ACM 常用算法模板

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1 数据结构

1.1 并查集

```
struct dsu {
1
2
      int n;
3
      vector<int> fa, sz;
4
      dsu(int _n) : n(_n), fa(n + 1), sz(n + 1, 1) {
5
          iota(fa.begin(), fa.end(), 0);
6
      }
7
      int find(int x) { return x == fa[x] ? x : fa[x] = find(fa[x]); }
8
      int merge(int x, int y) {
9
          int fax = find(x), fay = find(y);
10
          if (fax == fay) return 0; // 一个集合
11
          sz[fay] += fax;
12
          return fa[fax] = fay; // 合并到哪个集合了
13
14
      int size(int x) { return sz[find(x)]; }
15 };
```

1.2 树状数组

1.2.1 一维

```
template <class T>
2
  struct fenwick {
3
       int n;
4
       vector<T> t;
5
       fenwick(int _n) : n(_n), t(n + 1) {}
6
       T query(int 1, int r) {
7
           auto query = [&](int pos) {
8
                T res = 0;
9
                while (pos) {
10
                    res += t[pos];
                    pos -= lowbit(pos);
11
12
13
                return res;
14
15
           return query(r) - query(l - 1);
16
17
       void add(int pos, T num) {
           while (pos <= n) {</pre>
18
19
                t[pos] += num;
                pos += lowbit(pos);
20
21
           }
22
       }
23 };
```

1.2.2 二维

```
template <class T>
2
  struct Fenwick_tree_2 {
3
       Fenwick_tree_2(int n, int m) : n(n), m(m), tree(n + 1, vector<T>(m + 1)) {}
4
       T query(int 11, int r1, int 12, int r2) {
5
           auto query = [&](int 1, int r) {
6
               T res = 0;
7
               for (int i = 1; i; i -= lowbit(i))
8
                   for (int j = r; j; j -= lowbit(j)) res += tree[i][j];
9
               return res;
10
           };
           return query(12, r2) - query(12, r1 - 1) - query(11 - 1, r2) +
11
12
                  query(11 - 1, r1 - 1);
13
       void update(int x, int y, T num) {
14
15
           for (int i = x; i <= n; i += lowbit(i))</pre>
16
               for (int j = y; j <= m; j += lowbit(j)) tree[i][j] += num;</pre>
17
       }
18
  private:
19
       int n, m;
20
       vector<vector<T>> tree;
21 };
```

1.2.3 三维

```
template <class T>
2
   struct Fenwick_tree_3 {
3
       Fenwick_tree_3(int n, int m, int k)
 4
           : n(n),
5
             m(m),
6
             k(k),
7
             tree(n + 1, vector<vector<T>>(m + 1, vector<T>(k + 1))) {}
8
       T query(int a, int b, int c, int d, int e, int f) {
9
           auto query = [&](int x, int y, int z) {
10
               T res = 0;
               for (int i = x; i; i -= lowbit(i))
11
12
                    for (int j = y; j; j -= lowbit(j))
13
                        for (int p = z; p; p -= lowbit(p)) res += tree[i][j][p];
14
               return res;
15
           };
16
           T res = query(d, e, f);
17
           res -= query(a - 1, e, f) + query(d, b - 1, f) + query(d, e, c - 1);
           res += query(a - 1, b - 1, f) + query(a - 1, e, c - 1) +
18
19
                  query(d, b - 1, c - 1);
           res -= query(a - 1, b - 1, c - 1);
20
21
           return res;
22
23
       void update(int x, int y, int z, T num) {
           for (int i = x; i <= n; i += lowbit(i))</pre>
24
25
               for (int j = y; j <= m; j += lowbit(j))</pre>
26
                    for (int p = z; p \leftarrow k; p += lowbit(p)) tree[i][j][p] += num;
```

```
27     }
28     private:
29     int n, m, k;
30     vector<vector<T>>> tree;
31     };
```

1.3 线段树

```
template <class Data, class Num>
   struct Segment_Tree {
3
       inline void update(int 1, int r, Num x) { update(1, 1, r, x); }
 4
       inline Data query(int 1, int r) { return query(1, 1, r); }
5
       Segment_Tree(vector<Data>& a) {
6
           n = a.size();
7
           tree.assign(n * 4 + 1, {});
8
           build(a, 1, 1, n);
9
       }
10
   private:
11
       int n;
12
       struct Tree {
13
           int 1, r;
14
           Data data;
15
       };
16
       vector<Tree> tree;
17
       inline void pushup(int pos) {
18
           tree[pos].data = tree[pos << 1].data + tree[pos << 1 | 1].data;</pre>
19
20
       inline void pushdown(int pos) {
21
           tree[pos << 1].data = tree[pos << 1].data + tree[pos].data.lazytag;</pre>
22
           tree[pos << 1 | 1].data =
23
               tree[pos << 1 | 1].data + tree[pos].data.lazytag;</pre>
24
           tree[pos].data.lazytag = Num::zero();
25
       }
26
       void build(vector<Data>& a, int pos, int 1, int r) {
27
           tree[pos].l = 1;
28
           tree[pos].r = r;
29
           if (1 == r) {
30
               tree[pos].data = a[l - 1];
31
               return;
32
           }
33
           int mid = (tree[pos].l + tree[pos].r) >> 1;
34
           build(a, pos << 1, 1, mid);</pre>
35
           build(a, pos << 1 | 1, mid + 1, r);
36
           pushup(pos);
37
       void update(int pos, int& 1, int& r, Num& x) {
38
39
           if (1 > tree[pos].r || r < tree[pos].1) return;</pre>
40
           if (1 <= tree[pos].1 && tree[pos].r <= r) {</pre>
41
               tree[pos].data = tree[pos].data + x;
42
               return;
43
```

```
44
           pushdown(pos);
45
           update(pos << 1, 1, r, x);
46
           update(pos << 1 | 1, 1, r, x);
47
           pushup(pos);
48
       }
49
       Data query(int pos, int& 1, int& r) {
50
           if (1 > tree[pos].r || r < tree[pos].l) return Data::zero();</pre>
           if (1 <= tree[pos].1 && tree[pos].r <= r) return tree[pos].data;</pre>
51
52
           pushdown(pos);
           return query(pos << 1, 1, r) + query(pos << 1 | 1, 1, r);</pre>
53
54
       }
55
  };
56
  struct Num {
57
       ll add;
58
       inline static Num zero() { return {0}; }
59
       inline Num operator+(Num b) { return {add + b.add}; }
60
  };
61
  struct Data {
62
       11 sum, len;
63
       Num lazytag;
64
       inline static Data zero() { return {0, 0, Num::zero()}; }
65
       inline Data operator+(Num b) {
66
           return {sum + len * b.add, len, lazytag + b};
67
       }
68
       inline Data operator+(Data b) {
69
           return {sum + b.sum, len + b.len, Num::zero()};
70
       }
71
  };
```

