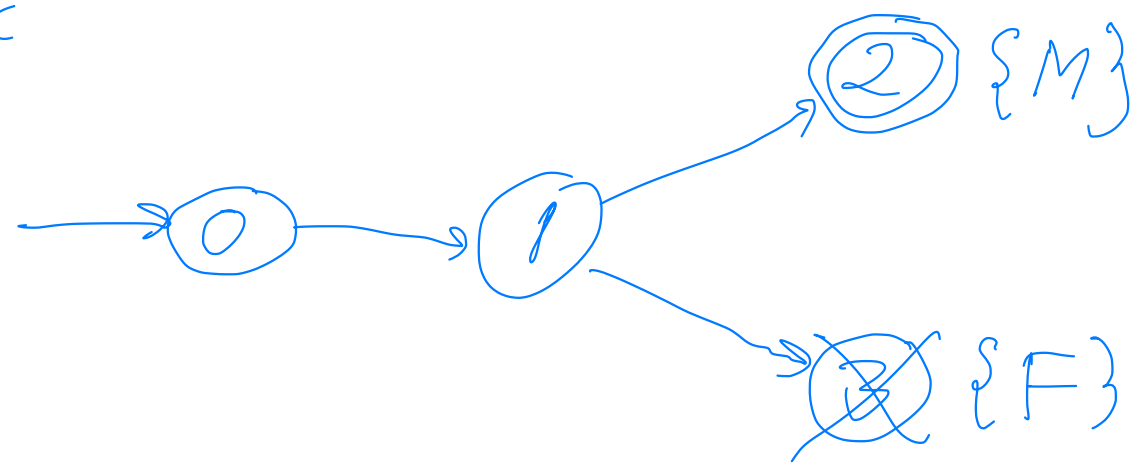


9 (cont) Temporal Logic

state labels in automata

An automaton including also state labels is called a transition system

Ex



AP = set of atomic propositions. In this example $AP = \{M, F\}$.

Labeling function $\lambda: \overline{X} \rightarrow 2^{AP}$
↑
set of states

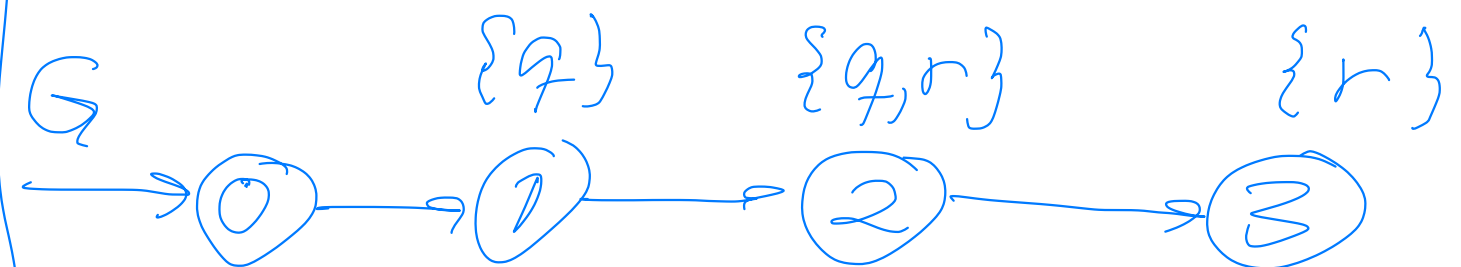
$$\lambda(2) = \{M\} \quad \lambda(3) = \{F\}$$

$$\lambda(0) = \lambda(1) = \emptyset \quad (\text{default})$$

More generally

$$AP = \{p, q, r, \dots\}$$

Ex



$[q] = \text{set of states including label } q = \{1, 2\}$

$$[r] = \{2, 3\}$$

An algorithm will now be presented where

$\llbracket \exists \Diamond q \rrbracket$ is computed

Note that $\llbracket \exists \Diamond q \rrbracket = \{0, 1, 2\}$

μ -calculus

A tool to generate algorithms for analyzing in which states a temporal logic formula is valid. This includes both LTL, CTL and CTL*.

