

# DISCUSSION 04/12

For definitions, refer Sipser.

$$\Sigma = \{a, b\}$$

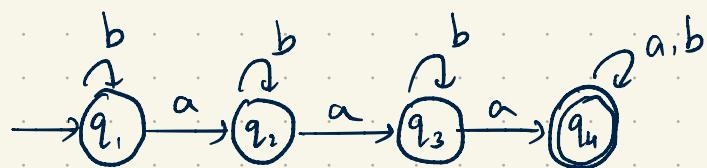
1.4 (a)  $\{w \mid w \text{ has at least 3 'a's \& at least 2 'b's}\}$

Sol: We want the intersection of 2 DFAs  $M_1$  &  $M_2$

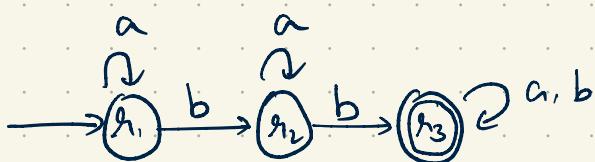
$M_1$  recognizes  $\{w \mid w \text{ has at least 3 'a's}\}$

$M_2$  recognizes  $\{w \mid w \text{ has at least 2 'b's}\}$

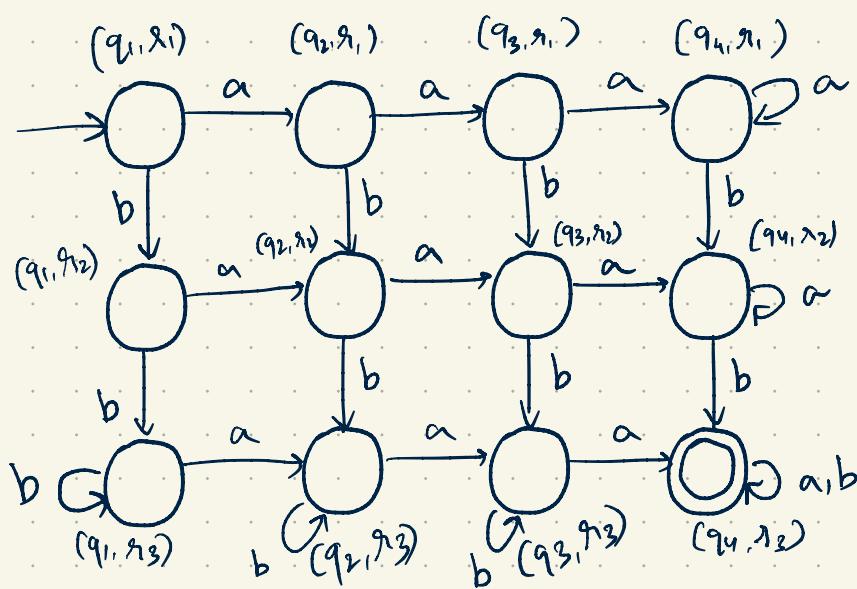
$M_1$



$M_2$



$M_1 \cap M_2$

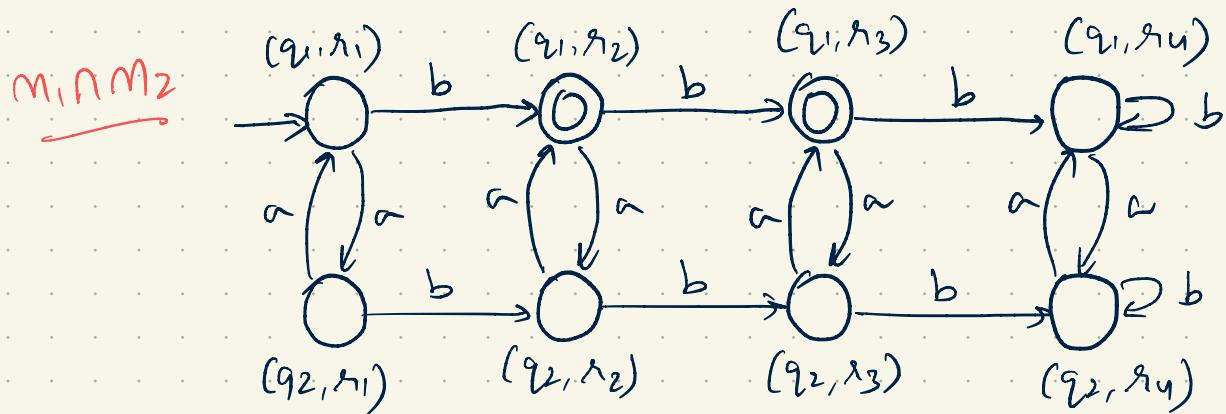
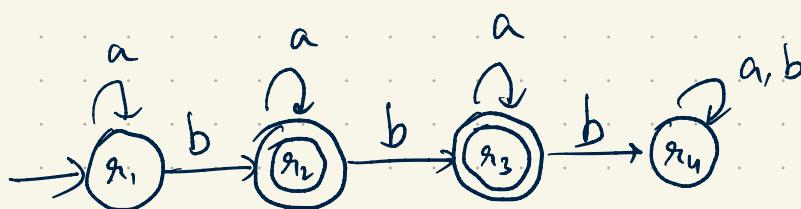
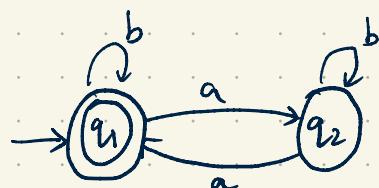


