

ACM 常用算法模板

therehello

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1 数据结构

1.1 并查集

```
1 struct dsu {
2     int n;
3     vector<int> fa, sz;
4     dsu(int \_n) : n(\_n), fa(n + 1), sz(n + 1, 1) {
5         iota(fa.begin(), fa.end(), 0);
6     }
7     int find(int x) { return x == fa[x] ? x : fa[x] = find(fa[x]); }
8     int merge(int x, int y) {
9         int fax = find(x), fay = find(y);
10        if (fax == fay) return 0; // 一个集合
11        sz[fay] += sz[fax];
12        return fa[fax] = fay; // 合并到哪个集合了
13    }
14    int size(int x) { return sz[find(x)]; }
15};
```

1.2 树状数组

1.2.1 一维

```
1 template <class T>
2 struct fenwick {
3     int n;
4     vector<T> t;
5     fenwick(int \_n) : n(\_n), t(n + 1) {}
6     T query(int l, int r) {
7         auto query = [&](int pos) {
8             T res = 0;
9             while (pos) {
10                res += t[pos];
11                pos -= lowbit(pos);
12            }
13            return res;
14        };
15        return query(r) - query(l - 1);
16    }
17    void add(int pos, T num) {
18        while (pos <= n) {
19            t[pos] += num;
20            pos += lowbit(pos);
21        }
22    }
23};
```

1.2.2 二维

```

1 template <class T>
2 struct Fenwick\_tree\_2 {
3     Fenwick\_tree\_2(int n, int m) : n(n), m(m), tree(n + 1, vector<T>(m + 1)) {}
4     T query(int l1, int r1, int l2, int r2) {
5         auto query = [&](int l, int r) {
6             T res = 0;
7             for (int i = l; i; i -= lowbit(i))
8                 for (int j = r; j; j -= lowbit(j)) res += tree[i][j];
9             return res;
10        };
11        return query(l2, r2) - query(l2, r1 - 1) - query(l1 - 1, r2) +
12            query(l1 - 1, r1 - 1);
13    }
14    void update(int x, int y, T num) {
15        for (int i = x; i <= n; i += lowbit(i))
16            for (int j = y; j <= m; j += lowbit(j)) tree[i][j] += num;
17    }
18 private:
19     int n, m;
20     vector<vector<T>> tree;
21 };

```

1.2.3 三维

```

1 template <class T>
2 struct Fenwick\_tree\_3 {
3     Fenwick\_tree\_3(int n, int m, int k)
4         : n(n),
5           m(m),
6           k(k),
7           tree(n + 1, vector<vector<T>>(m + 1, vector<T>(k + 1))) {}
8     T query(int a, int b, int c, int d, int e, int f) {
9         auto query = [&](int x, int y, int z) {
10             T res = 0;
11             for (int i = x; i; i -= lowbit(i))
12                 for (int j = y; j; j -= lowbit(j))
13                     for (int p = z; p; p -= lowbit(p)) res += tree[i][j][p];
14             return res;
15        };
16        T res = query(d, e, f);
17        res -= query(a - 1, e, f) + query(d, b - 1, f) + query(d, e, c - 1);
18        res += query(a - 1, b - 1, f) + query(a - 1, e, c - 1) +
19            query(d, b - 1, c - 1);
20        res -= query(a - 1, b - 1, c - 1);
21        return res;
22    }
23    void update(int x, int y, int z, T num) {
24        for (int i = x; i <= n; i += lowbit(i))
25            for (int j = y; j <= m; j += lowbit(j))
26                for (int p = z; p <= k; p += lowbit(p)) tree[i][j][p] += num;

```

```

27     }
28 private:
29     int n, m, k;
30     vector<vector<vector<T>>> tree;
31 };

```

1.3 线段树

```

1 template <class Data, class Num>
2 struct Segment_Tree {
3     inline void update(int l, int r, Num x) { update(1, l, r, x); }
4     inline Data query(int l, int r) { return query(1, l, r); }
5     Segment_Tree(vector<Data>& a) {
6         n = a.size();
7         tree.assign(n * 4 + 1, {});
8         build(a, 1, 1, n);
9     }
10 private:
11     int n;
12     struct Tree {
13         int l, r;
14         Data data;
15     };
16     vector<Tree> tree;
17     inline void pushup(int pos) {
18         tree[pos].data = tree[pos << 1].data + tree[pos << 1 | 1].data;
19     }
20     inline void pushdown(int pos) {
21         tree[pos << 1].data = tree[pos << 1].data + tree[pos].data.lazytag;
22         tree[pos << 1 | 1].data =
23             tree[pos << 1 | 1].data + tree[pos].data.lazytag;
24         tree[pos].data.lazytag = Num::zero();
25     }
26     void build(vector<Data>& a, int pos, int l, int r) {
27         tree[pos].l = l;
28         tree[pos].r = r;
29         if (l == r) {
30             tree[pos].data = a[l - 1];
31             return;
32         }
33         int mid = (tree[pos].l + tree[pos].r) >> 1;
34         build(a, pos << 1, l, mid);
35         build(a, pos << 1 | 1, mid + 1, r);
36         pushup(pos);
37     }
38     void update(int pos, int& l, int& r, Num& x) {
39         if (l > tree[pos].r || r < tree[pos].l) return;
40         if (l <= tree[pos].l && tree[pos].r <= r) {
41             tree[pos].data = tree[pos].data + x;
42             return;
43         }

```

```

44     pushdown(pos);
45     update(pos << 1, 1, r, x);
46     update(pos << 1 | 1, 1, r, x);
47     pushup(pos);
48 }
49 Data query(int pos, int& l, int& r) {
50     if (l > tree[pos].r || r < tree[pos].l) return Data::zero();
51     if (l <= tree[pos].l && tree[pos].r <= r) return tree[pos].data;
52     pushdown(pos);
53     return query(pos << 1, l, r) + query(pos << 1 | 1, l, r);
54 }
55 };
56 struct Num {
57     ll add;
58     inline static Num zero() { return {0}; }
59     inline Num operator+(Num b) { return {add + b.add}; }
60 };
61 struct Data {
62     ll sum, len;
63     Num lazytag;
64     inline static Data zero() { return {0, 0, Num::zero()}; }
65     inline Data operator+(Num b) {
66         return {sum + len * b.add, len, lazytag + b};
67     }
68     inline Data operator+(Data b) {
69         return {sum + b.sum, len + b.len, Num::zero()};
70     }
71 };

```

1.4 普通平衡树

1.4.1 树状数组实现

需要预先处理出来所有可能的数。

```

1 int lowbit(int x) { return x & -x; }
2
3 template <typename T>
4 struct treap {
5     int n, size;
6     vector<int> t;
7     vector<T> t2, S;
8     treap(const vector<T>& a) : S(a) {
9         sort(S.begin(), S.end());
10        S.erase(unique(S.begin(), S.end()), S.end());
11        n = S.size();
12        size = 0;
13        t = vector<int>(n + 1);
14        t2 = vector<T>(n + 1);
15    }
16    int pos(T x) { return lower\_bound(S.begin(), S.end(), x) - S.begin() + 1; }
17    int sum(int pos) {

```



```

18     int res = 0;
19     while (pos) {
20         res += t[pos];
21         pos -= lowbit(pos);
22     }
23     return res;
24 }
25
26 // 插入cnt个x
27 void insert(T x, int cnt) {
28     size += cnt;
29     int i = pos(x);
30     assert(i <= n && S[i - 1] == x);
31     for (; i <= n; i += lowbit(i)) {
32         t[i] += cnt;
33         t2[i] += cnt * x;
34     }
35 }
36
37 // 删除cnt个x
38 void erase(T x, int cnt) {
39     assert(cnt <= count(x));
40     insert(x, -cnt);
41 }
42
43 // x的排名
44 int rank(T x) {
45     assert(count(x));
46     return sum(pos(x) - 1) + 1;
47 }
48
49 // 统计出现次数
50 int count(T x) { return sum(pos(x)) - sum(pos(x) - 1); }
51
52 // 第k小
53 T kth(int k) {
54     assert(0 < k && k <= size);
55     int cnt = 0, x = 0;
56     for (int i = __lg(n); i >= 0; i--) {
57         x += 1 << i;
58         if (x >= n || cnt + t[x] >= k) x -= 1 << i;
59         else cnt += t[x];
60     }
61     return S[x];
62 }
63
64 // 前k小的数之和
65 T pre_sum(int k) {
66     assert(0 < k && k <= size);
67     int cnt = 0, x = 0;
68     T res = 0;
69     for (int i = __lg(n); i >= 0; i--) {

```

```

70         x += 1 << i;
71         if (x >= n || cnt + t[x] >= k) x -= 1 << i;
72         else {
73             cnt += t[x];
74             res += t2[x];
75         }
76     }
77     return res + (k - cnt) * S[x];
78 }
79
80 // 小于x, 最大的数
81 T prev(T x) { return kth(sum(pos(x) - 1)); }
82
83 // 大于x, 最小的数
84 T next(T x) { return kth(sum(pos(x)) + 1); }
85 };

```

1.5 可持久化线段树

```

1 constexpr int MAXN = 200000;
2 vector<int> root(MAXN << 5);
3 struct Persistent\_seg {
4     int n;
5     struct Data {
6         int ls, rs;
7         int val;
8     };
9     vector<Data> tree;
10    Persistent\_seg(int n, vector<int>& a) : n(n) { root[0] = build(1, n, a); }
11    int build(int l, int r, vector<int>& a) {
12        if (l == r) {
13            tree.push\_back({0, 0, a[l]});
14            return tree.size() - 1;
15        }
16        int mid = l + r >> 1;
17        int ls = build(l, mid, a), rs = build(mid + 1, r, a);
18        tree.push\_back({ls, rs, tree[ls].val + tree[rs].val});
19        return tree.size() - 1;
20    }
21    int update(int rt, const int& idx, const int& val, int l, int r) {
22        if (l == r) {
23            tree.push\_back({0, 0, tree[rt].val + val});
24            return tree.size() - 1;
25        }
26        int mid = l + r >> 1, ls = tree[rt].ls, rs = tree[rt].rs;
27        if (idx <= mid) ls = update(ls, idx, val, l, mid);
28        else rs = update(rs, idx, val, mid + 1, r);
29        tree.push\_back({ls, rs, tree[ls].val + tree[rs].val});
30        return tree.size() - 1;
31    }
32    int query(int rt1, int rt2, int k, int l, int r) {

```

```

33     if (l == r) return l;
34     int mid = l + r >> 1;
35     int lcnt = tree[tree[rt2].ls].val - tree[tree[rt1].ls].val;
36     if (k <= lcnt) return query(tree[rt1].ls, tree[rt2].ls, k, l, mid);
37     else return query(tree[rt1].rs, tree[rt2].rs, k - lcnt, mid + 1, r);
38 }
39 };

