Chapter 1

Oblig 1

Course “Compiler Construction”

Martin Steffen

Spring 2018

Section

Compila 18

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Oblig 1

- material (also for oblig 2) based on previous years, including contributions from Eyvind W. Axelsen, Henning Berg, Fredrik Sørensen, and others
- see also the course web-page, containing links to "resources"

Goal (of oblig 1)

Parsing

Determine if programs written in *Compila 18* are syntactically correct:

- scanner
- parser

- first part of a compiler, oblig 2 will add to it
- language spec provided separatly

Learning outcomes

- using tools for parser/scanner generation
  - JFlex
  - CUP
- variants of a grammar for the same languages
  - transforming one form (EBNF) to another (compatible with the used tools)
  - controlling precedence and associativity
- designing and implementing an AST data structure
  - using the parsing tools to build such trees
  - pretty-printing such trees

Compila language at a glance

```
program MyProgram
begin
    struct complex {      // record data type, but
      var    re: float;   // no subtyping, polymorphism...
      var    im: float;
    }
end;

proc add (a: complex, b: complex) : complex
begin
    var retval : complex;
    retval :=  new complex;
    retval.re := a.re + b.re;
    retval.im := a.im + b.im;
    return retval;
end;
proc main()               // exectution start here
begin
   var c1: complex;
   var c2: complex:
   var result := add (c1,c2);
   ...
   return;
end;
end;
```

Another glance

```
proc swap (a: ref(int), b: ref(int))  // passed a reference
begin
    var tmp: int;
    tmp := deref(a);                  // dereferencing
    deref(a) := deref(b);             // deref can be used both
    deref(b) := tmp;                  // left and right of
                                      //  an assignment.
end;
```

Grammar (1): declarations

```
PROGRAM          -> "program" NAME "begin" { DECL ";" }
"end" ";"

DECL             -> VAR_DECL | PROC_DECL | REC_DECL

VAR_DECL         -> "var" NAME ":" TYPE

PROC_DECL        -> "proc" NAME
                    "(" [ PARAM_DECL { "," PARAM_DECL } ] ")"
                    [ ":" TYPE ]
                    "begin" { DECL ";" } { STMT ";" } "end"

REC_DECL         -> "struct" NAME "{" { VAR_DECL ";" } "}"

PARAM_DECL       -> NAME ":" TYPE
```

Grammar (2): declarations

```
EXP          -> EXP LOG_OP EXP
              | "not" EXP
              | EXP REL_OP EXP
              | EXP ARIT_OP EXP
              | "(" EXP ")"
              | LITERAL
              | CALL_STMT
              | "new" NAME
              | VAR
              | REF_VAR
              | DEREF_VAR

REF_VAR      -> "ref" "(" VAR ")"
DEREF_VAR    -> "deref" "(" VAR ")" | "deref" "(" DEREF_VAR ")"

VAR          -> NAME | EXP "." NAME

LOG_OP       -> "&&" | "||"

REL_OP       -> "<" | "<=" | ">" | ">=" | "=" | "<>"

ARIT_OP      -> "+" | "-" | "*" | "/" | "^"

LITERAL      -> FLOAT_LITERAL | INT_LITERAL | STRING_LITERAL
              | "true" | "false" | "null"
```

Grammar (3): statements and types

```
STMT            -> ASSIGN_STMT
                 | IF_STMT
                 | WHILE_STMT
                 | RETURN_STMT
                 | CALL_STMT

ASSIGN_STMT     -> VAR ":=" EXP | DEREF_VAR ":=" EXP

IF_STMT         -> "if" EXP "then" "begin" { STMT ";" } "end"
                   [ "else" "begin" { STMT ";" } "end" ]
WHILE_STMT      -> "while" EXP "do" "begin" { STMT ";" } "end"

RETURN_STMT     -> "return" [ EXP ]

CALL_STMT       -> NAME "(" [ EXP { "," EXP } ] ")"

TYPE            -> "float" | "int" | "string" | "bool" | NAME
                 | "ref" "(" TYPE ")"
```