1.4 普通平衡树

1.4.1 树状数组实现

需要预先处理出来所有可能的数。

```
template <typename T>
 1
2
  struct treap {
3
       int n, size;
4
       vector<int> t;
5
       vector<T> t2, S;
6
       treap(const vector<T>& b) {
7
           S = b;
8
           sort(S.begin(), S.end());
9
           S.erase(unique(S.begin(), S.end()), S.end());
10
           n = S.size();
11
           size = 0;
12
           t = vector < int > (n + 1);
13
           t2 = vector < T > (n + 1);
14
15
       int pos(T x) { return lower_bound(S.begin(), S.end(), x) - S.begin() + 1; }
16
       int sum(int pos) {
17
           int res = 0;
```

```
18
           while (pos) {
19
               res += t[pos];
20
               pos -= lowbit(pos);
21
22
           return res;
23
       }
24
25
       // 插入cnt个x
26
       void insert(T x, int cnt) {
27
           size += cnt;
28
           for (int i = pos(x); i <= n; i += lowbit(i)) {</pre>
29
               t[i] += cnt;
30
               t2[i] += cnt * x;
31
           }
32
       }
33
34
       // 删除cnt个x
       void erase(T x, int cnt) { insert(x, -cnt); }
35
36
37
       // x的排名
38
       int rank(T x) { return sum(pos(x) - 1) + 1; }
39
40
       // 统计出现次数
41
       int count(T x) { return sum(pos(x)) - sum(pos(x) - 1); }
42
       // 第k小
43
       T kth(int k) {
44
           int cnt = 0, x = 0;
45
46
           for (int i = log2(n); i >= 0; i--) {
47
               x += 1 << i;
48
               if (x >= n \mid | cnt + t[x] >= k) x -= 1 << i;
49
               else cnt += t[x];
50
51
           return S[x];
52
       }
53
54
       // 前k小的数之和
55
       T pre_sum(int k) {
56
           int cnt = 0, x = 0;
57
           T res = 0;
58
           for (int i = log2(n); i >= 0; i--) {
59
               x += 1 << i;
60
               if (x >= n \mid | cnt + t[x] >= k) x -= 1 << i;
61
               else {
62
                   cnt += t[x];
63
                   res += t2[x];
64
               }
65
66
           return res + (k - cnt) * S[x];
67
       }
68
69
      // 小于x,最大的数
```

```
T prev(int x) { return kth(sum(pos(x) - 1)); }

// 大于x, 最小的数

T next(int x) { return kth(sum(pos(x)) + 1); }

};
```

1.5 可持久化线段树

```
constexpr int MAXN = 200000;
  vector<int> root(MAXN << 5);</pre>
  struct Persistent_seg {
3
 4
       int n;
5
       struct Data {
6
           int ls, rs;
7
           int val;
8
       };
9
       vector<Data> tree;
10
       Persistent_seg(int n, vector<int>& a) : n(n) { root[0] = build(1, n, a); }
11
       int build(int 1, int r, vector<int>& a) {
12
           if (1 == r) {
13
               tree.push_back({0, 0, a[1]});
14
               return tree.size() - 1;
15
           }
16
           int mid = 1 + r \gg 1;
17
           int ls = build(l, mid, a), rs = build(mid + 1, r, a);
18
           tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
19
           return tree.size() - 1;
20
       }
21
       int update(int rt, const int& idx, const int& val, int 1, int r) {
22
           if (1 == r) {
23
               tree.push_back({0, 0, tree[rt].val + val});
24
               return tree.size() - 1;
25
26
           int mid = l + r >> 1, ls = tree[rt].ls, rs = tree[rt].rs;
27
           if (idx <= mid) ls = update(ls, idx, val, l, mid);</pre>
28
           else rs = update(rs, idx, val, mid + 1, r);
29
           tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
30
           return tree.size() - 1;
31
       }
32
       int query(int rt1, int rt2, int k, int l, int r) {
33
           if (1 == r) return 1;
           int mid = 1 + r \gg 1;
34
           int lcnt = tree[tree[rt2].ls].val - tree[tree[rt1].ls].val;
35
36
           if (k <= lcnt) return query(tree[rt1].ls, tree[rt2].ls, k, l, mid);</pre>
37
           else return query(tree[rt1].rs, tree[rt2].rs, k - lcnt, mid + 1, r);
38
       }
39 };
```

```
1
  auto lg = []() {
2
       array<int, 10000001> lg;
3
       lg[1] = 0;
4
       for (int i = 2; i \leftarrow 10000000; i++) lg[i] = lg[i >> 1] + 1;
5
       return lg;
6 }();
7
  template <typename T>
8
  struct st {
9
       int n;
10
       vector<vector<T>> a;
11
       st(vector<T>& _a) : n(_a.size()) {
12
           a.assign(lg[n] + 1, vector<int>(n));
13
           for (int i = 0; i < n; i++) a[0][i] = _a[i];</pre>
           for (int j = 1; j <= lg[n]; j++)</pre>
14
15
               for (int i = 0; i + (1 << j) - 1 < n; i++)
16
                    a[j][i] = max(a[j - 1][i], a[j - 1][i + (1 << (j - 1))]);
17
       }
18
       T query(int 1, int r) {
           int k = lg[r - l + 1];
19
20
           return max(a[k][l], a[k][r - (1 << k) + 1]);</pre>
21
       }
22 };
```

2 图论

存图

```
1
  struct Graph {
2
       int n;
3
       struct Edge {
           int to, w;
 4
5
       };
 6
       vector<vector<Edge>> graph;
7
       Graph(int _n) {
8
           n = _n;
9
           graph.assign(n + 1, vector<Edge>());
10
11
       void add(int u, int v, int w) { graph[u].push_back({v, w}); }
12 };
```

2.1 最短路

2.1.1 dijkstra

```
void dij(Graph& graph, vector<int>& dis, int t) {
 1
2
       vector<int> visit(graph.n + 1, 0);
3
       priority_queue<pair<int, int>> que;
4
       dis[t] = 0;
5
       que.emplace(0, t);
6
       while (!que.empty()) {
7
           int u = que.top().second;
8
           que.pop();
9
           if (visit[u]) continue;
10
           visit[u] = 1;
           for (auto& [to, w] : graph.graph[u]) {
11
12
               if (dis[to] > dis[u] + w) {
13
                   dis[to] = dis[u] + w;
                   que.emplace(-dis[to], to);
14
15
               }
16
           }
17
       }
18
```

2.2 树上问题

2.2.1 最近公公祖先

倍增法

```
vector<int> dep;
vector<array<int, 21>> fa;
dep.assign(n + 1, 0);
fa.assign(n + 1, array<int, 21>{});
void binary_jump(int root) {
 function<void(int)> dfs = [&](int t) {
```

```
7
           dep[t] = dep[fa[t][0]] + 1;
8
           for (auto& [to] : graph[t]) {
9
               if (to == fa[t][0]) continue;
10
               fa[to][0] = t;
               dfs(to);
11
12
           }
13
       };
       dfs(root);
14
15
       for (int j = 1; j <= 20; j++)
16
           for (int i = 1; i <= n; i++) fa[i][j] = fa[fa[i][j - 1]][j - 1];</pre>
17
  int lca(int x, int y) {
18
19
       if (dep[x] < dep[y]) swap(x, y);</pre>
       for (int i = 20; i >= 0; i--)
20
21
           if (dep[fa[x][i]] >= dep[y]) x = fa[x][i];
       if (x == y) return x;
22
23
       for (int i = 20; i >= 0; i--) {
           if (fa[x][i] != fa[y][i]) {
24
25
               x = fa[x][i];
26
               y = fa[y][i];
27
           }
28
29
       return fa[x][0];
30 }
```

树剖

```
int lca(int x, int y) {
    while (top[x] != top[y]) {
        if (dep[top[x]] < dep[top[y]]) swap(x, y);
        x = fa[top[x]];
    }
    if (dep[x] < dep[y]) swap(x, y);
    return y;
}</pre>
```

2.2.2 树链剖分

```
1 vector<int> fa, siz, dep, son, dfn, rnk, top;
2 fa.assign(n + 1, 0);
3 | siz.assign(n + 1, 0);
 4 dep.assign(n + 1, 0);
5
  son.assign(n + 1, 0);
6 dfn.assign(n + 1, 0);
7
  rnk.assign(n + 1, 0);
8
  top.assign(n + 1, 0);
9
  void hld(int root) {
10
       function<void(int)> dfs1 = [&](int t) {
           dep[t] = dep[fa[t]] + 1;
11
12
           siz[t] = 1;
13
           for (auto& [to, w] : graph[t]) {
14
               if (to == fa[t]) continue;
```

```
15
                fa[to] = t;
16
                dfs1(to);
17
                if (siz[son[t]] < siz[to]) son[t] = to;</pre>
                siz[t] += siz[to];
18
19
           }
20
       };
       dfs1(root);
21
22
       int dfn_tail = 0;
23
       for (int i = 1; i <= n; i++) top[i] = i;</pre>
       function<void(int)> dfs2 = [&](int t) {
24
25
           dfn[t] = ++dfn_tail;
           rnk[dfn_tail] = t;
26
27
           if (!son[t]) return;
28
           top[son[t]] = top[t];
29
           dfs2(son[t]);
           for (auto& [to, w] : graph[t]) {
30
31
                if (to == fa[t] || to == son[t]) continue;
32
33
           }
34
       };
35
       dfs2(root);
36
```

