

# Introduction to Data Management

Databases in Theory and Practice

Alyssa Pittman

Based on slides by Jonathan Leang, Dan Suciu, et al

Paul G. Allen School of Computer Science and Engineering University of Washington, Seattle

# First, a story...

#### **UW's Databases**

# The Seattle Times

Education | Local News

# New UW payroll system behind schedule, more costly than expected

Originally published November 26, 2015 at 2:36 pm | Updated November 27, 2015 at 6:19 am

A project to modernize the University of Washington's payroll system is costing millions more and taking longer than expected.

## Outline

- 1. Administrivia
- The Relational Data Model
- 3. Databases, SQL, and RA

# What am I going to learn?

- Course Topics
  - Relational Databases
    - Queries
    - Database Design
    - Optimization
    - Transactions and Parallelism
  - NoSQL Databases
    - Semi-Structured Document Databases
    - Graph Databases
- Tools
  - Experimental to Enterprise Platforms
  - Cloud Services (AWS, Azure)

# What am I going to learn?

- After the course, you will be able to...
  - Explain how a query is processed end-to-end
  - Integrate a database into an application
  - Effectively manage data for long-term use
  - Create database constructs to provide speedups
  - Make design choices when selecting tools for a project

## 344 Staff

#### Instructor: Alyssa Pittman

- Email: <a href="mailto:smooo@cs.washington.edu">smooo@cs.washington.edu</a>
- Office hours: Wednesday 3:00-4:00 in CSE 216 or by appointment

#### **TAs**

- Andrew Guterman
- Jack Khuu
- Matthew Liu
- Phong Ly
- Samuel Oh
- Khang Phan
- Matt Tenczar
- Eric Yeh
- Stella Zhang

# Communication - you to us

#### Getting in touch with us

- Office hours daily
- Piazza: <a href="https://piazza.com/washington/winter2020/cse344/home">https://piazza.com/washington/winter2020/cse344/home</a>
  - Best place to ask questions
  - And to help other students!
- Mailing list: cse344-staff@cs.washington.edu
  - For personal inquiries

"Asking for help is not a sign of weakness, it's a sign of strength."

# Communication - us to you

- Course website: <a href="http://cs.uw.edu/344">http://cs.uw.edu/344</a>
  - Everything will be here
- Piazza announcements
- Class mailing list
  - Very low traffic, only important announcements

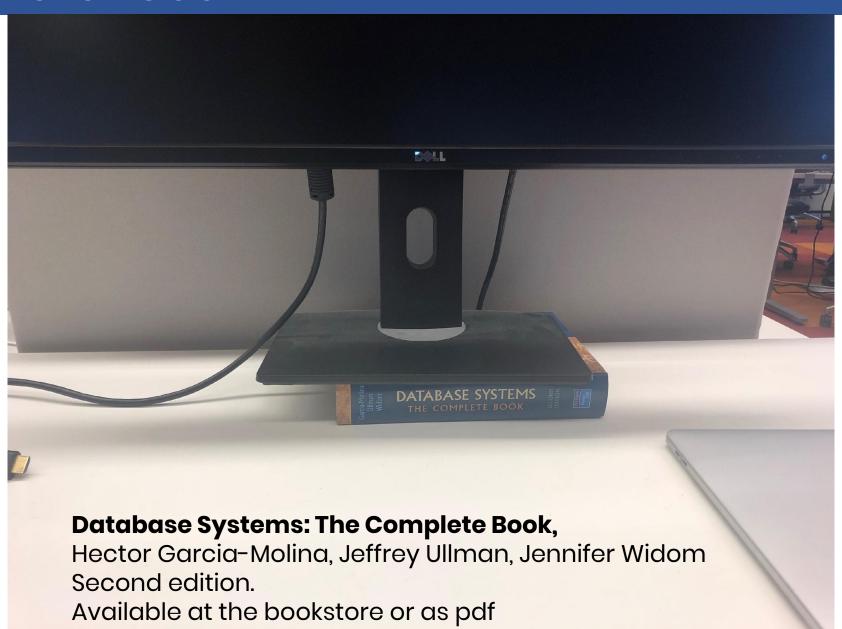
## Course Format

- Lectures: this room, please attend!
- Sections: Thursdays, see web for locations.
  - Bring your laptop!
- 8 homework assignments
- Midterm and final