1.4 (c)  $\{w \mid w \text{ has an even number of 'a's and 1 or 2 'b's}\}$

Sol: Want  $M_1 \cap M_2$ , where

$M_1$  recognizes  $\{w \mid w \text{ has an even number of 'a's}\}$

$M_2$  recognizes  $\{w \mid w \text{ has 1 or 2 'b's}\}$



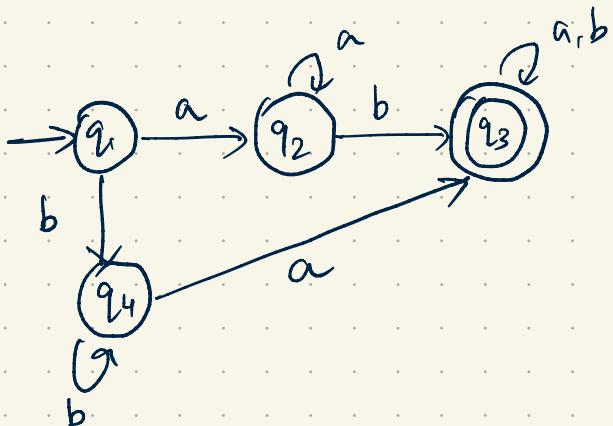
1.5 (c) Construct a DFA that recognizes

$L = \{w \mid w \text{ contains neither } ab \text{ nor } ba\}$

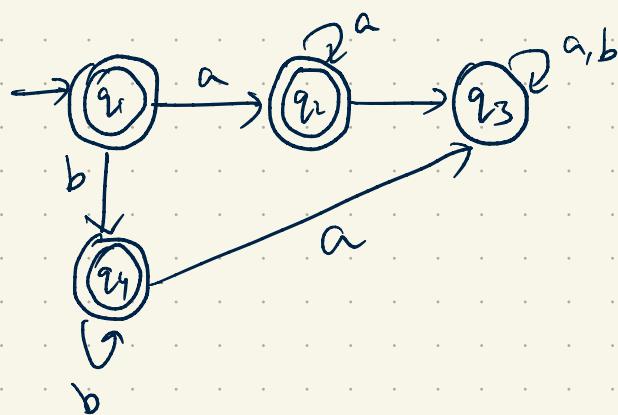
The complement of this language is

$\bar{L} = \{w \mid w \text{ contains } ab \text{ or } ba\}$

DFA that recognizes  $\bar{L}$ :



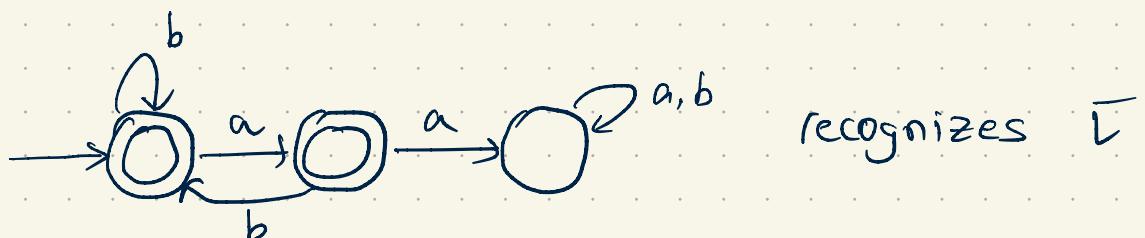
Complement this to get the DFA that recognizes  $L$ :



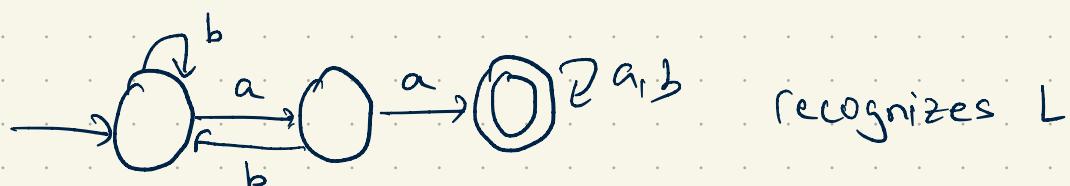
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1.5(e)  $\{w \mid w \text{ is any string NOT in } (ab^+)^*\} = \bar{L}$

$$\bar{L} = \{w \mid w \text{ is in } (ab^+)^*\}$$



Complement this:



1.6 (i) Construct a DFA that recognizes  
 $\{ w \mid \text{every odd position of } w \text{ is a } 1 \}$   
over  $\Sigma = \{0, 1\}$

Sol Important :  $\epsilon$  (empty string) BELONGS to this language.