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Tools

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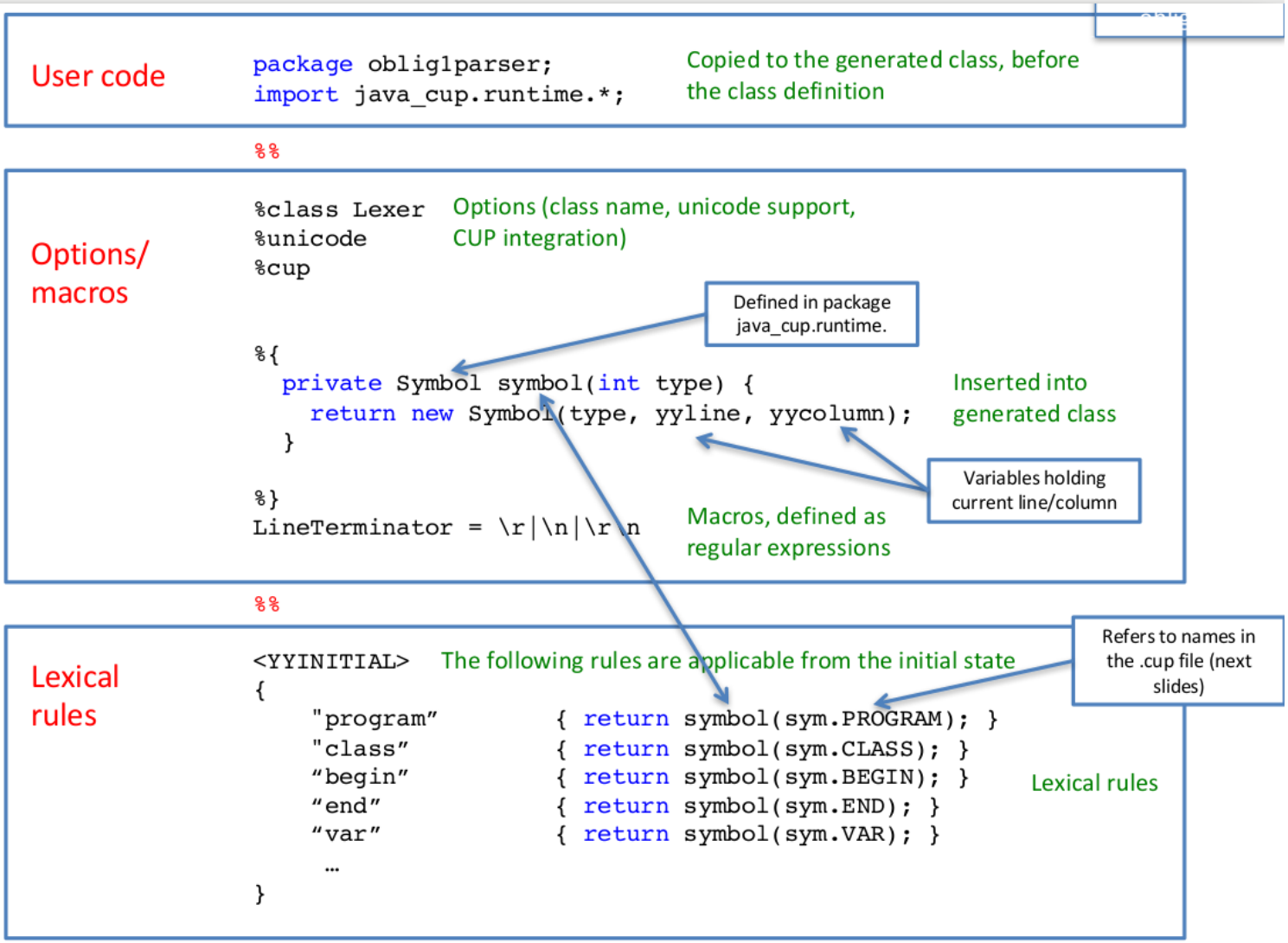
Tools: JFlex

- scanner generator (or lexer generator) tool
  - input: lexical specification
  - output: scanner program in Java
- lexical spec written as `.lex` file
- consists of 3 parts
  - user code
  - options and macros
  - lexical rules

Sample lex code



CUP: Construction of useful parsers (for Java)

