

$$= \text{CURS } g =$$

Complexitate : măsură de complexitatea calcedei
Relativ între clase.

P vs. NP - reducere

Complexitate : calculul lui Kolmogorov.
Descriptive

MaSura time (TIME)
space (SPACE)

Modelul de reț : 1. reț se opoarte

2. Ave la bezi
infinita la
ambele capete

3. Poate stativna

Calcul $C(M, \pi) = C_1 C_2 \dots C_n$

$\swarrow \quad \searrow \quad \swarrow \quad \searrow$

$n-1$ intrave configuration

$$\text{time}(c(m, x)) = n$$

$$\text{time}_M(\underline{x}) = \max \{ \text{time}(C(M, x)) \mid C(M, x) \}$$

calcolati tutti i x


given $(\frac{n}{2}) = \max \{ \text{time}_n(x) \mid \forall x \text{ a.s. } |x| = n \}$

$$\text{TIME}_k(f(n)) = \left\{ L \mid \exists M_0 \text{ s.t. } L = L(M_0) \text{ and } \forall n \geq n_0, \text{time}_{M_0}(n) \leq f(n) \right\}$$

Obs ~~not~~ ~~is~~ $f(n) \geq n+1$.

SPACE : Modelul : 1. UT se opreste
(off-line) 2. Are o banda de intrare-
ie intrare

disponibilită doar pentru citire

3.  active series

$\approx 2 =$

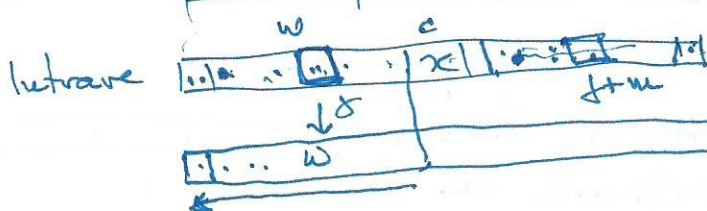
$$C(M, x) = C_1 \dots C_n$$

$\text{space}(M, x) =$ cel mai mare număr de celule folosite pe o bandă auxiliară (1-k) în calculul C .

$$\text{space}_M(n) = \max \{ \text{space}(M, x) \}$$

$$\text{SPACE}_k(f(n)) = \{ L \mid \exists M \text{ o m.t. cu } k \text{ benzi, } a.1. L = L(M) \text{ și } \text{space}_M(n) \leq f(n) \forall n \geq n_0 \}$$

Ex. $L = \{ waw \mid w \in \{a, b\}^+ \}$



$n/2$



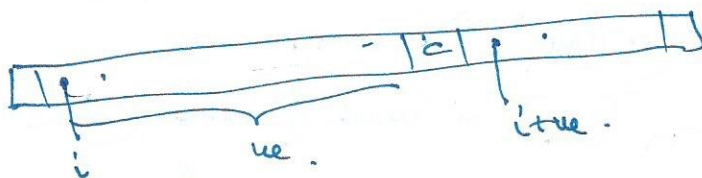
$\log u$

A, 2 p.

în timpul lungimea lui prefixului până la c .



$\log u$



u

itue.

$$O(f(n)) = O(cf(n))$$

Eliminarea constantelor

$$\textcircled{1} \textcircled{2} (N) \text{SPACE}_k(f(n)) = \textcircled{1} (N) \text{SPACE}_k(cf(n)) \quad \forall c \geq 0$$

Lemma Aleg N . Este suficient să avem.

$$\textcircled{*} \text{NSPACE}_k(f(n)) \subseteq \text{NSPACE}_k(cf(n)) \quad 0 < c < 1$$

$$c \geq 1 \quad \text{NSPACE}_k(f(n)) \subseteq \text{NSPACE}_k(cf(n))$$

$$= 3 =$$

$$0 < c \leq 1$$

Dem. M a.i. $\text{space}_M(u) \leq f(u) \implies M'$ a.i.

$$\text{space}_{M'}(u) \leq c f(u)$$

$$\underline{L(M)} = \underline{L(M')}.$$

M :

Aleg r a.i.
rc. ≥ 2

k

q

M' :

$\left. \begin{array}{l} \dots \\ \dots \end{array} \right\} r.$

$$\text{space}_{M'}(u) \leq \left\lceil \frac{\text{space}_M(u)}{r} \right\rceil \leq k$$

$$\leq \frac{\text{space}_M(u)}{r} + 1 \leq$$

$$\frac{c}{2} \cdot \text{space}_M(u) + 1 \leq c \text{space}_M(u) \quad q$$

$$1 \leq \left(\frac{c}{2} \text{space}_M(u) \right)$$

$$\text{Dacă } 1 \leq \frac{c}{2} \text{space}_M(u)$$

$$\begin{aligned} rc \geq 2 &\Rightarrow \\ c \geq \frac{2}{r} &\Rightarrow \\ \frac{1}{r} &\leq \frac{c}{2} \end{aligned}$$

$$\left(\frac{c}{2} \cdot f(u) + 1 \right) \leq c f(u) \iff \underline{1 \leq \frac{c}{2} f(u)}$$

Dacă $1 \leq \frac{c}{2} f(u)$, am terminat

$$\text{Dacă } \cancel{1 \leq \frac{c}{2} f(u)} \quad 1 > \frac{c}{2} f(u) \Rightarrow \underline{f(u) < \frac{2}{c}}$$

M' va continua k bazezi a.i. spațial nru este
mărit de 1.

$$f(u) < f'(u)$$

$$\longrightarrow f(u).$$

$$\longrightarrow f'(u).$$