```

1.6 st 表

```

1 auto lg = []() {
2     array<int, 10000001> lg;
3     lg[1] = 0;
4     for (int i = 2; i <= 100000000; i++) lg[i] = lg[i >> 1] + 1;
5     return lg;
6 }();
7 template <typename T>
8 struct st {
9     int n;
10    vector<vector<T>>> a;
11    st(vector<T>& _a) : n(_a.size()) {
12        a.assign(lg[n] + 1, vector<int>(n));
13        for (int i = 0; i < n; i++) a[0][i] = _a[i];
14        for (int j = 1; j <= lg[n]; j++)
15            for (int i = 0; i + (1 << j) - 1 < n; i++)
16                a[j][i] = max(a[j - 1][i], a[j - 1][i + (1 << (j - 1))]);
17    }
18    T query(int l, int r) {
19        int k = lg[r - l + 1];
20        return max(a[k][l], a[k][r - (1 << k) + 1]);
21    }
22 };

```

2 图论

存图

```

1 struct Graph {
2     int n;
3     struct Edge {
4         int to, w;
5     };
6     vector<vector<Edge>> graph;
7     Graph(int \_n) {
8         n = \_n;
9         graph.assign(n + 1, vector<Edge>());
10    };
11    void add(int u, int v, int w) { graph[u].push\_back({v, w}); }
12 };

```

2.1 最短路

2.1.1 dijkstra

```

1 void dij(Graph& graph, vector<int>& dis, int t) {
2     vector<int> visit(graph.n + 1, 0);
3     priority\_queue<pair<int, int>> que;
4     dis[t] = 0;
5     que.emplace(0, t);
6     while (!que.empty()) {
7         int u = que.top().second;
8         que.pop();
9         if (visit[u]) continue;
10        visit[u] = 1;
11        for (auto& [to, w] : graph.graph[u]) {
12            if (dis[to] > dis[u] + w) {
13                dis[to] = dis[u] + w;
14                que.emplace(-dis[to], to);
15            }
16        }
17    }
18 }

```

2.2 树上问题

2.2.1 最近公公祖先

倍增法

```

1 vector<int> dep;
2 vector<array<int, 21>> fa;
3 dep.assign(n + 1, 0);
4 fa.assign(n + 1, array<int, 21>{});
5 void binary\_jump(int root) {
6     function<void(int)> dfs = [&](int t) {

```

```

7     dep[t] = dep[fa[t][0]] + 1;
8     for (auto& [to] : graph[t]) {
9         if (to == fa[t][0]) continue;
10        fa[to][0] = t;
11        dfs(to);
12    }
13 };
14 dfs(root);
15 for (int j = 1; j <= 20; j++)
16     for (int i = 1; i <= n; i++) fa[i][j] = fa[fa[i][j - 1]][j - 1];
17 }
18 int lca(int x, int y) {
19     if (dep[x] < dep[y]) swap(x, y);
20     for (int i = 20; i >= 0; i--)
21         if (dep[fa[x][i]] >= dep[y]) x = fa[x][i];
22     if (x == y) return x;
23     for (int i = 20; i >= 0; i--) {
24         if (fa[x][i] != fa[y][i]) {
25             x = fa[x][i];
26             y = fa[y][i];
27         }
28     }
29     return fa[x][0];
30 }

```

树剖

```

1 int lca(int x, int y) {
2     while (top[x] != top[y]) {
3         if (dep[top[x]] < dep[top[y]]) swap(x, y);
4         x = fa[top[x]];
5     }
6     if (dep[x] < dep[y]) swap(x, y);
7     return y;
8 }

```

2.2.2 树链剖分

```

1 vector<int> fa, siz, dep, son, dfn, rnk, top;
2 fa.assign(n + 1, 0);
3 siz.assign(n + 1, 0);
4 dep.assign(n + 1, 0);
5 son.assign(n + 1, 0);
6 dfn.assign(n + 1, 0);
7 rnk.assign(n + 1, 0);
8 top.assign(n + 1, 0);
9 void hld(int root) {
10     function<void(int)> dfs1 = [&](int t) {
11         dep[t] = dep[fa[t]] + 1;
12         siz[t] = 1;
13         for (auto& [to, w] : graph[t]) {
14             if (to == fa[t]) continue;

```

```

15         fa[to] = t;
16         dfs1(to);
17         if (siz[son[t]] < siz[to]) son[t] = to;
18         siz[t] += siz[to];
19     }
20 };
21 dfs1(root);
22 int dfn\_tail = 0;
23 for (int i = 1; i <= n; i++) top[i] = i;
24 function<void(int)> dfs2 = [&](int t) {
25     dfn[t] = ++dfn\_tail;
26     rnk[dfn\_tail] = t;
27     if (!son[t]) return;
28     top[son[t]] = top[t];
29     dfs2(son[t]);
30     for (auto& [to, w] : graph[t]) {
31         if (to == fa[t] || to == son[t]) continue;
32         dfs2(to);
33     }
34 };
35 dfs2(root);
36 }

```

2.3 强连通分量

```

1 void tarjan(Graph& g1, Graph& g2) {
2     int dfn\_tail = 0, cnt = 0;
3     vector<int> dfn(g1.n + 1, 0), low(g1.n + 1, 0), exist(g1.n + 1, 0),
4         belong(g1.n + 1, 0);
5     stack<int> sta;
6     function<void(int)> dfs = [&](int t) {
7         dfn[t] = low[t] = ++dfn\_tail;
8         sta.push(t);
9         exist[t] = 1;
10        for (auto& [to] : g1.graph[t])
11            if (!dfn[to]) {
12                dfs(to);
13                low[t] = min(low[t], low[to]);
14            } else if (exist[to]) low[t] = min(low[t], dfn[to]);
15        if (dfn[t] == low[t]) {
16            cnt++;
17            while (int temp = sta.top()) {
18                belong[temp] = cnt;
19                exist[temp] = 0;
20                sta.pop();
21                if (temp == t) break;
22            }
23        }
24    };
25    for (int i = 1; i <= g1.n; i++)
26        if (!dfn[i]) dfs(i);

```

```
27 g2 = Graph(cnt);
28 for (int i = 1; i <= g1.n; i++) g2.w[belong[i]] += g1.w[i];
29 for (int i = 1; i <= g1.n; i++)
30     for (auto& [to] : g1.graph[i])
31         if (belong[i] != belong[to]) g2.add(belong[i], belong[to]);
32 }
```

2.4 拓扑排序

```
1 void toposort(Graph& g, vector<int>& dis) {
2     vector<int> in(g.n + 1, 0);
3     for (int i = 1; i <= g.n; i++)
4         for (auto& [to] : g.graph[i]) in[to]++;
5     queue<int> que;
6     for (int i = 1; i <= g.n; i++)
7         if (!in[i]) {
8             que.push(i);
9             dis[i] = g.w[i]; // dp
10        }
11    while (!que.empty()) {
12        int u = que.front();
13        que.pop();
14        for (auto& [to] : g.graph[u]) {
15            in[to]--;
16            dis[to] = max(dis[to], dis[u] + g.w[to]); // dp
17            if (!in[to]) que.push(to);
18        }
19    }
20 }
```

3 字符串

3.1 kmp

```

1 auto kmp(string& s) {
2     vector next(s.size(), -1);
3     for (int i = 1, j = -1; i < s.size(); i++) {
4         while (j >= 0 && s[i] != s[j + 1]) j = next[j];
5         if (s[i] == s[j + 1]) j++;
6         next[i] = j;
7     }
8     // next 意为长度
9     for (auto& i : next) i++;
10    return next;
11 }

```

3.2 哈希

```

1 constexpr int N = 1e6;
2 int pow\_base[N + 1][2];
3 constexpr ll mod[2] = {(int)2e9 + 11, (int)2e9 + 33},
4                     base[2] = {(int)2e5 + 11, (int)2e5 + 33};
5
6 struct Hash {
7     int size;
8     vector<array<int, 2>> a;
9     Hash() {}
10    Hash(const string& s) {
11        size = s.size();
12        a.resize(size);
13        a[0][0] = a[0][1] = s[0];
14        for (int i = 1; i < size; i++) {
15            a[i][0] = (a[i - 1][0] * base[0] + s[i]) % mod[0];
16            a[i][1] = (a[i - 1][1] * base[1] + s[i]) % mod[1];
17        }
18    }
19    array<int, 2> get(int l, int r) const {
20        if (l == 0) return a[r];
21        auto getone = [&](bool f) {
22            int x =
23                (a[r][f] - 11l * a[l - 1][f] * pow\_base[r - l + 1][f]) % mod[f];
24            if (x < 0) x += mod[f];
25            return x;
26        };
27        return {getone(0), getone(1)};
28    }
29 };
30
31 auto \_ = []() {
32     pow\_base[0][0] = pow\_base[0][1] = 1;
33     for (int i = 1; i <= N; i++) {

```



```
34     pow\_base[i][0] = pow\_base[i - 1][0] * base[0] % mod[0];
35     pow\_base[i][1] = pow\_base[i - 1][1] * base[1] % mod[1];
36 }
37 return true;
38 }();
```

3.3 manacher

```
1 auto manacher(const string& \_s) {
2     string s(\_s.size() * 2 + 1, '$');
3     for (int i = 0; i < \_s.size(); i++) s[2 * i + 1] = \_s[i];
4     vector r(s.size(), 0);
5     for (int i = 0, maxr = 0, mid = 0; i < s.size(); i++) {
6         if (i < maxr) r[i] = min(r[mid * 2 - i], maxr - i);
7         while (i - r[i] - 1 >= 0 && i + r[i] + 1 < s.size() &&
8             s[i - r[i] - 1] == s[i + r[i] + 1])
9             ++r[i];
10        if (i + r[i] > maxr) maxr = i + r[i], mid = i;
11    }
12    return r;
13 }
```

4 数学

4.1 扩展欧几里得

需保证 $a, b \geq 0$

$$x = x + k * dx, y = y - k * dy$$

若要求 $x \geq p, k \geq \lceil \frac{p-x}{dx} \rceil$

若要求 $x \leq q, k \leq \lfloor \frac{q-x}{dx} \rfloor$

若要求 $y \geq p, k \leq \lfloor \frac{y-p}{dy} \rfloor$

若要求 $y \leq q, k \geq \lceil \frac{y-q}{dy} \rceil$

```

1 int \_\_exgcd(int a, int b, int& x, int& y) {
2     if (!b) {
3         x = 1;
4         y = 0;
5         return a;
6     }
7     int g = \_\_exgcd(b, a % b, y, x);
8     y -= a / b * x;
9     return g;
10 }
11
12 array<int, 2> exgcd(int a, int b, int c) {
13     int x, y;
14     int g = \_\_exgcd(a, b, x, y);
15     if (c % g) return {INT\_MAX, INT\_MAX};
16     int dx = b / g;
17     int dy = a / g;
18     x = c / g % dx * x % dx;
19     if (x < 0) x += dx;
20     y = (c - a * x) / b;
21     return {x, y};
22 }

