1}) \; \text{max}_{k} \; a_{kl} \; \textbf{V}_{k}(\textbf{i}) \end{split}$$

# The Viterbi Algorithm - Review



Input: 
$$x = x_1 \dots x_N$$

#### **Initialization:**

$$V_0(0) = 1$$
 (0 is the imaginary first position)  
 $V_k(0) = 0$ , for all  $k > 0$ 

#### **Iteration:**

$$V_{j}(i) = e_{j}(x_{i}) \times \max_{k} a_{kj} V_{k}(i-1)$$

$$Ptr_{i}(i) = \operatorname{argmax}_{k} a_{ki} V_{k}(i-1)$$

#### **Termination:**

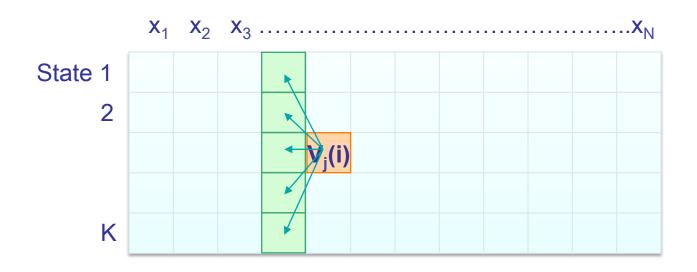
$$P(x, \pi^*) = \max_k V_k(N)$$

#### **Traceback:**

$$\pi_N^* = \operatorname{argmax}_k V_k(N)$$
  
 $\pi_{i-1}^* = \operatorname{Ptr}_{\pi_i}(i)$ 

### The Viterbi Algorithm - Review





Similar to "aligning" a set of states to a sequence

#### Time:

 $O(K^2N)$ 

#### Space:

O(KN)

### Viterbi Algorithm – Review



Underflows are a significant problem

$$P[x_1,...,x_i,\pi_1,...,\pi_i] = a_{0\pi 1} a_{\pi 1\pi 2}....a_{\pi i} e_{\pi 1}(x_1)....e_{\pi i}(x_i)$$

These numbers become extremely small – underflow

**Solution:** Take the logs of all values

$$V_l(i) = log e_k(x_i) + max_k [V_k(i-1) + log a_{kl}]$$

### **Example - Review**



Let x be a long sequence with a portion of  $\sim 1/6$  6's, followed by a portion of  $\sim \frac{1}{2}$  6's...

x = 123456123456...12345 6626364656...1626364656

Then, it is not hard to show that optimal parse is (exercise):

FFF.....L

6 characters "123456" parsed as F, contribute  $.95^6 \times (1/6)^6 = 1.6 \times 10^{-5}$ parsed as L, contribute  $.95^6 \times (1/2)^1 \times (1/10)^5 = 0.4 \times 10^{-5}$ 

"162636" parsed as F, contribute  $.95^6 \times (1/6)^6 = 1.6 \times 10^{-5}$  parsed as L, contribute  $.95^6 \times (1/2)^3 \times (1/10)^3 = 9.0 \times 10^{-5}$ 



### **Problem 2: Evaluation**

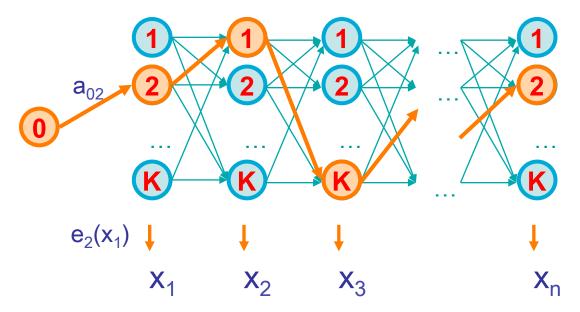
Find the likelihood a sequence is generated by the model





Given a HMM, we can generate a sequence of length n as follows:

- 1. Start at state  $\pi_1$  according to prob  $a_{0\pi 1}$
- 2. Emit letter  $x_1$  according to prob  $e_{\pi 1}(x_1)$
- 3. Go to state  $\pi_2$  according to prob  $a_{\pi 1\pi 2}$
- 4. ... until emitting x<sub>n</sub>



