

Theory of Algorithms

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Scheme

Turing machines

Computational complexity

NP completeness

Scheme



- John McCarthy while at MIT – late 1950s.
- Created Lisp.
- Lisp generally considered first functional programming language (not really though).
- Lots of dialects exist today, such as Scheme and Common Lisp.

Wasteful for loops

How do you parallelise this?

```
inds = list(range(10))
```

```
total = 0
```

```
for i in inds:
```

```
    total = total + (i * 7)
```

```
def byseven(i):
```

```
    return i * 7
```

```
inds = list(range(10))
```

```
sevens = map(byseven, inds)
```

```
total = sum(sevens)
```

State

Imperative programming is a programming paradigm where statements are used to change the *state*.

State is the name given to the current data/values related to an executing process, including internal stuff like the call stack.

Processes begin with an initial state and (possibly) have (a number of) halt states.

Statements change the state.

Functions in imperative programming languages might return different values for the same input at different times, because of the state.

Functional programming languages (try to) not depend on state.

Functions are said to have side effects if they modify the state (on top of returning a value).

Static and global variables are often good examples of side effects in action.

Functional programming tries to avoid side effects.

It's tricky to avoid them – such as when we need user input.

Basic operators

```
> (+ 3 4)
```

```
7
```

```
> (* 3 2)
```

```
6
```

```
> (- 5 3)
```

```
2
```

```
> (/ 6 3)
```

```
2
```


More arguments

```
> (+ 3 4 5)
```

```
(+ 3 4 5)
```

```
> (- 3 4 5)
```

```
-6
```

```
> (* 2 3 4)
```

```
24
```

```
> (/ 6 3 3)
```

```
2/3
```

```
> (/ 6 3 3 3)
```

```
2/9
```

Functions and values

; Define a value called foo with value 3.

```
>(define foo 3)
```

; Define a function f.

```
>(define (f x)
  (+ (* 3 x) 12))
```

```
>(define (g x)
  (* 3 (+ x 4)))
```

```
>(g 2)
18
```

Conditionals

```
> (if (< 1 2) '(y e s) '(n o))  
(y e s)
```

```
>(define (abs x)  
  (if (< x 0)  
      (- x)  
      x))
```

Lists

```
>(list 1 2 3)  
(1 2 3)
```

```
>(list 'a 'b 'c)  
(a b c)
```

```
> (length (list 1 2 3))  
3
```

car and cdr

```
> (car (list 1 2 3))  
1  
> (cdr (list 1 2 3))  
(2 3)  
> (define l (list 1 2 3))  
> (car l)  
1  
> (cdr l)  
(2 3)  
> (car (cdr l))  
2  
> (cadr l)  
2
```

Recursion

```
> (define (sum lv)
  (if (null? lv)
      0
      (+ (car lv) (sum (cdr lv)))))
> (sum (list 1 2 3))
6
> (define (derange n)
  (if (= 0 n)
      '()
      (cons n (derange (- n 1)))))
> (derange 12)
(12 11 10 9 8 7 6 5 4 3 2 1)
```

Looping recursively

```
> (let loop ((i 5))  
  (print "i is " i ".\n")  
  (if (> i 0) (loop (- i 1))))  
i is 5.  
i is 4.  
i is 3.  
i is 2.  
i is 1.  
i is 0.
```

```
> (define (swap3-1-2 x)
  (list (cadr x) (car x) (caddr x)))
```

```
> (swap3-1-2 (list 1 2 3))
(2 1 3)
```

```
> (define four-over-two (list 4 '/ 2))
```

```
> four-over-two
(4 / 2)
```

```
> (eval (swap3-1-2 four-over-two))
```


More on functions

; Printing stuff to terminal.

```
> (print "Ay" "-" "yo.\n")
```

; Proper way to define a function.

```
> (define foo (lambda (bar) (print "Bar is " bar ".\n")))
```