```

4.2 线性筛法

```

1 constexpr int N = 10000000;
2 array<int, N + 1> min\_prime;
3 vector<int> primes;
4 bool ok = []() {
5     for (int i = 2; i <= N; i++) {
6         if (min\_prime[i] == 0) {
7             min\_prime[i] = i;
8             primes.push\_back(i);
9         }
10        for (auto& j : primes) {
11            if (j > min\_prime[i] || j > N / i) break;
12            min\_prime[j * i] = j;
13        }
14    }
15    return 1;

```

```
16 }();
```

4.3 分解质因数

```
1 auto getprimes(int n) {
2     vector<array<int, 2>> res;
3     for (auto& i : primes) {
4         if (i > n / i) break;
5         if (n % i == 0) {
6             res.push_back({i, 0});
7             while (n % i == 0) {
8                 n /= i;
9                 res.back()[1]++;
10            }
11        }
12    }
13    if (n > 1) res.push_back({n, 1});
14    return res;
15 }
```

4.4 pollard rho

```
1 using LL = __int128_t;
2
3 random_device rd;
4 mt19937 seed(rd());
5
6 ll power(ll a, ll b, ll mod) {
7     ll res = 1;
8     while (b) {
9         if (b & 1) res = (LL)res * a % mod;
10        a = (LL)a * a % mod;
11        b >>= 1;
12    }
13    return res;
14 }
15
16 bool isprime(ll n) {
17     static array primes{2, 3, 5, 7, 11, 13, 17, 19, 23};
18     static unordered_map<ll, bool> S;
19     if (n < 2) return 0;
20     if (S.count(n)) return S[n];
21     ll d = n - 1, r = 0;
22     while (!(d & 1)) {
23         r++;
24         d >>= 1;
25     }
26     for (auto& a : primes) {
27         if (a == n) return S[n] = 1;
28         ll x = power(a, d, n);
```

```

29     if (x == 1 || x == n - 1) continue;
30     for (int i = 0; i < r - 1; i++) {
31         x = (LL)x * x % n;
32         if (x == n - 1) break;
33     }
34     if (x != n - 1) return S[n] = 0;
35 }
36 return S[n] = 1;
37 }
38
39 ll pollard_rho(ll n) {
40     ll s = 0, t = 0;
41     ll c = seed() % (n - 1) + 1;
42     ll val = 1;
43     for (int goal = 1;; goal *= 2, s = t, val = 1) {
44         for (int step = 1; step <= goal; step++) {
45             t = ((LL)t * t + c) % n;
46             val = (LL)val * abs(t - s) % n;
47             if (step % 127 == 0) {
48                 ll g = gcd(val, n);
49                 if (g > 1) return g;
50             }
51         }
52         ll g = gcd(val, n);
53         if (g > 1) return g;
54     }
55 }
56 auto getprimes(ll n) {
57     unordered_set<ll> S;
58     auto get = [&](auto self, ll n) {
59         if (n < 2) return;
60         if (isprime(n)) {
61             S.insert(n);
62             return;
63         }
64         ll mx = pollard_rho(n);
65         self(self, n / mx);
66         self(self, mx);
67     };
68     get(get, n);
69     return S;
70 }

```

4.5 组合数

```

1 constexpr int N = 1e6;
2 array<modint, N + 1> fac, ifac;
3
4 modint C(int n, int m) {
5     if (n < m) return 0;
6     if (n <= mod) return fac[n] * ifac[m] * ifac[n - m];

```

```

7 // n >= mod 时需要这个
8 return C(n % mod, m % mod) * C(n / mod, m / mod);
9 }
10
11 auto \_ = []() {
12     fac[0] = 1;
13     for (int i = 1; i <= N; i++) fac[i] = fac[i - 1] * i;
14     ifac[N] = fac[N].inv();
15     for (int i = N - 1; i >= 0; i--) ifac[i] = ifac[i + 1] * (i + 1);
16     return true;
17 }();

```

4.6 数论分块

求解形如 $\sum_{i=1}^n f(i)g(\lfloor \frac{n}{i} \rfloor)$ 的合式

$$s(n) = \sum_{i=1}^n f(i)$$

```

1 modint sqrt\_decomposition(int n) {
2     auto s = [&](int x) { return x; };
3     auto g = [&](int x) { return x; };
4     modint res = 0;
5     while (l <= R) {
6         int r = n / (n / l);
7         res = res + (s(r) - s(l - 1)) * g(n / l);
8         l = r + 1;
9     }
10    return res;
11 }

```

4.7 积性函数

4.7.1 定义

函数 $f(n)$ 满足 $f(1) = 1$ 且 $\forall x, y \in \mathbf{N}^*, \gcd(x, y) = 1$ 都有 $f(xy) = f(x)f(y)$, 则 $f(n)$ 为积性函数。
 函数 $f(n)$ 满足 $f(1) = 1$ 且 $\forall x, y \in \mathbf{N}^*$ 都有 $f(xy) = f(x)f(y)$, 则 $f(n)$ 为完全积性函数。

4.7.2 例子

- 单位函数: $\varepsilon(n) = [n = 1]$ 。(完全积性)
- 恒等函数: $\text{id}_k(n) = n^k$ 。(完全积性)
- 常数函数: $1(n) = 1$ 。(完全积性)
- 除数函数: $\sigma_k(n) = \sum_{d|n} d^k$ 。 $\sigma_0(n)$ 通常简记作 $d(n)$ 或 $\tau(n)$, $\sigma_1(n)$ 通常简记作 $\sigma(n)$ 。
- 欧拉函数: $\varphi(n) = \sum_{i=1}^n [\gcd(i, n) = 1]$ 。

- 莫比乌斯函数: $\mu(n) = \begin{cases} 1 & n = 1 \\ 0 & \exists d > 1, d^2 | n, \text{ 其中 } \omega(n) \text{ 表示 } n \text{ 的本质不同质因子个数, 它是} \\ (-1)^{\omega(n)} & \text{otherwise} \end{cases}$
 一个加性函数。

4.8 狄利克雷卷积

对于两个数论函数 $f(x)$ 和 $g(x)$ ，则它们的狄利克雷卷积得到的结果 $h(x)$ 定义为：

$$h(x) = \sum_{d|x} f(d)g\left(\frac{x}{d}\right) = \sum_{ab=x} f(a)g(b)$$

可以简记为： $h = f * g$ 。

4.8.1 性质

交换律： $f * g = g * f$ 。

结合律： $(f * g) * h = f * (g * h)$ 。

分配律： $(f + g) * h = f * h + g * h$ 。

等式的性质： $f = g$ 的充要条件是 $f * h = g * h$ ，其中数论函数 $h(x)$ 要满足 $h(1) \neq 0$ 。

4.8.2 例子

- $\varepsilon = \mu * 1 \iff \varepsilon(n) = \sum_{d|n} \mu(d)$
- $id = \varphi * 1 \iff id(n) = \sum_{d|n} \varphi(d)$
- $d = 1 * 1 \iff d(n) = \sum_{d|n} 1$
- $\sigma = id * 1 \iff \sigma(n) = \sum_{d|n} nd$
- $\varphi = \mu * id \iff \varphi(n) = \sum_{d|n} nd \cdot \mu\left(\frac{n}{d}\right)$

4.9 欧拉函数

```

1 constexpr int N = 1e6;
2 array<int, N + 1> phi;
3 auto \_ = []() {
4     iota(phi.begin() + 1, phi.end(), 1);
5     for (int i = 2; i <= N; i++) {
6         if (phi[i] == i)
7             for (int j = i; j <= N; j += i) phi[j] = phi[j] / i * (i - 1);
8     }
9     return true;
10 }();

```

4.10 莫比乌斯反演

4.10.1 莫比乌斯函数性质

- $\sum_{d|n} \mu(d) = \begin{cases} 1 & n = 1 \\ 0 & n \neq 1 \end{cases}$ ，即 $\sum_{d|n} \mu(d) = \varepsilon(n)$ ， $\mu * 1 = \varepsilon$
- $[\gcd(i, j) = 1] = \sum_{d|\gcd(i, j)} \mu(d)$

```

1 constexpr int N = 1e6;
2 array<int, N + 1> miu;
3 array<bool, N + 1> ispr;
4
5 auto \_ = []() {

```

```

6   miu.fill(1);
7   ispr.fill(1);
8   for (int i = 2; i <= N; i++) {
9       if (!ispr[i]) continue;
10      miu[i] = -1;
11      for (int j = 2 * i; j <= N; j += i) {
12          ispr[j] = 0;
13          if ((j / i) % i == 0) miu[j] = 0;
14          else miu[j] *= -1;
15      }
16  }
17  return true;
18 }();

```

4.10.2 莫比乌斯变换/反演

$f(n) = \sum_{d|n} g(d)$, 那么有 $g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) = \sum_{d|n} d \mu(\frac{d}{n}) f(d)$ 。

用狄利克雷卷积表示则为 $f = g * 1$, 有 $g = f * \mu$ 。

$f \rightarrow g$ 称为莫比乌斯反演, $g \rightarrow f$ 称为莫比乌斯反演。

4.11 杜教筛

杜教筛被用于处理一类数论函数的前缀和问题。对于数论函数 f , 杜教筛可以在低于线性时间的复杂度内计算 $S(n) = \sum_{i=1}^n f(i)$ 。

$$S(n) = \frac{\sum_{i=1}^n (f * g)(i) - \sum_{i=2}^n g(i) S(\lfloor \frac{n}{i} \rfloor)}{g(1)}$$

可以构造恰当的数论函数 g 使得:

- 可以快速计算 $\sum_{i=1}^n (f * g)(i)$ 。
- 可以快速计算 g 的单点值, 用数论分块求解 $\sum_{i=2}^n g(i) S(\lfloor \frac{n}{i} \rfloor)$ 。

4.11.1 示例

```

1 ll sum_phi(ll n) {
2     if (n <= N) return sp[n];
3     if (sp2.count(n)) return sp2[n];
4     ll res = 0, l = 2;
5     while (l <= n) {
6         ll r = n / (n / l);
7         res = res + (r - l + 1) * sum_phi(n / l);
8         l = r + 1;
9     }
10    return sp2[n] = (ll)n * (n + 1) / 2 - res;
11 }
12
13 ll sum_miu(ll n) {
14     if (n <= N) return sm[n];
15     if (sm2.count(n)) return sm2[n];
16     ll res = 0, l = 2;

