

# Quant. Comp. HW - 2

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Oct. 18, 2005

## 1 Simon's Problem

(a) There are 128 possible outputs

(b) See attached source code:

```
0011001100001110100110000011111110111101100000010000001001110011010100001
110001001000001110111000101001001011001010010010010001
```

340282366920938463463374607431768211455

(c) I reported an average of 129 trials

(d)

## 2 Modular Exponentiation

Here was my python code:

```
#Modular Exponentiation
import math

def modular_exponentiation(base, exponent, modulus):
    c = 1
    for e_prime in range(1, exponent+1):
        c = (c * base) % modulus
    return c

print modular_exponentiation(1234, 1234*1234, math.pow(10, 10))
```

And the output was:

3102217216.0

### 3 RSA Misuse

I first tried to solve this as strictly a math problem, but had little success, in large part because I thought that the  $\gcd(e_1, e_2)$  was somehow irrelevant to the problem. However, I noticed that normally the exponents are the same value when performing RSA, so I scoured google to see if there was some well known attack where you have a common modulus with different attacks. Turns out that it is fairly well documented, and Simmons wrote a paper on it a while back.

The basic idea is: Since

$$\gcd(e_1, e_2) = 1$$

, then

$$\exists u, v \text{ s.t. } e_1 * u + e_2 * v = 1$$

To solve for u and v, I used the extended Euclidean algorithm. I then raise each side to u and v

$$c1^u = (M^{e1})^u \bmod n$$

$$c2^v = (M^{e2})^v \bmod n$$

$$c1^u * c2^v = (M^{e1})^u * (M^{e2})^v \bmod n = M^{e1*u+e2*v} \bmod n = M \bmod n$$

From this I can multiply  $c1^d$  and  $(c2^f)^{-1}$ :

$$c1^d * (c2^f)^{-1} = M^{bd} M^{-ce} = M^{bd-ce} = M \bmod n$$

Because my modular exponentiation code can't handle negatives, note that  $c2^{-v} \bmod n = c2^{n-f} \bmod n$

### 4 Prime factorization

Problem: Consider  $n=121932632103337941464563328643500519$

(a) How many bits is n?

```
print len(str(121932632103337941464563328643500519))
```

Output:

36

(b) Find if  $n$  is prime with program that runs in less than one second.

```
def miller_rabin_pass(a, s, d, n):
    a_to_power = pow(a, d, n)
    if a_to_power == 1:
        return True
    for i in xrange(s-1):
        if a_to_power == n - 1:
            return True
        a_to_power = (a_to_power * a_to_power) % n
    return a_to_power == n - 1

def miller_rabin(n):
    #compute s and d
    d = n - 1
    s = 0
    while d % 2 == 0:
        d >>= 1
        s += 1

    #Run several miller_rabin passes
    for repeat in xrange(20):
        a = randint(2, n-1)
        if not miller_rabin_pass(a, s, d, n):
            return False
    return True

print miller_rabin(n)
```

- (b)
- (c)
- (d)
- (e)