Ex Expansion law for

$$\exists \Diamond q \equiv q \vee \exists O(\exists \Diamond q)$$

Introduce a logical variable $y = \exists \Diamond q \Rightarrow$

The expansion law can be expressed as

$$y = q \vee \exists O y$$

This expression will be iterated until a fixed point is reached, μ -calculus is a logic that includes fixed point operators.

Syntax for μ -calculus

$\psi ::= p \mid y \mid \neg \psi \mid \psi_1 \wedge \psi_2 \mid f(\psi) \mid \mu y. \psi$

$p \in AP$

f = function including the next modality + quantifier either

\exists or \forall .

$\mu y. \psi$ is the least fixed point

$\nu y. \psi$ is the greatest fixed point

Semantics of μ -calculus

Given a transition system

$$G = \langle \Sigma, \Sigma, T, I, AP, \rangle$$

the set of states $[[\psi]]$ where the μ -calculus formula ψ holds is defined as follows:

$$[[p]] = \{x \mid p \in \lambda(x)\}$$

$$[[y]] = Y \in 2^{\Sigma}$$

$$[[\neg \psi]] = \sim [[\psi]] = \Sigma \setminus [[\psi]]$$

$$[[\psi_1 \wedge \psi_2]] = [[\psi_1]] \cap [[\psi_2]]$$

$$[[f(\psi)]] = \text{Pre}^f([[\psi]])$$

$$\begin{aligned} [[\mu y. \psi]] &= \mu Y. \psi(Y) = \\ &\text{least fixed point for } Y = \\ &= \bigcap \{Y \in 2^{\Sigma} \mid Y = \psi(Y)\} \end{aligned}$$

$\llbracket \forall y. \psi \rrbracket = \nu Y. \psi(Y) =$
 $= \text{greatest fixed point}$

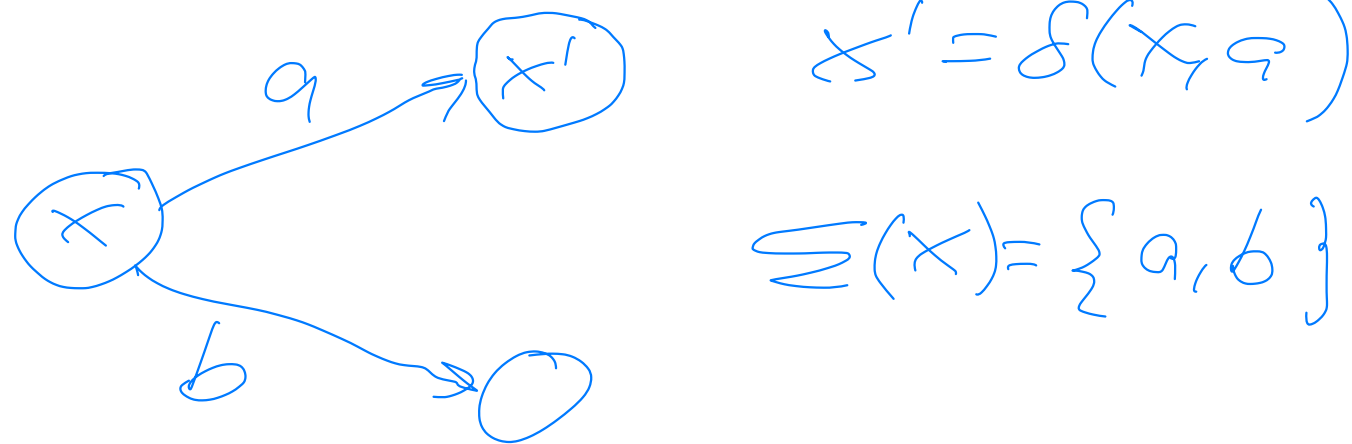
for $Y = \bigcup \{ Y \in 2^{\Sigma} \mid Y = \psi(Y) \}$

where $\psi(Y) = \psi(\llbracket y \rrbracket) =$
 $= \llbracket \psi(y) \rrbracket$

