# **DES Documentation**

## Made with Love, by Mousa Emarah

- 1. Key Generation Process:
- 1.1 Initial Permutation (PC1):
- The 64-bit key undergoes permutation using the PC1 table
- Results in a 56-bit key
- 1.2 Key Division:
- The 56-bit key is divided into two equal halves:
- Left half (C<sub>o</sub>): 28 bits
- Right half (D<sub>o</sub>): 28 bits
- 1.3 Subkey Creation:
- Sixteen subkeys are generated through 16 rounds
- Each round applies circular left shifts to both halves
- Shift amounts follow the predefined shift schedule
- Each shifted pair forms a 56-bit intermediate key

- 1.4 Final Subkey Permutation (PC2):
- Each 56-bit intermediate key is permuted using PC2 table
- Produces 16 final 48-bit subkeys (K<sub>1</sub> to K<sub>16</sub>)
- These subkeys are used in the Feistel network
- 2. Message Encryption Process:
- 2.1 Input Conversion:
- Plaintext message is converted to 64-bit binary representation
- Padding applied if message is not 64 bits
- 2.2 Initial Permutation (IP):
- 64-bit block undergoes permutation using IP table
- Reorders the bits according to IP table specification
- 2.3 Block Division:
- Permuted block is split into:
- Left half (L<sub>o</sub>): 32 bits
- Right half (R<sub>o</sub>): 32 bits

### 2.4 Feistel Rounds (16 iterations):

For each round n (1 to 16):

$$2.4.1 L_n = R_{n-1}$$

2.4.2 
$$R_n = L_{n-1} \oplus f(R_{n-1}, K_n)$$

#### Where:

- $\bigoplus = XOR$  operation
- f = Feistel function
- $K_n$  = Subkey for round n

## 2.5 Expansion (E):

- Right half  $(R_{n-1})$  expanded from 32 to 48 bits
- Uses expansion table E
- Allows combination with 48-bit subkey

#### 2.6 S-box Substitution:

- 48-bit result divided into eight 6-bit groups
- Each group processed through corresponding S-box (S1-S8)
- Each S-box outputs 4 bits
- Total output: 32 bits (8 × 4 bits)

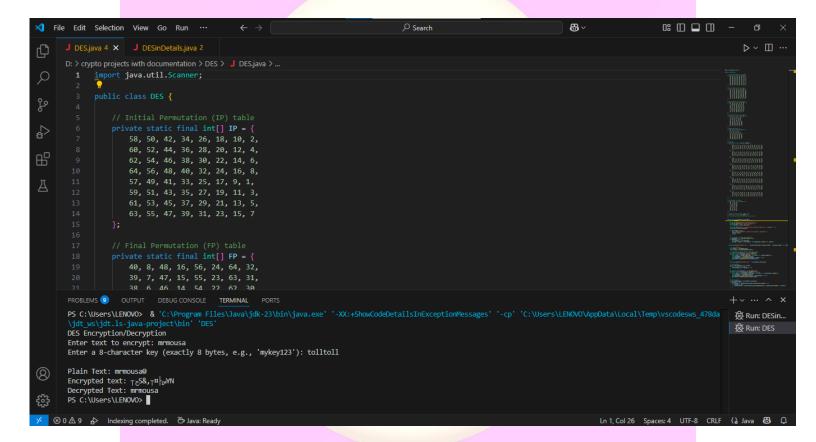
#### 2.7 P Permutation:

- 32-bit S-box output undergoes permutation
- Uses P table to rearrange bits
- Produces final output of Feistel function

## 2.8 XOR Operation:

- P-permuted result is XORed with L<sub>n-1</sub>
- Produces new right half R<sub>n</sub>
- Left half L<sub>n</sub> becomes previous R<sub>n-1</sub>
- 3. Final Steps:
- 3.1 After 16 rounds:
- Left and right halves are concatenated
- Note: No swap after final round
- 3.2 Final Permutation (FP):
- Combined block undergoes final permutation
- Uses inverse of initial permutation table
- Produces final 64-bit ciphertext

- 4. Decryption Process:
- 4.1 Same algorithm as encryption
- 4.2 Subkeys used in reverse order (K<sub>16</sub> to K<sub>1</sub>)
- 4.3 Final output is original plaintext



## <mark>Always remember:</mark>

#### 1) Key generation:

- 1- permutation using pc1 table.
- 2- divide each key into 2 each of 28 bits.
- 3- now we have 16 subkeys "56 bits"
- 4- therefore apply pc2 on contacanated subkeys to have 48 bits to work in fesitel network.

(These enter the function with RO then XORed with LO to get R1)

## 2) Message encryption:

- 1-change the message inputed by the user into binary
- 2- apply ip permutation
- 3-divide into 2 halves
- 4- Ln = Rn-1, Rn = Ln-1 + f(Rn-1, Kn), where n is the number of iteration
- 5- use expansion permutation no have 8 groups of 6 bits to have 48 bits with the 48 bits of the keys
- 6-Apply s boxes (now we have 8 groups of 4 bits)
- 7-Apply P permuation which gives us the ifnal value of f
- 8- xor with the Ln-1

Reference:

https://page.math.tu-berlin.de/``kant/teaching/hess/krypto-ws2006/des.htm