2.3 强连通分量

```
void tarjan(Graph& g1, Graph& g2) {
2
       int dfn_tail = 0, cnt = 0;
3
       vector\langle int \rangle dfn(g1.n + 1, 0), low(g1.n + 1, 0), exist(g1.n + 1, 0),
4
           belong(g1.n + 1, 0);
5
       stack<int> sta;
6
       function<void(int)> dfs = [&](int t) {
7
           dfn[t] = low[t] = ++dfn_tail;
8
           sta.push(t);
9
           exist[t] = 1;
10
           for (auto& [to] : g1.graph[t])
11
                if (!dfn[to]) {
12
                    dfs(to);
13
                    low[t] = min(low[t], low[to]);
                } else if (exist[to]) low[t] = min(low[t], dfn[to]);
14
15
           if (dfn[t] == low[t]) {
16
                cnt++;
17
                while (int temp = sta.top()) {
                    belong[temp] = cnt;
18
19
                    exist[temp] = 0;
20
                    sta.pop();
21
                    if (temp == t) break;
22
                }
23
           }
24
       };
25
       for (int i = 1; i <= g1.n; i++)</pre>
26
           if (!dfn[i]) dfs(i);
```

2.4 拓扑排序

```
void toposort(Graph& g, vector<int>& dis) {
2
      vector<int> in(g.n + 1, 0);
3
      for (int i = 1; i <= g.n; i++)
4
           for (auto& [to] : g.graph[i]) in[to]++;
5
      queue<int> que;
6
      for (int i = 1; i <= g.n; i++)
7
           if (!in[i]) {
8
               que.push(i);
9
               dis[i] = g.w[i]; // dp
10
11
      while (!que.empty()) {
12
           int u = que.front();
13
           que.pop();
14
           for (auto& [to] : g.graph[u]) {
15
               in[to]--;
16
               dis[to] = max(dis[to], dis[u] + g.w[to]); // dp
17
               if (!in[to]) que.push(to);
18
           }
19
      }
20 }
```

3 字符串 14

3 字符串

3.1 kmp

```
auto kmp(string& s) {
1
2
       vector next(s.size(), -1);
3
       for (int i = 1, j = -1; i < s.size(); i++) {</pre>
4
           while (j \ge 0 \&\& s[i] != s[j + 1]) j = next[j];
5
           if (s[i] == s[j + 1]) j++;
6
           next[i] = j;
7
       }
8
       // next 意为长度
9
       for (auto& i : next) i++;
10
       return next;
11
```

3.2 哈希

```
1 constexpr int N = 2e6;
2 constexpr 11 mod[2] = {20000000011, 2000000033}, base[2] = {20011, 20033};
3
  vector<array<11, 2>> pow_base(N);
5
  pow_base[0][0] = pow_base[0][1] = 1;
  for (int i = 1; i < N; i++) {</pre>
6
       pow_base[i][0] = pow_base[i - 1][0] * base[0] % mod[0];
7
       pow_base[i][1] = pow_base[i - 1][1] * base[1] % mod[1];
8
9
10
  struct Hash {
11
12
       int size;
13
       vector<array<11, 2>> hash;
14
       Hash() {}
15
       Hash(const string& s) {
           size = s.size();
16
17
           hash.resize(size);
18
           hash[0][0] = hash[0][1] = s[0];
19
           for (int i = 1; i < size; i++) {</pre>
               hash[i][0] = (hash[i - 1][0] * base[0] + s[i]) % mod[0];
20
21
               hash[i][1] = (hash[i - 1][1] * base[1] + s[i]) % mod[1];
           }
22
23
       }
24
       array<11, 2> operator[](const array<int, 2>& range) const {
25
           int 1 = range[0], r = range[1];
           if (1 == 0) return hash[r];
26
27
           auto single_hash = [&](bool flag) {
28
               return (hash[r][flag] -
29
                        hash[l - 1][flag] * pow_base[r - l + 1][flag] % mod[flag] +
30
                        mod[flag]) %
31
                       mod[flag];
32
           };
33
           return {single_hash(0), single_hash(1)};
```

3 字符串 15

```
34 };
```

3.3 manacher

```
void manacher(const string& _s, vector<int>& r) {
2
       string s(_s.size() * 2 + 1, '$');
3
       for (int i = 0; i < _s.size(); i++) s[2 * i + 1] = _s[i];
4
       r.resize(_s.size() * 2 + 1);
5
       for (int i = 0, maxr = 0, mid = 0; i < s.size(); i++) {</pre>
6
           if (i < maxr) r[i] = min(r[mid * 2 - i], maxr - i);</pre>
7
           while (i - r[i] - 1 >= 0 \&\& i + r[i] + 1 < s.size() \&\&
8
                  s[i - r[i] - 1] == s[i + r[i] + 1])
9
               ++r[i];
10
           if (i + r[i] > maxr) maxr = i + r[i], mid = i;
11
      }
12 }
```

4 数学

4.1 扩展欧几里得

```
需保证 a,b>=0 x=x+k*dx, y=y-k*dy 若要求 x\geq p,\ k\geq \left\lceil\frac{p-x}{dx}\right\rceil 若要求 x\leq q,\ k\leq \left\lfloor\frac{q-x}{dx}\right\rfloor 若要求 y\geq p,\ k\leq \left\lfloor\frac{y-p}{dy}\right\rfloor 若要求 y\leq q,\ k\geq \left\lceil\frac{y-q}{dy}\right\rceil
```

```
int __exgcd(int a, int b, int& x, int& y) {
2
       if (!b) {
3
           x = 1;
4
           y = 0;
5
           return a;
6
7
       int g = __exgcd(b, a % b, y, x);
8
       y -= a / b * x;
9
       return g;
10
11
12
  array<int, 2> exgcd(int a, int b, int c) {
13
      int x, y;
14
       int g = \__exgcd(a, b, x, y);
15
       if (c % g) return {INT_MAX, INT_MAX};
16
       int dx = b / g;
17
       int dy = a / g;
18
       x = c / g % dx * x % dx;
19
       if (x < 0) x += dx;
20
       y = (c - a * x) / b;
21
       return {x, y};
22 }
```

4.2 线性筛法

```
1 constexpr int N = 10000000;
2 array<int, N + 1> min_prime;
3
  vector<int> primes;
 4
  bool ok = []() {
5
       for (int i = 2; i <= N; i++) {</pre>
6
           if (min_prime[i] == 0) {
7
               min_prime[i] = i;
8
               primes.push_back(i);
9
10
           for (auto& j : primes) {
11
               if (j > min_prime[i] || j > N / i) break;
12
               min_prime[j * i] = j;
13
           }
14
15
       return 1;
```

```
16 }();
```

4.3 分解质因数

```
1
  auto getprimes(int n) {
2
       vector<array<int, 2>> res;
3
       for (auto& i : primes) {
4
           if (i > n / i) break;
5
           if (n % i == 0) {
6
               res.push_back({i, 0});
7
               while (n % i == 0) {
8
                    n /= i;
9
                    res.back()[1]++;
10
               }
11
           }
12
       }
13
       if (n > 1) res.push_back({n, 1});
14
       return res;
15 }
```