### A couple of questions



Given a sequence x,

- What is the probability that

Example: the dishonest ca 0.23-9

P(box: FFFFFFFFF) =  $(1/6)^{11} * 0.95^{12} =$  $2.76^{-9} * 0.54 =$ 1.49-9

```
Given a position i, what is the P(box: LLLLLLLLL) =
                                  [(1/2)^6 * (1/10)^5] * 0.95^{10} * 0.05^2 =
                                   1.56*10^{-7} * 1.5^{-3} =
```

Most likely path:  $\pi = FF.....F$ (too "unlikely" to transition  $F \rightarrow L \rightarrow F$ )

However: marked letters more likely to be L than unmarked letters

### **Evaluation**



We will develop algorithms that allow us to compute:

P(x) Probability of x given the model

 $P(x_i...x_i)$  Probability of a substring of x given the model

 $P(\pi_i = k \mid x)$  "Posterior" probability that the i<sup>th</sup> state is k, given x

A more refined measure of which states x may be in

### The Forward Algorithm



We want to calculate

P(x) = probability of x, given the HMM

Sum over all possible ways of generating x:

$$P(x) = \Sigma_{\pi} P(x, \pi) = \Sigma_{\pi} P(x \mid \pi) P(\pi)$$

To avoid summing over an exponential number of paths  $\pi$ , define

$$f_k(i) = P(x_1...x_i, \pi_i = k)$$
 (the forward probability)

"generate i first characters of x and end up in state k"

# The Forward Algorithm – derivation



Define the forward probability:

$$\begin{split} f_k(i) &= P(x_1...x_i, \, \pi_i = k) \\ &= \sum_{\pi_1...\pi_{i-1}} P(x_1...x_{i-1}, \, \pi_1, ..., \, \pi_{i-1}, \, \pi_i = k) \, e_k(x_i) \\ &= \sum_{l} \sum_{\pi_1...\pi_{l-2}} P(x_1...x_{i-1}, \, \pi_1, ..., \, \pi_{i-2}, \, \pi_{i-1} = l) \, a_{lk} \, e_k(x_i) \\ &= \sum_{l} P(x_1...x_{i-1}, \, \pi_{i-1} = l) \, a_{lk} \, e_k(x_i) \\ &= e_k(x_i) \, \sum_{l} \, f_l(i-1) \, a_{lk} \end{split}$$

# The Forward Algorithm



We can compute f<sub>k</sub>(i) for all k, i, using dynamic programming!

#### **Initialization:**

$$f_0(0) = 1$$
  
 $f_k(0) = 0$ , for all  $k > 0$ 

#### **Iteration:**

$$f_k(i) = e_k(x_i) \sum_i f_i(i-1) a_{ik}$$

#### **Termination:**

$$P(x) = \sum_{k} f_{k}(N)$$

### Relation between Forward and Viterbi



#### **VITERBI**

#### **FORWARD**

#### **Initialization:**

$$V_0(0) = 1$$
  
  $V_k(0) = 0$ , for all  $k > 0$ 

### Iteration:

$$V_j(i) = e_j(x_i) \max_k V_k(i-1) a_{kj}$$

#### **Termination:**

$$P(x, \pi^*) = \max_{k} V_k(N)$$

#### **Initialization:**

$$f_0(0) = 1$$
  
 $f_k(0) = 0$ , for all  $k > 0$ 

#### **Iteration:**

$$f_{l}(i) = e_{l}(x_{i}) \sum_{k} f_{k}(i-1) a_{kl}$$

#### **Termination:**

$$P(x) = \sum_{k} f_{k}(N)$$

### Motivation for the Backward Algorithm



We want to compute

$$P(\pi_i = k \mid x),$$

the probability distribution on the ith position, given x

We start by computing

$$P(\pi_{i} = k, x) = P(x_{1}...x_{i}, \pi_{i} = k, x_{i+1}...x_{N})$$

$$= P(x_{1}...x_{i}, \pi_{i} = k) P(x_{i+1}...x_{N} \mid x_{1}...x_{i}, \pi_{i} = k)$$

