

ACM 常用算法模板

therehello

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1 数据结构

1.1 并查集

```
1 struct dsu {
2     int n;
3     vector<int> fa, sz;
4     dsu(int _n) : n(_n), fa(n + 1), sz(n + 1, 1) {
5         iota(fa.begin(), fa.end(), 0);
6     }
7     int find(int x) { return x == fa[x] ? x : fa[x] = find(fa[x]); }
8     int merge(int x, int y) {
9         int fax = find(x), fay = find(y);
10        if (fax == fay) return 0; // 一个集合
11        sz[fay] += fax;
12        return fa[fax] = fay; // 合并到哪个集合了
13    }
14    int size(int x) { return sz[find(x)]; }
15};
```

1.2 树状数组

1.2.1 一维

```
1 template <class T>
2 struct fenwick {
3     int n;
4     vector<T> t;
5     fenwick(int _n) : n(_n), t(n + 1) {}
6     T query(int l, int r) {
7         auto query = [&](int pos) {
8             T res = 0;
9             while (pos) {
10                res += t[pos];
11                pos -= lowbit(pos);
12            }
13            return res;
14        };
15        return query(r) - query(l - 1);
16    }
17    void add(int pos, T num) {
18        while (pos <= n) {
19            t[pos] += num;
20            pos += lowbit(pos);
21        }
22    }
23};
```

1.2.2 二维

```

1 template <class T>
2 struct Fenwick_tree_2 {
3     Fenwick_tree_2(int n, int m) : n(n), m(m), tree(n + 1, vector<T>(m + 1)) {}
4     T query(int l1, int r1, int l2, int r2) {
5         auto query = [&](int l, int r) {
6             T res = 0;
7             for (int i = l; i; i -= lowbit(i))
8                 for (int j = r; j; j -= lowbit(j)) res += tree[i][j];
9             return res;
10        };
11        return query(l2, r2) - query(l2, r1 - 1) - query(l1 - 1, r2) +
12            query(l1 - 1, r1 - 1);
13    }
14    void update(int x, int y, T num) {
15        for (int i = x; i <= n; i += lowbit(i))
16            for (int j = y; j <= m; j += lowbit(j)) tree[i][j] += num;
17    }
18 private:
19     int n, m;
20     vector<vector<T>> tree;
21 };

```

1.2.3 三维

```

1 template <class T>
2 struct Fenwick_tree_3 {
3     Fenwick_tree_3(int n, int m, int k)
4         : n(n),
5           m(m),
6           k(k),
7           tree(n + 1, vector<vector<T>>(m + 1, vector<T>(k + 1))) {}
8     T query(int a, int b, int c, int d, int e, int f) {
9         auto query = [&](int x, int y, int z) {
10             T res = 0;
11             for (int i = x; i; i -= lowbit(i))
12                 for (int j = y; j; j -= lowbit(j))
13                     for (int p = z; p; p -= lowbit(p)) res += tree[i][j][p];
14             return res;
15        };
16        T res = query(d, e, f);
17        res -= query(a - 1, e, f) + query(d, b - 1, f) + query(d, e, c - 1);
18        res += query(a - 1, b - 1, f) + query(a - 1, e, c - 1) +
19            query(d, b - 1, c - 1);
20        res -= query(a - 1, b - 1, c - 1);
21        return res;
22    }
23    void update(int x, int y, int z, T num) {
24        for (int i = x; i <= n; i += lowbit(i))
25            for (int j = y; j <= m; j += lowbit(j))
26                for (int p = z; p <= k; p += lowbit(p)) tree[i][j][p] += num;

```

```

27     }
28 private:
29     int n, m, k;
30     vector<vector<vector<T>>> tree;
31 };