4.4 pollard rho

```
1 using LL = __int128_t;
3 random_device rd;
  mt19937 seed(rd());
4
5
6
  11 power(ll a, ll b, ll mod) {
7
       11 \text{ res} = 1;
8
       while (b) {
9
           if (b & 1) res = (LL)res * a % mod;
10
           a = (LL)a * a % mod;
11
           b >>= 1;
12
13
       return res;
14
  }
15
  bool isprime(ll n) {
16
17
       static array primes{2, 3, 5, 7, 11, 13, 17, 19, 23};
       static unordered_map<11, bool> S;
18
19
       if (n < 2) return 0;</pre>
20
       if (S.count(n)) return S[n];
21
       11 d = n - 1, r = 0;
22
       while (!(d & 1)) {
23
           r++;
24
           d >>= 1;
25
       }
26
       for (auto& a : primes) {
27
           if (a == n) return S[n] = 1;
           11 x = power(a, d, n);
28
```

```
if (x == 1 \mid \mid x == n - 1) continue;
29
30
           for (int i = 0; i < r - 1; i++) {
                x = (LL)x * x % n;
31
                if (x == n - 1) break;
32
33
           if (x != n - 1) return S[n] = 0;
34
35
36
       return S[n] = 1;
37
38
39
  11 pollard_rho(ll n) {
       11 s = 0, t = 0;
40
       11 c = seed() % (n - 1) + 1;
41
       11 \text{ val} = 1;
42
       for (int goal = 1;; goal *= 2, s = t, val = 1) {
43
44
           for (int step = 1; step <= goal; step++) {</pre>
45
                t = ((LL)t * t + c) % n;
                val = (LL)val * abs(t - s) % n;
46
                if (step % 127 == 0) {
47
                    ll g = gcd(val, n);
48
49
                    if (g > 1) return g;
                }
50
51
52
           ll g = gcd(val, n);
53
           if (g > 1) return g;
       }
54
55
56
  auto getprimes(ll n) {
57
       unordered_set<ll> S;
       auto get = [&](auto self, ll n) {
58
59
           if (n < 2) return;</pre>
60
           if (isprime(n)) {
61
                S.insert(n);
62
                return;
63
           }
64
           11 mx = pollard_rho(n);
65
           self(self, n / mx);
66
           self(self, mx);
67
       };
68
       get(get, n);
69
       return S;
70 }
```

4.5 组合数

```
constexpr int N = 2e5 + 1;
array<modint, N + 1> fac, ifac;
auto ok = []() {
  fac[0] = ifac[0] = 1;
  for (int i = 1; i <= N; i++) {
    fac[i] = fac[i - 1] * i;</pre>
```

```
ifac[i] = fac[i].inv();
8
9
       return true;
10 }();
11
12 modint C(int n, int m) {
      if (n < m) return 0;</pre>
13
       if (m == 0) return 1;
14
15
       if (n <= mod) return fac[n] * ifac[m] * ifac[n - m];</pre>
       // n >= mod 时需要这个
16
17
       return C(n % mod, m % mod) * C(n / mod, m / mod);
18
```

4.6 数论分块

求解形如 $\sum_{i=1}^n f(i)g(\lfloor \frac{n}{i} \rfloor)$ 的合式 $s(n) = \sum_{i=1}^n f(i)$

```
modint sqrt decomposition(int n) {
2
      auto s = [\&](int x) \{ return x; \};
3
      auto g = [&](int x) { return x; };
4
      modint res = 0;
5
      while (1 <= R) {
6
           int r = n / (n / 1);
7
           res = res + (s(r) - s(1 - 1)) * g(n / 1);
8
9
       }
10
       return res;
11 }
```

4.7 积性函数

4.7.1 定义

函数 f(n) 满足 f(1) = 1 且 $\forall x, y \in \mathbf{N}^*$, gcd(x, y) = 1 都有 f(xy) = f(x)f(y),则 f(n) 为积性函数。 函数 f(n) 满足 f(1) = 1 且 $\forall x, y \in \mathbf{N}^*$ 都有 f(xy) = f(x)f(y),则 f(n) 为完全积性函数。

4.7.2 例子

- 单位函数: $\varepsilon(n) = [n = 1]$ 。(完全积性)
- 恒等函数: $id_k(n) = n^k$ 。(完全积性)
- 常数函数: 1(n) = 1。(完全积性)
- 除数函数: $\sigma_k(n) = \sum_{d|n} d^k$ 。 $\sigma_0(n)$ 通常简记作 d(n) 或 $\tau(n)$, $\sigma_1(n)$ 通常简记作 $\sigma(n)$ 。
- 欧拉函数: $\varphi(n) = \sum_{i=1}^{n} [\gcd(i, n) = 1]$.
- 莫比乌斯函数: $\mu(n) = \begin{cases} 1 & n=1 \\ 0 & \exists d>1, d^2 \mid n \text{, 其中 } \omega(n) \text{ 表示 } n \text{ 的本质不同质因子个数, 它是} \\ (-1)^{\omega(n)} & \text{otherwise} \end{cases}$

一个加性函数。

4.8 狄利克雷卷积

对于两个数论函数 f(x) 和 g(x),则它们的狄利克雷卷积得到的结果 h(x) 定义为: $h(x) = \sum_{d|x} f(d)g\left(\frac{x}{d}\right) = \sum_{ab=x} f(a)g(b)$ 可以简记为: h = f * g。

4.8.1 性质

交換律: f*g=g*f。 结合律: (f*g)*h=f*(g*h)。 分配律: (f+g)*h=f*h+g*h。 等式的性质: f=g 的充要条件是 f*h=g*h,其中数论函数 h(x) 要满足 $h(1)\neq 0$ 。

4.8.2 例子

- $\varepsilon = \mu * 1 \iff \varepsilon(n) = \sum_{d|n} \mu(d)$
- $id = \varphi * 1 \iff id(n) = \sum_{d|n} \varphi(d)$
- $d=1*1 \iff d(n)=\sum_{d|n}1$
- $\sigma = id *1 \iff \sigma(n) = \sum_{d|n} d$
- $\varphi = \mu * id \iff \varphi(n) = \sum_{d|n} d \cdot \mu(\frac{n}{d})$

4.9 欧拉函数

```
1 array<int, N + 1> phi;
2 auto _ = []() {
3    iota(phi.begin() + 1, phi.end(), 1);
4    for (int i = 2; i <= N; i++) {
5        if (phi[i] == i)
6             for (int j = i; j <= N; j += i) phi[j] = phi[j] / i * (i - 1);
7    }
8    return true;
9 }();</pre>
```

4.10 莫比乌斯反演

4.10.1 莫比乌斯函数性质

- $\sum_{d|n} \mu(d) = \begin{cases} 1 & n=1 \\ 0 & n \neq 1 \end{cases}$, $\text{EF} \sum_{d|n} \mu(d) = \varepsilon(n), \ \mu * 1 = \varepsilon$
- $[\gcd(i,j) = 1] = \sum_{d \mid \gcd(i,j)} \mu(d)$

```
1 array<int, N + 1> miu;
2 array<bool, N + 1> ispr;
3 auto _ = []() {
    miu.fill(1);
    ispr.fill(1);
```

```
for (int i = 2; i <= N; i++) {
7
           if (!ispr[i]) continue;
8
           miu[i] = -1;
9
           for (int j = 2 * i; j <= N; j += i) {
10
               ispr[j] = 0;
11
               if ((j / i) % i == 0) miu[j] = 0;
12
               else miu[j] *= -1;
13
           }
14
15
       return true;
16 }();
```

4.10.2 莫比乌斯变换/反演

```
f(n) = \sum_{d|n} g(d),那么有 g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) = \sum_{n|d} \mu(\frac{d}{n}) f(d)。
用狄利克雷卷积表示则为 f = g * 1,有 g = f * \mu。
f \to g 称为莫比乌斯反演,g \to f 称为莫比乌斯反演。
```

4.11 杜教筛

杜教筛被用于处理一类数论函数的前缀和问题。对于数论函数 f,杜教筛可以在低于线性时间的复杂 度内计算 $S(n) = \sum_{i=1}^{n} f(i)$ 。

$$S(n) = \frac{\sum_{i=1}^{n} (f * g)(i) - \sum_{i=2}^{n} g(i)S\left(\left\lfloor \frac{n}{i} \right\rfloor\right)}{g(1)}$$

可以构造恰当的数论函数 g 使得:

- 可以快速计算 $\sum_{i=1}^{n} (f * g)(i)$ 。
- 可以快速计算 g 的单点值,用数论分块求解 $\sum_{i=2}^{n} g(i) S\left(\left|\frac{n}{i}\right|\right)$ 。

