

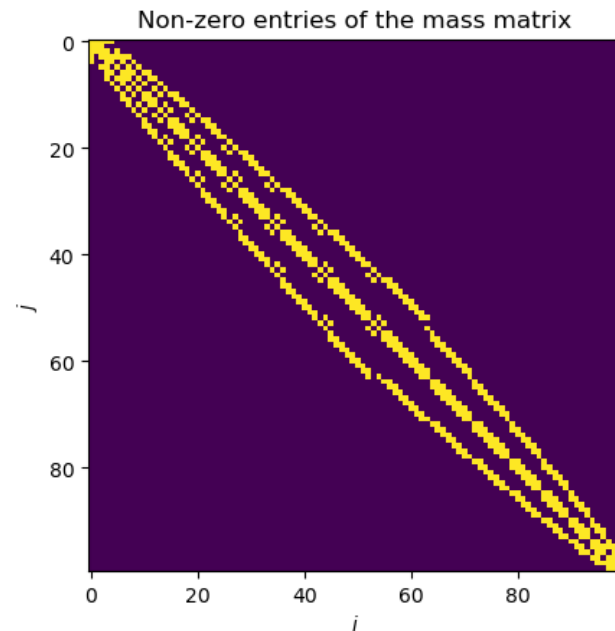
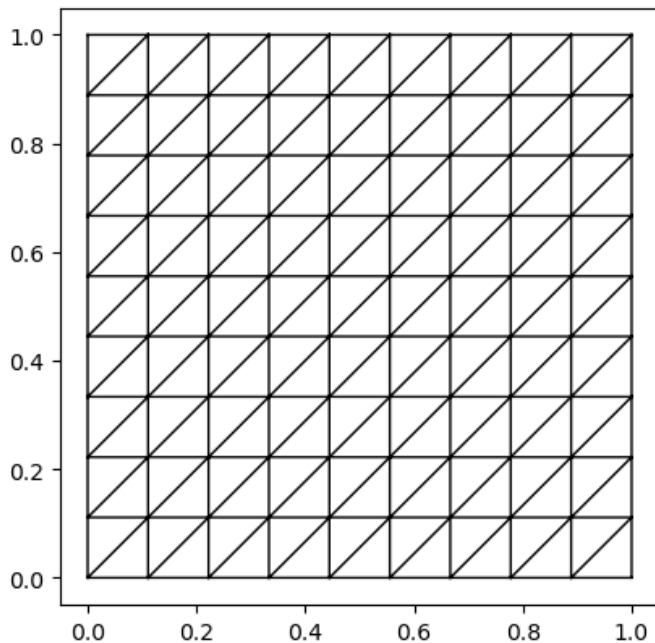
# Parallel assembly of mass matrix

PHYS-743 Parallel  
programming

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Novembre 20th 2025

- Compute the entries of the mass matrix  $M_{ij} = \int_{\Omega} \phi_i \phi_j$ .
- $V = \{v \in H^1(\Omega), v|_K \in \mathbb{P}^1(K), \forall K \in \mathcal{T}_h\}$ , where  $\mathcal{T}_h$  is a conformal mesh of  $\Omega$ .
- $V = \text{span}(\phi_1, \dots, \phi_M)$ ,  $M$  number of mesh vertices.