$\delta(x, a) = \text{transition function}$

$\Sigma(x) = \text{active event set}$

$= \text{possible events defined}$
 $\text{in state } x$



$\text{Pre}^{\exists}(Y) = \llbracket \exists o y \rrbracket =$

$= \{x \mid \exists a \in \Sigma(x) : \delta(x, a) \subseteq Y\}$

For at least one event

$a, x' = \delta(x, a) \in Y \Leftrightarrow$

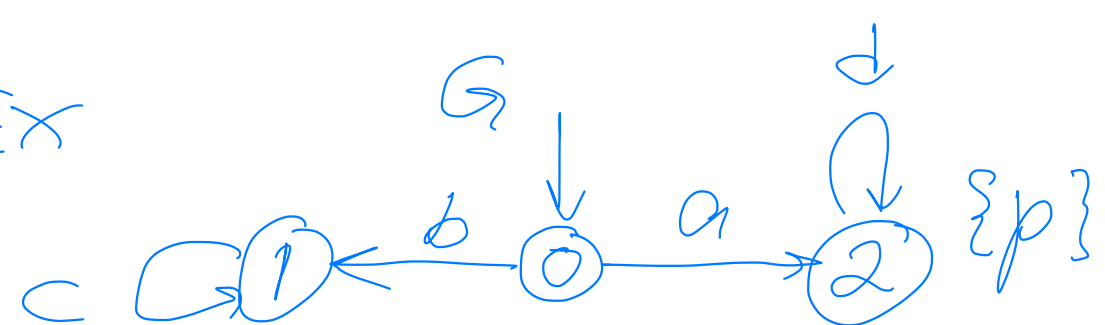
coreachability one
 step backward.

$\text{Pre}^{\forall}(Y) = \llbracket \forall o y \rrbracket =$

$= \{x \mid \forall a \in \Sigma(x) : \delta(x, a) \subseteq Y\}$

All target states from
 x must belong to Y

Ex

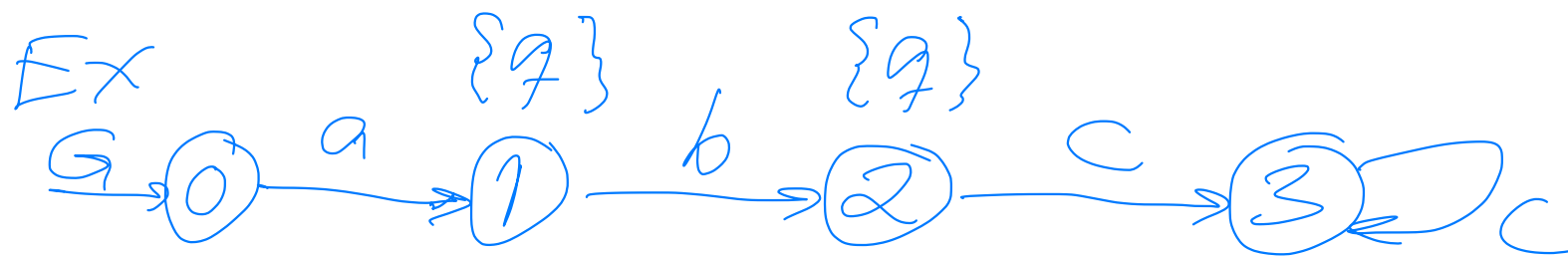


$$\text{Pre}^{\exists}(\llbracket p \rrbracket) = \text{Pre}^{\exists}(\{2\}) = \{0, 2\}$$

$$\text{Pre}^{\forall}(\llbracket p \rrbracket) = \text{Pre}^{\forall}(\{2\}) = \{2\}$$

state 0 is not included
since state 1 is missing
in $\text{Pre}^{\forall}(\{2\})$