4.11.1 示例

```
11 sum phi(11 n) {
       if (n <= N) return sp[n];</pre>
3
       if (sp2.count(n)) return sp2[n];
4
       11 \text{ res} = 0, 1 = 2;
5
       while (1 <= n) {
6
            ll r = n / (n / 1);
7
            res = res + (r - 1 + 1) * sum_phi(n / 1);
8
            1 = r + 1;
9
       }
10
       return sp2[n] = (11)n * (n + 1) / 2 - res;
11 }
12 | 11 sum miu(11 n) {
13
       if (n <= N) return sm[n];</pre>
       if (sm2.count(n)) return sm2[n];
14
15
       11 \text{ res} = 0, 1 = 2;
16
       while (1 <= n) {
17
            ll r = n / (n / 1);
            res = res + (r - l + 1) * sum_miu(n / l);
18
```

4.12 盒子与球

n 个球, m 个盒

球同	盒同	可空	公式
✓	✓	✓	$f_{n,m} = f_{n-1,m-1} + f_{n-m,m}$
✓	✓	×	$f_{n-m,m}$
×	✓	✓	$\sum_{i=1}^{m} g_{n,i}$
×	✓	×	$g_{n,m} = g_{n-1,m-1} + m * g_{n-1,m}$
✓	×	✓	C_{n+m-1}^{m-1}
✓	×	×	C_{n-1}^{m-1}
×	×	✓	m^n
×	×	×	$m!*g_{n,m}$

扩展:

n 个相同的球,m 个不同的盒,每个盒子超过 k 个球,问方案数? 可以考虑容斥,f(d) 表示至少有 d 个盒子装了 > k 个球方案数,总方案数则为 $f(0) - f(1) + f(2) - \dots$

4.13 线性基

```
1 // 线性基
2
  struct basis {
3
      int rnk = 0;
      array<ull, 64> p{};
4
5
6
      // 将×插入此线性基中
7
      void insert(ull x) {
8
           for (int i = 63; i >= 0; i--) {
9
               if ((x >> i) & 1) {
10
                   if (p[i]) x ^= p[i];
11
                   else {
12
                       for (int j = 0; j < i; j++)
13
                           if (x >> j & 1) x ^= p[j];
14
                       for (int j = i + 1; j <= 63; j++)
15
                           if (p[j] >> i & 1) p[j] ^= x;
16
                       p[i] = x;
17
                       rnk++;
18
                       break;
19
                   }
20
```

```
21
          }
22
      }
23
      // 将另一个线性基插入此线性基中
24
25
      void insert(basis other) {
26
           for (int i = 0; i <= 63; i++) {
27
               if (!other.p[i]) continue;
               insert(other.p[i]);
28
29
          }
      }
30
31
32
      // 最大异或值
33
      ull max_basis() {
          ull res = 0;
34
          for (int i = 63; i >= 0; i--)
35
               if ((res ^ p[i]) > res) res ^= p[i];
36
37
          return res;
38
      }
39 };
```