# Serial algorithm

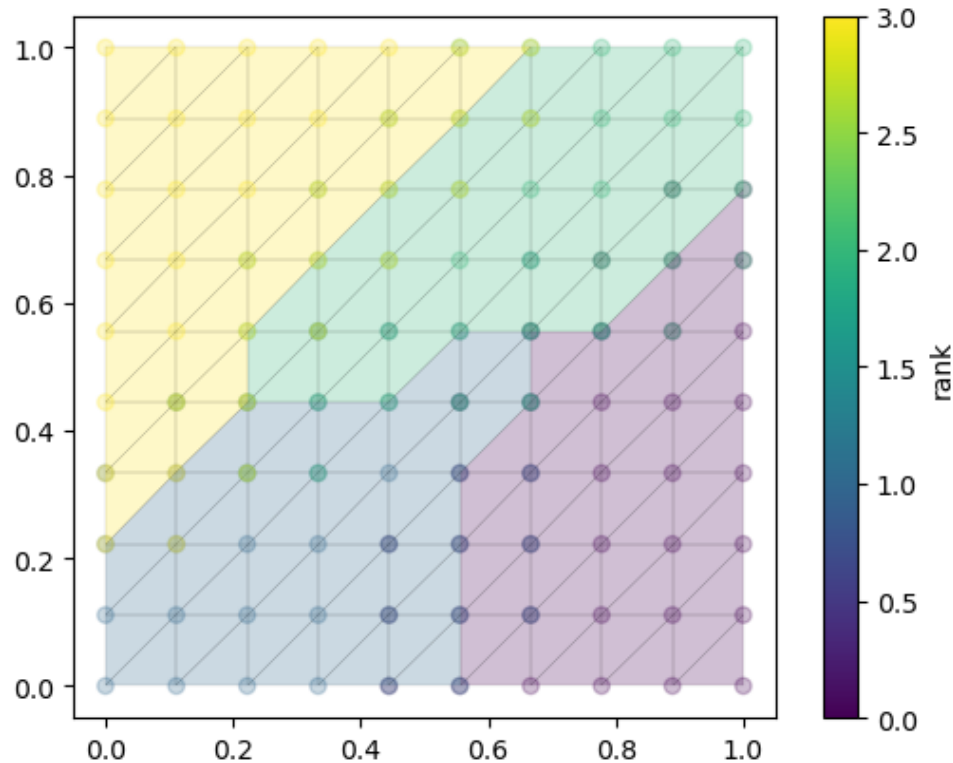
- Loop over the elements  $K$ .
- For each  $K$ , we have a mapping  $\left| \begin{array}{l} \{1,2,3\} \rightarrow \{a_1, a_2, \dots, a_M\} \\ i \mapsto \mathcal{N}_i \end{array} \right.$ .

For  $K \in \mathcal{T}_h$   
 For  $i = 1,2,3$   
 For  $j = 1,2,3$   
 $A_{\mathcal{N}_i \mathcal{N}_j} += \int_K \phi_{i,K} \phi_{j,K}$

- Complexity  $O(9N_{\text{elems}})$ .

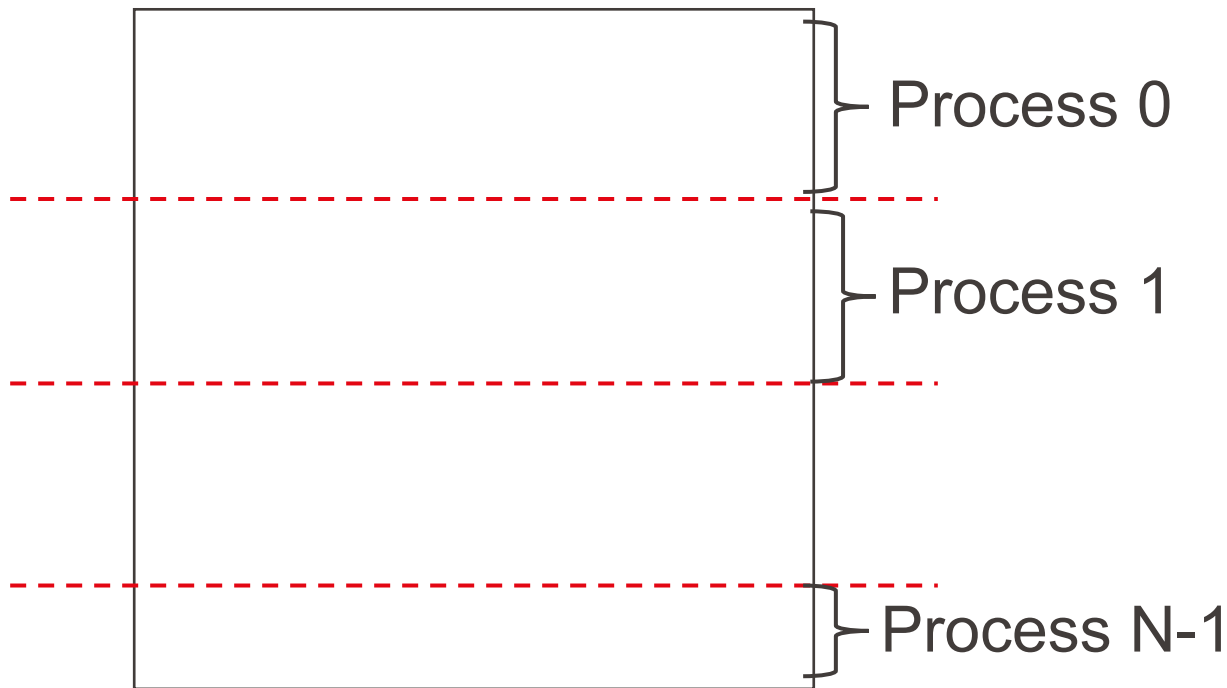
- Partition of the mesh.
- Compute mass matrix contribution on sub-mesh, using serial algorithm.
- Communicate the entries belonging to other processes.
- Build a part of the mass matrix on each process.

