VE281 Sorting Algorithms Comparison Report

Xinhao Liao 516370910037

September 2018

1 Introduction

There are various algorithms used to sort an array of integers. Among them, Bubble Sorting, Insertion Sorting, Selection Sorting, Merge Sorting and Quicksorting are all commonly used comparison sorting algorithms. In this report, these algorithms are to be compared with randomly generated integer arrays. (Quicksoting is respectively implemented with and without $\Omega(n)$ cost of space.)

2 Background

Consider an array A with n integers as elements. The array can be sorted with following algorithms.

Bubble Sorting

Compares two adjacent items and swap them to keep them in ascending order. From the beginning to the end, compare two adjacent items and swap them to keep them in ascending order. The last item will be the largest one. Then similarly find the second largest element, and so on. The time complexity should be $\Theta(n^2)$ for the best case, the average case, and the worst case. The C++ code for this algorithm can be found in Appendix A as the function BubbleSort.

Insertion Sorting

Notice that A[0] alone is a sorted array. For i=1 to N-1, insert A[i] into the appropriate location in the sorted array with the first i-1 elements, so that the first i elements are sorted. To do so, save A[i] in a temporary variable temp, shift sorted elements greater than temp right, and then insert temp in the gap. The time complexity should be $\Theta(n)$ for the best case, and $\Theta(n^2)$ for the average case and the worst case. The C++ code for this algorithm can be found in Appendix A as the function InsertionSort.

Selection Sorting

For i=0 to n-2, find the smallest item in the array $A[i], \ldots, A[n-1]$. Then, swap that item with A[i]. Finding the smallest item requires linear scan. The time complexity should be $\Theta(n^2)$ for the best case, the average case, and the worst case. The C++ code for this algorithm can be found in Appendix A as the function SelectionSort.

Merge Sorting

Spilt array into two (roughly) equal subarrays. Merge sort each subarray recursively. Then the two subarrays will be sorted. Finally, recursively merge the two sorted subarrays back into a sorted array. The time complexity should be $\Theta(nlogn)$ for the best case, the average case, and the worst case. The C++ code for this algorithm can be found in Appendix A as the function MergeSort, which uses a function MergeSortHelper to help make it recursive, and a function Merge to help merge subarrays back.

Quick Sorting

First, choose an array element as pivot. Put all elements less than pivot to the left of pivot. Put all elements greater than or equal to pivot to the right of pivot. Move pivot to its correct place in the array. Finally, sort left and right subarrays recursively (not including pivot).

The time complexity should be $\Theta(nlogn)$ for the best case and the average case, and $\Theta(n^2)$ for the worst case. The C++ code for this algorithm can be found in Appendix A as the function QuickSort1 and QuickSort2, which uses a function QuickSortHelper to help make it recursive and a function getpivot to generate a pivot index. For QuickSort1, it used a function $partition_extra_place$ to sort with $\Omega(n)$ space cost. And for QuickSort2, it used a function $partition_in_place$ to sort, which costs on average only $\Theta(logn)$ stack space, which makes it weakly in-place.

Summary

	Worst Case Time	Average Case Time	In Place	Stable
Insertion	$O(N^2)$	$O(N^2)$	Yes	Yes
Selection	$O(N^2)$	$O(N^2)$	Yes	No
Bubble	$O(N^2)$	$O(N^2)$	Yes	Yes
Merge Sort	$O(N \log N)$	$O(N \log N)$	No	Yes
Quick Sort	$O(N^2)$	$O(N \log N)$	Weakly	No

Figure 1: Comparison sorting algorithms summary.

3 Procedures

- 1. Set i to be the size of the array, which is initially 10 in this report.
- 2. Generate i random integers in the array. Save a copy in an auxiliary array.
- 3. Sort the array with the algorithms in order and measure the time cost. Write the measured time cost to an output file. Remember to copy the integers in order back to keep the same order for every algorithm.
- 4. Multiply i with 10 and repeat steps 2 to 3 until i reaches 100000.