4.14 矩阵快速幂

```
constexpr 11 mod = 2147493647;
2
   struct Mat {
3
       int n, m;
 4
       vector<vector<ll>> mat;
5
       Mat(int n, int m) : n(n), m(n), mat(n, vector<11>(m, 0)) {}
6
       Mat(vector<vector<1l>> mat) : n(mat.size()), m(mat[0].size()), mat(mat) {}
7
       Mat operator*(const Mat& other) {
8
           assert(m == other.n);
9
           Mat res(n, other.m);
10
           for (int i = 0; i < res.n; i++)</pre>
11
               for (int j = 0; j < res.m; j++)</pre>
12
                    for (int k = 0; k < m; k++)
13
                        res.mat[i][j] =
14
                             (res.mat[i][j] + mat[i][k] * other.mat[k][j] % mod) %
15
                            mod;
16
           return res;
17
       }
18
  };
19
  Mat ksm(Mat a, ll b) {
20
       assert(a.n == a.m);
21
       Mat res(a.n, a.m);
22
       for (int i = 0; i < res.n; i++) res.mat[i][i] = 1;</pre>
23
       while (b) {
24
           if (b & 1) res = res * a;
25
           b >>= 1;
26
           a = a * a;
27
28
       return res;
29
```

```
1 double eps = 1e-8;
  const double PI = acos(-1);
3 using T = 11;
 4
5
  template <typename T>
6
  int cmp(T a, T b) {
7
       return a != b ? a < b ? -1 : 1 : 0;
8
9
  int cmp(double a, double b) {
10
11
       double c = a - b;
12
       if (abs(c) < eps) return 0;</pre>
13
       return c < 0 ? -1 : 1;
14
  }
15
  // 向量
16
17
  struct vec {
18
      T x, y;
19
       vec(T_x = 0, T_y = 0) : x(_x), y(_y) {}
20
21
22
       double length2() const { return x * x + y * y; }
       double length() const { return sqrt(x * x + y * y); }
23
24
25
       // 与×轴正方向的夹角
26
       double angle() const {
27
           double angle = atan2(y, x);
           if (angle < 0) angle += 2 * PI;</pre>
28
29
           return angle;
30
       }
31
32
       // 逆时针旋转
33
       vec &rotate(const double &theta) {
34
           double tmp = x;
           x = x * cos(theta) - y * sin(theta);
35
           y = y * cos(theta) + tmp * sin(theta);
36
37
           return *this;
38
       }
39
       bool operator==(const vec &other) const {
40
41
           return !cmp(x, other.x) && !cmp(y, other.y);
42
       bool operator<(const vec &other) const {</pre>
43
44
           int tmp = cmp(angle(), other.angle());
           if (tmp) return tmp == -1 ? 0 : 1;
45
46
           tmp = cmp(x, other.x);
47
           return tmp == -1 ? 0 : 1;
48
       }
49
50
       vec operator+(const vec &other) const { return {x + other.x, y + other.y}; }
```

```
51
       vec operator-() const { return {-x, -y}; }
52
       vec operator-(const vec &other) const { return -other + (*this); }
53
       vec operator*(const T &other) const { return {x * other, y * other}; }
       vec operator/(const T &other) const { return {x / other, y / other}; }
54
55
       T operator*(const vec &other) const { return x * other.x + y * other.y; }
56
57
       // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
       // 等于0共线,可能同向或反向,结果绝对值表示 a b 形成的平行四边行的面积
58
       T operator^(const vec &other) const { return x * other.y - y * other.x; }
59
60
61
       friend istream &operator>>(istream &input, vec &data) {
62
           input >> data.x >> data.y;
63
           return input;
64
65
       friend ostream &operator<<(ostream &output, const vec &data) {</pre>
66
           output << fixed << setprecision(6);
67
           output << data.x << " " << data.y;</pre>
68
           return output;
69
       }
70
   };
71
72 bool xycmp(const vec &a, const vec &b) {
73
       int tmp = cmp(a.x, b.x);
74
       if (tmp) return tmp == -1 ? 0 : 1;
75
       tmp = cmp(a.y, b.y);
76
       return tmp == -1 ? 0 : 1;
77
78
79 T cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
80
   // 两点间的距离
81
82
   T distance(const vec &a, const vec &b) {
83
       return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
84
85
86
   // 两向量夹角
87 double angle(const vec &a, const vec &b) {
88
       double theta = abs(a.angle() - b.angle());
89
       if (theta > PI) theta = 2 * PI - theta;
90
       return theta;
91 }
92
93 // 判断点是否在凸包内
94
   bool in_polygon(const vec &a, vector<vec> &p) {
95
       int n = p.size();
96
       if (n == 1) return a == p[0];
97
       if (cross(a, p[1], p[0]) > 0 || cross(p.back(), a, p[0]) > 0) return 0;
       auto cmp = [&p](vec &x, const vec &y) { return ((x - p[0]) ^ y) >= 0; };
98
99
       int i = lower_bound(p.begin() + 2, p.end(), a - p[0], cmp) - p.begin() - 1;
100
       return cross(p[(i + 1) % n], a, p[i]) >= 0;
101
   }
102
```

```
103 // 多边形的面积
104
   double polygon_area(vector<vec> &p) {
105
       T area = 0;
       for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];</pre>
106
107
       area += p.back() ^ p[0];
108
       return abs(area / 2.0);
109
110
   // 多边形的周长
111
   double polygon_length(vector<vec> &p) {
112
       double length = 0;
113
       for (int i = 1; i < p.size(); i++) length += (p[i - 1] - p[i]).length();</pre>
114
115
       length += (p.back() - p[0]).length();
116
       return length;
117
118
119
   // 以整点为顶点的线段上的整点个数
120
   T count(const vec &a, const vec &b) {
121
       vec c = a - b;
       return gcd(abs(c.x), abs(c.y)) + 1;
122
123 }
124
125
   // 以整点为顶点的多边形边上整点个数
126
   T count(vector<vec> &p) {
127
       T cnt = 0;
128
       for (int i = 1; i < p.size(); i++) cnt += count(p[i - 1], p[i]);</pre>
129
       cnt += count(p.back(), p[0]);
130
       return cnt - p.size();
131
132
133
   // 凸包直径的两个端点
134
   auto polygon dia(vector<vec> &p) {
135
       int n = p.size();
136
       array<vec, 2> res{};
137
       if (n == 1) return res;
138
       if (n == 2) return res = {p[0], p[1]};
139
       T mx = 0;
       for (int i = 0, j = 2; i < n; i++) {
140
141
            while (abs(cross(p[i], p[(i + 1) % n], p[j])) <=
142
                   abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n]))
143
                j = (j + 1) \% n;
            if (T tmp = distance(p[i], p[j]); tmp > mx) {
144
145
                mx = tmp;
146
                res = \{p[i], p[j]\};
147
            if (T tmp = distance(p[(i + 1) % n], p[j]); tmp > mx) {
148
149
                mx = tmp;
150
                res = \{p[(i + 1) \% n], p[j]\};
151
            }
152
153
       return res;
154
```

```
155
156
   // 凸包
   auto convex_hull(vector<vec> &p) {
157
158
        sort(p.begin(), p.end(), xycmp);
159
        int n = p.size();
160
        vector sta(n + 1, 0);
161
        vector v(n, false);
162
        int tp = -1;
163
        sta[++tp] = 0;
        auto update = [&](int lim, int i) {
164
            while (tp > lim &&
165
                   ((p[sta[tp]] - p[sta[tp - 1]]) ^ (p[i] - p[sta[tp]])) <= 0)
166
167
                v[sta[tp--]] = 0;
168
            sta[++tp] = i;
169
            v[i] = 1;
170
        };
171
       for (int i = 1; i < n; i++) update(0, i);</pre>
172
        int cnt = tp;
173
        for (int i = n - 1; i >= 0; i --) {
            if (v[i]) continue;
174
175
            update(cnt, i);
176
        }
177
       vector<vec> res(tp);
178
        for (int i = 0; i < tp; i++) res[i] = p[sta[i]];</pre>
179
        return res;
180
181
   // 闵可夫斯基和 两个点集的和构成一个凸包
182
183
   auto minkowski(vector<vec> &a, vector<vec> &b) {
184
        rotate(a.begin(), min_element(a.begin(), a.end(), xycmp), a.end());
185
        rotate(b.begin(), min_element(b.begin(), b.end(), xycmp), b.end());
        int n = a.size(), m = b.size();
186
187
       vector<vec> c{a[0] + b[0]};
188
        c.reserve(n + m);
189
        int i = 0, j = 0;
190
        while (i < n && j < m) {</pre>
191
            vec x = a[(i + 1) \% n] - a[i];
192
            vec y = b[(j + 1) \% m] - b[j];
193
            c.push back(c.back() + ((x ^ y) >= 0 ? (i++, x) : (j++, y)));
194
        }
195
        while (i + 1 < n) {
196
            c.push_back(c.back() + a[(i + 1) % n] - a[i]);
197
            i++;
198
199
        while (j + 1 < m) {
200
            c.push_back(c.back() + b[(j + 1) \% m] - b[j]);
201
            j++;
202
203
       return c;
204
   }
205
206 // 过凸多边形外一点求凸多边形的切线,返回切点下标
```

```
auto tangent(const vec &a, vector<vec> &p) {
207
        int n = p.size();
208
       int 1 = -1, r = -1;
209
       for (int i = 0; i < n; i++) {</pre>
210
211
           T \text{ tmp1} = cross(p[i], p[(i - 1 + n) \% n], a);
212
           T \text{ tmp2} = cross(p[i], p[(i + 1) % n], a);
213
           if (1 == -1 \&\& tmp1 <= 0 \&\& tmp2 <= 0) 1 = i;
           else if (r == -1 \&\& tmp1 >= 0 \&\& tmp2 >= 0) r = i;
214
215
216
       return array{1, r};
217
218
219
   // 直线
220
   struct line {
221
       vec point, direction;
222
       line(const vec &p = vec(), const vec &d = vec()) : point(p), direction(d) {}
223
   };
224
225
   // 点到直线距离
226
   double distance(const vec &a, const line &b) {
227
       return abs((b.point - a) ^ (b.point + b.direction - a)) /
228
               b.direction.length();
229
230
231
   // 判断点在直线哪边,大于0在左边,等于0在线上,小于0在右边
232
   T side_line(const vec &a, const line &b) { return b.direction ^ (a - b.point); }
233
234
   // 两直线是否垂直
235
   bool perpendicular(const line &a, const line &b) {
236
       return !cmp(a.direction * b.direction, 0);
237
   }
238
239
   // 两直线是否平行
240
   bool parallel(const line &a, const line &b) {
241
       return !cmp(a.direction ^ b.direction, 0);
242
243
   // 点的垂线是否与线段有交点
244
245
   bool perpendicular(const vec &a, const line &b) {
       vec perpen(-b.direction.y, b.direction.x);
246
247
       bool cross1 = (perpen ^ (b.point - a)) > 0;
       bool cross2 = (perpen ^ (b.point + b.direction - a)) > 0;
248
249
       return cross1 != cross2;
250
251
252
   // 点到线段距离
253
   double distance seg(const vec &a, const line &b) {
       if (perpendicular(a, b)) return distance(a, b);
254
255
       return min(distance(a, b.point), distance(a, b.point + b.direction));
256
   }
257
258 // 两直线交点
```

```
vec intersection(T A, T B, T C, T D, T E, T F) {
259
       return {(B * F - C * E) / (A * E - B * D),
260
                (C * D - A * F) / (A * E - B * D);
261
262
263
264
   // 两直线交点
265
   vec intersection(const line &a, const line &b) {
266
       return intersection(a.direction.y, -a.direction.x,
267
                            a.direction.x * a.point.y - a.direction.y * a.point.x,
268
                            b.direction.y, -b.direction.x,
                            b.direction.x * b.point.y - b.direction.y * b.point.x);
269
270
271
272
   struct circle {
273
       vec o;
274
       double r;
275
       circle(const vec &_o, T _r) : o(_o), r(_r){};
276
       // 点与圆的关系 -1在圆内, 0在圆上, 1在圆外
277
       int relation(const vec &other) const {
            double len = (other - o).length();
278
279
            return cmp(len, r);
280
       }
281
       double area() { return PI * r * r; }
282
   };
283
284
   // 圆与直线交点
285
   auto intersection(const circle &c, const line &l) {
286
       double d = distance(c.o, 1);
287
       vector<vec> res;
288
       double len = 1.direction.length();
289
       vec mid = 1.point + 1.direction * ((c.o - 1.point) * 1.direction / len);
290
       if (!cmp(d, c.r)) res.push back(mid);
291
       else if (d < c.r) {</pre>
292
            d = sqrt(c.r * c.r - d * d) / len;
293
            res.push_back(mid + 1.direction * d);
294
            res.push_back(mid - 1.direction * d);
295
       }
       return res;
296
297
   }
298
299
   // oab三角形与圆相交的面积
   double area(const circle &c, const vec &a, const vec &b) {
300
301
       vec oa = a - c.o, ob = b - c.o;
302
       T cab = oa ^ ob;
303
       if (!cmp(cab, 0)) return 0;
304
       if (c.relation(a) != 1 && c.relation(b) != 1) return cab / 2.0;
305
       vec ba = a - b, bo = -ob;
306
       vec ab = -ba, ao = -oa;
307
       auto r = c.r;
308
       double ang;
309
       double loa = oa.length(), lob = ob.length(), lab = ab.length();
310
       double x =
```

```
311
            (ba * bo + sqrt(r * r * lab * lab - (ba ^ bo) * (ba ^ bo))) / lab;
312
        double y =
            (ab * ao + sqrt(r * r * lab * lab - (ab ^ ao) * (ab ^ ao))) / lab;
313
314
        if (cmp(lob, r) == -1 \&\& cmp(loa, r) != -1) {
315
            ang = cab * (1 - x / lab) / (r * loa);
316
            ang = min(max((double)-1, ang), (double)1);
317
            return (asin(ang) * r * r + cab * x / lab) / 2;
318
        }
        if (cmp(lob, r) != -1 \&\& cmp(loa, r) == -1) {
319
            ang = cab * (1 - y / lab) / (r * lob);
320
321
            ang = min(max((double)-1, ang), (double)1);
            return (asin(ang) * r * r + cab * y / lab) / 2;
322
323
       }
       if (cmp(abs(cab), r * lab) != -1 || cmp(ab * ao, 0) != 1 ||
324
325
            cmp(ba * bo, 0) != 1) {
326
            ang = cab / (loa * lob);
327
            ang = min(max((double)-1, ang), (double)1);
328
            double tmp = -asin(ang);
329
            if (cmp(oa * ob, 0) == -1)
                if (cmp(cab, 0) == -1) tmp -= PI;
330
331
                else tmp += PI;
332
            else tmp = -tmp;
333
            return tmp * r * r / 2;
334
        }
335
        ang = cab * (1 - x / lab) / (r * loa);
336
        ang = min(max((double)-1, ang), (double)1);
337
        double ang2 = cab * (1 - y / lab) / (r * lob);
338
        ang2 = min(max((double)-1, ang2), (double)1);
339
        return ((a\sin(ang) + a\sin(ang2)) * r * r + cab * ((x + y) / lab - 1)) / 2;
340
   }
341
342
   // 多边形与圆相交的面积
343 double area(vector<vec> &p, circle c) {
344
        double res = 0;
345
        for (int i = 1; i < p.size(); i++) res += area(c, p[i - 1], p[i]);</pre>
346
        res += area(c, p.back(), p[0]);
347
        return abs(res);
348
   }
```