- This is done using dolfinx [1], which calls PT-Scotch [2], ParMETIS [3] or KaHIP [4].
- Process  $n \in [0: N - 1]$  owns its elements  $\mathcal{E}_n$ , its vertices  $\mathcal{N}_n$  and needs ghost vertices  $\tilde{\mathcal{N}}_n$  which are vertices on other sub-meshes.



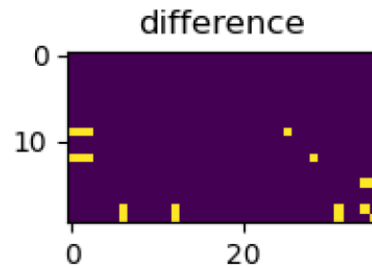
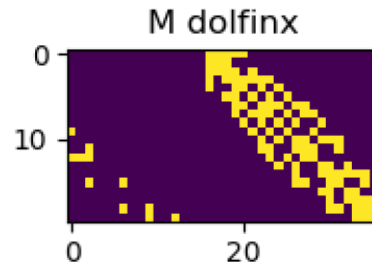
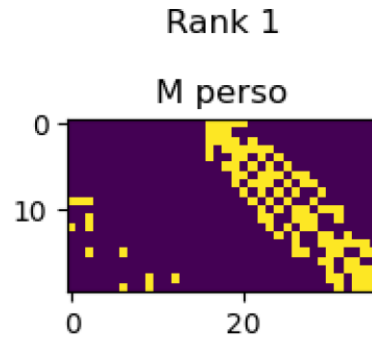
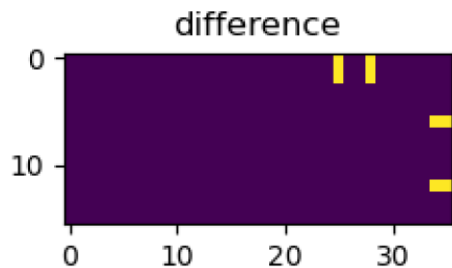
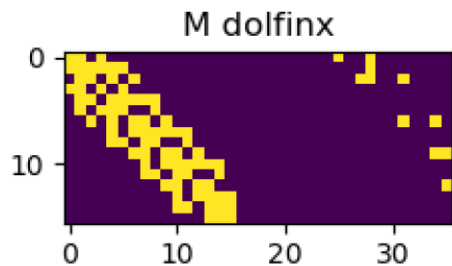
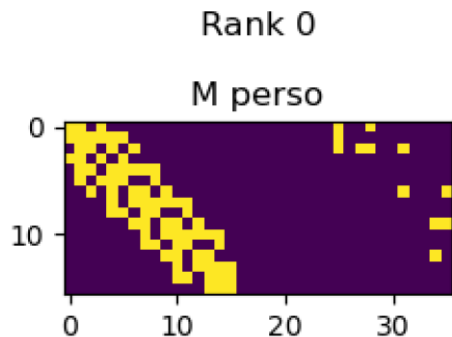
# Matrix partitioning

- Each process owns only a part of the mass matrix  $M_{ij}^n$ , where  $i \in R_n$  and  $j \in [1:M]$ .



