

Ph.D. Comprehensive Examination, Spring, 2007

Design and Analysis of Algorithms

1 Short Questions

Answer 3 of 4 questions

[S1] Prove $h(n) \triangleq \sum_{i=1}^n \frac{1}{i} = \Theta(\log_2 n)$.

[S2] Given the recursive equation

$$\begin{cases} T(n) = 5T(n-1) - 6T(n-2) + 7^n \\ T(0) = 1, \quad T(1) = 6 \end{cases}$$

Obtain a closed-form formula for $T(n)$ and determine its growth rate (Θ).

[S3] Formally define

(S3a) deterministic finite automaton (DFA);

(S3b) pushdown automaton (PDA);

(S3c) context free grammar (CFG);

Give the relations between the classes of languages of these devices (containment or equality).

[S4] Construct

(S4a) a finite automaton or a regular expression for the language

$$\{x \in \{0, 1\}^* : \text{any substring } 000 \text{ in } x \text{ is followed immediately by } 1\};$$

(S4b) a context free grammar or pushdown automaton for the language

$$\{a^n b^m c^{2n} : n, m > 0\}.$$

2 Long Questions

Answer 3 of 4 questions

[L1] Assume that the 3-dimensional matching problem (3-DM) has been proved *NP*-complete. Prove that the sub-set sum problem is *NP*-complete.

[L2] Consider the use of branch and bound method to solve the traveling salesman problem.

(L2a) Given a cost matrix M , how to calculate the value $V = V(M)$ of the matrix M ?

(L2b) Consider a graph on 4 vertices corresponding to the following cost matrix M

$$M = \begin{bmatrix} \infty & 8 & 6 & 7 \\ 8 & \infty & 7 & 4 \\ 6 & 7 & \infty & 6 \\ 7 & 4 & 6 & \infty \end{bmatrix}.$$

Using a branch-and-bound method, find the minimum-cost Hamiltonian circuit.

[L3] Let A and B be in *NP*, and A be polynomial-time reducible to B . Briefly prove:

(L3a) If B is in P , then A is in P .

(L3b) If A is *NP*-complete, then B is *NP*-complete.

(L3c) If A is *NP*-complete and B is in P , then $P = NP$.

[L4] Classify each of the following languages as regular, context free but not regular, or decidable but not context free. Prove your answers.

(L4a) $\{a^n b^m c^m : n, m \geq 0\}$;

(L4b) $\{a^{n^2} : n \geq 0\}$;

(L4c) $\{a^{2n+1} : n \geq 0\}$.