```

1.3 线段树

```

1 template <class Data, class Num>
2 struct Segment_Tree {
3     inline void update(int l, int r, Num x) { update(1, l, r, x); }
4     inline Data query(int l, int r) { return query(1, l, r); }
5     Segment_Tree(vector<Data>& a) {
6         n = a.size();
7         tree.assign(n * 4 + 1, {});
8         build(a, 1, 1, n);
9     }
10 private:
11     int n;
12     struct Tree {
13         int l, r;
14         Data data;
15     };
16     vector<Tree> tree;
17     inline void pushup(int pos) {
18         tree[pos].data = tree[pos << 1].data + tree[pos << 1 | 1].data;
19     }
20     inline void pushdown(int pos) {
21         tree[pos << 1].data = tree[pos << 1].data + tree[pos].data.lazytag;
22         tree[pos << 1 | 1].data =
23             tree[pos << 1 | 1].data + tree[pos].data.lazytag;
24         tree[pos].data.lazytag = Num::zero();
25     }
26     void build(vector<Data>& a, int pos, int l, int r) {
27         tree[pos].l = l;
28         tree[pos].r = r;
29         if (l == r) {
30             tree[pos].data = a[l - 1];
31             return;
32         }
33         int mid = (tree[pos].l + tree[pos].r) >> 1;
34         build(a, pos << 1, l, mid);
35         build(a, pos << 1 | 1, mid + 1, r);
36         pushup(pos);
37     }
38     void update(int pos, int& l, int& r, Num& x) {
39         if (l > tree[pos].r || r < tree[pos].l) return;
40         if (l <= tree[pos].l && tree[pos].r <= r) {
41             tree[pos].data = tree[pos].data + x;
42             return;
43         }

```

```

44     pushdown(pos);
45     update(pos << 1, 1, r, x);
46     update(pos << 1 | 1, 1, r, x);
47     pushup(pos);
48 }
49 Data query(int pos, int& l, int& r) {
50     if (l > tree[pos].r || r < tree[pos].l) return Data::zero();
51     if (l <= tree[pos].l && tree[pos].r <= r) return tree[pos].data;
52     pushdown(pos);
53     return query(pos << 1, l, r) + query(pos << 1 | 1, l, r);
54 }
55 };
56 struct Num {
57     ll add;
58     inline static Num zero() { return {0}; }
59     inline Num operator+(Num b) { return {add + b.add}; }
60 };
61 struct Data {
62     ll sum, len;
63     Num lazytag;
64     inline static Data zero() { return {0, 0, Num::zero()}; }
65     inline Data operator+(Num b) {
66         return {sum + len * b.add, len, lazytag + b};
67     }
68     inline Data operator+(Data b) {
69         return {sum + b.sum, len + b.len, Num::zero()};
70     }
71 };

```

1.4 普通平衡树

1.4.1 树状数组实现

需要预先处理出来所有可能的数。

```

1  template <typename T>
2  struct treap {
3      int n, size;
4      vector<int> t;
5      vector<T> t2, S;
6      treap(const vector<T>& b) {
7          S = b;
8          sort(S.begin(), S.end());
9          S.erase(unique(S.begin(), S.end()), S.end());
10         n = S.size();
11         size = 0;
12         t = vector<int>(n + 1);
13         t2 = vector<T>(n + 1);
14     }
15     int pos(T x) { return lower_bound(S.begin(), S.end(), x) - S.begin() + 1; }
16     int sum(int pos) {
17         int res = 0;

```



```
18     while (pos) {
19         res += t[pos];
20         pos -= lowbit(pos);
21     }
22     return res;
23 }
24
25 // 插入cnt个x
26 void insert(T x, int cnt) {
27     size += cnt;
28     for (int i = pos(x); i <= n; i += lowbit(i)) {
29         t[i] += cnt;
30         t2[i] += cnt * x;
31     }
32 }
33
34 // 删除cnt个x
35 void erase(T x, int cnt) { insert(x, -cnt); }
36
37 // x的排名
38 int rank(T x) { return sum(pos(x) - 1) + 1; }
39
40 // 统计出现次数
41 int count(T x) { return sum(pos(x)) - sum(pos(x) - 1); }
42
43 // 第k小
44 T kth(int k) {
45     int cnt = 0, x = 0;
46     for (int i = log2(n); i >= 0; i--) {
47         x += 1 << i;
48         if (x >= n || cnt + t[x] >= k) x -= 1 << i;
49         else cnt += t[x];
50     }
51     return S[x];
52 }
53
54 // 前k小的数之和
55 T pre_sum(int k) {
56     int cnt = 0, x = 0;
57     T res = 0;
58     for (int i = log2(n); i >= 0; i--) {
59         x += 1 << i;
60         if (x >= n || cnt + t[x] >= k) x -= 1 << i;
61         else {
62             cnt += t[x];
63             res += t2[x];
64         }
65     }
66     return res + (k - cnt) * S[x];
67 }
68
69 // 小于x, 最大的数
```

```

70 T prev(int x) { return kth(sum(pos(x) - 1)); }
71
72 // 大于x, 最小的数
73 T next(int x) { return kth(sum(pos(x)) + 1); }
74 };