- a tool to easily (ymmv) generate *parsers*
- reads tokes from the scanner using `next_token()`
- the `%cup` option (previous slide) makes that work

Input

grammar in BNF with action code

```
var_decl ::= VAR ID:name COLON type:vtype
 {: RESULT =  new VarDecl(name, vtype); :};
```

- output: parser program (in Java)

Sample CUP code

Package/ imports

```
package oblig1parser;
import java_cup.runtime.*;
import syntaxtree.*;
```

Package name for generated code and imports of packages we need

The syntaxtree package contains our own AST classes

User code

```
parser code {: :};
```

Code between {: and :} is inserted directly into the generated class (parser.java)

Symbol list

```
terminal            PROGRAM, CLASS;
terminal            BEGIN, END;
…
terminal            String      ID;
terminal            String      STRING_LITERAL;

non terminal        Program           program;
non terminal        List<ClassDecl>   decl_list;
non terminal        ClassDecl         class_decl, decl;
```

Terminals and non-terminals are defined here. They can also be given a Java type for the "value" that they carry, e.g. a node in the AST

Precedence

```
precedence left     AND;
```

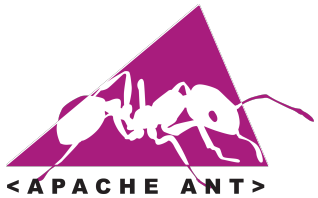
Precedence declarations are listed in ascending order, last = highest

Grammar

```
program     := PROGRAM BEGIN decl_list:dl END SEMI {: RESULT = new Program(dl); :} ;

decl_list   ::= decl:d
      {: List<ClassDecl> l = new LinkedList<ClassDecl>(); l.add(d); RESULT = l; :} ;

decl        ::= class_decl:sd {: RESULT = sd; :}    ;

class_decl  ::= CLASS ID:name BEGIN END
                    {: RESULT = new ClassDecl(name); :}   ;
```

AST is built during parsing. The left hand side of each production is implicitly labeled RESULT.

Build tool: ant



- Java-based build tool (think “make”)
- config in `build.xml`
- can contain different targets