## Exams

- Midterm (February 10th in class) and Final (March 16th 2:30 pm)
- You may bring letter-size piece of paper with notes
  - Handwritten
  - May write on both sides
  - o Midterm: 1 sheet, Final: 2 sheets
- Closed book. No computers, phones, watches,...
- Location: this classroom

## References



## Administrivia

#### Grading:

- 40% HW, 20% Midterm, 30% Final, 10% Participation
- 4 late days, 2 days max per assignment

#### Collaboration:

- HW must be done and typed up individually, though you can talk to other students about your approach
- We will run cheating detection

## Administrivia



# Let's get started!

## Database

What is a database?

A collection of files storing related data

Examples?

## Database

#### What is a database?

A collection of files storing related data

#### Examples?

- Payroll database
- UW student database
- Amazon products database
- Airline reservation database

What is a DBMS?

#### What is a DBMS?

 A big program written by someone else that allows us to manage efficiently a large database and allows it to persist over long periods of time

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Examples?

#### What is a DBMS?

 A big program written by someone else that allows us to manage efficiently a large database and allows it to persist over long periods of time

#### Examples?

- Oracle, Microsoft SQL Server, Teradata
- Open source: MySQL (Sun/Oracle), PostgreSQL, CouchDB
- Open source library: SQLite

We will focus on relational DBMSs most quarter

## Think About This

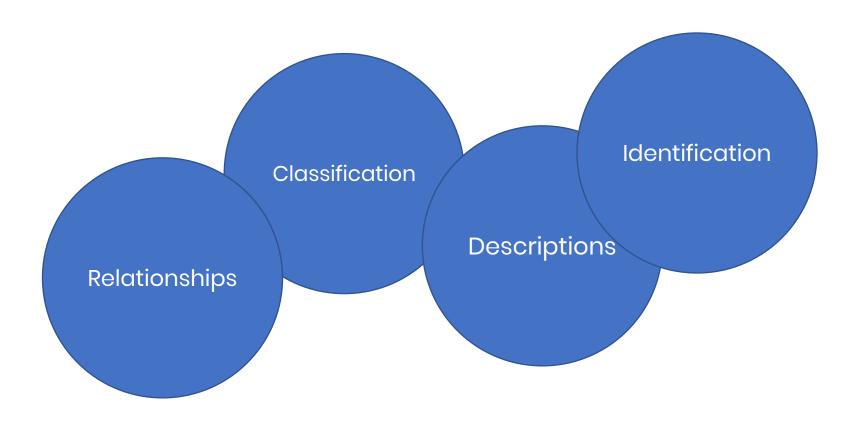


How do we describe information?

## Think About This



#### How do we describe information?



## Think About This



#### How do we describe information?

#### **Data Model**

A **Data Model** is a mathematical formalism to describe data. It is how we can talk about data conceptually without having to think about implementation.

## 3 Parts of a Data Model

## The 3 parts of any data model

- Instance
  - The actual data
- Schema
  - A description of what data is being stored
- Query Language
  - How to retrieve and manipulate data

## Data Model Zoo

#### There are lots of models out there!

#### Relational

- Semi-structured
- Key-value pairs
- Graph
- Object-oriented

**-** ...

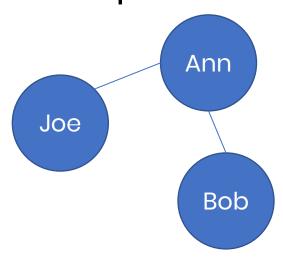
Jan 2019	Rank Dec 2018	Jan 2018	DBMS	Database Model
1.	1.	1.	Oracle 🚹	Relational DBMS
2.	2.	2.	MySQL 🖽	Relational DBMS
3.	3.	3.	Microsoft SQL Server 🖽	Relational DBMS
4.	4.	4.	PostgreSQL 🖽	Relational DBMS
5.	5.	5.	MongoDB 🚦	Document store
6.	6.	6.	IBM Db2 □	Relational DBMS
7.	7.	<b>1</b> 9.	Redis 🗄	Key-value store
8.	8.	<b>1</b> 0.	Elasticsearch 🗄	Search engine
9.	9.	<b>4</b> 7.	Microsoft Access	Relational DBMS
10.	10.	<b>1</b> 11.	SQLite 🛅	Relational DBMS