$$= P(x_{1}...x_{i}, \pi_{i} = k) P(x_{i+1}...x_{N} \mid \pi_{i} = k)$$

Forward,  $f_k(i)$  Backward,  $b_k(i)$ 

Then, 
$$P(\pi_i = k \mid x) = P(\pi_i = k, x) / P(x)$$

### The Backward Algorithm – derivation



Define the backward probability:

$$\begin{split} \mathbf{b}_{k}(\mathbf{i}) &= \mathbf{P}(\mathbf{x}_{i+1}...\mathbf{x}_{N} \mid \pi_{i} = \mathbf{k}) \qquad \text{"starting from $i^{th}$ state = $k$, generate rest of $x$"} \\ &= \sum_{\pi_{i+1}...\pi_{N}} \mathbf{P}(\mathbf{x}_{i+1}, \mathbf{x}_{i+2}, \, ..., \, \mathbf{x}_{N}, \, \pi_{i+1}, \, ..., \, \pi_{N} \mid \pi_{i} = \mathbf{k}) \\ &= \sum_{l} \sum_{\pi_{i+1}...\pi_{N}} \mathbf{P}(\mathbf{x}_{i+1}, \mathbf{x}_{i+2}, \, ..., \, \mathbf{x}_{N}, \, \pi_{i+1} = \mathbf{I}, \, \pi_{i+2}, \, ..., \, \pi_{N} \mid \pi_{i} = \mathbf{k}) \\ &= \sum_{l} \mathbf{e}_{l}(\mathbf{x}_{i+1}) \, \mathbf{a}_{kl} \, \sum_{\pi_{i+1}...\pi_{N}} \mathbf{P}(\mathbf{x}_{i+2}, \, ..., \, \mathbf{x}_{N}, \, \pi_{i+2}, \, ..., \, \pi_{N} \mid \pi_{i+1} = \mathbf{I}) \\ &= \sum_{l} \mathbf{e}_{l}(\mathbf{x}_{i+1}) \, \mathbf{a}_{kl} \, \mathbf{b}_{l}(\mathbf{i+1}) \end{split}$$

### The Backward Algorithm



We can compute  $b_k(i)$  for all k, i, using dynamic programming

#### **Initialization:**

$$b_k(N) = 1$$
, for all k

#### **Iteration:**

$$b_k(i) = \sum_i e_i(x_{i+1}) a_{ki} b_i(i+1)$$

#### **Termination:**

$$P(x) = \sum_{i} a_{0i} e_{i}(x_{1}) b_{i}(1)$$

# **Computational Complexity**



What is the running time, and space required, for Forward, and Backward?

Time:  $O(K^2N)$ 

Space: O(KN)

Useful implementation technique to avoid underflows

Viterbi: sum of logs

Forward/Backward: rescaling at each few positions by multiplying by a

constant

### **Posterior Decoding**



We can now calculate

$$P(\pi_i = k \mid x) = \frac{f_k(i) b_k(i)}{P(x)}$$

$$\begin{split} &P(\pi_i = k \mid x) = \\ &P(\pi_i = k \;,\; x)/P(x) = \\ &P(x_1,\; \dots,\; x_i,\; \pi_i = k,\; x_{i+1},\; \dots\; x_n) \; / \; P(x) = \\ &P(x_1,\; \dots,\; x_i,\; \pi_i = k) \; P(x_{i+1},\; \dots\; x_n \mid \pi_i = k) \; / \; P(x) = \\ &f_k(i) \; b_k(i) \; / \; P(x) \end{split}$$

Then, we can ask

What is the most likely state at position i of sequence x:

Define  $\pi$  by Posterior Decoding:

$$\pi_i$$
 = argmax<sub>k</sub> P( $\pi_i$  = k | x)