```

1.5 可持久化线段树

```

1 constexpr int MAXN = 200000;
2 vector<int> root(MAXN << 5);
3 struct Persistent_seg {
4     int n;
5     struct Data {
6         int ls, rs;
7         int val;
8     };
9     vector<Data> tree;
10    Persistent_seg(int n, vector<int>& a) : n(n) { root[0] = build(1, n, a); }
11    int build(int l, int r, vector<int>& a) {
12        if (l == r) {
13            tree.push_back({0, 0, a[l]});
14            return tree.size() - 1;
15        }
16        int mid = l + r >> 1;
17        int ls = build(l, mid, a), rs = build(mid + 1, r, a);
18        tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
19        return tree.size() - 1;
20    }
21    int update(int rt, const int& idx, const int& val, int l, int r) {
22        if (l == r) {
23            tree.push_back({0, 0, tree[rt].val + val});
24            return tree.size() - 1;
25        }
26        int mid = l + r >> 1, ls = tree[rt].ls, rs = tree[rt].rs;
27        if (idx <= mid) ls = update(ls, idx, val, l, mid);
28        else rs = update(rs, idx, val, mid + 1, r);
29        tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
30        return tree.size() - 1;
31    }
32    int query(int rt1, int rt2, int k, int l, int r) {
33        if (l == r) return l;
34        int mid = l + r >> 1;
35        int lcnt = tree[tree[rt2].ls].val - tree[tree[rt1].ls].val;
36        if (k <= lcnt) return query(tree[rt1].ls, tree[rt2].ls, k, l, mid);
37        else return query(tree[rt1].rs, tree[rt2].rs, k - lcnt, mid + 1, r);
38    }
39 };

```

1.6 st 表

```
1 auto lg = []() {
2     array<int, 10000001> lg;
3     lg[1] = 0;
4     for (int i = 2; i <= 10000000; i++) lg[i] = lg[i >> 1] + 1;
5     return lg;
6 }();
7 template <typename T>
8 struct st {
9     int n;
10    vector<vector<T>>> a;
11    st(vector<T>& _a) : n(_a.size()) {
12        a.assign(lg[n] + 1, vector<int>(n));
13        for (int i = 0; i < n; i++) a[0][i] = _a[i];
14        for (int j = 1; j <= lg[n]; j++)
15            for (int i = 0; i + (1 << j) - 1 < n; i++)
16                a[j][i] = max(a[j - 1][i], a[j - 1][i + (1 << (j - 1))]);
17    }
18    T query(int l, int r) {
19        int k = lg[r - l + 1];
20        return max(a[k][l], a[k][r - (1 << k) + 1]);
21    }
22 };
```

2 图论

存图

```

1 struct Graph {
2     int n;
3     struct Edge {
4         int to, w;
5     };
6     vector<vector<Edge>> graph;
7     Graph(int _n) {
8         n = _n;
9         graph.assign(n + 1, vector<Edge>());
10    };
11    void add(int u, int v, int w) { graph[u].push_back({v, w}); }
12 };

```

2.1 最短路

2.1.1 dijkstra

```

1 void dij(Graph& graph, vector<int>& dis, int t) {
2     vector<int> visit(graph.n + 1, 0);
3     priority_queue<pair<int, int>> que;
4     dis[t] = 0;
5     que.emplace(0, t);
6     while (!que.empty()) {
7         int u = que.top().second;
8         que.pop();
9         if (visit[u]) continue;
10        visit[u] = 1;
11        for (auto& [to, w] : graph.graph[u]) {
12            if (dis[to] > dis[u] + w) {
13                dis[to] = dis[u] + w;
14                que.emplace(-dis[to], to);
15            }
16        }
17    }
18 }