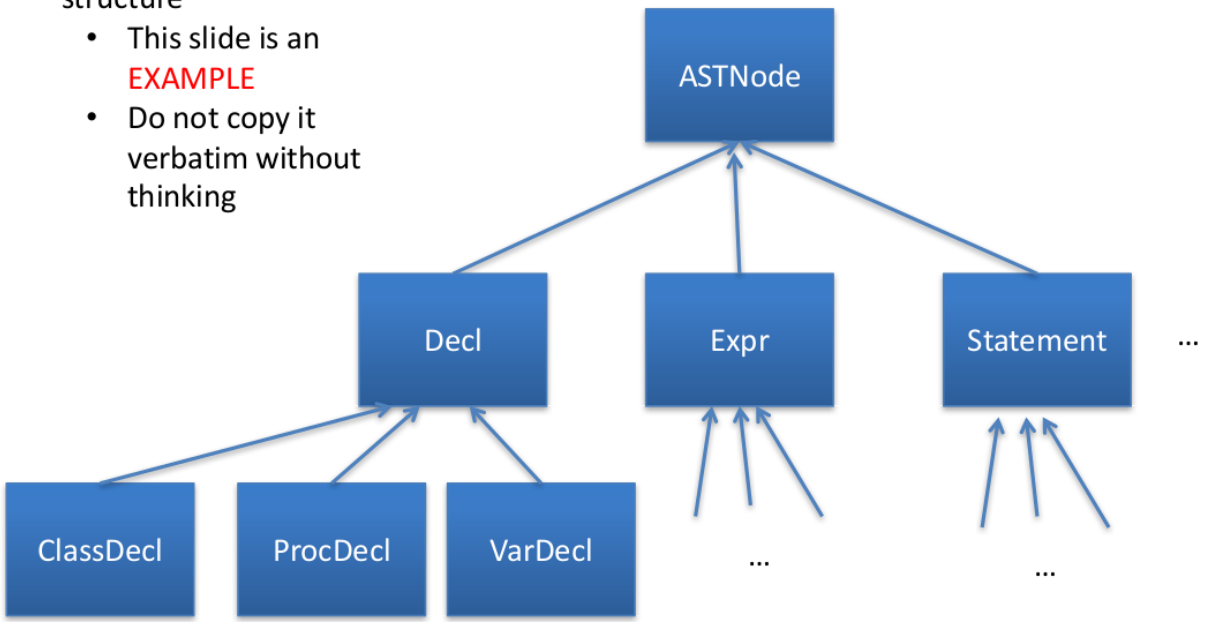
typical general targets

- test
- clean
- build
- run

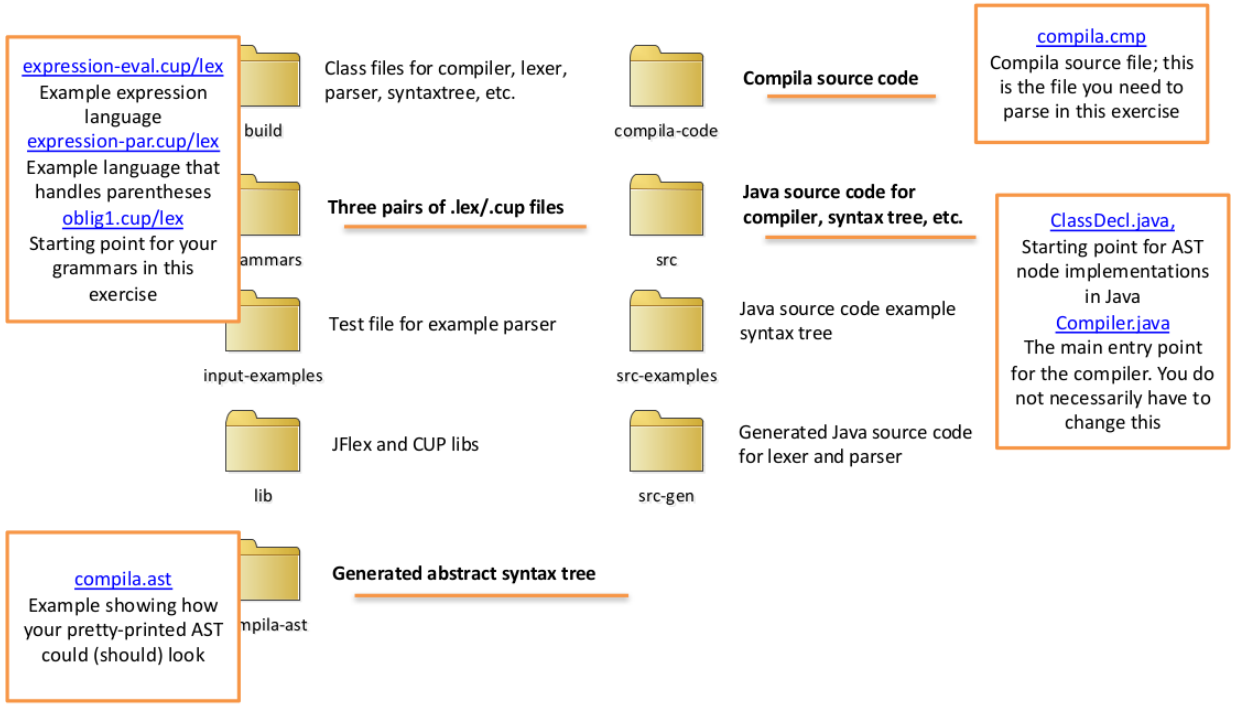
- supplied configuration should take care of calling `jflex`, `cup`, and `javadoc` for you