- Mass matrix computation returns local indices  $I_{\text{local}}, J_{\text{local}}$  and values  $V$  on the sub-mesh.
- Some are ghosts : they are sent to other process using `MPI.COMM_WORLD.allgather`.
- The current process verify which of the node it recieved belongs to it.
- The new values are added to the mass matrix on process

# Issue with indexing... (comparison with dolfinx)





# Solution to run on Helvetios

- I did not manage to use fenicsx on Helvetios
- The solution was to partition the mesh on my computer, and save the partition on separate h5 files using h5py.
- The code on the cluster reads the h5 files to recover the partition of the mesh.

- $N_V$  number of vertices,  $N_c$  number of cells,  $N_V/n \sim N_{\text{loc}} + N_g$  number of vertices on current process.
- Reading mesh :
 

<code>with h5py.File(filename, 'r') as f:</code>	$O(N_V/n)$
<code>nodes = f['nodes'][:]</code>	$O(N_c/n)$
<code>cells = f['cells'][:]</code>	$O(N_V/n)$
<code>local_to_global = f['local_to_global'][:]</code>	$O(N_V)!$
<code>global_to_local = f['global_to_local'][:]</code>	
<code>size_local = f['size_local'][(0)]</code>	
- Computing mass matrix :  $O(N_c/n)$
- Find ghosts, etc... : at most  $O(N_c/n)$
- All gather :  $O(N_g \times n)!$
- Find accepted vertices :  $O(\max(N_{\text{loc}}, N_g \times n))$
- Conclusion  $p \sim 0.95$  (experimentally...)

n	1	2	4	8	16	32	64
N triangles	~2.1M	~2.1M	~2.1M	~2.1M	~2.1M	~2.1M	~2.1M
T	2.095 s	1.087 s	0.597 s	0.366 s	0.270 s	0.169 s	0.118 s
S	1	1.93	3.51	5.72	7.56	12.37	17.75
E	1	0.96	0.88	0.72	0.48	0.39	0.28

n	1	2	4	8	16	32	64
Load mesh	2.57%	1.60%	4.55%	5.29%	9.93%	18.67%	31.29%
Compute matrix	54.40%	54.90%	52.13%	44.80%	33.90%	25.66%	20.68%
Communicate	43.03%	43.50%	43.32%	49.91%	56.16%	55.67%	48.03%

n	1	2	4	8	16	32	64
N triangles	~2.1M	~4.2M	~8.4M	~16.8M	~33.6 M	~67.1M	~134.2M
T	2.095 s	2.264 s	2.298 s	2.504 s	2.570 s	4.057 s	6.731 s
S	1	1.85	3.65	6.70	13.04	16.52	19.92
E	1	0.93	0.91	0.84	0.82	0.52	0.31