The code implementing the above procedures is shown in Appendix B.

4 Result

Size	10	100	1000	10000	100000
Bubble Sort	$0.004 \mathrm{ms}$	$0.123 \mathrm{ms}$	$5.721 \mathrm{ms}$	777.441 ms	$80180.9 \mathrm{ms}$
Insertion Sort	$0.001 \mathrm{ms}$	$0.017 \mathrm{ms}$	$1.918 \mathrm{ms}$	$146.042 { m ms}$	$14213.4\mathrm{ms}$
Selection Sort	$0.002 \mathrm{ms}$	$0.08 \mathrm{ms}$	$3.105 \mathrm{ms}$	289.307 ms	$27408 \mathrm{ms}$
Merge Sort	$0.006 \mathrm{ms}$	$0.023 \mathrm{ms}$	$0.343 \mathrm{ms}$	$3.57 \mathrm{ms}$	$42.832 \mathrm{ms}$
Quicksort not in place	$0.004 \mathrm{ms}$	$0.048 \mathrm{ms}$	$0.392 \mathrm{ms}$	$4.19 \mathrm{ms}$	$50 \mathrm{ms}$
Quicksort in place	$0.002 \mathrm{ms}$	$0.019 \mathrm{ms}$	$0.238 \mathrm{ms}$	$3.167 \mathrm{ms}$	$36.442 \mathrm{ms}$

Table 1: Result for sorting with different algorithms.

The results for sorting an array of 10, 100, 1000, 10000, and 100000 random integers are shown in Table 1. The integers are randomly generated in the range $[-2^{31}, 2^{31} - 1]$ and saved as *long int* in 32 bits. And we can plot all six curves corresponding to the six sorting algorithms in the same figure in Figure 1. The curves are plotted in log scales.

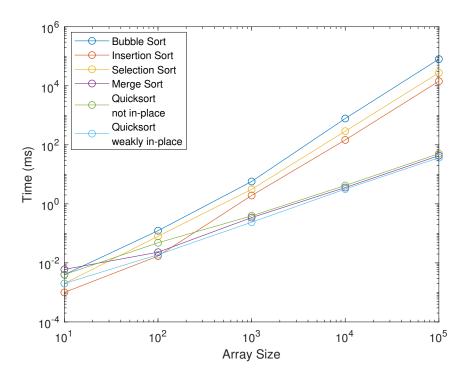


Figure 2: Curves corresponding to the six sorting algorithms.

The results are porduced in Linux Ubuntu18.04 running in VMware Workstation 14.x virtual machine distributed with 4GB memory and 3 processors, with Intel CORE i9.

According to the results shown above, Bubble Sorting, Insertion Sorting, and Selection Sorting on average have far worse performance than Merge Sorting and Quick Sorting. In general, Quicksorting in-place has the best performance. Quicksoting weakly-in-place is the second best, Merge Sorting the third, Insertion Sorting the fourth, Selection the fifth, and Bubble Sorting the last one. And when the size is small enough, Insertion sorting shows the best performance.

5 Discussion and analysis

The performance of Bubble Sorting, Insertion Sorting, and Selection Sorting is generally worse than that of Merge Sorting and Quick Sorting. This conforms to the theoretical analysis that Bubble Sorting, Insertion Sorting, and Selection Sorting have the time cost of $\Theta(n^2)$ in the average case, which are far worse than that of Merge Sorting and Quick Sorting, which are $\Theta(n\log n)$.