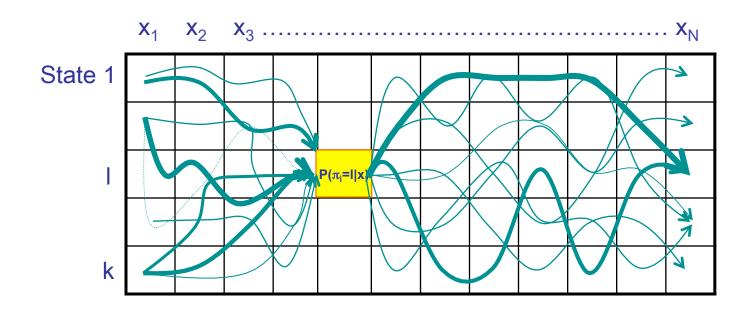
### **Posterior Decoding**



- For each state,
  - Posterior Decoding gives us a curve of likelihood of state for each position
  - That is sometimes more informative than Viterbi path π\*
- Posterior Decoding may give an invalid sequence of states (of prob 0)
  - Why?

### **Posterior Decoding**





• 
$$P(\pi_i = k \mid x) = \sum_{\pi} P(\pi \mid x) \mathbf{1}(\pi_i = k)$$
  
=  $\sum_{\pi:\pi[i] = k} P(\pi \mid x)$ 

$$\mathbf{1}(\psi)$$
 = 1, if  $\psi$  is true 0, otherwise

### Viterbi, Forward, Backward



#### **VITERBI**

#### **FORWARD**

#### **BACKWARD**

#### **Initialization:**

$$V_0(0) = 1$$
  
 $V_k(0) = 0$ , for all  $k > 0$ 

#### **Initialization:**

$$f_0(0) = 1$$
  
 $f_k(0) = 0$ , for all  $k > 0$ 

#### **Initialization:**

$$b_k(N) = 1$$
, for all k

#### **Iteration:**

$$V_l(i) = e_l(x_i) \max_k V_k(i-1) a_{kl}$$

#### **Iteration:**

$$f_{i}(i) = e_{i}(x_{i}) \sum_{k} f_{k}(i-1) a_{ki}$$

#### <u>Iteration:</u>

$$b_{l}(i) = \sum_{k} e_{l}(x_{i}+1) a_{kl} b_{k}(i+1)$$

#### **Termination:**

$$P(x, \pi^*) = \max_{k} V_k(N)$$

#### **Termination:**

$$P(x) = \sum_{k} f_{k}(N)$$

#### **Termination:**

$$P(x) = \sum_{k} a_{0k} e_{k}(x_{1}) b_{k}(1)$$

### Viterbi, Forward, Backward



#### **VITERBI**

#### **FORWARD**

#### **BACKWARD**

#### **Initialization:**

$$V_0(0) = 1$$
  
 $V_k(0) = 0$ , for all  $k > 0$ 

#### **Initialization:**

$$f_0(0) = 1$$
  
 $f_k(0) = 0$ , for all  $k > 0$ 

#### **Initialization:**

$$b_k(N) = 1$$
, for all k

#### **Iteration:**

$$V_l(i) = e_l(x_i) \max_k V_k(i-1) a_{kl}$$

#### **Iteration:**

$$f_{i}(i) = e_{i}(x_{i}) \sum_{k} f_{k}(i-1) a_{ki}$$

#### <u>Iteration:</u>

$$b_i(i) = \sum_k e_i(x_i+1) a_{ki} b_k(i+1)$$

#### **Termination:**

$$P(x, \pi^*) = \max_k V_k(N)$$

#### **Termination:**

$$P(x) = \sum_{k} f_{k}(N)$$

#### **Termination:**

$$P(x) = \sum_{k} a_{0k} e_{k}(x_{1}) b_{k}(1)$$



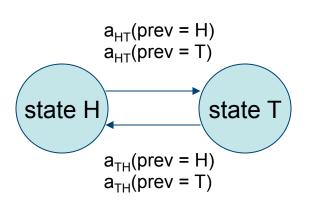
# **Variants of HMMs**

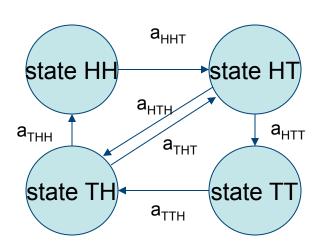


### **Higher-order HMMs**



- How do we model "memory" larger than one time point?
- $P(\pi_{i+1} = I \mid \pi_i = k)$   $a_{kl}$
- $P(\pi_{i+1} = I \mid \pi_i = k, \pi_{i-1} = j)$   $a_{ikl}$
- •
- A second order HMM with K states is equivalent to a first order HMM with K<sup>2</sup> states