```

2.2 树上问题

2.2.1 最近公公祖先

倍增法

```

1 vector<int> dep;
2 vector<array<int, 21>> fa;
3 dep.assign(n + 1, 0);
4 fa.assign(n + 1, array<int, 21>{});
5 void binary_jump(int root) {
6     function<void(int)> dfs = [&](int t) {

```

```

7     dep[t] = dep[fa[t][0]] + 1;
8     for (auto& [to] : graph[t]) {
9         if (to == fa[t][0]) continue;
10        fa[to][0] = t;
11        dfs(to);
12    }
13 };
14 dfs(root);
15 for (int j = 1; j <= 20; j++)
16     for (int i = 1; i <= n; i++) fa[i][j] = fa[fa[i][j - 1]][j - 1];
17 }
18 int lca(int x, int y) {
19     if (dep[x] < dep[y]) swap(x, y);
20     for (int i = 20; i >= 0; i--)
21         if (dep[fa[x][i]] >= dep[y]) x = fa[x][i];
22     if (x == y) return x;
23     for (int i = 20; i >= 0; i--) {
24         if (fa[x][i] != fa[y][i]) {
25             x = fa[x][i];
26             y = fa[y][i];
27         }
28     }
29     return fa[x][0];
30 }

```

树剖

```

1 int lca(int x, int y) {
2     while (top[x] != top[y]) {
3         if (dep[top[x]] < dep[top[y]]) swap(x, y);
4         x = fa[top[x]];
5     }
6     if (dep[x] < dep[y]) swap(x, y);
7     return y;
8 }

```

2.2.2 树链剖分

```

1 vector<int> fa, siz, dep, son, dfn, rnk, top;
2 fa.assign(n + 1, 0);
3 siz.assign(n + 1, 0);
4 dep.assign(n + 1, 0);
5 son.assign(n + 1, 0);
6 dfn.assign(n + 1, 0);
7 rnk.assign(n + 1, 0);
8 top.assign(n + 1, 0);
9 void hld(int root) {
10     function<void(int)> dfs1 = [&](int t) {
11         dep[t] = dep[fa[t]] + 1;
12         siz[t] = 1;
13         for (auto& [to, w] : graph[t]) {
14             if (to == fa[t]) continue;

```

```

15         fa[to] = t;
16         dfs1(to);
17         if (siz[son[t]] < siz[to]) son[t] = to;
18         siz[t] += siz[to];
19     }
20 };
21 dfs1(root);
22 int dfn_tail = 0;
23 for (int i = 1; i <= n; i++) top[i] = i;
24 function<void(int)> dfs2 = [&](int t) {
25     dfn[t] = ++dfn_tail;
26     rnk[dfn_tail] = t;
27     if (!son[t]) return;
28     top[son[t]] = top[t];
29     dfs2(son[t]);
30     for (auto& [to, w] : graph[t]) {
31         if (to == fa[t] || to == son[t]) continue;
32         dfs2(to);
33     }
34 };
35 dfs2(root);
36 }

```

2.3 强连通分量

```

1 void tarjan(Graph& g1, Graph& g2) {
2     int dfn_tail = 0, cnt = 0;
3     vector<int> dfn(g1.n + 1, 0), low(g1.n + 1, 0), exist(g1.n + 1, 0),
4         belong(g1.n + 1, 0);
5     stack<int> sta;
6     function<void(int)> dfs = [&](int t) {
7         dfn[t] = low[t] = ++dfn_tail;
8         sta.push(t);
9         exist[t] = 1;
10        for (auto& [to] : g1.graph[t])
11            if (!dfn[to]) {
12                dfs(to);
13                low[t] = min(low[t], low[to]);
14            } else if (exist[to]) low[t] = min(low[t], dfn[to]);
15        if (dfn[t] == low[t]) {
16            cnt++;
17            while (int temp = sta.top()) {
18                belong[temp] = cnt;
19                exist[temp] = 0;
20                sta.pop();
21                if (temp == t) break;
22            }
23        }
24    };
25    for (int i = 1; i <= g1.n; i++)
26        if (!dfn[i]) dfs(i);

```

```
27 g2 = Graph(cnt);
28 for (int i = 1; i <= g1.n; i++) g2.w[belong[i]] += g1.w[i];
29 for (int i = 1; i <= g1.n; i++)
30     for (auto& [to] : g1.graph[i])
31         if (belong[i] != belong[to]) g2.add(belong[i], belong[to]);
32 }
```