Among the 3 algorithms with average time cost of $\Theta(n^2)$, Bubble Sorting has the worst performance. Bubble Sorting, and Insertion Sorting has the best performance. This is reasonable, since in the best case, the time complexity of Insertion Sorting is $\Omega(n)$ which is better than all the other algorithms here. Insertion Sorting only needs to compare and swap when some order is not ascending. Once a comparison gives a correct answer, all the elements before doesn't need to be checked or swapped any more. And for Bubble Sorting, whatever order is given, it needs $\Theta(n^2)$ comparisons and can only swap with the neighboring element, which implies swapping $O(n^2)$ times. And for Selection sorting, it's similar to Bubble Sorting except that it only needs $\Theta(n)$ swaps. Comparing Insertion Sorting with Selection Sorting, Insertion Sorting may terminate in advance, and that may be the reason why it has better performance. However, notice that if Insertion Sorting is implemented by swapping the compared neighbors, rather than just copying the former ones as shown in Appendix B, Insertion Sorting may have worse performance than Selection Sorting. This is because a swap implies three assignments, which can make the Insertion Sorting cost much more.

Among the algorithms with average time cost of $\Theta(nlogn)$, as we can see, Quicksorting in place has the best performance, and Quicksorting not in place has the worst performance. The reason may be that reading

and writing in memory are quite slow. And compared with Merge Sorting, though they both require $\Omega(n)$ extra place, Merge Sorting avoids the worst case of Quicksorting, which partitions only 1 element, by always partitioning in half.

Also notice that when the size is small enough, the Insertion Sorting algorithm shows the best performance, even better than the algorithms with $\Theta(nlogn)$ average time complexity. This is reasonable since in the best case, the time complexity of Insertion Sorting is only $\Theta(n)$, which is the best time complexity, and Insertion Sorting may terminate in advance in some cases. And when the size is small, there is more chance that it's the case that the algorithm can terminate in advance. There is always 1 best case, and there are n! cases for an array of size n. The probability of the best case is $\frac{1}{n!}$, which increases quickly as n decreases.