https://db-engines.com/en/ranking

# Multiple Representation

Most data can be represented in different ways

An example of Facebook friends



**Graph model** 

Person 1	Person 2	Friend
Joe	Ann	1
Ann	Bob	1
Bob	Joe	0

**Relational Model** 

## The Relational Model

Information Retrieval

P. BAXENDALE, Editor

#### A Relational Model of Data for Large Shared Data Banks

E. F. Codd IBM Research Laboratory, San Jose, California

Future users of large data banks must be protected from having to know how the data is organized in the machine (the internal representation). A prompting service which supplies such information is not a satisfactory solution. Activities of users at terminals and most application programs should remain

The relational view (or model) of data described in Section 1 appears to be superior in several respects to the graph or network model [3, 4] presently in vogue for non-inferential systems. It provides a means of describing data with its natural structure only—that is, without superimposing any additional structure for machine representation purposes. Accordingly, it provides a basis for a high level data language which will yield maximal independence between programs on the one hand and machine representation and organization of data on the other.

A further advantage of the relational view is that it forms a sound basis for treating derivability, redundancy, and consistency of relations—these are discussed in Section

Levein and Maron [2] provide numerous references to work in this area.

In contrast, the problems treated here are those of data independence—the independence of application programs and terminal activities from growth in data types and changes in data representation—and certain kinds of data inconsistency which are expected to become troublesome even in nondeductive systems.

Volume 13 / Number / June, 1970

/ June, 1970

those existing systems which either require or permit data elements to be stored in at least one total ordering which is closely associated with the hardware-determined ordering of addresses. For example, the records of a file concerning parts might be stored in ascending order by part serial number. Such systems normally permit application programs to assume that the order of presentation of records from such a file is identical to (or is a subordering of) the

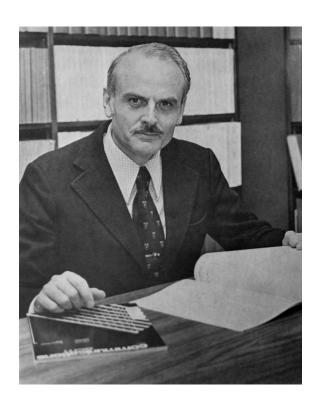
Communications of the ACM

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## The Relational Model

Most common way people describe information.





Payroll(Userld, Name, Job, Salary)

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000*
345	Allison	TA	60000*
567	Magda	Prof	90000
789	Dan	Prof	100000

\* salaries may not reflect reality

# Table/Relation

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

# Table/Relation

Rows/ Tuples/ Records

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

Table/Relation

Columns/Attributes/Fields

Rows/ Tuples/ Records

×				
	UserID	Name	Job	Salary
	123	Jack	TA	50000
	345	Allison	TA	60000
	567	Magda	Prof	90000
	789	Dan	Prof	100000

## Characteristics of the Relational Model

- Set semantics
  - No duplicate tuples
- Attributes are typed and static
  - INTEGER, FLOAT, VARCHAR(n), DATETIME, ...
- Tables are flat

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#### Set semantics

- No duplicate tuples
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UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

UserID	name	Jon	Salary
567	Magda	Prof	90000
123	Jack	TA	50000
789	Dan	Prof	100000
345	Allison	ТΔ	60000

Order doesn't matter

### Set semantics

- No duplicate tuples
- Attributes are typed and static
  - INTEGER, FLOAT, VARCHAR(n), DATETIME, ...
- Tables are flat

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000
789	Dan	Prof	100000

Violates set semantics!