Ex



$$\llbracket \exists \Diamond q \rrbracket \neq \llbracket \Diamond q \rrbracket$$

$$\llbracket \exists \Diamond q \rrbracket \leftrightarrow \mu\text{-calculus formula}$$

$$\mu y. q \vee \exists o y$$

$\underbrace{\hspace{10em}}_{\psi(y)}$

$$\begin{aligned} \llbracket \exists \Diamond q \rrbracket &= \llbracket \mu y. \psi(y) \rrbracket = \\ &= \mu Y. \underline{\psi}(Y) \quad \text{where} \end{aligned}$$

$$\underline{\psi}(Y) = \llbracket \psi(y) \rrbracket = \llbracket q \vee \exists o y \rrbracket$$

$$= \llbracket q \rrbracket \cup \llbracket \exists o y \rrbracket = \llbracket q \rrbracket \cup \text{Pre}^{\exists}(Y)$$

$$\text{where } Y = \llbracket y \rrbracket$$

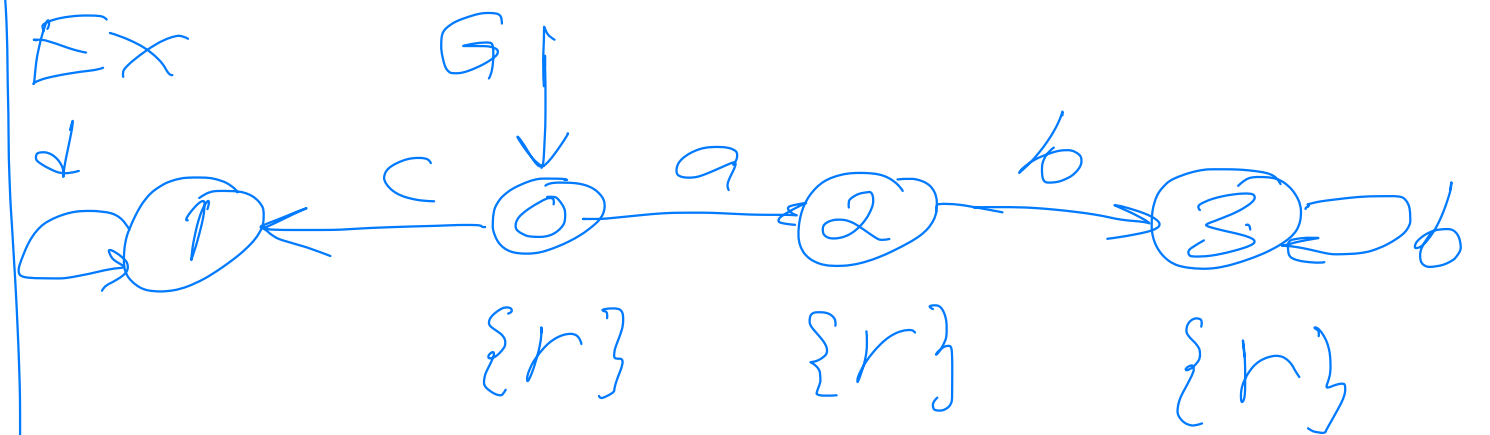
The least fixed point of $Y = \Psi(Y)$ is obtained by iterating $Y_{i+1} = \Psi(Y_i)$ for $i = 0, 1, 2, \dots$ with $Y_0 = \emptyset$ until $Y_{i+1} = Y_i$

$$Y_1 = \Psi(Y_0) = [q] \cup \text{Pre}^\exists(Y_0) \\ = \{1, 2\} \cup \underbrace{\text{Pre}^\exists(\emptyset)}_{\emptyset} = \{1, 2\}$$

$$Y_2 = \Psi(Y_1) = [q] \cup \text{Pre}^\exists(Y_1) \\ = \{1, 2\} \cup \{0, 1\} = \{0, 1, 2\}$$

$$Y_3 = \Psi(Y_2) = \{1, 2\} \cup \underbrace{\text{Pre}^\exists(\{0, 1, 2\})}_{\{0, 1\}} \\ = \{0, 1, 2\} = Y_2 = Y_\omega$$