; Shorthand.

```
> (define (foo bar) (print "Bar is " bar ".\n"))
```

; Local variables.

```
> (define (foo bar) (let ((thing "Bar"))
```

```
  (print thing " is " bar ".\n")))
```

```
> (foo "open")
```

```
Bar is open.
```

Function example

```
> (define l
  (let
    ((d 4) (e 5))
    (lambda (a b c) (list a b c d e))
  )
)
> (1 1 2 3)
(1 2 3 4 5)
```

```
> (cons 1 '())  
(1)  
> (cons 1 (cons 2 null))  
(1 2)  
> (cons 1 (cons 2 (cons 3 null)))  
(1 2 3)  
> (define mylist (cons 1 (cons 2 (cons 3 null))))  
> mylist  
(1 2 3)  
> (car mylist)  
1  
> (cdr mylist)  
(2 3)
```

More on lists

```
> (list "a" "b" "c")
("a" "b" "c")
> (list a b c)
reference to undefined identifier: a
> (list 'a 'b 'c)
(a b c)

> (equal?
  (list 1 2 3)
  (cons 1 (cons 2 (cons 3 '()))))
#t
```

Quoting

```
> (list a b c)
*** ERROR IN (console)@1.7--Unbound variable: a
> (quote (a b c))
(a b c)
> (quote a b c)
*** ERROR IN (console)@2.1--Ill-formed special form: quote
> '(a b c)
(a b c)
> (define forty-two '(* 6 9))
> forty-two
(* 6 9)
> (eval forty-two)
```

54

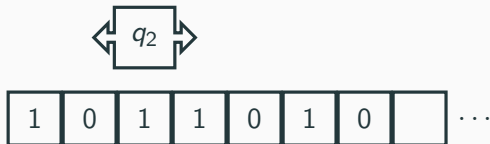
Null list

```
> ()  
missing procedure expression  
> (list)  
( )  
> '()  
( )  
> null  
( )  
> 'null  
null
```

Closures

```
> (define (container value)
  (lambda ()
    (string-append "This container contains " value ".")))
> (define apple (container "an apple"))
> (define pie (container "a pie"))
> (apple)
"This container contains an apple."
> (apple)
"This container contains an apple."
> (pie)
"This container contains a pie."
```

Turing machines



Notation

Q Set of states (finite).

G Tape alphabet (finite).

B Blank symbol, element of G .

S Input alphabet, subset of $G \setminus \{B\}$.

δ Transition function.

q_0 Initial state, $\in Q$.

q_a Accept state, $\in Q$.

q_r Reject state, $\in Q$.

M Turing Machine: $[Q, G, B, S, \delta, q_0, q_a, q_r]$.

State Table

State	Input	Write	Move	Next
0	B	B		Accept
	0	0	L	0
	1	1	L	1
1	B	B		Fail
	0	0	L	1
	1	1	L	0

$$\delta(q_i, g_n) \rightarrow (q_j, g_m, L/R)$$

Turing's Second Example

A slightly more difficult example

We can construct a machine to compute the sequence

$$00101101110111101111\dots$$

The machine is to be capable of five m -configurations, viz. o , q , p , f , b and of printing a , x , 0 , 1 . The first three symbols on the tape will be aaO ; the other figures follow on alternate squares. On the intermediate squares we never print anything but x . These letters serve to keep the place for us and are erased when we have finished with them. We also arrange that in the sequence of figures on alternate squares there shall be no blanks.

Turing's Second Example: Table

<i>Configuration</i>		<i>Behaviour</i>	
<i>m-config.</i>	<i>symbol</i>	<i>operations</i>	<i>final m-config.</i>
b		$P\emptyset, R, P\emptyset, R, P0, R, R, P0, L, L$	o
o	{ 1	R, Px, L, L, L	o
	{ 0		q
q	{ Any (0 or 1)	R, R	q
	{ None	$P1, L$	p
p	{ x	E, R	q
	{ \emptyset	R	f
	{ None	L, L	p
f	{ Any	R, R	f
	{ None	$P0, L, L$	o

