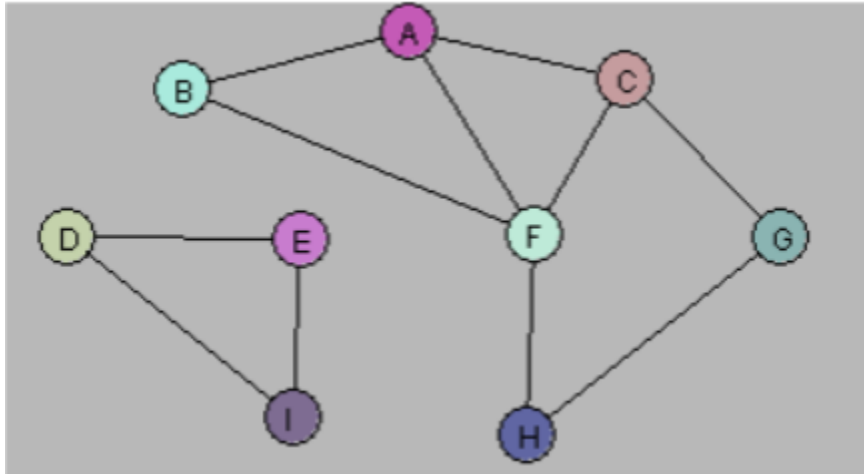
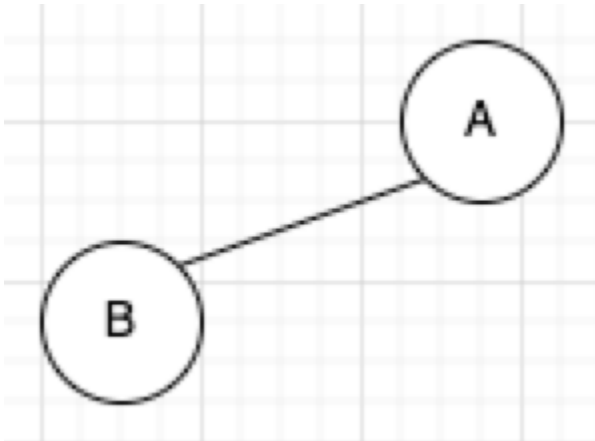


Lab 8

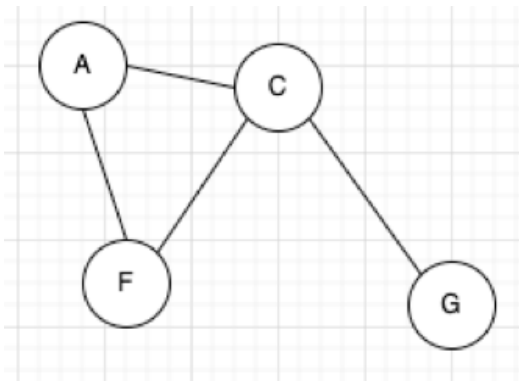
1. Induced Graphs. Answer questions about the graph $G = (V, E)$ displayed below.



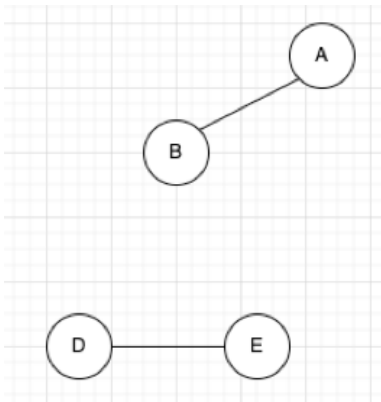
A. Let $U = \{A, B\}$. Draw $G[U]$.



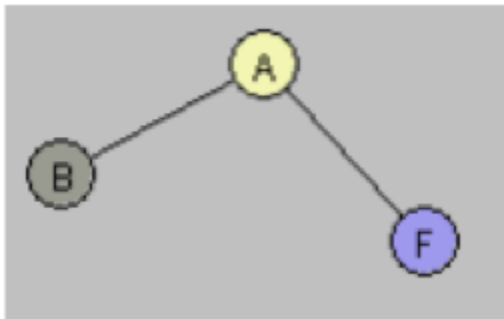
B. Let $W = \{A, C, G, F\}$. Draw $G[W]$.



C. Let $Y = \{A, B, D, E\}$. Draw $G[Y]$



D. Consider the following subgraph H of G:



Is there a subset X of the vertex set V so that $H = G[X]$? Explain.

-> yes. $X = \{A, B, F\} \subseteq V$ & $H = G[X]$

E. Find a way to partition the vertex set V into two subsets V_1, V_2 so that each of the induced graphs $G[V_1]$ and $G[V_2]$ is connected and $G = G[V_1] \cup G[V_2]$.

- $V_1 = \{A, B, C, F, G, H\}$
- $V_2 = \{D, E, I\}$

$G[V_1]$ is connected.

$G[V_2]$ is connected.

$G = G[V_1] \cup G[V_2]$

3. Graph Exercises.

A. Suppose $G = (V, E)$ is a connected simple graph. Suppose V_1, V_2, \dots, V_k are disjoint subsets of V and that $V_1 \cup V_2 \cup \dots \cup V_k = V$. Show that there is an edge (x, y) in E such that for some $i \neq j$, x is in V_i and y is in V_j .

- Assume that there is no such edge in E . This means that for every edge (x, y) in E , either x and y belong to the same subset V_i or they both belong to different subsets, say V_i and V_j , where $i = j$.

Since G is a connected graph, there exists a path between any two vertices in V . Let's consider two arbitrary vertices u and v in V , where u is in V_i and v is in V_j , with $i \neq j$.

Since there is no edge (x, y) in E such that x is in V_i and y is in V_j , there must be a path P between u and v that consists only of vertices from V_i . However, since v is not in V_i , the path P cannot reach v , which contradicts the assumption that G is connected.

Therefore, our assumption that there is no edge (x, y) in E such that for some i different from j , x is in V_i and y is in V_j must be false. Hence, there must exist at least one edge (x, y) in E such that x is in V_i and y is in V_j , with $i \neq j$.

- B. In class it was shown that a graph $G = (V, E)$ is connected whenever the following is true,

$$(*) \quad m > C(n-1, 2)$$

where n is the number of vertices and m is the number of edges. Is the following true or false?

Every connected graph satisfies the inequality $(*)$.

Prove your answer.

If G is not connected then $m \leq C(n-1, 2) = (n-1)(n-2)/2$

It can be shown that the graph having n vertices with the maximum possible number edge is: $G = K_{n-1} + K_1$
 $\Rightarrow m > C(n-1, 2)$

- C. Suppose G is a graph with two vertices. What is the minimum number of edges it must have in order to be a connected graph? Suppose instead G has three vertices; what is the minimum number of edges it must have in order to be connected? Fill in the blank with a reasonable conjecture:

If G has n vertices, G must have at least $\frac{(n-1)(n-2)}{2} + 1$ edges in order to be connected.