### Similar Algorithms to 1<sup>st</sup> Order



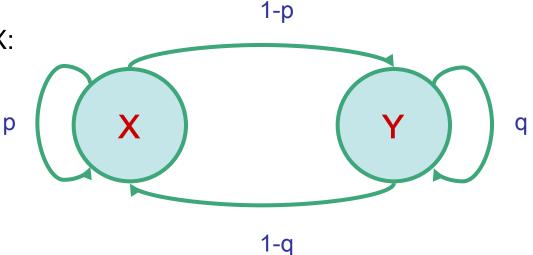
- $P(\pi_{i+1} = I \mid \pi_i = k, \pi_{i-1} = j)$ 
  - $V_{lk}(i) = max_{i} \{ V_{kj}(i-1) + ... \}$
  - Time? Space?

### **Modeling the Duration of States**



Length distribution of region X:

$$E[I_X] = 1/(1-p)$$



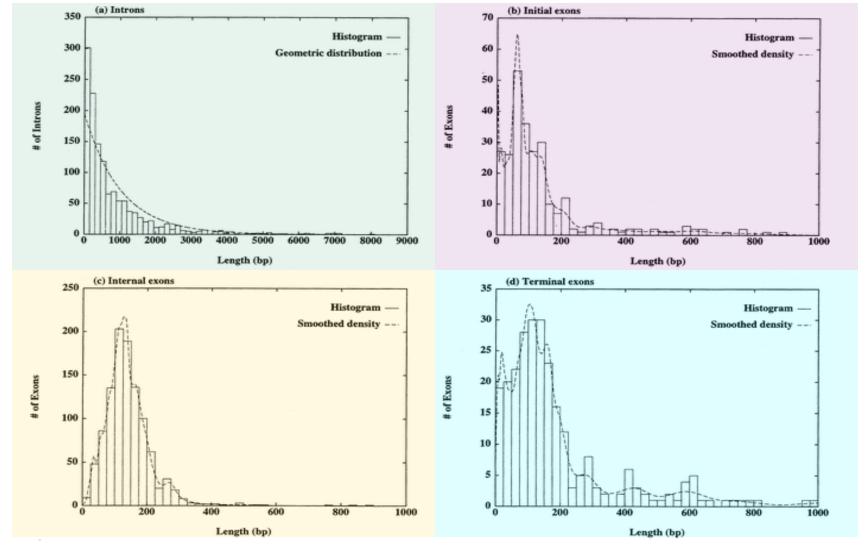
Geometric distribution, with mean 1/(1-p)