2.4 拓扑排序

```
1 void toposort(Graph& g, vector<int>& dis) {
2     vector<int> in(g.n + 1, 0);
3     for (int i = 1; i <= g.n; i++)
4         for (auto& [to] : g.graph[i]) in[to]++;
5     queue<int> que;
6     for (int i = 1; i <= g.n; i++)
7         if (!in[i]) {
8             que.push(i);
9             dis[i] = g.w[i]; // dp
10        }
11    while (!que.empty()) {
12        int u = que.front();
13        que.pop();
14        for (auto& [to] : g.graph[u]) {
15            in[to]--;
16            dis[to] = max(dis[to], dis[u] + g.w[to]); // dp
17            if (!in[to]) que.push(to);
18        }
19    }
20 }
```

3 字符串

3.1 kmp

```

1 auto kmp(string& s) {
2     vector next(s.size(), -1);
3     for (int i = 1, j = -1; i < s.size(); i++) {
4         while (j >= 0 && s[i] != s[j + 1]) j = next[j];
5         if (s[i] == s[j + 1]) j++;
6         next[i] = j;
7     }
8     // next 意为长度
9     for (auto& i : next) i++;
10    return next;
11 }

```

3.2 哈希

```

1 constexpr int N = 2e6;
2 constexpr ll mod[2] = {2000000011, 2000000033}, base[2] = {20011, 20033};
3 vector<array<ll, 2>> pow_base(N);
4
5 pow_base[0][0] = pow_base[0][1] = 1;
6 for (int i = 1; i < N; i++) {
7     pow_base[i][0] = pow_base[i - 1][0] * base[0] % mod[0];
8     pow_base[i][1] = pow_base[i - 1][1] * base[1] % mod[1];
9 }
10
11 struct Hash {
12     int size;
13     vector<array<ll, 2>> hash;
14     Hash() {}
15     Hash(const string& s) {
16         size = s.size();
17         hash.resize(size);
18         hash[0][0] = hash[0][1] = s[0];
19         for (int i = 1; i < size; i++) {
20             hash[i][0] = (hash[i - 1][0] * base[0] + s[i]) % mod[0];
21             hash[i][1] = (hash[i - 1][1] * base[1] + s[i]) % mod[1];
22         }
23     }
24     array<ll, 2> operator[] (const array<int, 2>& range) const {
25         int l = range[0], r = range[1];
26         if (l == 0) return hash[r];
27         auto single_hash = [&](bool flag) {
28             return (hash[r][flag] -
29                     hash[l - 1][flag] * pow_base[r - l + 1][flag] % mod[flag] +
30                     mod[flag]) %
31                     mod[flag];
32         };
33         return {single_hash(0), single_hash(1)};

```



```
34     }  
35 };
```

3.3 manacher

```
1 void manacher(const string& _s, vector<int>& r) {  
2     string s(_s.size() * 2 + 1, '$');  
3     for (int i = 0; i < _s.size(); i++) s[2 * i + 1] = _s[i];  
4     r.resize(_s.size() * 2 + 1);  
5     for (int i = 0, maxr = 0, mid = 0; i < s.size(); i++) {  
6         if (i < maxr) r[i] = min(r[mid * 2 - i], maxr - i);  
7         while (i - r[i] - 1 >= 0 && i + r[i] + 1 < s.size() &&  
8             s[i - r[i] - 1] == s[i + r[i] + 1])  
9             ++r[i];  
10        if (i + r[i] > maxr) maxr = i + r[i], mid = i;  
11    }  
12 }
```