```

```

17 while (l <= n) {
18     ll r = n / (n / l);
19     res = res + (r - l + 1) * sum\_miu(n / l);
20     l = r + 1;
21 }
22 return sm2[n] = 1 - res;
23 }

```

4.12 多项式

```

1 #define countr\_zero(n) \_\_builtin\_ctz(n)
2 constexpr int N = 1e6;
3 array<int, N + 1> inv;
4
5 int power(int a, int b) {
6     int res = 1;
7     while (b) {
8         if (b & 1) res = 1ll * res * a % mod;
9         a = 1ll * a * a % mod;
10        b >>= 1;
11    }
12    return res;
13 }
14
15 namespace NFTS {
16     int g = 3;
17     vector<int> rev, roots{0, 1};
18     void dft(vector<int> &a) {
19         int n = a.size();
20         if (rev.size() != n) {
21             int k = countr\_zero(n) - 1;
22             rev.resize(n);
23             for (int i = 0; i < n; ++i) rev[i] = rev[i >> 1] >> 1 | (i & 1) << k;
24         }
25         if (roots.size() < n) {
26             int k = countr\_zero(roots.size());
27             roots.resize(n);
28             while ((1 << k) < n) {
29                 int e = power(g, (mod - 1) >> (k + 1));
30                 for (int i = 1 << (k - 1); i < (1 << k); ++i) {
31                     roots[2 * i] = roots[i];
32                     roots[2 * i + 1] = 1ll * roots[i] * e % mod;
33                 }
34                 ++k;
35             }
36         }
37         for (int i = 0; i < n; ++i)
38             if (rev[i] < i) swap(a[i], a[rev[i]]);
39         for (int k = 1; k < n; k *= 2) {
40             for (int i = 0; i < n; i += 2 * k) {
41                 for (int j = 0; j < k; ++j) {

```



```

42         int u = a[i + j];
43         int v = 111 * a[i + j + k] * roots[k + j] % mod;
44         int x = u + v, y = u - v;
45         if (x >= mod) x -= mod;
46         if (y < 0) y += mod;
47         a[i + j] = x;
48         a[i + j + k] = y;
49     }
50 }
51 }
52 }
53 void idft(vector<int> &a) {
54     int n = a.size();
55     reverse(a.begin() + 1, a.end());
56     dft(a);
57     int inv_n = power(n, mod - 2);
58     for (int i = 0; i < n; ++i) a[i] = 111 * a[i] * inv_n % mod;
59 }
60 } // namespace NFTS
61
62 struct poly {
63     poly &format() {
64         while (!a.empty() && a.back() == 0) a.pop_back();
65         return *this;
66     }
67     poly &reverse() {
68         ::reverse(a.begin(), a.end());
69         return *this;
70     }
71     vector<int> a;
72     poly() {}
73     poly(int x) {
74         if (x) a = {x};
75     }
76     poly(const vector<int> &\_a) : a(\_a) {}
77     int size() const { return a.size(); }
78     int &operator[](int id) { return a[id]; }
79     int at(int id) const {
80         if (id < 0 || id >= (int)a.size()) return 0;
81         return a[id];
82     }
83     poly operator-() const {
84         auto A = *this;
85         for (auto &x : A.a) x = (x == 0 ? 0 : mod - x);
86         return A;
87     }
88     poly mulXn(int n) const {
89         auto b = a;
90         b.insert(b.begin(), n, 0);
91         return poly(b);
92     }
93     poly modXn(int n) const {

```

```

94     if (n > size()) return *this;
95     return poly({a.begin(), a.begin() + n});
96 }
97 poly divXn(int n) const {
98     if (size() <= n) return poly();
99     return poly({a.begin() + n, a.end()});
100 }
101 poly &operator+=(const poly &rhs) {
102     if (size() < rhs.size()) a.resize(rhs.size());
103     for (int i = 0; i < rhs.size(); ++i)
104         if ((a[i] += rhs.a[i]) >= mod) a[i] -= mod;
105     return *this;
106 }
107 poly &operator-=(const poly &rhs) {
108     if (size() < rhs.size()) a.resize(rhs.size());
109     for (int i = 0; i < rhs.size(); ++i)
110         if ((a[i] -= rhs.a[i]) < 0) a[i] += mod;
111     return *this;
112 }
113 poly &operator*=(poly rhs) {
114     int n = size(), m = rhs.size(), tot = max(1, n + m - 1);
115     int sz = 1 << _lg(tot * 2 - 1);
116     a.resize(sz);
117     rhs.a.resize(sz);
118     NFTS::dft(a);
119     NFTS::dft(rhs.a);
120     for (int i = 0; i < sz; ++i) a[i] = 1ll * a[i] * rhs.a[i] % mod;
121     NFTS::idft(a);
122     return *this;
123 }
124 poly &operator/=(poly rhs) {
125     int n = size(), m = rhs.size();
126     if (n < m) return (*this) = poly();
127     reverse();
128     rhs.reverse();
129     (*this) *= rhs.inv(n - m + 1);
130     a.resize(n - m + 1);
131     reverse();
132     return *this;
133 }
134 poly &operator%=(poly rhs) { return (*this) -= (*this) / rhs * rhs; }
135 poly operator+(const poly &rhs) const { return poly(*this) += rhs; }
136 poly operator-(const poly &rhs) const { return poly(*this) -= rhs; }
137 poly operator*(poly rhs) const { return poly(*this) *= rhs; }
138 poly operator/(poly rhs) const { return poly(*this) /= rhs; }
139 poly operator%(poly rhs) const { return poly(*this) %= rhs; }
140 poly powModPoly(int n, poly p) {
141     poly r(1), x(*this);
142     while (n) {
143         if (n & 1) (r *= x) %= p;
144         (x *= x) %= p;
145         n >>= 1;

```

```

146     }
147     return r;
148 }
149 int inner(const poly &rhs) {
150     int r = 0, n = min(size(), rhs.size());
151     for (int i = 0; i < n; ++i) r = (r + 111 * a[i] * rhs.a[i]) % mod;
152     return r;
153 }
154 poly derivation() const {
155     if (a.empty()) return poly();
156     int n = size();
157     vector<int> r(n - 1);
158     for (int i = 1; i < n; ++i) r[i - 1] = 111 * a[i] * i % mod;
159     return poly(r);
160 }
161 poly integral() const {
162     if (a.empty()) return poly();
163     int n = size();
164     vector<int> r(n + 1);
165     for (int i = 0; i < n; ++i) r[i + 1] = 111 * a[i] * ::inv[i + 1] % mod;
166     return poly(r);
167 }
168 poly inv(int n) const {
169     assert(a[0] != 0);
170     poly x(power(a[0], mod - 2));
171     int k = 1;
172     while (k < n) {
173         k *= 2;
174         x = (poly(2) - modXn(k) * x).modXn(k);
175     }
176     return x.modXn(n);
177 }
178 // 需要保证首项为 1
179 poly log(int n) const {
180     return (derivation() * inv(n)).integral().modXn(n);
181 }
182 // 需要保证首项为 0
183 poly exp(int n) const {
184     poly x(1);
185     int k = 1;
186     while (k < n) {
187         k *= 2;
188         x = (x * (poly(1) - x.log(k) + modXn(k))).modXn(k);
189     }
190     return x.modXn(n);
191 }
192 // 需要保证首项为 1, 开任意次方可以先 ln 再 exp 实现。
193 poly sqrt(int n) const {
194     poly x(1);
195     int k = 1;
196     while (k < n) {
197         k *= 2;

```

```

198         x += modXn(k) * x.inv(k);
199         x = x.modXn(k) * inv2;
200     }
201     return x.modXn(n);
202 }
203 // 减法卷积, 也称转置卷积  $\{ \rm MULT \} (F(x), G(x)) = \sum_{i \geq 0} (\sum_{j \geq 0} f_{i+j} g_j) x^i$ 
204 //  $0 \} f_{i+j} g_j x^i$ 
205 poly mulT(poly rhs) const {
206     if (rhs.size() == 0) return poly();
207     int n = rhs.size();
208     ::reverse(rhs.a.begin(), rhs.a.end());
209     return ((*this) * rhs).divXn(n - 1);
210 }
211 int eval(int x) {
212     int r = 0, t = 1;
213     for (int i = 0, n = size(); i < n; ++i) {
214         r = (r + 1ll * a[i] * t) % mod;
215         t = 1ll * t * x % mod;
216     }
217     return r;
218 }
219 // 多点求值新科技: https://jklloverdcoi.github.io/2020/08/04/转置原理及其应用/
220 // 模板例题: https://www.luogu.com.cn/problem/P5050
221 auto evals(vector<int> &x) const {
222     if (size() == 0) return vector(x.size(), 0);
223     int n = x.size();
224     vector ans(n, 0);
225     vector<poly> g(4 * n);
226     auto build = [&](auto self, int l, int r, int p) -> void {
227         if (r - l == 1) {
228             g[p] = poly({1, x[l] ? mod - x[l] : 0});
229         } else {
230             int m = (l + r) / 2;
231             self(self, l, m, 2 * p);
232             self(self, m, r, 2 * p + 1);
233             g[p] = g[2 * p] * g[2 * p + 1];
234         }
235     };
236     build(build, 0, n, 1);
237     auto solve = [&](auto self, int l, int r, int p, poly f) -> void {
238         if (r - l == 1) {
239             ans[l] = f[0];
240         } else {
241             int m = (l + r) / 2;
242             self(self, l, m, 2 * p, f.mulT(g[2 * p + 1]).modXn(m - l));
243             self(self, m, r, 2 * p + 1, f.mulT(g[2 * p]).modXn(r - m));
244         }
245     };
246     solve(solve, 0, n, 1, mulT(g[1].inv(size())).modXn(n));
247     return ans;
248 }
249 }; // 全家桶测试: https://www.luogu.com.cn/training/3015#information

```

```

250
251 auto \_ = []() {
252     inv[0] = inv[1] = 1;
253     for (int i = 2; i < inv.size(); i++)
254         inv[i] = 1ll * (mod - mod / i) * inv[mod % i] % mod;
255     return true;
256 }();