This is a significant disadvantage of HMMs

Several solutions exist for modeling different length distributions

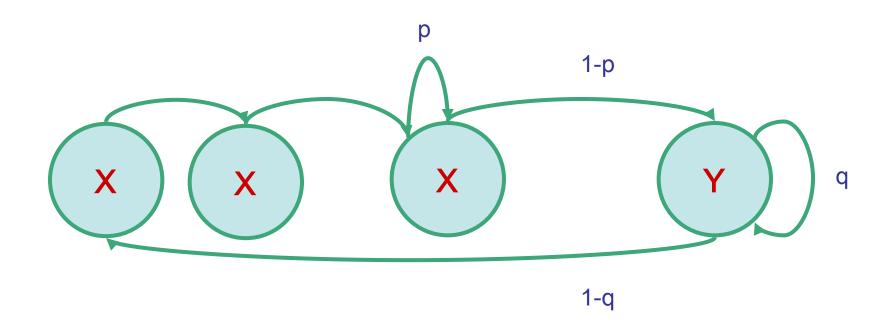
# Example: exon lengths in genes





### **Solution 1: Chain several states**



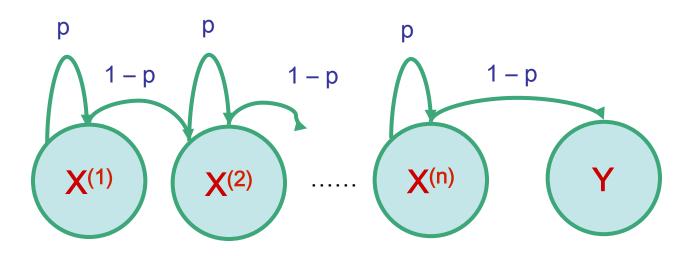


Disadvantage: Still very inflexible

 $I_X = C + geometric with mean 1/(1-p)$ 

# Solution 2: Negative binomial distribution





Duration in X: m turns, where

- During first m 1 turns, exactly n 1 arrows to next state are followed
- During m<sup>th</sup> turn, an arrow to next state is followed

$$P(I_X = m) = {m-1 \choose n-1} (1-p)^{n-1+1} p^{(m-1)-(n-1)} = {m-1 \choose n-1} (1-p)^n p^{m-n}$$

## **Example:** genes in prokaryotes



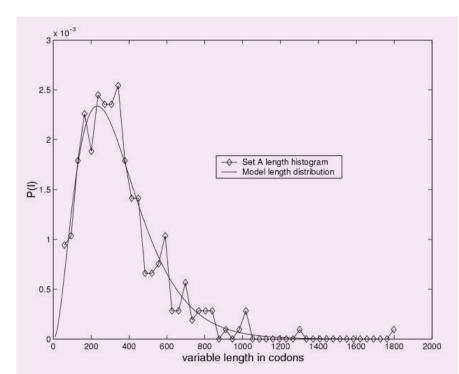
EasyGene: Prokaryotic gene-finder Shadows

Background

n=1 n=2 n=3

Larsen TS, Krogh A

Negative binomial with n = 3

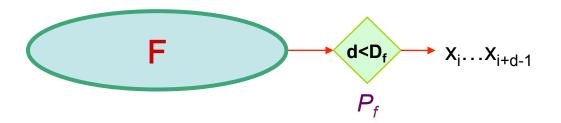


## **Solution 3: Duration modeling**



### Upon entering a state:

- 1. Choose duration d, according to probability distribution
- 2. Generate d letters according to emission probs
- 3. Take a transition to next state according to transition probs



Disadvantage: Increase in complexity of Viterbi:

Time: O(D) o O

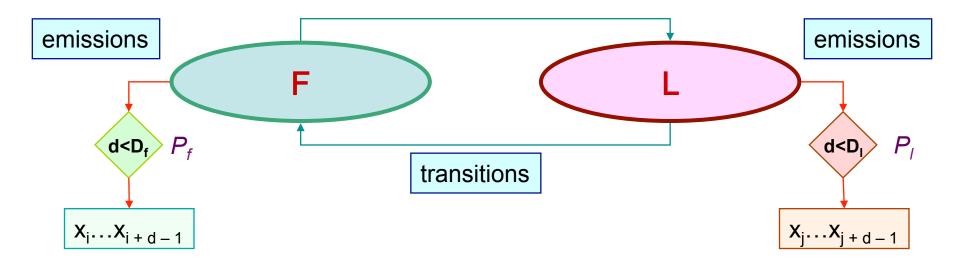
Space: O(1)

Warning, Rabiner's tutorial claims O(D²) & O(D) increases

where D = maximum duration of state

## Viterbi with duration modeling





Recall original iteration:

$$VI(i) = max_k V_k(i-1) a_{kl} \times e_l(x_i)$$

New iteration:

$$V_{l}(i) = \max_{k} \max_{d=1...Dl} V_{k}(i-d) \times P_{l}(d) \times a_{kl} \times \prod_{j=i-d+1...i} e_{l}(x_{j})$$