4 数学

4.1 扩展欧几里得

需保证 $a, b \geq 0$

$$x = x + k * dx, y = y - k * dy$$

若要求 $x \geq p$, $k \geq \lceil \frac{p-x}{dx} \rceil$

若要求 $x \leq q$, $k \leq \lfloor \frac{q-x}{dx} \rfloor$

若要求 $y \geq p$, $k \leq \lfloor \frac{y-p}{dy} \rfloor$

若要求 $y \leq q$, $k \geq \lceil \frac{y-q}{dy} \rceil$

```

1 int __exgcd(int a, int b, int& x, int& y) {
2     if (!b) {
3         x = 1;
4         y = 0;
5         return a;
6     }
7     int g = __exgcd(b, a % b, y, x);
8     y -= a / b * x;
9     return g;
10 }
11
12 array<int, 2> exgcd(int a, int b, int c) {
13     int x, y;
14     int g = __exgcd(a, b, x, y);
15     if (c % g) return {INT_MAX, INT_MAX};
16     int dx = b / g;
17     int dy = a / g;
18     x = c / g % dx * x % dx;
19     if (x < 0) x += dx;
20     y = (c - a * x) / b;
21     return {x, y};
22 }

```

4.2 线性筛法

```

1 constexpr int N = 10000000;
2 array<int, N + 1> min_prime;
3 vector<int> primes;
4 bool ok = []() {
5     for (int i = 2; i <= N; i++) {
6         if (min_prime[i] == 0) {
7             min_prime[i] = i;
8             primes.push_back(i);
9         }
10        for (auto& j : primes) {
11            if (j > min_prime[i] || j > N / i) break;
12            min_prime[j * i] = j;
13        }
14    }
15    return 1;

```

```
16 }();
```

4.3 分解质因数

```
1 auto getprimes(int n) {
2     vector<array<int, 2>> res;
3     for (auto& i : primes) {
4         if (i > n / i) break;
5         if (n % i == 0) {
6             res.push_back({i, 0});
7             while (n % i == 0) {
8                 n /= i;
9                 res.back()[1]++;
10            }
11        }
12    }
13    if (n > 1) res.push_back({n, 1});
14    return res;
15 }
```

4.4 pollard rho

```
1 using LL = __int128_t;
2
3 random_device rd;
4 mt19937 seed(rd());
5
6 ll power(ll a, ll b, ll mod) {
7     ll res = 1;
8     while (b) {
9         if (b & 1) res = (LL)res * a % mod;
10        a = (LL)a * a % mod;
11        b >>= 1;
12    }
13    return res;
14 }
15
16 bool isprime(ll n) {
17     static array primes{2, 3, 5, 7, 11, 13, 17, 19, 23};
18     static unordered_map<ll, bool> S;
19     if (n < 2) return 0;
20     if (S.count(n)) return S[n];
21     ll d = n - 1, r = 0;
22     while (!(d & 1)) {
23         r++;
24         d >>= 1;
25     }
26     for (auto& a : primes) {
27         if (a == n) return S[n] = 1;
28         ll x = power(a, d, n);
```

```

29     if (x == 1 || x == n - 1) continue;
30     for (int i = 0; i < r - 1; i++) {
31         x = (LL)x * x % n;
32         if (x == n - 1) break;
33     }
34     if (x != n - 1) return S[n] = 0;
35 }
36 return S[n] = 1;
37 }
38
39 ll pollard_rho(ll n) {
40     ll s = 0, t = 0;
41     ll c = seed() % (n - 1) + 1;
42     ll val = 1;
43     for (int goal = 1;; goal *= 2, s = t, val = 1) {
44         for (int step = 1; step <= goal; step++) {
45             t = ((LL)t * t + c) % n;
46             val = (LL)val * abs(t - s) % n;
47             if (step % 127 == 0) {
48                 ll g = gcd(val, n);
49                 if (g > 1) return g;
50             }
51         }
52         ll g = gcd(val, n);
53         if (g > 1) return g;
54     }
55 }
56 auto getprimes(ll n) {
57     unordered_set<ll> S;
58     auto get = [&](auto self, ll n) {
59         if (n < 2) return;
60         if (isprime(n)) {
61             S.insert(n);
62             return;
63         }
64         ll mx = pollard_rho(n);
65         self(self, n / mx);
66         self(self, mx);
67     };
68     get(get, n);
69     return S;
70 }

```

4.5 组合数

```

1 constexpr int N = 1e7;
2 array<modint, N + 1> fac, ifac;
3 auto _ = []() {
4     fac[0] = 1;
5     for (int i = 1; i <= N; i++) fac[i] = fac[i - 1] * i;
6     ifac[N] = fac[N].inv();