```

4.13 盒子与球

n 个球, m 个盒

球同	盒同	可空	公式
✓	✓	✓	$f_{n,m} = f_{n,m-1} + f_{n-m,m}$ 或 $[x^n]e^{\sum_{i=1}^m \sum_{j=1}^{\infty} \frac{x^{ij}}{j}}$
✓	✓	✗	$f_{n-m,m}$
✗	✓	✓	$\sum_{i=1}^m g_{n,i}$ 或 $\sum_{i=1}^m \sum_{j=0}^i \frac{j^n}{j!} \frac{(-1)^{i-j}}{(i-j)!}$
✗	✓	✗	$g_{n,m} = g_{n-1,m-1} + m * g_{n-1,m}$ 或 $\frac{1}{m!} \sum_{i=0}^m (-1)^i \binom{m}{i} (m-i)^n$
✓	✗	✓	C_{n+m-1}^{m-1}
✓	✗	✗	C_{n-1}^{m-1}
✗	✗	✓	m^n
✗	✗	✗	$m! * g_{n,m}$ 或 $\sum_{i=0}^m (-1)^i \binom{m}{i} (m-i)^n$

4.13.1 球同, 盒同, 可空

```

1 int solve(int n, int m) {
2     vector a(n + 1, 0);
3     for (int i = 1; i <= m; i++)
4         for (int j = i, k = 1; j <= n; j += i, k++)
5             a[j] = (a[j] + inv[k]) % mod;
6     auto p = poly(a).exp(n + 1);
7     return (p.a[n] + mod) % mod;
8 }

```

若要求不超过 k 个, 答案为 $[x^n y^m] \prod_{i=0}^k (1 + \sum_{j=0}^m x^{ij} y^j)$ 。

4.13.2 球不同, 盒同, 可空

```

1 int solve(int n, int m) {
2     vector a(n + 1, 0);
3     vector b(n + 1, 0);
4     for (int i = 0; i <= n; i++) {
5         a[i] = ifac[i];
6         if (i & 1) a[i] = -a[i];
7         b[i] = 1ll * power(i, n) * ifac[i] % mod;
8     }
9     auto p = poly(a) * poly(b);
10    int ans = 0;
11    for (int i = 1; i <= min(n, m); i++) ans = (ans + p.a[i]) % mod;
12    return (ans + mod) % mod;
13 }

```

若要求不超过 k 个, 答案为 $m! \cdot [x^n y^m] \prod_{i=0}^k (\sum_{j=0}^n \frac{1}{i! j!} x^{ij} y^j)$ 。

4.13.3 球同, 盒不同, 可空

若要求不超过 k 个, 答案为 $[x^n] (\sum_{i=0}^k x^i)^m = [x^n] \frac{(x^{k+1}-1)^m}{(x-1)^m}$ 。

也可以考虑容斥, 令 $f(i)$ 表示至少有 i 个盒子装了 $> k$ 个球方案数, $f(i) = \binom{m}{i} \binom{n-(k+1)i+m-1}{m-1}$ 。

总方案数则为 $\sum_{i=0}^m (-1)^i f(i)$ 。

4.13.4 球同, 盒不同, 不可空

若要求不超过 k 个, 答案为 $[x^n] (\sum_{i=1}^k x^i)^m = [x^n] \frac{(x^k-1)^m x^m}{(x-1)^m}$ 。

也可以考虑容斥, 令 $f(i)$ 表示至少有 i 个盒子装了 $> k$ 个球方案数, $f(i) = \binom{m}{i} \binom{n-ki-1}{m-i}$ 。

总方案数则为 $\sum_{i=0}^m (-1)^i f(i)$ 。

4.13.5 球不同, 盒不同, 可空

若要求不超过 k 个, 答案为 $m! \cdot [x^n] (\sum_{i=0}^k \frac{1}{i!} x^i)^m$ 。

4.13.6 球不同, 盒不同, 不可空

若要求不超过 k 个, 答案为 $m! \cdot [x^n] (\sum_{i=1}^k \frac{1}{i!} x^i)^m$ 。

4.14 线性基

```

1 // 线性基
2 struct basis {
3     int rnk = 0;
4     array<ull, 64> p{};
5
6     // 将x插入此线性基中
7     void insert(ull x) {
8         for (int i = 63; i >= 0; i--) {
9             if (!(x >> i & 1)) continue;
10            if (p[i] * x ^= p[i];
11            else {
12                p[i] = x;
13                rnk++;

```

```

14         break;
15     }
16 }
17 }
18
19 // 将另一个线性基插入此线性基中
20 void insert(basis other) {
21     for (int i = 0; i <= 63; i++) {
22         if (!other.p[i]) continue;
23         insert(other.p[i]);
24     }
25 }
26
27 // 最大异或值
28 ull max\_basis() {
29     ull res = 0;
30     for (int i = 63; i >= 0; i--)
31         if ((res ^ p[i]) > res) res ^= p[i];
32     return res;
33 }
34 };

```

4.15 矩阵快速幂

```

1 constexpr ll mod = 2147493647;
2 struct Mat {
3     int n, m;
4     vector<vector<ll>> mat;
5     Mat(int n, int m) : n(n), m(m), mat(n, vector<ll>(m, 0)) {}
6     Mat(vector<vector<ll>> mat) : n(mat.size()), m(mat[0].size()), mat(mat) {}
7     Mat operator*(const Mat& other) {
8         assert(m == other.n);
9         Mat res(n, other.m);
10        for (int i = 0; i < res.n; i++)
11            for (int j = 0; j < res.m; j++)
12                for (int k = 0; k < m; k++)
13                    res.mat[i][j] =
14                        (res.mat[i][j] + mat[i][k] * other.mat[k][j] % mod) %
15                        mod;
16        return res;
17    }
18 };
19 Mat ksm(Mat a, ll b) {
20     assert(a.n == a.m);
21     Mat res(a.n, a.m);
22     for (int i = 0; i < res.n; i++) res.mat[i][i] = 1;
23     while (b) {
24         if (b & 1) res = res * a;
25         b >>= 1;
26         a = a * a;
27     }

```

```
28 |     return res;  
29 | }
```


5 计算几何

5.1 整数

```

1 constexpr double inf = 1e100;
2
3 // 向量
4 struct vec {
5     static bool cmp(const vec &a, const vec &b) {
6         return tie(a.x, a.y) < tie(b.x, b.y);
7     }
8
9     ll x, y;
10    vec() : x(0), y(0) {}
11    vec(ll _x, ll _y) : x(_x), y(_y) {}
12
13    // 模
14    ll len2() const { return x * x + y * y; }
15    double len() const { return sqrt(x * x + y * y); }
16
17    // 是否在上半轴
18    bool up() const { return y > 0 || y == 0 && x >= 0; }
19
20    bool operator==(const vec &b) const { return tie(x, y) == tie(b.x, b.y); }
21    // 极角排序
22    bool operator<(const vec &b) const {
23        if (up() != b.up()) return up() > b.up();
24        ll tmp = (*this) ^ b;
25        return tmp ? tmp > 0 : cmp(*this, b);
26    }
27
28    vec operator+(const vec &b) const { return {x + b.x, y + b.y}; }
29    vec operator-() const { return {-x, -y}; }
30    vec operator-(const vec &b) const { return -b + (*this); }
31    vec operator*(ll b) const { return {x * b, y * b}; }
32    ll operator*(const vec &b) const { return x * b.x + y * b.y; }
33
34    // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
35    // 等于0共线, 可能同向或反向, 结果绝对值表示 a b 形成的平行四边形的面积
36    ll operator^(const vec &b) const { return x * b.y - y * b.x; }
37
38    friend istream &operator>>(istream &in, vec &data) {
39        in >> data.x >> data.y;
40        return in;
41    }
42    friend ostream &operator<<(ostream &out, const vec &data) {
43        out << fixed << setprecision(6);
44        out << data.x << " " << data.y;
45        return out;
46    }
47 };
48

```

```

49 ll cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
50
51 // 多边形的面积a
52 double polygon\_area(vector<vec> &p) {
53     ll area = 0;
54     for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];
55     area += p.back() ^ p[0];
56     return abs(area / 2.0);
57 }
58
59 // 多边形的周长
60 double polygon\_len(vector<vec> &p) {
61     double len = 0;
62     for (int i = 1; i < p.size(); i++) len += (p[i - 1] - p[i]).len();
63     len += (p.back() - p[0]).len();
64     return len;
65 }
66
67 // 以整点为顶点的线段上的整点个数
68 ll count(const vec &a, const vec &b) {
69     vec c = a - b;
70     return gcd(abs(c.x), abs(c.y)) + 1;
71 }
72
73 // 以整点为顶点的多边形边上整点个数
74 ll count(vector<vec> &p) {
75     ll cnt = 0;
76     for (int i = 1; i < p.size(); i++) cnt += count(p[i - 1], p[i]);
77     cnt += count(p.back(), p[0]);
78     return cnt - p.size();
79 }
80
81 // 判断点是否在凸包内，凸包必须为逆时针顺序
82 bool in\_polygon(const vec &a, vector<vec> &p) {
83     int n = p.size();
84     if (n == 0) return 0;
85     if (n == 1) return a == p[0];
86     if (n == 2)
87         return cross(a, p[1], p[0]) == 0 && (p[0] - a) * (p[1] - a) <= 0;
88     if (cross(a, p[1], p[0]) > 0 || cross(p.back(), a, p[0]) > 0) return 0;
89     auto cmp = [&](vec &x, const vec &y) { return ((x - p[0]) ^ y) >= 0; };
90     int i =
91         lower\_bound(p.begin() + 2, p.end() - 1, a - p[0], cmp) - p.begin() - 1;
92     return cross(p[(i + 1) % n], a, p[i]) >= 0;
93 }
94
95 // 凸包直径的两个端点
96 auto polygon\_dia(vector<vec> &p) {
97     int n = p.size();
98     array<vec, 2> res{};
99     if (n == 1) return res;
100    if (n == 2) return res = {p[0], p[1]};

```

```

101     ll mx = 0;
102     for (int i = 0, j = 2; i < n; i++) {
103         while (abs(cross(p[i], p[(i + 1) % n], p[j])) <=
104             abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n])))
105             j = (j + 1) % n;
106         ll tmp = (p[i] - p[j]).len2();
107         if (tmp > mx) {
108             mx = tmp;
109             res = {p[i], p[j]};
110         }
111         tmp = (p[(i + 1) % n] - p[j]).len2();
112         if (tmp > mx) {
113             mx = tmp;
114             res = {p[(i + 1) % n], p[j]};
115         }
116     }
117     return res;
118 }
119
120 // 凸包
121 auto convex_hull(vector<vec> &p) {
122     sort(p.begin(), p.end(), vec::cmp);
123     int n = p.size();
124     vector sta(n + 1, 0);
125     vector v(n, false);
126     int tp = -1;
127     sta[++tp] = 0;
128     auto update = [&](int lim, int i) {
129         while (tp > lim && cross(p[i], p[sta[tp]], p[sta[tp - 1]]) >= 0)
130             v[sta[tp--]] = 0;
131         sta[++tp] = i;
132         v[i] = 1;
133     };
134     for (int i = 1; i < n; i++) update(0, i);
135     int cnt = tp;
136     for (int i = n - 1; i >= 0; i--) {
137         if (v[i]) continue;
138         update(cnt, i);
139     }
140     vector<vec> res(tp);
141     for (int i = 0; i < tp; i++) res[i] = p[sta[i]];
142     return res;
143 }
144
145 // 闵可夫斯基和，两个点集的和构成一个凸包
146 auto minkowski(vector<vec> &a, vector<vec> &b) {
147     rotate(a.begin(), min_element(a.begin(), a.end(), vec::cmp), a.end());
148     rotate(b.begin(), min_element(b.begin(), b.end(), vec::cmp), b.end());
149     int n = a.size(), m = b.size();
150     vector<vec> c{a[0] + b[0]};
151     c.reserve(n + m);
152     int i = 0, j = 0;