- **Set semantics** □ not in most DBMS implementations
  - No duplicate tuples
- Attributes are typed and static
  - INTEGER, FLOAT, VARCHAR(n), DATETIME, ...
- Tables are flat

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000
789	Dan	Prof	100000

Violates set semantics!

- **Set semantics** □ not in most DBMS implementations
  - No duplicate tuples
- Attributes are typed and static
  - INTEGER, FLOAT, VARCHAR(n), DATETIME, ...
- Tables are flat

UserID	Name	Job	Salary
123	Jack	TA	banana
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

Violates attribute type assuming INT

- **Set semantics** □ not in most DBMS implementations
  - No duplicate tuples
- Attributes are typed and static
  - INTEGER, FLOAT, VARCHAR(n), DATETIME, ...
- Tables are flat

#### No sub-tables allowed!

UserID	Name	Job		Salary
123	Jack	JobName	HasBananas	50000
		TA	0	
		banana picker	1	
345	Allison	TA		60000
567	Magda	Prof		90000
789	Dan	Prof		100000

### But how is this data ACTUALLY stored?

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

### But how is this data ACTUALLY stored?

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

Don't know. Don't care.

# Physical Data Independence

# Structured Query Language - SQL

Alright, I have data and a schema. How do I access it?

# Structured Query Language - SQL

"SQL (standing for Structured Query Language) is the standard language for relational database management systems. When it originated back in the 1970s, the domain-specific language was intended to fulfill the need of conducting a database query that could navigate through a network of pointers to find the desired location. Its application in handling structured data has fostered in the Digital Age. In fact, the powerful database manipulation and definition capabilities of SQL and its intuitive tabular view have become available in some form on virtually every important computer platform in the world.

Some notable features of SQL include the ability to process sets of data as groups instead of individual units, automatic navigation to data, and the use of statements that are complex and powerful individually. Used for a variety of tasks, such as querying data, controlling access to the database and its objects, guaranteeing database consistency, updating rows in a table, and creating, replacing, altering and dropping objects, SQL lets users work with data at the logical level.

Read more at the ANSI Blog: The SQL Standard – ISO/IEC 9075:2016 <a href="https://blog.ansi.org/?p=158690">https://blog.ansi.org/?p=158690</a>

# Structured Query Language - SQL

- Key points about SQL:
  - A domain-specific language
    - SQL only works on relational databases
    - Not for general purpose programming (Java, C/C++, ...)
  - A declarative language
  - Logical level of interaction with data

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

```
SELECT P.Name, P.UserID
FROM Payroll AS P
WHERE P.Job = 'TA';
```

