20/08/2017 MIPS Quick Tutorial

MIPS Architecture and Assembly Language Overview

 $Adapted\ from:\ http://edge.mcs.dre.g.el.edu/GICL/people/sevy/architecture/MIPSRef(SPIM).html (SPIM) and (SPIM) architecture/MIPSRef(SPIM) and (SPIM) architecture/MIPSRef(SPIM) arch$

[Register Description] [I/O Description]

Data Types and Literals

Data types:

- Instructions are all 32 bits
- byte(8 bits), halfword (2 bytes), word (4 bytes)
- · a character requires 1 byte of storage
- an integer requires 1 word (4 bytes) of storage

Literals:

- numbers entered as is. e.g. 4
- characters enclosed in single quotes. e.g. 'b'
- strings enclosed in double quotes. e.g. "A string"

Registers

- 32 general-purpose registers
- register preceded by \$\\$\\$ in assembly language instruction two formats for addressing:
 - o using register number <u>e.g.</u> \$0 through \$31
 - o using equivalent names e.g. \$t1, \$sp
- · special registers Lo and Hi used to store result of multiplication and division
 - o not directly addressable; contents accessed with special instruction mfhi ("move from Hi") and mflo ("move from Lo")
- stack grows from high memory to low memory

This is from Figure 9.9 in the Goodman&Miller text

Register Number	Alternative Name	Description	
0	zero	the value 0	
1	\$at	(assembler temporary) reserved by the assembler	
2-3	\$v0 - \$v1	(values) from expression evaluation and function results	
4-7	\$a0 - \$a3	(arguments) First four parameters for subroutine. Not preserved across procedure calls	
8-15	\$t0 - \$t7	(temporaries) Caller saved if needed. Subroutines can use w/out saving. Not preserved across procedure calls	
16-23	\$s0 - \$s7	(saved values) - Callee saved. A subroutine using one of these must save original and restore it before exiting. Preserved across procedure calls	
24-25	\$t8 - \$t9	(temporaries) Caller saved if needed. Subroutines can use w/out saving. These are in addition to \$t0 - \$t7 above. Not preserved across procedure calls.	
26-27	\$k0 - \$k1	reserved for use by the interrupt/trap handler	
28	\$gp	global pointer. Points to the middle of the 64K block of memory in the static data segment.	
29	\$sp	stack pointer Points to last location on the stack.	
30	\$s8/\$fp	saved value / frame pointer Preserved across procedure calls	
31	\$ra	return address	

See also Britton section 1.9, Sweetman section 2.21, Larus Appendix section A.6

Program Structure

- just plain text file with data declarations, program code (name of file should end in suffix .s to be used with SPIM simulator)
- data declaration section followed by program code section

Data Declarations

- placed in section of program identified with assembler directive $\boldsymbol{.data}$
- declares variable names used in program; storage allocated in main memory (RAM) $\,$

Code

- placed in section of text identified with assembler directive .text
- contains program code (instructions)
- starting point for code e.g.ecution given label main:
- ending point of main code should use exit system call (see below under System Calls)

Comments

```
• anything following # on a line
# This stuff would be considered a comment
```

```
• Template for a MIPS assembly language program:
```

```
# Comment giving name of program and description of function
# Template.s
# Bare-bones outline of MIPS assembly language program

.data  # variable declarations follow this line
# ...

.text  # instructions follow this line
main:  # indicates start of code (first instruction to execute)
# ...
# End of program, leave a blank line afterwards to make SPIM happy
```

Data Declarations

```
format for declarations:
```

```
name: storage_type value(s)
```

- o create storage for variable of specified type with given name and specified value
- value(s) usually gives initial value(s); for storage type .space, gives number of spaces to be allocated

Note: labels always followed by colon (:)

Load / Store Instructions

- · RAM access only allowed with load and store instructions
- all other instructions use register operands

load:

```
lw register_destination, RAM_source
#copy word (4 bytes) at source RAM location to destination register.
lb register_destination, RAM_source
#copy byte at source RAM location to low-order byte of destination register,
# and sign-e.g.tend to higher-order bytes
```

store word:

```
sw register_source, RAM_destination
#store word in source register into RAM destination
sb register_source, RAM_destination
#store byte (low-order) in source register into RAM destination
```

load immediate:

```
li register_destination, value
#load immediate value into destination register
```

Indirect and Based Addressing

• Used only with load and store instructions

load address:

```
la $t0, var1
```

• copy RAM address of var1 (presumably a label defined in the program) into register \$t0

indirect addressing:

```
lw $t2, ($t0)
```

• load word at RAM address contained in \$t0 into \$t2

```
sw $t2, ($t0)
```

· store word in register \$t2 into RAM at address contained in \$t0

based or indexed addressing:

```
lw $t2, 4($t0)
```

- load word at RAM address (\$t0+4) into register \$t2
- "4" gives offset from address in register \$t0

```
sw $t2, -12($t0)
```

- store word in register \$t2 into RAM at address (\$t0 12)
- negative offsets are fine

Note: based addressing is especially useful for:

- · arrays; access elements as offset from base address
- · stacks; easy to access elements at offset from stack pointer or frame pointer

example

```
array1:
                 .space
                        12
                                         # declare 12 bytes of storage to hold array of 3 integers
                 .text
start:
                la
                         $t0, array1
                                                    load base address of array into register $t0
                                            $t1 = 5
                li
                         $±1.5
                                                       ("load immediate"
                                            first array element set to 5; indirect addressing
                sw $t1, ($t0)
                li $t1, 13
                sw $t1, 4($t0)
li $t1, -7
                                         #
                                            second array element set to 13
                                              $t1 = -7
                sw $t1, 8($t0)
                                            third array element set to -7
                done
```

Arithmetic Instructions

- · most use 3 operands
- · all operands are registers; no RAM or indirect addressing
- operand size is word (4 bytes)

```
add
            $t0.$t1.$t2
                                         $t0 = $t1 + $t2;
$t2 = $t3 D $t4
                                                                      add as signed (2's complement) integers
            $t2,$t3,$t4
sub
addi
            $t2,$t3,5
                                     #
                                         $t2 = $t3 + 5;
                                                                    "add immediate" (no sub immediate)
                                         $t1 = $t6 + $t7;
$t1 = $t6 + $t7;
addu
            $t1.$t6.$t7
                                                                       \hbox{add as unsigned integers} \\
                                                                      subtract as unsigned integers
subu
            $t1.$t6.$t7
                                         multiply 32-bit quantities in $t3 and $t4, and store 64-bit result in special registers Lo and Hi: (Hi,Lo) = $t3 * $t4 Lo = $t5 / $t6 (integer quotient)
mult
            $t3.$t4
                                        Ho = $t5 / $t6 (integer quotient)
Hi = $t5 mod $t6 (remainder)
move quantity in special register Hi to $t0:
move quantity in special register Lo to $t1:
div
            $t5,$t6
mfhi
            $t0
                                                                                                                  $t0 = Hi
                                         used to get at result of product or quotient
            $t2,$t3 # $t2 = $t3
move
```

Control Structures

Branches

Jumps

• comparison for conditional branches is built into instruction

```
b
          target
                                 unconditional branch to program label target
                                 branch to target if $t0 = $t1
branch to target if $t0 < $t1
beq
          $t0,$t1,target
          $t0,$t1,target #
$t0,$t1,target #
b1t
                                 branch to target if
ble
                                                           $t0 <= $t1
bgt
          $t0,$t1,target
                             #
                                 branch to target if
                                                           $t0 > $t1
          $t0,$t1,target #
$t0,$t1,target #
                                 branch to target if branch to target if
                                                           $t0 >= $t1
$t0 <> $t1
bge
bne
          target # unconditional jump to program label target
                             # jump to address contained in $t3 ("jump register")
          $t3
```

Subroutine Calls

subroutine call: "jump and link" instruction

```
jal sub_label # "jump and link"
```

- copy program counter (return address) to register \$ra (return address register)
- · jump to program statement at sub_label

subroutine return: "jump register" instruction

```
jr $ra # "jump register"
```

• jump to return address in \$ra (stored by jal instruction)

Note: return address stored in register \$ra; if subroutine will call other subroutines, or is recursive, return address should be copied from \$ra onto stack to preserve it, since jal always places return address in this register and hence will overwrite previous value

System Calls and I/O (SPIM Simulator)

- · used to read or print values or strings from input/output window, and indicate program end
- · use syscall operating system routine call
- first supply appropriate values in registers \$v0 and \$a0-\$a1
- result value (if any) returned in register \$v0

The following table lists the possible **syscall** services.

Service	Code in \$v0	Arguments	Results
print_int	1	\$a0 = integer to be printed	
print_float	2	\$f12 = float to be printed	
print_double	3	\$f12 = double to be printed	
print_string	4	\$a0 = address of string in memory	
read_int	5		integer returned in \$v0
read_float	6		float returned in \$v0
read_double	7		double returned in \$v0
read_string	8	\$a0 = memory address of string input buffer \$a1 = length of string buffer (n)	
sbrk	9	\$a0 = amount	address in \$v0
exit	10		

- o The print_string service expects the address to start a null-terminated character string. The directive .asciiz creates a null-terminated character string.
- o The read_int, read_float and read_double services read an entire line of input up to and including the newline character.
- The read_string service has the same semantices as the UNIX library routine fgets.
 - It reads up to n-1 characters into a buffer and terminates the string with a null character.
 - If fewer than n-1 characters are in the current line, it reads up to and including the newline and terminates the string with a null character.
- The sbrk service returns the address to a block of memory containing n additional bytes. This would be used for dynamic memory allocation.
- The exit service stops a program from running.
- e.g. Print out integer value contained in register \$t2

e.g. Read integer value, store in RAM location with label int_value (presumably declared in data section)

e.g. Print out string (useful for prompts)

 $\underline{\text{e.g.}}$ To indicate end of program, use $\underline{\text{exit}}$ system call; thus last lines of program should be:

```
li $v0, 10  # system call code for exit = 10
syscall  # call operating sys
```