5.1 扫描线

```
1 #define ls (pos << 1)
  #define rs (ls | 1)
3 #define mid ((tree[pos].l + tree[pos].r) >> 1)
4
  struct Rectangle {
5
      ll x_l, y_l, x_r, y_r;
6
  };
7
  11 area(vector<Rectangle>& rec) {
8
       struct Line {
9
           11 x, y_up, y_down;
10
           int pd;
```

```
11
       };
12
       vector<Line> line(rec.size() * 2);
13
       vector<ll> y_set(rec.size() * 2);
       for (int i = 0; i < rec.size(); i++) {</pre>
14
15
           y_set[i * 2] = rec[i].y_l;
16
           y_{set[i * 2 + 1] = rec[i].y_r;}
17
           line[i * 2] = {rec[i].x_l, rec[i].y_r, rec[i].y_l, 1};
           line[i * 2 + 1] = {rec[i].x_r, rec[i].y_r, rec[i].y_l, -1};
18
19
20
       sort(y_set.begin(), y_set.end());
21
       y_set.erase(unique(y_set.begin(), y_set.end()), y_set.end());
22
       sort(line.begin(), line.end(), [](Line a, Line b) { return a.x < b.x; });</pre>
23
       struct Data {
24
           int 1, r;
25
           11 len, cnt, raw_len;
26
       };
27
       vector<Data> tree(4 * y_set.size());
28
       function<void(int, int, int)> build = [&](int pos, int 1, int r) {
29
           tree[pos].l = 1;
30
           tree[pos].r = r;
31
           if (1 == r) {
32
               tree[pos].raw_len = y_set[r + 1] - y_set[l];
33
               tree[pos].cnt = tree[pos].len = 0;
34
               return;
35
36
           build(ls, l, mid);
37
           build(rs, mid + 1, r);
38
           tree[pos].raw_len = tree[ls].raw_len + tree[rs].raw_len;
39
       };
       function < void(int, int, int, int) > update = [&](int pos, int 1, int r,
40
41
                                                          int num) {
42
           if (1 <= tree[pos].1 && tree[pos].r <= r) {</pre>
43
               tree[pos].cnt += num;
44
               tree[pos].len = tree[pos].cnt ? tree[pos].raw_len
45
                                 : tree[pos].1 == tree[pos].r
46
47
                                     : tree[ls].len + tree[rs].len;
48
               return;
49
50
           if (1 <= mid) update(ls, 1, r, num);</pre>
51
           if (r > mid) update(rs, l, r, num);
52
           tree[pos].len =
53
               tree[pos].cnt ? tree[pos].raw_len : tree[ls].len + tree[rs].len;
54
       };
55
       build(1, 0, y_set.size() - 2);
56
       auto find_pos = [&](ll num) {
57
           return lower_bound(y_set.begin(), y_set.end(), num) - y_set.begin();
58
       };
59
       11 \text{ res} = 0;
60
       for (int i = 0; i < line.size() - 1; i++) {</pre>
61
           update(1, find_pos(line[i].y_down), find_pos(line[i].y_up) - 1,
62
                   line[i].pd);
```

6 杂项

6.1 高精度

```
1
  struct bignum {
2
       string num;
3
4
       bignum() : num("0") {}
5
       bignum(const string& num) : num(num) {
6
           reverse(this->num.begin(), this->num.end());
7
8
       bignum(ll num) : num(to_string(num)) {
9
           reverse(this->num.begin(), this->num.end());
10
       }
11
12
       bignum operator+(const bignum& other) {
13
           bignum res;
14
           res.num.pop_back();
15
           res.num.reserve(max(num.size(), other.num.size()) + 1);
16
           for (int i = 0, j = 0, x; i < num.size() || i < other.num.size() || j;
17
                i++) {
18
               x = j;
19
               j = 0;
20
               if (i < num.size()) x += num[i] - '0';</pre>
21
               if (i < other.num.size()) x += other.num[i] - '0';</pre>
22
               if (x >= 10) j = 1, x -= 10;
23
               res.num.push_back(x + '0');
24
25
           res.num.capacity();
26
           return res;
27
       }
28
29
       bignum operator*(const bignum& other) {
30
           vector<int> res(num.size() + other.num.size() - 1, 0);
           for (int i = 0; i < num.size(); i++)</pre>
31
32
               for (int j = 0; j < other.num.size(); j++)</pre>
33
                    res[i + j] += (num[i] - '0') * (other.num[j] - '0');
34
           int g = 0;
           for (int i = 0; i < res.size(); i++) {</pre>
35
36
               res[i] += g;
37
               g = res[i] / 10;
38
               res[i] %= 10;
39
40
           while (g) {
               res.push_back(g % 10);
41
42
               g /= 10;
43
44
           int lim = res.size();
           while (lim > 1 && res[lim - 1] == 0) lim--;
45
46
           bignum res2;
47
           res2.num.resize(lim);
48
           for (int i = 0; i < lim; i++) res2.num[i] = res[i] + '0';
```

```
49
            return res2;
50
       }
51
52
       bool operator<(const bignum& other) {</pre>
53
           if (num.size() == other.num.size())
54
                for (int i = num.size() - 1; i >= 0; i--)
55
                    if (num[i] == other.num[i]) continue;
                    else return num[i] < other.num[i];</pre>
56
            return num.size() < other.num.size();</pre>
57
       }
58
59
60
       friend istream& operator>>(istream& in, bignum& a) {
61
           in >> a.num;
62
            reverse(a.num.begin(), a.num.end());
63
           return in;
64
       }
65
       friend ostream& operator<<(ostream& out, bignum a) {</pre>
66
            reverse(a.num.begin(), a.num.end());
67
           return out << a.num;</pre>
68
       }
69 };
```