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

SELECT What kind

What kind of data I want

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```
SELECT P.Name, P.UserID
```

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

**SELECT** 

What kind of data I want

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**FROM** 

Where the data coming from

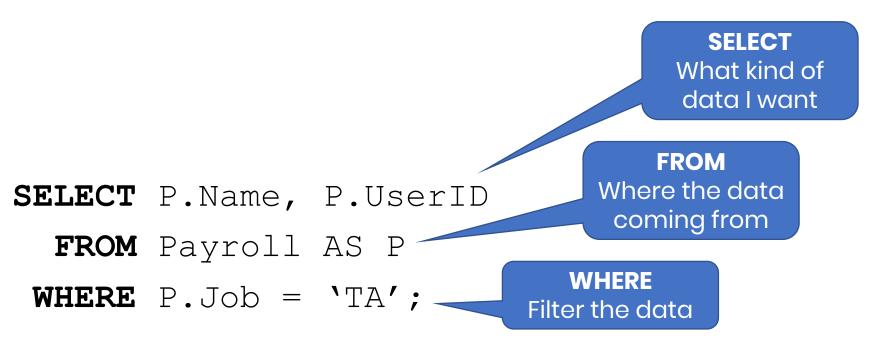
SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = 'TA';

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000

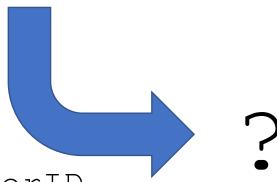


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#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



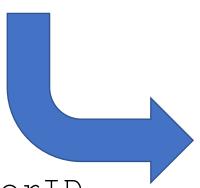
SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = TA';

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



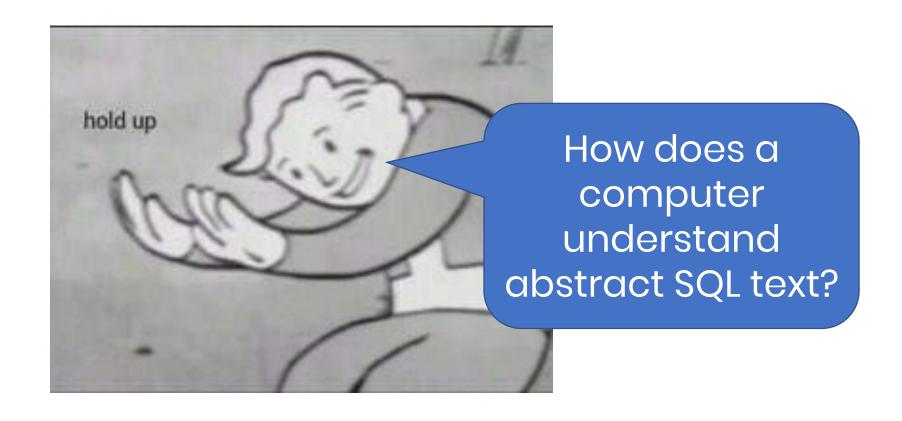
Name	UserID
Jack	123
Allison	345

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SELECT P.Name, P.UserID

FROM Payroll AS P

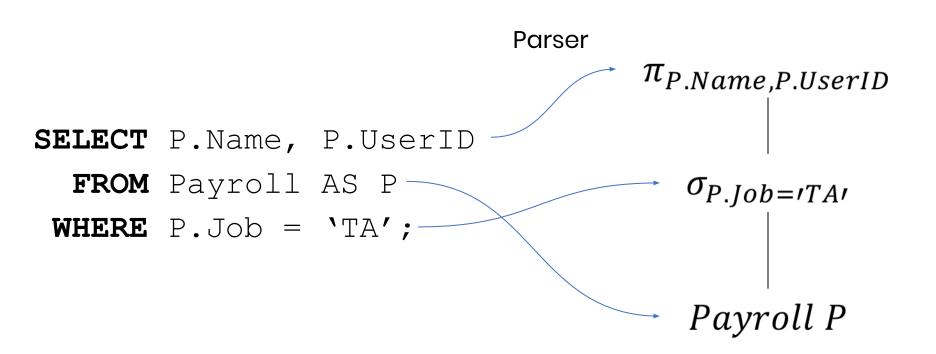
WHERE P.Job = TA';



- Code has to boil down to instructions at some point
- Relational Database Management Systems (RDBMSs) use Relational Algebra (RA)

```
SELECT P.Name, P.UserID
FROM Payroll AS P
WHERE P.Job = 'TA';
```

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 $\pi_{P.Name,P.UserID}$   $\sigma_{P.Job=\prime TA\prime}$  Payroll~P

For-each semantics

- Code has to boil down to instructions at some point
- Relational Database Management Systems (RDBMSs) use Relational Algebra (RA).

### For-each semantics

```
\pi_{P.Name,P.UserID}

for each row in P:

\sigma_{P.Job='TA'}

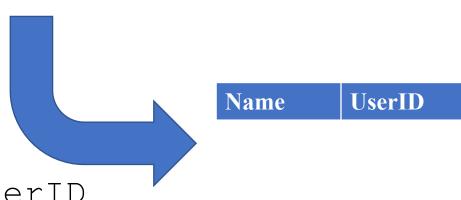
if (row.Job == 'TA'):

output (row.Name, row.UserID)