```

```

153     while (i < n && j < m) {
154         vec x = a[(i + 1) % n] - a[i];
155         vec y = b[(j + 1) % m] - b[j];
156         c.push_back(c.back() + ((x ^ y) >= 0 ? (i++, x) : (j++, y)));
157     }
158     while (i + 1 < n) {
159         c.push_back(c.back() + a[(i + 1) % n] - a[i]);
160         i++;
161     }
162     while (j + 1 < m) {
163         c.push_back(c.back() + b[(j + 1) % m] - b[j]);
164         j++;
165     }
166     return c;
167 }
168
169 // 过凸多边形外一点求凸多边形的切线，返回切点下标
170 auto tangent(const vec &a, vector<vec> &p) {
171     int n = p.size();
172     int l = -1, r = -1;
173     for (int i = 0; i < n; i++) {
174         ll tmp1 = cross(p[i], p[(i - 1 + n) % n], a);
175         ll tmp2 = cross(p[i], p[(i + 1) % n], a);
176         if (l == -1 && tmp1 <= 0 && tmp2 <= 0) l = i;
177         else if (r == -1 && tmp1 >= 0 && tmp2 >= 0) r = i;
178     }
179     return array{l, r};
180 }
181
182 // 直线
183 struct line {
184     vec p, d;
185     line() : p(vec()), d(vec()) {}
186     line(const vec &\_p, const vec &\_d) : p(\_p), d(\_d) {}
187 };
188
189 // 点到直线距离
190 double dis(const vec &a, const line &b) {
191     return abs((b.p - a) ^ (b.p + b.d - a)) / b.d.len();
192 }
193
194 // 点在直线哪边，大于0在左边，等于0在线上，小于0在右边
195 ll side_line(const vec &a, const line &b) { return b.d ^ (a - b.p); }
196
197 // 两直线是否垂直
198 bool perpen(const line &a, const line &b) { return a.d * b.d == 0; }
199
200 // 两直线是否平行
201 bool parallel(const line &a, const line &b) { return (a.d ^ b.d) == 0; }
202
203 // 点的垂线是否与线段有交点
204 bool perpen(const vec &a, const line &b) {

```

```

205     vec p(-b.d.y, b.d.x);
206     bool cross1 = (p ^ (b.p - a)) > 0;
207     bool cross2 = (p ^ (b.p + b.d - a)) > 0;
208     return cross1 != cross2;
209 }
210
211 // 点到线段距离
212 double dis_seg(const vec &a, const line &b) {
213     if (perpen(a, b)) return dis(a, b);
214     return min((b.p - a).len(), (b.p + b.d - a).len());
215 }
216
217 // 点到凸包距离
218 double dis(const vec &a, vector<vec> &p) {
219     double res = inf;
220     for (int i = 1; i < p.size(); i++)
221         res = min(dis_seg(a, line(p[i - 1], p[i] - p[i - 1])), res);
222     res = min(dis_seg(a, line(p.back(), p[0] - p.back())), res);
223     return res;
224 }
225
226 // 两直线交点
227 vec intersection(ll A, ll B, ll C, ll D, ll E, ll F) {
228     return {(B * F - C * E) / (A * E - B * D),
229             (C * D - A * F) / (A * E - B * D)};
230 }
231
232 // 两直线交点
233 vec intersection(const line &a, const line &b) {
234     return intersection(a.d.y, -a.d.x, a.d.x * a.p.y - a.d.y * a.p.x, b.d.y,
235                        -b.d.x, b.d.x * b.p.y - b.d.y * b.p.x);
236 }

```

5.2 浮点数

```

1  using lf = double;
2
3  constexpr lf eps = 1e-8;
4  constexpr lf inf = 1e100;
5  const lf PI = acos(-1);
6
7  int sgn(lf a, lf b) {
8      lf c = a - b;
9      return c < -eps ? -1 : c < eps ? 0 : 1;
10 }
11
12 // 向量
13 struct vec {
14     static bool cmp(const vec &a, const vec &b) {
15         return sgn(a.x, b.x) ? a.x < b.x : sgn(a.y, b.y) < 0;
16     }

```

```

17
18     lf x, y;
19     vec() : x(0), y(0) {}
20     vec(lf \_x, lf \_y) : x(\_x), y(\_y) {}
21
22     // 模
23     lf len2() const { return x * x + y * y; }
24     lf len() const { return sqrt(x * x + y * y); }
25
26     // 与x轴正方向的夹角
27     lf angle() const {
28         lf angle = atan2(y, x);
29         if (angle < 0) angle += 2 * PI;
30         return angle;
31     }
32
33     // 逆时针旋转
34     vec rotate(const lf &theta) const {
35         return {x * cos(theta) - y * sin(theta),
36                y * cos(theta) + x * sin(theta)};
37     }
38
39     vec e() const {
40         lf tmp = len();
41         return {x / tmp, y / tmp};
42     }
43
44     // 是否在上半轴
45     bool up() const {
46         return sgn(y, 0) > 0 || sgn(y, 0) == 0 && sgn(x, 0) >= 0;
47     }
48
49     bool operator==(const vec &other) const {
50         return sgn(x, other.x) == 0 && sgn(y, other.y) == 0;
51     }
52     // 极角排序
53     bool operator<(const vec &b) const {
54         if (up() != b.up()) return up() > b.up();
55         lf tmp = (*this) ^ b;
56         return sgn(tmp, 0) ? tmp > 0 : cmp(*this, b);
57     }
58
59     vec operator+(const vec &b) const { return {x + b.x, y + b.y}; }
60     vec operator-() const { return {-x, -y}; }
61     vec operator-(const vec &b) const { return -b + (*this); }
62     vec operator*(lf b) const { return {x * b, y * b}; }
63     vec operator/(lf b) const { return {x / b, y / b}; }
64     lf operator*(const vec &b) const { return x * b.x + y * b.y; }
65
66     // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
67     // 等于0共线, 可能同向或反向, 结果绝对值表示 a b 形成的平行四边形的面积
68     lf operator^(const vec &b) const { return x * b.y - y * b.x; }

```

```

69
70     friend istream &operator>>(istream &in, vec &data) {
71         in >> data.x >> data.y;
72         return in;
73     }
74     friend ostream &operator<<(ostream &out, const vec &data) {
75         out << fixed << setprecision(6);
76         out << data.x << " " << data.y;
77         return out;
78     }
79 };
80
81 lf cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
82
83 lf angle(const vec &a, const vec &b) { return atan2(abs(a ^ b), a * b); }
84
85 // 多边形的面积
86 lf polygon\_area(vector<vec> &p) {
87     lf area = 0;
88     for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];
89     area += p.back() ^ p[0];
90     return abs(area / 2.0);
91 }
92
93 // 多边形的周长
94 lf polygon\_len(vector<vec> &p) {
95     lf len = 0;
96     for (int i = 1; i < p.size(); i++) len += (p[i - 1] - p[i]).len();
97     len += (p.back() - p[0]).len();
98     return len;
99 }
100
101 // 判断点是否在凸包内，凸包必须为逆时针顺序
102 bool in\_polygon(const vec &a, vector<vec> &p) {
103     int n = p.size();
104     if (n == 0) return 0;
105     if (n == 1) return a == p[0];
106     if (n == 2)
107         return sgn(cross(a, p[1], p[0]), 0) == 0 &&
108             sgn((p[0] - a) * (p[1] - a), 0) <= 0;
109     if (sgn(cross(a, p[1], p[0]), 0) > 0 ||
110         sgn(cross(p.back(), a, p[0]), 0) > 0)
111         return 0;
112     auto cmp = [&](vec &x, const vec &y) {
113         return sgn((x - p[0]) ^ y, 0) >= 0;
114     };
115     int i =
116         lower\_bound(p.begin() + 2, p.end() - 1, a - p[0], cmp) - p.begin() - 1;
117     return sgn(cross(p[(i + 1) % n], a, p[i]), 0) >= 0;
118 }
119
120 // 凸包直径的两个端点

```

```

121 auto polygon\_dia(vector<vec> &p) {
122     int n = p.size();
123     array<vec, 2> res{};
124     if (n == 1) return res;
125     if (n == 2) return res = {p[0], p[1]};
126     if mx = 0;
127     for (int i = 0, j = 2; i < n; i++) {
128         while (sgn(abs(cross(p[i], p[(i + 1) % n], p[j])),
129                     abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n])))) <= 0)
130             j = (j + 1) % n;
131         if tmp = (p[i] - p[j]).len();
132         if (tmp > mx) {
133             mx = tmp;
134             res = {p[i], p[j]};
135         }
136         tmp = (p[(i + 1) % n] - p[j]).len();
137         if (tmp > mx) {
138             mx = tmp;
139             res = {p[(i + 1) % n], p[j]};
140         }
141     }
142     return res;
143 }
144
145 // 凸包
146 auto convex\_hull(vector<vec> &p) {
147     sort(p.begin(), p.end(), vec::cmp);
148     int n = p.size();
149     vector sta(n + 1, 0);
150     vector v(n, false);
151     int tp = -1;
152     sta[++tp] = 0;
153     auto update = [&](int lim, int i) {
154         while (tp > lim && sgn(cross(p[i], p[sta[tp]], p[sta[tp - 1]]), 0) >= 0)
155             v[sta[tp--]] = 0;
156         sta[++tp] = i;
157         v[i] = 1;
158     };
159     for (int i = 1; i < n; i++) update(0, i);
160     int cnt = tp;
161     for (int i = n - 1; i >= 0; i--) {
162         if (v[i]) continue;
163         update(cnt, i);
164     }
165     vector<vec> res(tp);
166     for (int i = 0; i < tp; i++) res[i] = p[sta[i]];
167     return res;
168 }
169
170 // 闵可夫斯基和，两个点集的和构成一个凸包
171 auto minkowski(vector<vec> &a, vector<vec> &b) {
172     rotate(a.begin(), min\_element(a.begin(), a.end(), vec::cmp), a.end());