AST data structure

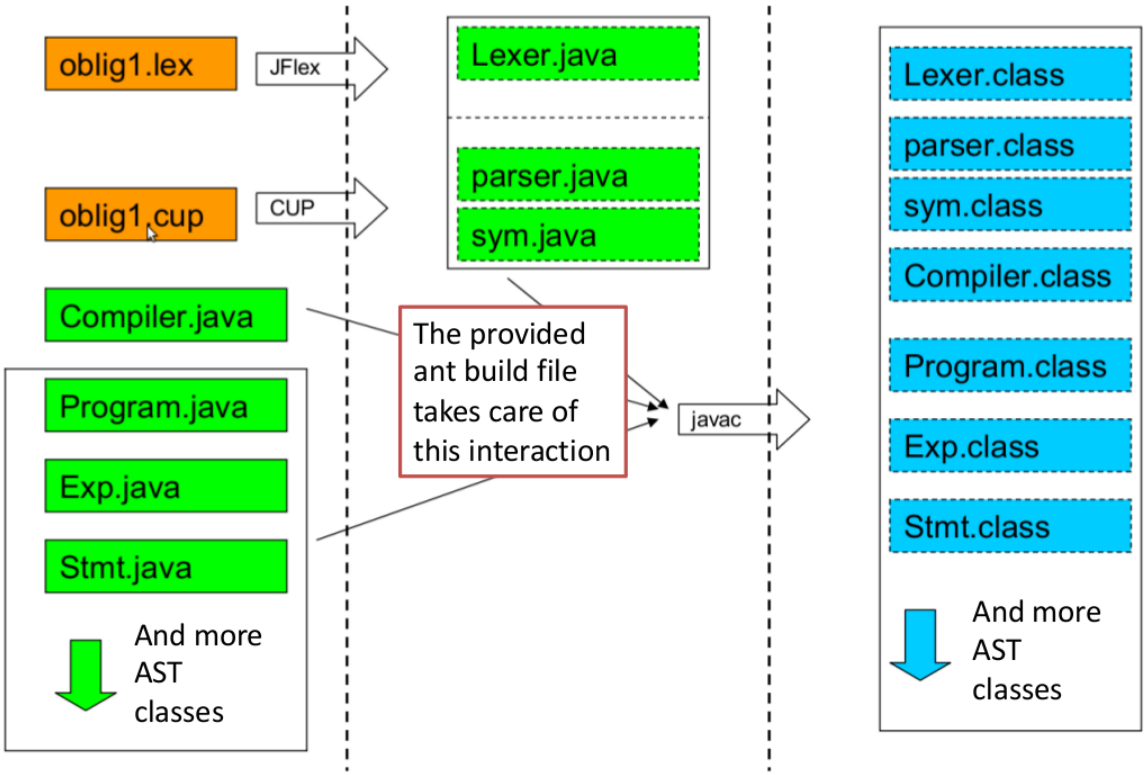
- Make a reasonable structure
  - This slide is an EXAMPLE
  - Do not copy it verbatim without thinking



Provides source code



Building: putting it together



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Official

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Deadline

Deadline

Friday 23. 03. 2018, 23:59

- don't miss the deadline
- for extensions, administration needs to agree (`studadm`), contact them if sick etc
- even if not 100% finished
  - deliver what you have
  - contact early when problems arise

Deliverables

- see also the “handout”

Deliverables (1)

- working parser
  - parse the supplied sample programs
  - printout the resulting AST
- two grammars (two `.cup`-files)
  - one unambiguious
  - one ambiguous, where ambibuities resolved through precedence declations in *CUP*, e.g.

```
precendence left AND;
```

Deliverables

Deliverables (2)

- report (with name(s) and UiO user name(s)
- discussion of the solution (see handout for questions)
- in particular: comparison of the two grammars
- "Readme"

- the code must *build* (with ant) and run
- test it on the UiO RHEL platform

Ask

If problems, ask in time (NOT Friday at the deadline)

Hand-in procedure

- this year we try *git*
- `https://github.uio.no` resp. `https://github.uio.no/msteffen/compila`
- you need
  - a login
  - send me emails that you want to do oblig (+ potential partner) ⇒ I tell you group number
  - create a project `compila<n>` (n = group number)
  - add collaborator + (at some point me)
- see also the handout
- code ready *tomorrow*

Section

Oblig 2

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Chapter 2

Oblig2

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Goal

1. semantic analysis, as far as
   - typing is concerned ("static semantics")
   - other coditions (no duplicate declaration etc)
2. code generation for `compila18` (ish) programs

Last time (O1)

Syntactic analysis

- lexer (scanner)
- parser
- abstract syntax tree

this time: continue with you previous deliv (and repos)

Learning outcome

- understand type checking, implementing a simmple variant
- undertand (simple form of) bytecode and how to generate it from "source code" (as AST)
- extend an existing compiler code base with new functionality

Semantic analysis & type checking

- parser / context-free grammars
  - not powerfull enough
  - cannot check all (static) properties of a language spec
- => extend the front-end by a type checker
  - use the AST classes of last time
  - add type checking code
  - allowed to make changes or adaptations if advantagous.