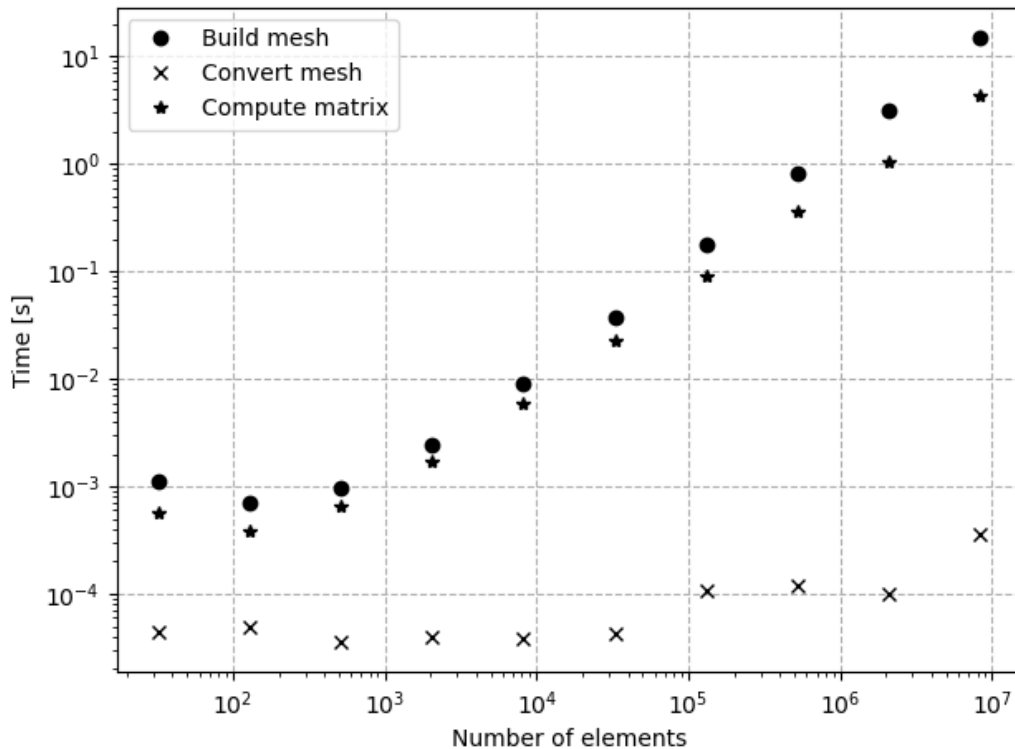
n	1	2	4	8	16	32	64
Load mesh	2.57%	2.39%	1.92%	3.73%	7.16%	8.05%	34.63%
Compute matrix	54.40%	52.32%	50.78%	51.94%	46.20%	37.68%	25.84%
Communicate	43.03%	45.29%	47.30%	44.32%	46.63%	54.26%	39.54%

- Implemented mass matrix computation in parallel with MPI.
- Helped to understand better how fenicsx works.

- [1] <https://github.com/FEniCS/dolfinx>
- [2] C. Chevalier, F. Pellegrini, PT-Scotch: A tool for efficient parallel graph ordering, *Parallel Computing*, Volume 34, Issues 6–8, 2008, Pages 318–331, ISSN 0167-8191, <https://doi.org/10.1016/j.parco.2007.12.001>
- [3] George Karypis and Vipin Kumar. 1996. Parallel multilevel k-way partitioning scheme for irregular graphs. In *Proceedings of the 1996 ACM/IEEE conference on Supercomputing (Supercomputing '96)*. IEEE Computer Society, USA, 35–es. <https://doi.org/10.1145/369028.369103>
- [4] Sanders, P., Schulz, C. (2013). Think Locally, Act Globally: Highly Balanced Graph Partitioning. In: Bonifaci, V., Demetrescu, C., Marchetti-Spaccamela, A. (eds) *Experimental Algorithms. SEA 2013. Lecture Notes in Computer Science*, vol 7933. Springer, Berlin, Heidelberg. [https://doi.org/10.1007/978-3-642-38527-8\\_16](https://doi.org/10.1007/978-3-642-38527-8_16)

# Appendix – Profile in serial

- In serial, it takes  $O(N_{\text{elems}})$  operations for dolfinx to create a mesh



# Find $p$ experimentally

- From Amdahl's law :  $S = \frac{1}{1-p+p/n} \Rightarrow p = \frac{1-1/S}{1-1/n}$ .
- For the strong scaling, we observe

$$p = [-, 96\%, 95\%, 94\%, 93\%, 95\%, 96\%]$$