$$[\exists \Diamond q] = \{0, 1, 2\}$$



$$\forall \Box r \leftrightarrow \forall y. r \wedge \forall \phi y$$

$\underbrace{\quad}_{\Psi(y)}$

$$[\forall \Box r] = \nu Y. \Psi(Y) \text{ where}$$

$$\Psi(Y) = [r \wedge \forall \phi y] = [r] \cap \text{Pre}^\forall(Y)$$

Greatest fixed point is obtained by iterating $Y_{i+1} = \Psi(Y_i)$ with $Y_0 = \Sigma = \{0, 1, 2, 3\}$

$$Y_1 = \Psi(Y_0) = \underbrace{\{0, 2, 3\}}_{[r]} \cap \underbrace{\text{Pre}^\forall(\Sigma)}_{\Sigma} =$$

$$= \{0, 2, 3\}$$

$$Y_2 = \Psi(Y_1) = \llbracket r \rrbracket \cap \text{Pre}^\forall(\{0, 2, 3\}) = \{2, 3\}$$

state 0 not included since the
target state 1 from state 0
is not included in Y in

$$\text{Pre}^\forall(Y) \quad Y = \{0, 2, 3\}$$

$$Y_3 = \llbracket r \rrbracket \cap \underbrace{\text{Pre}^\forall(\{2, 3\})}_{\{2, 3\}} = Y_2 = Y_\omega$$

$$\llbracket \forall \Box r \rrbracket = Y_\omega = \{2, 3\}$$

↑
fixed point

9. Temporal logic

For a transition system G with state set Σ and initial state set I , G satisfies a temporal logic formula Ψ , written as $G \models \Psi$, if Ψ holds in all initial states of G , i.e. if

$$I \subseteq \llbracket \Psi \rrbracket = \{x \in \Sigma \mid x \models \Psi\}$$

The set $\llbracket \Psi \rrbracket$ can be determined by μ -calculus.

Ex



Evaluate the nonblocking condition $\llbracket \forall \square \exists \Diamond p \rrbracket$, where the state label $\{p\}$ determines a marked state.

$$\text{since } \exists \Diamond p \equiv \mu z. p \vee \exists 0 z$$

$$\forall \square q \equiv \nu y. q \wedge \forall 0 y$$

$$\forall \square \exists \Diamond p \equiv \nu y. (\mu z. p \vee \exists 0 z) \wedge \forall 0 y$$

$$\llbracket \mu z. p \vee \exists 0 z \rrbracket = \mu z. \Psi(z)$$

where $\Psi(z) = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(z) = z$

Least fixed point:

$$z_{i+1} = \Psi(z_i), \quad z_0 = \emptyset$$

$$Z_1 = \Psi(Z_0) = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(Z_0) = \{2\} \cup \emptyset = \{2\}$$

$$Z_2 = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(Z_1) = \{2\} \cup \{0\} = \{0, 2\}$$

$$Z_3 = \{2\} \cup \text{Pre}^{\exists}(\{0, 2\}) = \{2\} \cup \{0\} = \{0, 2\} = Z_2 = \text{fixed point } Z^w$$

$$\llbracket \forall \Diamond \exists \Diamond p \rrbracket = \bigvee Y. \Psi(Y)$$

$$\text{where } \Psi(Y) = \llbracket M \models p \vee \exists 0 \in Z \rrbracket$$

$$\cap \text{Pre}^{\forall}(Y) = Z^w \cap \text{Pre}^{\forall}(Y) = Y$$

Greatest fixed point:

$$Y_{i+1} = \Psi(Y_i), \quad Y_0 = \overline{X}$$

$$Y_1 = Z^w \cap \text{Pre}^{\forall}(\overline{X}) = \{0, 2\} \cap \{0, 1, 2\} = \{0, 2\}$$

$$Y_2 = Z^w \cap \text{Pre}^{\forall}(Y_1) = \{0, 2\} \cap \{2\} \leftarrow (0 \text{ is excluded since its target state } 1 \text{ is not included in } Y_1 = \{0, 2\}) = \{2\}$$

$$Y_3 = Y_2 = \text{fixed point } Y^w$$

$$I = \{0\} \not\subseteq \llbracket \forall \Diamond \exists \Diamond p \rrbracket = Y^w = \{2\} \Rightarrow G \not\models \forall \Diamond \exists \Diamond p$$