Another glance at compila18

```
program MyProgram begin

 class Complex begin
   var Real : float;
   var Imag : float;
 end;


 proc Add(a : Complex, b : Complex) : Complex
 begin
   var retval : Complex;
   retval := new Complex;
   retval.Real := a.Real + b.Real;
   retval.Imag := a.Imag + b.Imag;

   return retval;
 end;

 proc Main()
 begin
  var c1 : Complex;
  var c2 : Complex;
  var result : Complex;
  ...
  result := Add ( c1, c2 );
  ...
  return;
end; end;
```

Real and Imag are of the (built-in) float type.
Complex defines a new (user-defined) type.

Check that the + operator is compatible with its operands' types, and that the assignment is legal.

Check that the actual parameters to Add(...) are of the correct type, according to the formal parameters, and that the assignment to result is legal.

NB: 2018: structs, not classes

Type checking for conditionals

- as “inspiration”, details may vary

```
class IfStatement extends Statement {
...
  public void typeCheck(){
    String condType = condition.get.Type ();
    if (condType != "bool") {
      throw new TypeException("condition in an if
        statement must be of type bool")
    }
}
```

Type checking: assignments

```
class Assignment extends Statement {
...
  public void typeCheck() {
    String varType = var.getType();
    String expType = exp.getType();
    if (varType != expType &&
       !isAssigmentCompatible(varType,expType){
                  throw new TypeException("Cannot assi
                  " from " + expType);
  }
}
```

Code generation

- lecture(s) of code gen start right now (so it might look puzzling, but hopefully will become clearer)
- byte code API and operations are described in the document "Interpreter and bytecode for INF5110"
- Task: add bytecode generation methods to your AST classes for instance

```
Ast.Node.GenerateCode(...)
```

- again: if adaptations of the AST are called for or useful, go for it...

Code generation: limitations

- interpreter and byte code library somewhat limited
  - cannot express full compila 18
  - no block structure
  - no reference types
- your delivery should support generating correct bytecode for the `compila` 18 source code file `runme.cmp`

Code generation: creating a procedure

```
CodeFile codeFile = new CodeFile();
// add the procedure by name first
codeFile.addProcedure("Main")
// then define it
CodeProcedure main = new
   CodeProcedure("Main", VoidType,TYPE, codeFile);
main.addInstruction( new RETURN());
//then update it in the code file
codeFile.updateProcedure(main);
```

Code generation: assignment


```
//1: proc add(a: int, b : int ) : int {
//2: var res : int;
//3: res := a + b; // only bytecode for this line
//4: return res;
//5: }

// push a onto the stack
proc.addInstruction(new LOADLOCAL(proc.variableNumber("a")));
// push b onto the stack
proc.addInstruction(new LOADLOCAL(proc.variableNumber("b")));
// perform addition with arguments on the stack
proc.addInstruction(new ADD());
// pop result from stack, and store it in variable res
proc.addInstruction(new
        STORELOCAL(proc.variableNumber("res")));
```

Testing

- bunch of test files, for testing the *type checker*
- preferable: make `ant test` workable
- test files ending with `fail` containt a syntactically correct but erronous program (erroneous as the type system or generally the semantic phase is concerned)
- => compiler returns error code 2 for semantic failure