```

```

7   for (int i = N - 1; i >= 0; i--) ifac[i] = ifac[i + 1] * (i + 1);
8   return true;
9 }();
10
11 modint C(int n, int m) {
12     if (n < m) return 0;
13     if (n <= mod) return fac[n] * ifac[m] * ifac[n - m];
14     // n >= mod 时需要这个
15     return C(n % mod, m % mod) * C(n / mod, m / mod);
16 }

```

4.6 数论分块

求解形如 $\sum_{i=1}^n f(i)g(\lfloor \frac{n}{i} \rfloor)$ 的合式

$$s(n) = \sum_{i=1}^n f(i)$$

```

1 modint sqrt_decomposition(int n) {
2     auto s = [&](int x) { return x; };
3     auto g = [&](int x) { return x; };
4     modint res = 0;
5     while (l <= R) {
6         int r = n / (n / l);
7         res = res + (s(r) - s(l - 1)) * g(n / l);
8         l = r + 1;
9     }
10    return res;
11 }

```

4.7 积性函数

4.7.1 定义

函数 $f(n)$ 满足 $f(1) = 1$ 且 $\forall x, y \in \mathbf{N}^*, \gcd(x, y) = 1$ 都有 $f(xy) = f(x)f(y)$, 则 $f(n)$ 为积性函数。

函数 $f(n)$ 满足 $f(1) = 1$ 且 $\forall x, y \in \mathbf{N}^*$ 都有 $f(xy) = f(x)f(y)$, 则 $f(n)$ 为完全积性函数。

4.7.2 例子

- 单位函数: $\varepsilon(n) = [n = 1]$ 。(完全积性)
- 恒等函数: $\text{id}_k(n) = n^k$ 。(完全积性)
- 常数函数: $1(n) = 1$ 。(完全积性)
- 除数函数: $\sigma_k(n) = \sum_{d|n} d^k$ 。 $\sigma_0(n)$ 通常简记作 $d(n)$ 或 $\tau(n)$, $\sigma_1(n)$ 通常简记作 $\sigma(n)$ 。
- 欧拉函数: $\varphi(n) = \sum_{i=1}^n [\gcd(i, n) = 1]$ 。
- 莫比乌斯函数: $\mu(n) = \begin{cases} 1 & n = 1 \\ 0 & \exists d > 1, d^2 | n, \text{ 其中 } \omega(n) \text{ 表示 } n \text{ 的本质不同质因子个数, 它是} \\ (-1)^{\omega(n)} & \text{otherwise} \end{cases}$
一个加性函数。

4.8 狄利克雷卷积

对于两个数论函数 $f(x)$ 和 $g(x)$ ，则它们的狄利克雷卷积得到的结果 $h(x)$ 定义为：

$$h(x) = \sum_{d|x} f(d)g\left(\frac{x}{d}\right) = \sum_{ab=x} f(a)g(b)$$

可以简记为： $h = f * g$ 。

4.8.1 性质

交换律： $f * g = g * f$ 。

结合律： $(f * g) * h = f * (g * h)$ 。

分配律： $(f + g) * h = f * h + g * h$ 。

等式的性质： $f = g$ 的充要条件是 $f * h = g * h$ ，其中数论函数 $h(x)$ 要满足 $h(1) \neq 0$ 。

4.8.2 例子

- $\varepsilon = \mu * 1 \iff \varepsilon(n) = \sum_{d|n} \mu(d)$
- $id = \varphi * 1 \iff id(n) = \sum_{d|n} \varphi(d)$
- $d = 1 * 1 \iff d(n) = \sum_{d|n} 1$
- $\sigma = id * 1 \iff \sigma(n) = \sum_{d|n} d$
- $\varphi = \mu * id \iff \varphi(n) = \sum_{d|n} d \cdot \mu\left(\frac{n}{d}\right)$

4.9 欧拉函数

```
1 array<int, N + 1> phi;
2 auto _ = []() {
3     iota(phi.begin() + 1, phi.end(), 1);
4     for (int i = 2; i <= N; i++) {
5         if (phi[i] == i)
6             for (int j = i; j <= N; j += i) phi[j] = phi[j] / i * (i - 1);
7     }
8     return true;
9 }();
```