Turing's Second Example: JavaScript 1

```
// The contents of the tape.  
var tape = []  
s// The current position of the machine on the tape.  
var pos = 0  
// The current state;  
var state = b;
```

Turing's Second Example: JavaScript 2

// Writes a symbol to the current cell on the tape.

```
function write(sym) {  
    tape[pos] = sym;  
}
```

// Returns true iff the current cell contains sym.

```
function read(sym) {  
    return sym == tape[pos] ? true : false;  
}
```

Turing's Second Example: JavaScript 3

// Erases the symbol in the current cell of the tape.

```
function erase() {  
    delete tape[pos];  
}
```

// Returns true iff the current cell is blank.

// Returns true iff the current cell is blank.

```
function blank() {  
    return typeof(tape[pos])  
        == 'undefined' ? true : false;  
}
```


Turing's Second Example: JavaScript 4

```
function b() {  
  write('e');  
  pos++;  
  write('e');  
  pos++;  
  write('0');  
  pos++;  
  pos++;  
  write('0');  
  pos--;  
  pos--;  
  state = o;  
}
```

Computational complexity

Alphabet Finite set of symbols, denoted Σ .

Word Sequence of symbols, w from Σ .

Language Set of words, denoted L .

Length Of a word, denoted $|w|$.

Empty word Unique word of length 0, denoted λ .

Words $w_1 w_2$ is the concatenation of words w_1 and w_2 .

Languages $L_1 L_2$ is the language resulting from the concatenation of all words in L_1 and all words in L_2 , in that order.

Powers $L^0 = \{\lambda\}$, $L^1 = L$ and $L^{n+1} = L^n L$ for all $n > 1$.

Kleene Star

$$L^* = \bigcup_{i=0}^{\infty} L^i$$

Note that treating the alphabet Σ as a language in itself, we get that Σ^* is the set of all words over Σ .

Example

$$\Sigma \{0, 1\}$$

$$L \{00, 01, 10, 11\}$$

$$w_1 \ 01$$

$$w_3 \ 11$$

$$w_1 w_3 \ 0111$$

$$\Sigma^* \{\lambda, 0, 1, 00, 01, 10, 11, 001, 010, \dots\}$$

$$L^* \{\lambda, 00, 01, 10, 11, 0000, 0001, \dots\}$$

$$L^+ \{00, 01, 10, 11, 0000, 0001, \dots\}$$

How many numbers can we represent with k bits?

- Consider the case of four bits.
- Imagine four placeholders:

?	?	?	?
---	---	---	---
- Each placeholder can contain either 0 or 1.
- There are $2 \times 2 \times 2 \times 2 = 2^4 = 16$ different numbers.
- Add another bit, how many numbers is it now?
- It's $2 \times 2 \times 2 \times 2 \times 2 = 2^5 = 32$.
- Generally k bits can represent 2^k numbers.
- This grows **exponentially** relative to k .

Definition

An algorithm is said to be solvable in *polynomial time* if the number of steps required to complete the algorithm for a given input is $O(n^k)$ for some nonnegative integer k , where n is the complexity of the input.

P complexity class

The P complexity class is the set of problems for which there exists, for each such problem, at least one algorithm to solve that problem in polynomial time.

Non-deterministic polynomial time

Definition

A problem is in the NP complexity class if it is solvable by a non-deterministic Turing Machine in polynomial time. A non-deterministic Turing Machine is one which may not have a unique action to take for some or all states and inputs.

P is a subset of NP

The P complexity class is a subset of P because all polynomial time solvable problems can be modelled using Nondeterministic Turing Machines.

Decision problems

Decision problems are problems where the answer is 0 or 1.

Restricting ourselves to decision problems is convenient and fair.

$f : \{0, 1\}^n \rightarrow \{0, 1\}$ is useful notation for considering decision problems.

Other problems can be easily (polynomial time) adapted into decision problems.

NP completeness

Definition

An problem is *NP*-hard if each problem in *NP* can be reduced to it in polynomial time.