Precompute cumulative values



# Learning

Re-estimate the parameters of the model based on training data

## Two learning scenarios



1. Estimation when the "right answer" is known

**Examples:** 

**GIVEN:** a genomic region  $x = x_1...x_{1,000,000}$  where we have good (experimental) annotations of the CpG islands

**GIVEN:** the casino player allows us to observe him one evening,

as he changes dice and produces 10,000 rolls

2. Estimation when the "right answer" is unknown

**Examples:** 

the porcupine genome; we don't know how frequent are the **GIVEN:** 

CpG islands there, neither do we know their composition

GIVEN: 10,000 rolls of the casino player, but we don't see when he

changes dice

**QUESTION:** Update the parameters  $\theta$  of the model to maximize  $P(x|\theta)$ 

### 1. When the states are known



Given 
$$x = x_1...x_N$$
  
for which the true  $\pi = \pi_1...\pi_N$  is known,

#### **Define:**

$$A_{kl}$$
 = # times k→l transition occurs in π  
 $E_k(b)$  = # times state k in π emits b in x

We can show that the maximum likelihood parameters  $\theta$  (maximize  $P(x|\theta)$ ) are:

$$a_{kl} = \frac{A_{kl}}{\Sigma_{i} A_{ki}}$$

$$e_k(b) = \frac{E_k(b)}{\sum_c E_k(c)}$$

### 1. When the states are known



Intuition: When we know the underlying states,

Best estimate is the normalized frequency of transitions & emissions that occur in the training data

#### **Drawback:**

Given little data, there may be **overfitting**:  $P(x|\theta)$  is maximized, but  $\theta$  is unreasonable **0 probabilities – BAD** 

#### **Example:**

Given 10 casino rolls, we observe

$$x = 2$$
, 1, 5, 6, 1, 2, 3, 6, 2, 3  
 $\pi = F$ ,  $F$ 

Then:

$$a_{FF} = 1;$$
  $a_{FL} = 0$   
 $e_{F}(1) = e_{F}(3) = .2;$   
 $e_{F}(2) = .3;$   $e_{F}(4) = 0;$   $e_{F}(5) = e_{F}(6) = .1$ 

### **Pseudocounts**



Solution for small training sets:

Add pseudocounts

$$A_{kl}$$
 = # times k→l transition occurs in π +  $r_{kl}$   
 $E_k(b)$  = # times state k in π emits b in x +  $r_k(b)$ 

 $r_{kl}$ ,  $r_{k}$ (b) are pseudocounts representing our prior belief

Larger pseudocounts ⇒ Strong priof belief

Small pseudocounts ( $\epsilon$  < 1): just to avoid 0 probabilities

### **Pseudocounts**



### **Example:** dishonest casino

We will observe player for one day, 600 rolls

Reasonable pseudocounts:

$$\begin{split} r_{0F} &= r_{0L} = r_{F0} = r_{L0} = 1; \\ r_{FL} &= r_{LF} = r_{FF} = r_{LL} = 1; \\ r_{F}(1) &= r_{F}(2) = \dots = r_{F}(6) = 20 \\ r_{L}(1) &= r_{L}(2) = \dots = r_{L}(6) = 5 \end{split} \qquad \text{(strong belief fair is fair)}$$

Above #s are arbitrary – assigning priors is an art

### 2. When the states are hidden



We don't know the true  $A_{kl}$ ,  $E_k(b)$ 

#### Idea:

- We estimate our "best guess" on what  $A_{kl}$ ,  $E_k(b)$  are
  - Or, we start with random / uniform values
- We update the parameters of the model, based on our guess
- We repeat

### 2. When the states are hidden



Starting with our best guess of a model M, parameters  $\theta$ :

Given 
$$x = x_1...x_N$$
  
for which the true  $\pi = \pi_1...\pi_N$  is unknown,

We can get to a provably more likely parameter set  $\theta$  *i.e.*,  $\theta$  that increases the probability  $P(x \mid \theta)$ 

Principle: EXPECTATION MAXIMIZATION

- 1. Estimate  $A_{kl}$ ,  $E_{k}(b)$  in the training data
- 2. Update  $\theta$  according to  $A_{kl}$ ,  $E_k(b)$
- 3. Repeat 1 & 2, until convergence

## **Estimating new parameters**



To estimate  $A_{kl}$ : (assume " $\mid \theta_{CURRENT}$ ", in all formulas below)

