

1. Solve the following recurrences:

$$T(n) = 2 \cdot T(n/4) + T(n/2) + n^{1/3} \quad (14\%)$$

$$T(n) = T(n^{1/2}) + \lg_2 n \quad (14\%)$$

$$T(n) = 2 \cdot T(n/2) + n \cdot \lg_2 n \quad (8\%)$$

2. Validate the following non-asymptotically tight bounds:

$$n^2 / \lg_2 n = o(n^2) \quad \text{and} \quad n^{0.9999} \cdot \lg_2 n = \omega(n) \quad (6\%)$$

3. Merge sort follows the divide-and-conquer approach and involves a key procedure, as follows. Fill the missing statements to complete the procedure. (12%)

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MERGE(A, p, q, r)
  n1 = q - p + 1
  n2 = r - q
  let L[1 .. n1 + 1] and R[1 .. n2 + 1] be new arrays
  for i = 1 to n1
    L[i] = A[p + i - 1]
  for j = 1 to n2
    R[j] = A[q + j]
  L[n1 + 1] = ∞
  R[n2 + 1] = ∞
  i = 1
  j = 1
  for k = p to r
    if L[i] ≤ R[j]
      
      
    else
      
      

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4. The Master Method aims to solve $T(n) = a \cdot T(n/b) + f(n)$ with constants $a \geq 1$ and n being an exact power of b (> 1). Show by the recursion tree technique that $T(n)$ equals (14%)

$$\Theta(n^{\log_b a}) + \sum_{j=0}^{\log_b n - 1} a^j \cdot f(n/b^j)$$

5. A hash function realized by the division method takes the form of $h(k) = k \bmod m$, where m is the number of table entries.

In general, what is a good choice for m , and in particular, what is a good m for the table to yield its load factor ≤ 0.5 when holding 120 distinct keys? (6%)

6. The utilization efficiency of a hash table depends heavily on its hashing function(s) employed. Describe with a diagram to illustrate how a multiplication method of hashing works on a machine with the word size of w bits for a hash table with 2^p entries, $p < w$. (16%)
7. Prove that for an open-address hash table with load factor $\alpha (< 1)$, the expected number of probes in unsuccessful search under uniform hashing is at most $1/(1-\alpha)$. (Hint: consider the probability of i probes to learn key absence.) (10%)

Good Luck!