6 Appendix A. Algorithms Implementation

sort.h

```
//sort.h
  //Implementation of the algorithms
  #ifndef __SORT_H__
  #define __SORT_H__
  #include <cstdlib>
  static inline void swap(long int Array[],long int x,long int y){
    //REQUIRES: Array initialized with index x and y,
    //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
    //MODIFIES: Array[x] and Array[y]
    //EFFECTS: swap Array[x] with Array[y]
    long int temp=Array[x];
11
    Array[x] = Array[y];
12
    Array [y]=temp;
14
15
  //Bubble Sorting
  void BubbleSort(long int Array[],long int number){
    //REQUIRES: Array initialized with at least 'number' elements
19
20
    //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements,
    //\text{number in } [0,2^{31}-1]
21
    //MODIFIES: Array
23
    //EFFECTS: sort the Array
24
    for (long int i=number-1; i>0; i--)
      for (long int j=0; j< i; j++){}
25
26
        if(Array[j]>Array[j+1])
27
          swap\left(\,Array\;,\;\;j\;,\;\;j+1\right);
28
    }
  }
30
31
  //Insertion Sorting
  void InsertionSort(long int Array[],long int number){
    //REQUIRES: Array initialized with at least 'number' elements
    //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
36
    //\text{number in } [0,2^{31}-1]
37
    //MODIFIES: Array
38
39
    //EFFECTS: sort the Array
40
    long int j;
    for (long int i=1; i<number; i++){
41
42
      long int temp=Array[i];
43
      //since all elements before have been sorted
        Array[j] = Array[j-1];
        //copy Array[j-1] to Array[j] instead of swap everytime to save time
46
47
      Array [j]=temp;
49
  }
50
51
  //Selection Sorting
```

```
54 void SelectionSort(long int Array[], long int number){
      //REQUIRES: Array initialized with at least 'number' elements
      //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31}-1 elements
56
      //\text{number in } [0,2^{31}-1]
57
      //MODIFIES: Array
58
59
      //EFFECTS: sort the Array
60
      long int min;
      for(long int i=0; i<(number-1); i++){}
61
62
        min = i:
        for (long int j=i; j<number; j++){
63
           if (Array[j]<Array[min])</pre>
65
             min=j;
66
        swap\left(\left.Array\,,\ i\,,\ min\right)\,;//\ similar\ to\ Bubble\ Sorting\,,
67
        //but instead of swaping every time only swap the minimum one
68
69
   }
70
71
72
   //Merge Sorting
   static void Merge(long int Array[], long int left, long int mid, long int right){
//REQUIRES: Array initialized with left mid and right as legal index
74
75
      //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
76
      //MODIFIES: Array
77
      //EFFECTS: Merge the array with elements [left:mid]
78
      //and the array with elements [mid+1:right]
79
      long int i = left;
80
      long int j = mid+1;
81
      long int k = 0;
82
      long int * Auxiliary = new long int [right - left + 1];
83
84
      //initizalize with new since there is a size limit for a static array
      while(i <= mid && j<= right){
85
86
        if (Array[i] <= Array[j])</pre>
           Auxiliary[k++]=Array[i++];
87
88
        else
89
           Auxiliary [k++]=Array [j++];
90
      while (k <= (right - left)) {
91
        if (i>mid)
92
           Auxiliary[k++]=Array[j++];
93
94
        else
95
           Auxiliary[k++]=Array[i++];
96
97
      for (i=left; i \le right; i++)
        Array[i]=Auxiliary[i-left];
98
99
      delete[] Auxiliary;
100
   }
102
   static void MergeSortHelper(long int Array[], long int left, long int right){
      //REQUIRES: Array initialized with left and right as legal index
      //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
      //MODIFIES: Array
106
      //EFFECTS: recursively merge sort the array
107
      if (left >= right) return;
108
      long int mid = (left+right)/2;
109
110
      MergeSortHelper(Array, left, mid);
      MergeSortHelper(Array, mid+1, right);
111
      {\it Merge}\,(\,{\it Array}\,\,,\,\,\,\,{\it left}\,\,,\,\,\,{\it mid}\,,\,\,\,\,{\it right}\,)\;;
112
113
114
   void MergeSort(long int Array[], long int number){
   //REQUIRES: Array initialized with at least 'number' elements
115
116
      //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31}-1 elements,
117
118
      //\text{number in } [0,2^{31}-1]
      //MODIFIES: Array
119
      //EFFECTS: sort the Array
120
121
      MergeSortHelper(Array, 0, number-1);
   }
122
```

```
125 //Quick Sorting
   static long int getpivot(long int left, long int right){ //REQUIRES: left and right in [0,2^{31}-1]
126
     //EFFECTS: return a number in the range [left, right]
     long int result = (rand() % (right - left) ) + left;
129
     return result; //here returns a random index in the range as the pivot index,
130
     //this can be modified for a better performance
131
   //Partition with extra place
134
   static long int partition_extra_place(long int Array[], long int left, long int right){
136
     //REQUIRES: Array initialized with with left and right as legal index
     //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
137
     //MODIFIES: Array
138
     //EFFECTS: rearrange the array to make it partitioned by a pivot
     long int partition_index = getpivot(left, right);
140
     long int * Auxiliary = new long int [right - left + 1];
141
     long int i,j,k;
142
     long int pivot=Array[partition_index];
143
144
     for (i=left, j=0, k=right-left; i <= right; i++){}
       if(i==partition_index) continue;
145
       else if(Array[i]<=pivot)</pre>
146
          Auxiliary [j++]=Array[i];
147
       else
148
149
          Auxiliary [k--]=Array [i];
150
     Auxiliary [j]=pivot;
     for (i=left; i \le right; i++){
       Array[i]=Auxiliary[i-left];
154
     delete [] Auxiliary;
     return left+j;
157
158
159
   //Partition in place
   static long int partition_in_place(long int Array[], long int left, long int right){
     //REQUIRES: Array initialized with with left and right as legal index
     //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
162
     //MODIFIES: Array
163
     //EFFECTS: rearrange the array to make it partitioned by a pivot
164
     long int pivot_index=getpivot(left, right);
165
     swap(Array, left, pivot_index);
     long int i=left;
167
168
     long int j=right;
     long int pivot=Array[left];
169
     while (i < j)
            while(i<j && pivot<=Array[j])</pre>
172
           while (i < j && Array [i] <= pivot)
173
174
           swap(Array, i, j);
176
177
     swap(Array,left,i);
178
     return i:
   }
179
180
   static void QuickSortHelper(long int Array[], long int left, long int right, long int (*
181
       partition)(long int[], long int, long int)) {
     //REQUIRES: Array initialized with left and right as legal index
182
     //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31} elements
183
     //MODIFIES: Array
184
     //EFFECTS: recursively quicksort the array
185
     long int pivotat; // index of the pivot
186
     if(left >= right) return;
187
188
     pivotat = partition(Array, left, right);
     QuickSortHelper(Array, left, pivotat-1, partition);
189
     QuickSortHelper\left(Array\,,\ pivotat+1,\ right\,,\ partition\,\right);
190
191
193
void QuickSort1(long int Array[], long int number){
```

```
//REQUIRES: Array initialized with at least 'number' elements
195
     //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31}-1 elements,
196
     ^{''}_{//\text{number in }[0,2^{\hat{}}\{31\}-1]}
197
     //MODIFIES: Array
198
     //EFFECTS: sort the Array
199
     QuickSortHelper(Array, 0, number-1, partition_extra_place);
200
201
202
   void QuickSort2(long int Array[], long int number){
203
     //REQUIRES: Array initialized with at least 'number' elements
204
     //with elements in [-2^{31}, 2^{31}-1] and no more than 2^{31}-1 elements,
205
     //\text{number in } [0,2^{31}-1]
     //MODIFIES: Array
207
     //EFFECTS: sort the Array
208
     QuickSortHelper(Array, 0, number-1, partition_in_place);
209
   }
210
211
   #endif
212
```