6.2 模运算

```
1
  struct modint {
 2
       int x;
3
       modint(11 _x = 0) : x(_x \% mod) {}
 4
       modint inv() const { return power(*this, mod - 2); }
5
       modint operator+(const modint& b) { return {x + b.x}; }
 6
       modint operator-() const { return {-x}; }
 7
       modint operator-(const modint& b) { return {-b + *this}; }
8
       modint operator*(const modint& b) { return {(11)x * b.x}; }
9
       modint operator/(const modint& b) { return *this * b.inv(); }
10
       friend istream& operator>>(istream& is, modint& other) {
11
           11 _x;
12
           is >> _x;
13
           other = modint(_x);
14
           return is;
15
       }
16
       friend ostream& operator<<(ostream& os, modint other) {</pre>
17
           other.x = (other.x + mod) % mod;
18
           return os << other.x;</pre>
19
       }
20 };
```

6.3 分数

```
1 struct frac {
2      11 a, b;
3      frac() : a(0), b(1) {}
```

```
frac(ll _a, ll _b) : a(_a), b(_b) {
 4
5
           assert(b);
6
           if (a) {
7
               int tmp = gcd(a, b);
8
               a /= tmp;
9
               b /= tmp;
10
           } else *this = frac();
11
12
       frac operator+(const frac& other) {
13
           return frac(a * other.b + other.a * b, b * other.b);
14
15
       frac operator-() const {
16
           frac res = *this;
17
           res.a = -res.a;
18
           return res;
19
       }
20
       frac operator-(const frac& other) const { return -other + *this; }
21
       frac operator*(const frac& other) const {
22
           return frac(a * other.a, b * other.b);
23
       }
24
       frac operator/(const frac& other) const {
25
           assert(other.a);
           return *this * frac(other.b, other.a);
26
27
       }
28
       bool operator<(const frac& other) const { return (*this - other).a < 0; }</pre>
       bool operator<=(const frac& other) const { return (*this - other).a <= 0; }</pre>
29
30
       bool operator>=(const frac& other) const { return (*this - other).a >= 0; }
31
       bool operator>(const frac& other) const { return (*this - other).a > 0; }
32
       bool operator==(const frac& other) const {
33
           return a == other.a && b == other.b;
34
35
       bool operator!=(const frac& other) const { return !(*this == other); }
36 };
```

6.4 表达式求值

```
1
  // 格式化表达式
  string format(const string& s1) {
3
       stringstream ss(s1);
4
       string s2;
5
       char ch;
6
       while ((ch = ss.get()) != EOF) {
           if (ch == ' ') continue;
7
8
           if (isdigit(ch)) s2 += ch;
9
           else {
10
               if (s2.back() != ' ') s2 += ' ';
11
               s2 += ch;
12
               s2 += ' ';
13
           }
14
       }
15
       return s2;
```

```
16 }
17
18
  // 中缀表达式转后缀表达式
19
  string convert(const string& s1) {
20
       unordered_map<char, int> rank{
21
           {'+', 2}, {'-', 2}, {'*', 1}, {'/', 1}, {'^', 0}};
22
       stringstream ss(s1);
23
       string s2, temp;
24
       stack<char> op;
       while (ss >> temp) {
25
26
           if (isdigit(temp[0])) s2 += temp + ' ';
           else if (temp[0] == '(') op.push('(');
27
           else if (temp[0] == ')') {
28
               while (op.top() != '(') {
29
30
                    s2 += op.top();
                    s2 += ' ';
31
32
                    op.pop();
33
               }
34
               op.pop();
35
           } else {
36
               while (!op.empty() && op.top() != '(' &&
                       (temp[0] != '^' && rank[op.top()] <= rank[temp[0]] ||</pre>
37
38
                        rank[op.top()] < rank[temp[0]])) {</pre>
39
                    s2 += op.top();
                    s2 += ' ';
40
41
                    op.pop();
42
43
               op.push(temp[0]);
44
           }
45
46
       while (!op.empty()) {
47
           s2 += op.top();
           s2 += ' ';
48
49
           op.pop();
50
51
       return s2;
52
53
54
  // 计算后缀表达式
55
  int calc(const string& s) {
56
       stack<int> num;
57
       stringstream ss(s);
58
       string temp;
59
       while (ss >> temp) {
60
           if (isdigit(temp[0])) num.push(stoi(temp));
61
           else {
62
               int b = num.top();
63
               num.pop();
64
               int a = num.top();
65
               num.pop();
66
               if (temp[0] == '+') a += b;
67
               else if (temp[0] == '-') a -= b;
```

```
else if (temp[0] == '*') a *= b;
else if (temp[0] == '/') a /= b;
else if (temp[0] == '^') a = ksm(a, b);
num.push(a);
}
return num.top();
}
```

6.5 日期

```
1 int month[] = {0, 31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31};
  int pre[13];
3
  vector<int> leap;
  struct Date {
4
5
       int y, m, d;
6
       bool operator<(const Date& other) const {</pre>
7
           return array<int, 3>{y, m, d} <</pre>
8
                  array<int, 3>{other.y, other.m, other.d};
9
10
       Date(const string& s) {
11
           stringstream ss(s);
12
           char ch;
13
           ss >> y >> ch >> m >> ch >> d;
14
15
       int dis() const {
16
           int yd = (y - 1) * 365 +
17
                     (upper_bound(leap.begin(), leap.end(), y - 1) - leap.begin());
18
           int md =
               pre[m - 1] + (m > 2 && (y % 4 == 0 && y % 100 || y % 400 == 0));
19
20
           return yd + md + d;
21
22
       int dis(const Date& other) const { return other.dis() - dis(); }
23 };
24
  for (int i = 1; i <= 12; i++) pre[i] = pre[i - 1] + month[2];</pre>
  for (int i = 1; i <= 1000000; i++)
26
       if (i % 4 == 0 && i % 100 || i % 400 == 0) leap.push_back(i);
```

6.6 对拍

linux/Mac

```
g++ a.cpp -o program/a -02 -std=c++17
g++ b.cpp -o program/b -02 -std=c++17
g++ suiji.cpp -o program/suiji -02 -std=c++17

cnt=0

while true; do
let cnt++
echo TEST:$cnt
```

```
10
11    ./program/suiji > in
12    ./program/a < in > out.a
13    ./program/b < in > out.b
14
15    diff out.a out.b
16    if [ $? -ne 0 ]; then break; fi
17 done
```

windows

```
@echo off
 2
3
  g++ a.cpp -o program/a -02 -std=c++17
  g++ b.cpp -o program/b -02 -std=c++17
5
  g++ suiji.cpp -o program/suiji -O2 -std=c++17
6
7
  set cnt=0
8
9
  :again
10
       set /a cnt=cnt+1
11
       echo TEST:%cnt%
12
       .\program\suiji > in
13
       .\program\a < in > out.a
       .\program\b < in > out.b
14
15
16
      fc output.a output.b
17 if not errorlevel 1 goto again
```

6.7 编译常用选项

```
1 -Wall -Woverflow -Wextra -Wpedantic -Wfloat-equal -Wshadow -fsanitize=address,undefined
```

6.8 开栈

不同的编译器可能命令不一样

```
1 -Wl,--stack=0x10000000
2 -Wl,-stack_size -Wl,0x10000000
3 -Wl,-z,stack-size=0x10000000
```

6.9 clang-format

```
BasedOnStyle: Google
IndentWidth: 4
ColumnLimit: 80
AllowShortIfStatementsOnASingleLine: AllIfsAndElse
AccessModifierOffset: -4
EmptyLineBeforeAccessModifier: Leave
RemoveBracesLLVM: true
```