At each position i of sequence x, find probability transition  $k\rightarrow l$  is used:

$$P(\pi_i = k, \pi_{i+1} = l \mid x) =$$

$$[1/P(x)] \times P(\pi_i = k, \pi_{i+1} = l, x_1...x_N) = Q/P(x)$$

where Q = P(
$$x_1...x_i$$
,  $\pi_i$  = k,  $\pi_{i+1}$  = I,  $x_{i+1}...x_N$ ) =   
= P( $\pi_{i+1}$  = I,  $x_{i+1}...x_N$  |  $\pi_i$  = k) P( $x_1...x_i$ ,  $\pi_i$  = k) =   
= P( $\pi_{i+1}$  = I,  $x_{i+1}x_{i+2}...x_N$  |  $\pi_i$  = k) f<sub>k</sub>(i) =   
= P( $x_{i+2}...x_N$  |  $\pi_{i+1}$  = I) P( $x_{i+1}$  |  $\pi_{i+1}$  = I) P( $\pi_{i+1}$  = I |  $\pi_i$  = k) f<sub>k</sub>(i) =   
= b<sub>l</sub>(i+1) e<sub>l</sub>( $x_{i+1}$ ) a<sub>kl</sub> f<sub>k</sub>(i)

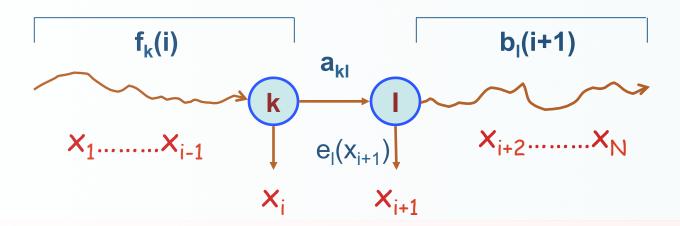
So: 
$$P(\pi_i = k, \pi_{i+1} = l \mid x, \theta) = \frac{f_k(i) a_{kl} e_l(x_{i+1}) b_l(i+1)}{P(x \mid \theta_{CURRENT})}$$

## **Estimating new parameters**



• So,  $A_{kl}$  is the E[# times transition  $k\rightarrow I$ , given current  $\theta$ ]

$$A_{kl} = \sum_{j} P(\pi_{i} = k, \ \pi_{i+1} = l \mid x, \ \theta) = \sum_{j} \frac{f_{k}(i) \ a_{kl} \ e_{l}(x_{i+1}) \ b_{l}(i+1)}{P(x \mid \theta)}$$



Similarly,

$$E_k(b) = [1/P(x \mid \theta)] \sum_{\{i \mid x_i = b\}} f_k(i) b_k(i)$$

## The Baum-Welch Algorithm



#### <u>Initialization:</u>

Pick the best-guess for model parameters (or arbitrary)

### **Iteration:**

- Forward
- Backward

3. Calculate  $A_{kl}$ ,  $E_k(b)$ , given  $\theta_{CURRENT}$ 

4. Calculate new model parameters  $\theta_{NEW}$ :  $a_{kl}$ ,  $e_{k}$ (b)

5. Calculate new log-likelihood  $P(x \mid \theta_{NEW})$ 

#### **GUARANTEED TO BE HIGHER BY EXPECTATION-MAXIMIZATION**

Until  $P(x \mid \theta)$  does not change much

## The Baum-Welch Algorithm



### Time Complexity:

# iterations  $\times$  O(K<sup>2</sup>N)

Guaranteed to increase the log likelihood P(x | θ)

Not guaranteed to find globally best parameters

Converges to local optimum, depending on initial conditions

Too many parameters / too large model: Overtraining

## **Alternative: Viterbi Training**



### **Initialization:** Same

### **Iteration:**

- 1. Perform Viterbi, to find  $\pi^*$
- 2. Calculate  $A_{kl}$ ,  $E_k(b)$  according to  $\pi^*$  + pseudocounts
- 3. Calculate the new parameters  $a_{kl}$ ,  $e_{k}$ (b)

### Until convergence

#### **Notes:**

- Not guaranteed to increase  $P(x | \theta)$
- Guaranteed to increase  $P(x | \theta, \pi^*)$
- In general, worse performance than Baum-Welch