```



```

173 rotate(b.begin(), min_element(b.begin(), b.end(), vec::cmp), b.end());
174 int n = a.size(), m = b.size();
175 vector<vec> c{a[0] + b[0]};
176 c.reserve(n + m);
177 int i = 0, j = 0;
178 while (i < n && j < m) {
179     vec x = a[(i + 1) % n] - a[i];
180     vec y = b[(j + 1) % m] - b[j];
181     c.push_back(c.back() + (sgn(x ^ y, 0) >= 0 ? (i++, x) : (j++, y)));
182 }
183 while (i + 1 < n) {
184     c.push_back(c.back() + a[(i + 1) % n] - a[i]);
185     i++;
186 }
187 while (j + 1 < m) {
188     c.push_back(c.back() + b[(j + 1) % m] - b[j]);
189     j++;
190 }
191 return c;
192 }
193
194 // 过凸多边形外一点求凸多边形的切线，返回切点下标
195 auto tangent(const vec &a, vector<vec> &p) {
196     int n = p.size();
197     int l = -1, r = -1;
198     for (int i = 0; i < n; i++) {
199         if tmp1 = cross(p[i], p[(i - 1 + n) % n], a);
200         if tmp2 = cross(p[i], p[(i + 1) % n], a);
201         if (l == -1 && sgn(tmp1, 0) <= 0 && sgn(tmp2, 0) <= 0) l = i;
202         else if (r == -1 && sgn(tmp1, 0) >= 0 && sgn(tmp2, 0) >= 0) r = i;
203     }
204     return array{l, r};
205 }
206
207 // 直线
208 struct line {
209     vec p, d;
210     line() : p(vec()), d(vec()) {}
211     line(const vec &\_p, const vec &\_d) : p(\_p), d(\_d) {}
212 };
213
214 // 点到直线距离
215 lf dis(const vec &a, const line &b) {
216     return abs((b.p - a) ^ (b.p + b.d - a)) / b.d.len();
217 }
218
219 // 点在直线哪边，大于0在左边，等于0在线上，小于0在右边
220 int side_line(const vec &a, const line &b) { return sgn(b.d ^ (a - b.p), 0); }
221
222 // 两直线是否垂直
223 bool perpen(const line &a, const line &b) { return sgn(a.d * b.d, 0) == 0; }
224

```

```

225 // 两直线是否平行
226 bool parallel(const line &a, const line &b) { return sgn(a.d ^ b.d, 0) == 0; }
227
228 // 点的垂线是否与线段有交点
229 bool perpen(const vec &a, const line &b) {
230     vec p(-b.d.y, b.d.x);
231     bool cross1 = sgn(p ^ (b.p - a), 0) > 0;
232     bool cross2 = sgn(p ^ (b.p + b.d - a), 0) > 0;
233     return cross1 != cross2;
234 }
235
236 // 点到线段距离
237 lf dis\_seg(const vec &a, const line &b) {
238     if (perpen(a, b)) return dis(a, b);
239     return min((b.p - a).len(), (b.p + b.d - a).len());
240 }
241
242 // 点到凸包距离
243 lf dis(const vec &a, vector<vec> &p) {
244     lf res = inf;
245     for (int i = 1; i < p.size(); i++)
246         res = min(dis\_seg(a, line(p[i - 1], p[i] - p[i - 1])), res);
247     res = min(dis\_seg(a, line(p.back(), p[0] - p.back())), res);
248     return res;
249 }
250
251 // 两直线交点
252 vec intersection(lf A, lf B, lf C, lf D, lf E, lf F) {
253     return {(B * F - C * E) / (A * E - B * D),
254             (C * D - A * F) / (A * E - B * D)};
255 }
256
257 // 两直线交点
258 vec intersection(const line &a, const line &b) {
259     return intersection(a.d.y, -a.d.x, a.d.x * a.p.y - a.d.y * a.p.x, b.d.y,
260                        -b.d.x, b.d.x * b.p.y - b.d.y * b.p.x);
261 }
262
263 struct circle {
264     vec o;
265     lf r;
266     circle(const vec &\_o, lf \_r) : o(\_o), r(\_r){};
267
268     // 点与圆的关系 -1在圆内, 0在圆上, 1在圆外
269     int relation(const vec &a) const { return sgn((a - o).len(), r); }
270
271     // 圆与圆的关系 -3包含, -2内切, -1相交, 0外切, 1相离
272     int relation(const circle &a) const {
273         lf l = (a.o - o).len();
274         if (sgn(l, abs(r - a.r)) < 0) return -3;
275         if (sgn(l, abs(r - a.r)) == 0) return -2;
276         if (sgn(l, abs(r + a.r)) < 0) return -1;

```

```

277     if (sgn(1, abs(r + a.r)) == 0) return 0;
278     return 1;
279 }
280
281 lf area() { return PI * r * r; }
282 };
283
284 // 圆与直线交点
285 auto intersection(const circle &c, const line &l) {
286     lf d = dis(c.o, l);
287     vector<vec> res;
288     vec mid = l.p + l.d.e() * ((c.o - l.p) * l.d / l.d.len());
289     if (sgn(d, c.r) == 0) res.push_back(mid);
290     else if (sgn(d, c.r) < 0) {
291         d = sqrt(c.r * c.r - d * d);
292         res.push_back(mid + l.d.e() * d);
293         res.push_back(mid - l.d.e() * d);
294     }
295     return res;
296 }
297
298 // oab三角形与圆相交的面积
299 lf area(const circle &c, const vec &a, const vec &b) {
300     if (sgn(cross(a, b, c.o), 0) == 0) return 0;
301     vector<vec> p;
302     p.push_back(a);
303     line l(a, b - a);
304     auto tmp = intersection(c, l);
305     if (tmp.size() == 2) {
306         for (auto &i : tmp)
307             if (sgn((a - i) * (b - i), 0) < 0) p.push_back(i);
308     }
309     p.push_back(b);
310     if (p.size() == 4 && sgn((p[0] - p[1]) * (p[2] - p[1]), 0) > 0)
311         swap(p[1], p[2]);
312     lf res = 0;
313     for (int i = 1; i < p.size(); i++)
314         if (c.relation(p[i - 1]) == 1 || c.relation(p[i]) == 1) {
315             lf ang = angle(p[i - 1] - c.o, p[i] - c.o);
316             res += c.r * c.r * ang / 2;
317         } else res += abs(cross(p[i - 1], p[i], c.o)) / 2.0;
318     return res;
319 }
320
321 // 多边形与圆相交的面积
322 lf area(vector<vec> &p, circle c) {
323     lf res = 0;
324     for (int i = 0; i < p.size(); i++) {
325         int j = i + 1 == p.size() ? 0 : i + 1;
326         if (sgn(cross(p[i], p[j], c.o), 0) <= 0) res += area(c, p[i], p[j]);
327         else res -= area(c, p[i], p[j]);
328     }

```

```

329     return abs(res);
330 }

```

5.3 扫描线

```

1  #define ls (pos << 1)
2  #define rs (ls | 1)
3  #define mid ((tree[pos].l + tree[pos].r) >> 1)
4  struct Rectangle {
5      ll x\_l, y\_l, x\_r, y\_r;
6  };
7  ll area(vector<Rectangle>& rec) {
8      struct Line {
9          ll x, y\_up, y\_down;
10         int pd;
11     };
12     vector<Line> line(rec.size() * 2);
13     vector<ll> y\_set(rec.size() * 2);
14     for (int i = 0; i < rec.size(); i++) {
15         y\_set[i * 2] = rec[i].y\_l;
16         y\_set[i * 2 + 1] = rec[i].y\_r;
17         line[i * 2] = {rec[i].x\_l, rec[i].y\_r, rec[i].y\_l, 1};
18         line[i * 2 + 1] = {rec[i].x\_r, rec[i].y\_r, rec[i].y\_l, -1};
19     }
20     sort(y\_set.begin(), y\_set.end());
21     y\_set.erase(unique(y\_set.begin(), y\_set.end()), y\_set.end());
22     sort(line.begin(), line.end(), [](Line a, Line b) { return a.x < b.x; });
23     struct Data {
24         int l, r;
25         ll len, cnt, raw\_len;
26     };
27     vector<Data> tree(4 * y\_set.size());
28     function<void(int, int, int)> build = [&](int pos, int l, int r) {
29         tree[pos].l = l;
30         tree[pos].r = r;
31         if (l == r) {
32             tree[pos].raw\_len = y\_set[r + 1] - y\_set[l];
33             tree[pos].cnt = tree[pos].len = 0;
34             return;
35         }
36         build(ls, l, mid);
37         build(rs, mid + 1, r);
38         tree[pos].raw\_len = tree[ls].raw\_len + tree[rs].raw\_len;
39     };
40     function<void(int, int, int, int)> update = [&](int pos, int l, int r,
41                                                     int num) {
42         if (l <= tree[pos].l && tree[pos].r <= r) {
43             tree[pos].cnt += num;
44             tree[pos].len = tree[pos].cnt ? tree[pos].raw\_len
45                             : tree[pos].l == tree[pos].r
46                             ? 0

```

```
47         : tree[ls].len + tree[rs].len;
48     return;
49 }
50 if (l <= mid) update(ls, l, r, num);
51 if (r > mid) update(rs, l, r, num);
52 tree[pos].len =
53     tree[pos].cnt ? tree[pos].raw\_len : tree[ls].len + tree[rs].len;
54 };
55 build(1, 0, y\_set.size() - 2);
56 auto find\_pos = [&](ll num) {
57     return lower\_bound(y\_set.begin(), y\_set.end(), num) - y\_set.begin();
58 };
59 ll res = 0;
60 for (int i = 0; i < line.size() - 1; i++) {
61     update(1, find\_pos(line[i].y\_down), find\_pos(line[i].y\_up) - 1,
62         line[i].pd);
63     res += (line[i + 1].x - line[i].x) * tree[1].len;
64 }
65 return res;
66 }
```

6 杂项

6.1 快读

```
1 namespace IO {
2 constexpr int N = (1 << 20) + 1;
3 char Buffer[N];
4 int p = N;
5
6 char& get() {
7     if (p == N) {
8         fread(Buffer, 1, N, stdin);
9         p = 0;
10    }
11    return Buffer[p++];
12 }
13
14 template <typename T = int>
15 T read() {
16     T x = 0;
17     bool f = 1;
18     char c = get();
19     while (!isdigit(c)) {
20         f = c != '-';
21         c = get();
22     }
23     while (isdigit(c)) {
24         x = x * 10 + c - '0';
25         c = get();
26     }
27     return f ? x : -x;
28 }
29 } // namespace IO
30 using IO::read;
```