Provided source code

https://github.uio.no/msteffen/compila

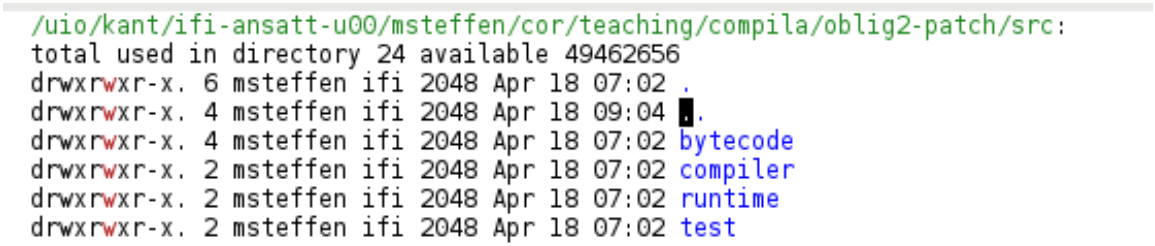
```
/uio/kant/ifi-ansatt-u00/msteffen/cor/teaching/compila:
total used in directory 48 available 49462784
drwxrwxr-x.  9 msteffen ifi 2048 Apr 18 10:24 .
drwxrwxr-x. 10 msteffen ifi 2048 Apr  9 13:33 ..
drwxrwxr-x.  8 msteffen ifi 2048 Apr 18 10:25 .git
-rw-rw-r--.  1 msteffen ifi    2 Feb 26 06:09 .gitignore
drwxrwxr-x.  3 msteffen ifi 2048 Feb 26 13:26 2017start
-rw-rw-r--.  1 msteffen ifi 4278 Apr 18 07:02 Readme.org
drwxrwxr-x.  5 msteffen ifi 2048 Feb 26 09:24 doc
drwxrwxr-x.  3 msteffen ifi 2048 Apr 17 16:04 material
drwxrwxr-x. 11 msteffen ifi 2048 Mar 26 06:58 oblig1-starting-point
drwxrwxr-x.  4 msteffen ifi 2048 Apr 18 09:04 oblig2-patch
drwxrwxr-x.  4 msteffen ifi 2048 Apr  5 09:29 src

/uio/kant/ifi-ansatt-u00/msteffen/cor/teaching/compila/oblig2-patch:
total used in directory 24 available 49462784
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 09:04 .
drwxrwxr-x. 9 msteffen ifi 2048 Apr 18 10:24 ..
-rwxrwxr-x. 1 msteffen ifi  592 Apr 18 09:03 Readme-patch
-rwxrwxr-x. 1 msteffen ifi 1203 Apr 18 09:04 build.xml
drwxrwxr-x. 6 msteffen ifi 2048 Apr 18 07:02 src
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 09:02 src-compila

/uio/kant/ifi-ansatt-u00/msteffen/cor/teaching/compila/oblig2-patch/src:
total used in directory 24 available 49462784
drwxrwxr-x. 6 msteffen ifi 2048 Apr 18 07:02 .
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 09:04 ..
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 07:02 bytecode
drwxrwxr-x. 2 msteffen ifi 2048 Apr 18 07:02 compiler
drwxrwxr-x. 2 msteffen ifi 2048 Apr 18 07:02 runtime
drwxrwxr-x. 2 msteffen ifi 2048 Apr 18 07:02 test

/uio/kant/ifi-ansatt-u00/msteffen/cor/teaching/compila/oblig2-patch/src-compila:
total used in directory 20 available 49462784
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 09:02 .
drwxrwxr-x. 4 msteffen ifi 2048 Apr 18 09:04 ..
drwxrwxr-x. 2 msteffen ifi 2048 Apr 18 09:02 examples
drwxrwxr-x. 2 msteffen ifi 6144 Apr 18 08:53 test-examples
```

Provided source code



- `compiler`: updated compiler class
- `test`: some code for performing tests
- `bytecode`: classes for constructing bytecode
- `runtime`: rte for executing the byte code

Deadline

Deadline

11th May 2018

Note: end of semester, and I need to report the ones passing the oblig some time before the exam.

delivs

- working type checker
- code generator (test with `runme.cmp`)
- report (including your name(s) etc.
  - discussion of your solution, choices you made, assumptions you rely on
  - printout of a test run (can be also checked in into the repos, but it n needs to be mentioned where it is)
  - printout of the bytecode from `runme.cmp` (with a target like `ant list-runme`)
  - solution must "build" and be "testable" (typically via `ant`)

References I