NP-hard problems are hard

NP-hard problems are at least as hard to solve as the hardest *NP* problems. Note that *NP*-hard problems don't have to be *NP*.

Definition

An problem is *NP*-complete if it's in *NP* and in *NP*-hard.

Subset sum problem

Problem

Given a set of integers S , is there a non-empty subset whose elements sum to zero?

Example

Does $\{1, 3, 7, -5, -13, 2, 9, -8\}$ have such a subset?

Note

If somebody suggests a solution, it is very quick to check it. Being able to quickly verify a solution is a characteristic of *NP* problems.

Propositional logic

Literals are Boolean variables (can be True or False), and their negations. They are represented by lower case letters like a and x_i .

Clause are expressions based on literals, that evaluate as True or False based on the literals. We use not, and and or on the literals. We'll sometimes call them expressions.

Not is depicted by \neg . So “not a ” is denoted by $\neg a$.

Or is depicted by \vee . So “ a or b ” is denoted by $a \vee b$.

And is depicted by \wedge . So “ a and b ” is denoted by $a \wedge b$.

CNF A clause is in Conjunctive Normal Form if it is a
“conjunction of disjunctions”: $(a \vee b) \wedge (\neg a \vee c) \wedge d$.

DNF A clause is in Disjunctive Normal Form if it is a
“disjunction of conjunctions”: $(a \wedge b) \vee (\neg a \wedge c) \vee d$.

Converting to CNF and DNF

Every Boolean expression can be converted to CNF, and every Boolean expression can also be converted to DNF.

Four laws

The following four laws can be used to convert expressions to CNF and DNF. The first two are known as De Morgan's laws, and the latter two are called the distributivity laws.

Conversion laws

$$\neg(a \vee b) = \neg a \wedge \neg b$$

$$\neg(a \wedge b) = \neg a \vee \neg b$$

$$c \wedge (a \vee b) = (c \wedge a) \vee (c \wedge b)$$

$$c \vee (a \wedge b) = (c \vee a) \wedge (c \vee b)$$

Boolean Satisfiability Problem (SAT)

The problem is the prototypical NP-complete problem. Note that it's quick to check the correctness of a solution for a given expression.

Only using CNF

All Boolean expressions can be converted in CNF. However, it's not always possible to do that in polynomial time. We can though, in polynomial time, create new expressions using some extra(neous) variables that are satisfiable if and only if the original expression is.

***k*-SAT**

k-SAT is like SAT except that all expressions must be in CNF and each clause must be a disjunction of *k* literals.

2-SAT $(a \vee b) \wedge (\neg c \vee d) \wedge \dots$

3-SAT $(a \vee b \vee c) \wedge (\neg c \vee d \vee \neg a) \wedge \dots$

2-SAT is not NP-complete. There are polynomial time algorithms that solve it.

3-SAT is NP complete.

3-SAT is NP-Complete

Reduction

3-SAT is a special case of SAT, so 3-SAT must be in NP. We can reduce SAT to 3-SAT in polynomial time. First take the expression and convert it to a CNF expression (in polynomial time). Then we just need to convert each clause into a CNF expression with 3 literals per clause.

Suppose we have a clause with 1 literal: a . Convert this to $(a \vee u_1 \vee u_2) \wedge (a \vee u_1 \vee \neg u_2) \wedge (a \vee \neg u_1 \vee u_2) \wedge (a \vee \neg u_1 \vee \neg u_2)$.
Suppose we have a clause with 2 literals: $a \vee b$. Convert this to $(a \vee b \vee u_1) \wedge (a \vee b \vee \neg u_1)$.

Now suppose we have a clause with n literals: $a \vee b \vee c \vee \dots$
Convert this to

$(a \vee b \vee u_1) \wedge (c \vee \neg u_1 \vee u_2) \wedge \dots \wedge (i \vee \neg u_{n-4} \vee u_{n-3}) \wedge (j \vee k \vee \neg u_{n-3})$.