6.2 高精度

```
1 struct bignum {
2     string num;
3
4     bignum() : num("0") {}
5     bignum(const string& num) : num(num) {
6         reverse(this->num.begin(), this->num.end());
7     }
8     bignum(ll num) : num(to\_string(num)) {
9         reverse(this->num.begin(), this->num.end());
10    }
11
12    bignum operator+(const bignum& other) {
13        bignum res;
14        res.num.pop\_back();
```

```

15     res.num.reserve(max(num.size(), other.num.size()) + 1);
16     for (int i = 0, j = 0, x; i < num.size() || i < other.num.size() || j;
17         i++) {
18         x = j;
19         j = 0;
20         if (i < num.size()) x += num[i] - '0';
21         if (i < other.num.size()) x += other.num[i] - '0';
22         if (x >= 10) j = 1, x -= 10;
23         res.num.push_back(x + '0');
24     }
25     res.num.capacity();
26     return res;
27 }
28
29 bignum operator*(const bignum& other) {
30     vector<int> res(num.size() + other.num.size() - 1, 0);
31     for (int i = 0; i < num.size(); i++)
32         for (int j = 0; j < other.num.size(); j++)
33             res[i + j] += (num[i] - '0') * (other.num[j] - '0');
34     int g = 0;
35     for (int i = 0; i < res.size(); i++) {
36         res[i] += g;
37         g = res[i] / 10;
38         res[i] %= 10;
39     }
40     while (g) {
41         res.push_back(g % 10);
42         g /= 10;
43     }
44     int lim = res.size();
45     while (lim > 1 && res[lim - 1] == 0) lim--;
46     bignum res2;
47     res2.num.resize(lim);
48     for (int i = 0; i < lim; i++) res2.num[i] = res[i] + '0';
49     return res2;
50 }
51
52 bool operator<(const bignum& other) {
53     if (num.size() == other.num.size())
54         for (int i = num.size() - 1; i >= 0; i--)
55             if (num[i] == other.num[i]) continue;
56             else return num[i] < other.num[i];
57     return num.size() < other.num.size();
58 }
59
60 friend istream& operator>>(istream& in, bignum& a) {
61     in >> a.num;
62     reverse(a.num.begin(), a.num.end());
63     return in;
64 }
65 friend ostream& operator<<(ostream& out, bignum a) {
66     reverse(a.num.begin(), a.num.end());

```

```

67     return out << a.num;
68 }
69 };

```

6.3 离散化

```

1  template <typename T>
2  struct Hash {
3      vector<int> S;
4      vector<T> a;
5      Hash(const vector<int>& b) : S(b) {
6          sort(S.begin(), S.end());
7          S.erase(unique(S.begin(), S.end()), S.end());
8          a = vector<T>(S.size());
9      }
10     T& operator[](int i) const {
11         auto pos = lower\_bound(S.begin(), S.end(), i) - S.begin();
12         assert(pos != S.size() && S[pos] == i);
13         return a[pos];
14     }
15 };

```

6.4 模运算

```

1  constexpr int mod = 998244353;
2
3  template <typename T>
4  T power(T a, int b) {
5      T res = 1;
6      while (b) {
7          if (b & 1) res = res * a;
8          a = a * a;
9          b >>= 1;
10     }
11     return res;
12 }
13
14 struct modint {
15     int x;
16     modint(int _x = 0) : x(_x) {
17         if (x < 0) x += mod;
18         else if (x >= mod) x -= mod;
19     }
20     modint inv() const { return power(*this, mod - 2); }
21     modint operator+(const modint& b) { return x + b.x; }
22     modint operator-(const modint& b) { return x - b.x; }
23     modint operator-() const { return -x; }
24     modint operator*(const modint& b) { return int((ll)x * b.x % mod); }
25     modint operator/(const modint& b) { return *this * b.inv(); }
26     friend istream& operator>>(istream& is, modint& other) {

```



```

27     ll _x;
28     is >> _x;
29     other = modint(_x);
30     return is;
31 }
32 friend ostream& operator<<(ostream& os, modint other) {
33     return os << other.x;
34 }
35 };

```

6.5 分数

```

1 struct frac {
2     ll a, b;
3     frac() : a(0), b(1) {}
4     frac(ll _a, ll _b) : a(_a), b(_b) {
5         assert(b);
6         if (a) {
7             int tmp = gcd(a, b);
8             a /= tmp;
9             b /= tmp;
10        } else *this = frac();
11    }
12    frac operator+(const frac& other) {
13        return frac(a * other.b + other.a * b, b * other.b);
14    }
15    frac operator-() const {
16        frac res = *this;
17        res.a = -res.a;
18        return res;
19    }
20    frac operator-(const frac& other) const { return -other + *this; }
21    frac operator*(const frac& other) const {
22        return frac(a * other.a, b * other.b);
23    }
24    frac operator/(const frac& other) const {
25        assert(other.a);
26        return *this * frac(other.b, other.a);
27    }
28    bool operator<(const frac& other) const { return (*this - other).a < 0; }
29    bool operator<=(const frac& other) const { return (*this - other).a <= 0; }
30    bool operator>=(const frac& other) const { return (*this - other).a >= 0; }
31    bool operator>(const frac& other) const { return (*this - other).a > 0; }
32    bool operator==(const frac& other) const {
33        return a == other.a && b == other.b;
34    }
35    bool operator!=(const frac& other) const { return !(*this == other); }
36 };

```

6.6 表达式求值

```
1 // 格式化表达式
2 string format(const string& s1) {
3     stringstream ss(s1);
4     string s2;
5     char ch;
6     while ((ch = ss.get()) != EOF) {
7         if (ch == ' ') continue;
8         if (isdigit(ch)) s2 += ch;
9         else {
10             if (s2.back() != ' ') s2 += ' ';
11             s2 += ch;
12             s2 += ' ';
13         }
14     }
15     return s2;
16 }
17
18 // 中缀表达式转后缀表达式
19 string convert(const string& s1) {
20     unordered_map<char, int> rank{
21         {'+', 2}, {'-', 2}, {'*', 1}, {'/', 1}, {'^', 0}};
22     stringstream ss(s1);
23     string s2, temp;
24     stack<char> op;
25     while (ss >> temp) {
26         if (isdigit(temp[0])) s2 += temp + ' ';
27         else if (temp[0] == '(') op.push('(');
28         else if (temp[0] == ')') {
29             while (op.top() != '(') {
30                 s2 += op.top();
31                 s2 += ' ';
32                 op.pop();
33             }
34             op.pop();
35         } else {
36             while (!op.empty() && op.top() != '(' &&
37                 (temp[0] != '^' && rank[op.top()] <= rank[temp[0]] ||
38                 rank[op.top()] < rank[temp[0]])) {
39                 s2 += op.top();
40                 s2 += ' ';
41                 op.pop();
42             }
43             op.push(temp[0]);
44         }
45     }
46     while (!op.empty()) {
47         s2 += op.top();
48         s2 += ' ';
49         op.pop();
50     }
```

```

51     return s2;
52 }
53
54 // 计算后缀表达式
55 int calc(const string& s) {
56     stack<int> num;
57     stringstream ss(s);
58     string temp;
59     while (ss >> temp) {
60         if (isdigit(temp[0])) num.push(stoi(temp));
61         else {
62             int b = num.top();
63             num.pop();
64             int a = num.top();
65             num.pop();
66             if (temp[0] == '+') a += b;
67             else if (temp[0] == '-') a -= b;
68             else if (temp[0] == '*') a *= b;
69             else if (temp[0] == '/') a /= b;
70             else if (temp[0] == '^') a = ksm(a, b);
71             num.push(a);
72         }
73     }
74     return num.top();
75 }

```

6.7 日期

```

1  int month[] = {0, 31, 28, 31, 30, 31, 30, 31, 31, 30, 31, 30, 31};
2  int pre[13];
3  vector<int> leap;
4  struct Date {
5      int y, m, d;
6      bool operator<(const Date& other) const {
7          return array<int, 3>{y, m, d} <
8              array<int, 3>{other.y, other.m, other.d};
9      }
10     Date(const string& s) {
11         stringstream ss(s);
12         char ch;
13         ss >> y >> ch >> m >> ch >> d;
14     }
15     int dis() const {
16         int yd = (y - 1) * 365 +
17             (upper\_bound(leap.begin(), leap.end(), y - 1) - leap.begin());
18         int md =
19             pre[m - 1] + (m > 2 && (y % 4 == 0 && y % 100 || y % 400 == 0));
20         return yd + md + d;
21     }
22     int dis(const Date& other) const { return other.dis() - dis(); }
23 };

```

```

24 for (int i = 1; i <= 12; i++) pre[i] = pre[i - 1] + month[2];
25 for (int i = 1; i <= 1000000; i++)
26     if (i % 4 == 0 && i % 100 || i % 400 == 0) leap.push_back(i);

```

6.8 __builtin 函数

如果是 long long 型，记得函数后多加个 ll。

- __builtin_ctz，从最低位连续的 0 的个数，如果传入 0 则行为未定义。
- __builtin_clz，从最高位连续的 0 的个数，如果传入 0 则行为未定义。
- __builtin_popcount，二进制 1 的个数。
- __builtin_parity，二进制 1 的个数奇偶性。

6.9 对拍

linux/Mac

```

1 #!/bin/bash
2
3 g++ $1 -o a -O2
4 g++ $2 -o b -O2
5 g++ random.cpp -o random -O2
6
7 cnt=0
8 while true; do
9     let cnt++
10    echo TEST:$cnt
11    ./random > in
12    ./a < in > out.a
13    ./b < in > out.b
14    if ! diff out.a out.b; then break; fi
15 done

```

windows

```

1 @echo off
2
3 g++ a.cpp -o a -O2 -std=c++17
4 g++ b.cpp -o b -O2 -std=c++17
5 g++ random.cpp -o random -O2 -std=c++17
6
7 set cnt=0
8
9 :again
10    set /a cnt=cnt+1
11    echo TEST:%cnt%
12    .\random > in
13    .\a < in > out.a
14    .\b < in > out.b
15    fc out.a out.b
16 if not errorlevel 1 goto again

```

6.10 编译常用选项

```
1 -Wall -Woverflow -Wextra -Wpedantic -Wfloat-equal -Wshadow -fsanitize=address,undefined
```

6.11 开栈

不同的系统/编译器可能命令不一样

```
1 ulimit -s
2 -Wl,--stack=0x10000000
3 -Wl,-stack\_size -Wl,0x10000000
4 -Wl,-z,stack-size=0x10000000
```

6.12 clang-format

转储配置

```
1 clang-format -style=Google -dump-config > ./clang-format
```

.clang-format

```
1 BasedOnStyle: Google
2 IndentWidth: 4
3 AllowShortIfStatementsOnASingleLine: AllIfsAndElse
```