Remove state  \Rightarrow

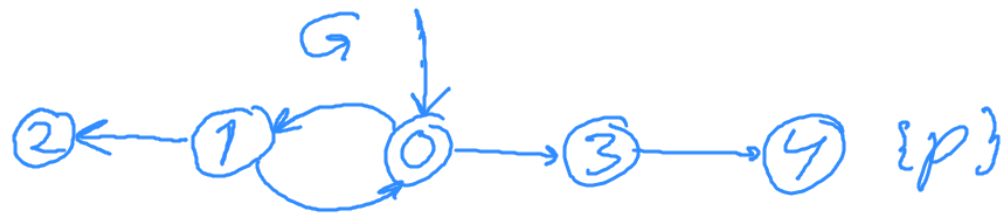
$$Z^w = \{0, 2\} \quad Y_1 = Y_2 = Y^w = \{0, 2\}$$

$$\Rightarrow I \subseteq Y^w \text{ and } G \models \forall \Diamond \exists \Diamond p.$$

Correct since state 2 is a blocking state.

Removing this blocking state
 $\Rightarrow G$ is nonblocking and satisfies
 the nonblocking condition $\forall \Box \exists \Diamond p$

Ex



Evaluate again the nonblocking
 condition $\forall \Box \exists \Diamond p$

$\llbracket \exists \Diamond p \rrbracket = \mu Z. \Psi(Z)$ where

$$\Psi(Z) = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(Z) = Z$$

Least fixed point iteration:

$$Z_0 = \emptyset, \quad Z_1 = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(\emptyset) = \{4\} \cup \emptyset$$

$$Z_2 = \llbracket p \rrbracket \cup \text{Pre}^{\exists}(\{4\}) = \{4\} \cup \{3\} = \{3, 4\}$$

$$Z_3 = \llbracket p \rrbracket \cup \text{Pre}^{\exists}\{3, 4\} = \{4\} \cup \{0, 3\} = \{0, 3, 4\}$$

$$Z_4 = \{4\} \cup \{1, 0, 3\} = \{0, 1, 3, 4\}$$

$$Z_5 = Z_4 = \text{fixed point } Z^w = \{0, 1, 3, 4\}$$

$$\llbracket \forall \Box \exists \Diamond p \rrbracket = \bigvee Y. Z^w \cap \text{Pre}^{\forall}(Y)$$

Greatest fixed point iteration:

$$Y_0 = \Sigma, \quad Y_1 = Z^w \cap \text{Pre}^{\forall}(\Sigma) = Z^w \cap \Sigma = Z^w = \{0, 1, 3, 4\}$$

$$Y_2 = Z^w \cap \text{Pre}^{\forall}(\underbrace{\{0, 1, 3, 4\}}_{Y_1}) =$$

$= Z^w \cap \{0, 3, 4\} = (\text{state 1 excluded since its target state is not included in } Y_1 \text{ in } \text{Pre}^V(Y_1)) = \{0, 3, 4\}$

$$Y_3 = Z^w \cap \text{Pre}^V(\underbrace{\{0, 3, 4\}}_{Y_2}) = \{3, 4\}$$

since the target state 1 of state 0 not included in Y_2

$$Y_4 = Y_3 = \text{fixed point } Y^w = \{3, 4\}$$

$$I = \{0\} \notin Y^w = \llbracket \forall \Diamond \exists \Diamond p \rrbracket \Rightarrow$$

$$G \not\models \forall \Diamond \exists \Diamond p$$

10. Reinforcement learning (RL)

RL = optimization method where actions are sent to the plant and resulting rewards are evaluated such that optimal actions are selected after an initial learning phase. Currently a very popular method within modern AI and machine learning.

AlphaGo is a popular program where its success is based on RL.