Let's attempt stable marriage.

```
[MAN, WOMAN]
  Reflexive[T] == \{r : T \leftrightarrow T \mid id T \subseteq r\}
  Transitive[T] == \{r : T \leftrightarrow T \mid r \ \ \ r \subseteq r\}
  Preorder[T] == Reflexive[T] \cap Transitive[T]
  Antisymmetric[T] == \{r : T \leftrightarrow T \mid r \cap r^{\sim} \subseteq id T\}
  PartialOrder[T] == Preorder[T] \cap Antisymmetric[T]
  TotalOrder[T] == \{r : PartialOrder[T] \mid r \cup r^{\sim} = T \times T\}
    \_StableMarriage \_
      male\_pref: MAN \rightarrow TotalOrder[WOMAN]
     female\_pref: WOMAN \rightarrow TotalOrder[MAN]
      MAN \in \mathbb{F} MAN
      WOMAN \in \mathbb{F} WOMAN
      \#MAN = \#WOMAN
Initially, no preferences are expressed
     Stable Marriage Init
      Stable Marriage'
      male\_pref' = \varnothing
     female\_pref' = \emptyset
    \_EnterWomanPref\_
      \Delta Stable Marriage
      name?: WOMAN
      inpref?: TotalOrder[MAN]
      name? \not\in dom female\_pref
     female\_pref' = female\_pref \cup \{name? \mapsto inpref?\}
      male\_pref' = male\_pref
     EnterManPref _
      \Delta Stable Marriage
      name?:MAN
      inpref?: TotalOrder[WOMAN]
      name? \not\in dom \; male\_pref
      male\_pref' = male\_pref \cup \{name? \mapsto inpref?\}
     female\_pref' = female\_pref
```

\_Marry\_

 $\Xi Stable Marriage$ 

 $\mathit{wife!} : \mathit{MAN} \rightarrowtail \mathit{WOMAN}$ 

 $\begin{array}{l} \operatorname{dom} male\_pref = MAN \\ \operatorname{dom} female\_pref = WOMAN \end{array}$ 

 $\forall \, m, n: \mathit{MAN} \mid m \neq n \, \bullet \, (\mathit{wife}!m \mapsto \mathit{wife}!n \in \mathit{male\_pref} \, m \, \lor \, n \mapsto m \in \mathit{female\_pref} \, (\mathit{wife}!n))$