7 Appendix B. Testing different algorothms

compare.cpp

```
//Generate 10, 100, 1000, 10000, and 100000 random integers in the range [-2^{3}], 2^{3}
  //Sort them with different algorithms and measure the time cost.
  //The output is written to a file called "CompareResult.txt".
5 #include <fstream>
 #include <cstdlib>
 #include <ctime>
s #include <iostream>
 #include <string>
#include "sort.h
11 using namespace std;
const long int MAXSIZE=100000; //the maximum size of the arrays to be tested
                                //the minimum size of the arrays to be tested
  const long int MINSIZE=10;
  const long int STEP=10;
15
  const int METHODSNUMBER=6;
16
  const string METHOSNAMES[6]={"Bubble Sort", "Insertion Sort", "Selection Sort",
               'Merge Sort", "Quick Sort Not-in-place", "Quick Sort In-place"};
18
19
  int main(){
    ofstream output;
21
    output.open("CompareResult.txt");
    long int * integers= new long int [MAXSIZE];
23
    long int * auxiliary= new long int [MAXSIZE];
25
    clock_t start, end;
    26
              MergeSort, QuickSort1, QuickSort2};
27
28
    for(long int i=MINSIZE; i<=MAXSIZE; i*=STEP){</pre>
      for (long int j=0; j < i; j++)
29
        integers[j]=auxiliary[j]=mrand48();//save a copy in auxiliary
30
31
      for(int choice=0; choice<METHODSNUMBER; choice++){</pre>
32
        start = clock();
33
34
        sort[choice](integers, (long)i);
        end = clock();
35
        output << METHOSNAMES [choice] << " with array size of "<<i << " costs"
36
            <<(double) ((end-start) *1000.0/CLOCKS_PER_SEC)<<"ms."<<endl;
37
38
        for (long int j=0; j < i; j++){
        //write back the saved numbers in order
39
        //to make it in the same order for every algorithm
40
          integers [j] = auxiliary [j];
41
42
      }
43
44
    output.close();
```