Payroll\ P
```

### **Payroll**

UserID	Name	Job	Salary	
123	Jack	TA	50000	<b>—</b>
345	Allison	TA	60000	
567	Magda	Prof	90000	
789	Dan	Prof	100000	



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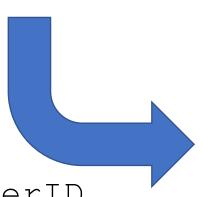
SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = TA';

### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



Name	UserID
Jack	123

SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = 'TA';

### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



Name	UserID
Jack	123

SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = 'TA';

### **Payroll**

UserID	Name	Job	Salary	
123	Jack	TA	50000	
345	Allison	TA	60000 ←	
567	Magda	Prof	90000	
789	Dan	Prof	100000	



Name	UserID
Jack	123
Allison	345

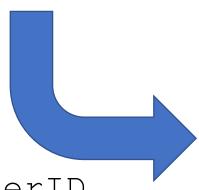
SELECT P.Name, P.UserID

FROM Payroll AS P

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### **Payroll**

UserID	Name	Job	Salary	
123	Jack	TA	50000	
345	Allison	TA	60000	
567	Magda	Prof	90000	<b>(</b>
789	Dan	Prof	100000	



Name	UserID
Jack	123
Allison	345

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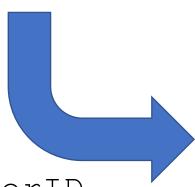
SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = TA';

#### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



Name	UserID
Jack	123
Allison	345

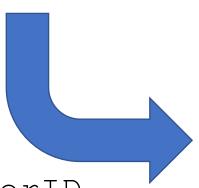
SELECT P.Name, P.UserID

FROM Payroll AS P

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### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
345	Allison	TA	60000
567	Magda	Prof	90000
789	Dan	Prof	100000



Name	UserID
Jack	123
Allison	345

SELECT P.Name, P.UserID

FROM Payroll AS P

WHERE P.Job = 'TA';

### Create Table Statement

# Payroll(Userld, Name, Job, Salary)

```
CREATE TABLE Payroll (
UserId INT,

Name VARCHAR(100),

Job VARCHAR(100),

Salary INT);
```

\* Case insensitive, but useful for readability

### Insert Statement

### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000



### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000
912	Alyssa	Lecturer	45000

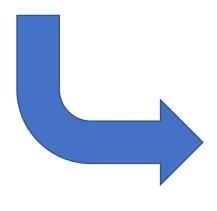
#### INSERT INTO

Payroll(UserId, Name, Job, Salary)
VALUES(912, 'Alyssa', 'Lecturer', 45000);

### Insert Statement

### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000



### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000
912	Alyssa	Lecturer	45000

Can omit column names if inserting to all columns in order

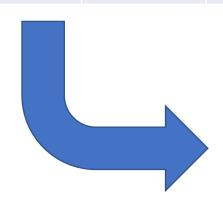
INSERT INTO

Payroll (UserId, Name, Job, Salary)
VALUES (912, 'Alyssa', 'Lecturer', 45000);

### Insert Statement

### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000



### **Payroll**

UserID	Name	Job	Salary
123	Jack	TA	50000
789	Dan	Prof	100000
912	Alyssa	Lecturer	45000

Can omit column names if inserting to all columns in order

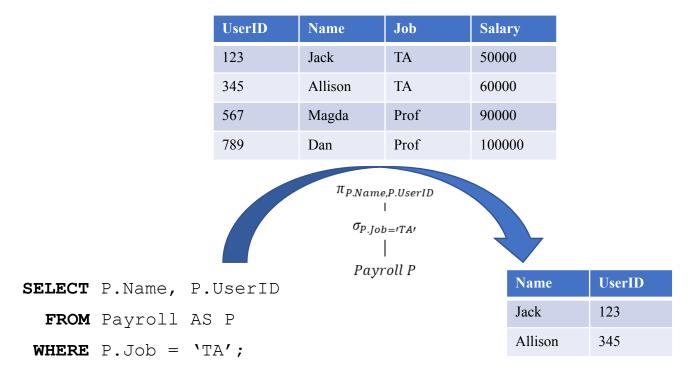
INSERT INTO

Payroll

VALUES (912, 'Alyssa', 'Lecturer', 45000);

# Today's Takeaways!

- The Relational Model concept
- How a basic SELECT-FROM-WHERE query works
- Basic execution process (RA) inside a RDBMS



### Homework

- HW1 released today via GitLab
  - Check the website!
- TAs will cover relevant tools in section