4.10 莫比乌斯反演

4.10.1 莫比乌斯函数性质

- $\sum_{d|n} \mu(d) = \begin{cases} 1 & n = 1 \\ 0 & n \neq 1 \end{cases}$ ，即 $\sum_{d|n} \mu(d) = \varepsilon(n)$ ， $\mu * 1 = \varepsilon$
- $[\gcd(i, j) = 1] = \sum_{d|\gcd(i, j)} \mu(d)$

```
1 array<int, N + 1> miu;
2 array<bool, N + 1> ispr;
3 auto _ = []() {
4     miu.fill(1);
5     ispr.fill(1);
```

```

6   for (int i = 2; i <= N; i++) {
7       if (!ispr[i]) continue;
8       miu[i] = -1;
9       for (int j = 2 * i; j <= N; j += i) {
10          ispr[j] = 0;
11          if ((j / i) % i == 0) miu[j] = 0;
12          else miu[j] *= -1;
13      }
14  }
15  return true;
16 }();

```

4.10.2 莫比乌斯变换/反演

$f(n) = \sum_{d|n} g(d)$, 那么有 $g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 。

用狄利克雷卷积表示则为 $f = g * 1$, 有 $g = f * \mu$ 。

$f \rightarrow g$ 称为莫比乌斯反演, $g \rightarrow f$ 称为莫比乌斯反演。

4.11 杜教筛

杜教筛被用于处理一类数论函数的前缀和问题。对于数论函数 f , 杜教筛可以在低于线性时间的复杂度内计算 $S(n) = \sum_{i=1}^n f(i)$ 。

$$S(n) = \frac{\sum_{i=1}^n (f * g)(i) - \sum_{i=2}^n g(i) S(\lfloor \frac{n}{i} \rfloor)}{g(1)}$$

可以构造恰当的数论函数 g 使得:

- 可以快速计算 $\sum_{i=1}^n (f * g)(i)$ 。
- 可以快速计算 g 的单点值, 用数论分块求解 $\sum_{i=2}^n g(i) S(\lfloor \frac{n}{i} \rfloor)$ 。

4.11.1 示例

```

1 ll sum_phi(ll n) {
2     if (n <= N) return sp[n];
3     if (sp2.count(n)) return sp2[n];
4     ll res = 0, l = 2;
5     while (l <= n) {
6         ll r = n / (n / l);
7         res = res + (r - l + 1) * sum_phi(n / l);
8         l = r + 1;
9     }
10    return sp2[n] = (ll)n * (n + 1) / 2 - res;
11 }
12 ll sum_miu(ll n) {
13     if (n <= N) return sm[n];
14     if (sm2.count(n)) return sm2[n];
15     ll res = 0, l = 2;
16     while (l <= n) {
17         ll r = n / (n / l);
18         res = res + (r - l + 1) * sum_miu(n / l);

```

```
19     l = r + 1;
20 }
21 return sm2[n] = 1 - res;
22 }
```

4.12 盒子与球

n 个球, m 个盒

球同	盒同	可空	公式
✓	✓	✓	$f_{n,m} = f_{n-1,m-1} + f_{n-m,m}$
✓	✓	✗	$f_{n-m,m}$
✗	✓	✓	$\sum_{i=1}^m g_{n,i}$
✗	✓	✗	$g_{n,m} = g_{n-1,m-1} + m * g_{n-1,m}$
✓	✗	✓	C_{n+m-1}^{m-1}
✓	✗	✗	C_{n-1}^{m-1}
✗	✗	✓	m^n
✗	✗	✗	$m! * g_{n,m}$

扩展:

n 个相同的球, m 个不同的盒, 每个盒子超过 k 个球, 问方案数?

可以考虑容斥, $f(d)$ 表示至少有 d 个盒子装了 $> k$ 个球方案数, 总方案数则为 $f(0)-f(1)+f(2)-\dots$

4.13 线性基

```
1 // 线性基
2 struct basis {
3     int rnk = 0;
4     array<ull, 64> p{};
5
6     // 将x插入此线性基中
7     void insert(ull x) {
8         for (int i = 63; i >= 0; i--) {
9             if ((x >> i) & 1) {
10                 if (p[i]) x ^= p[i];
11                 else {
12                     for (int j = 0; j < i; j++)
13                         if (x >> j & 1) x ^= p[j];
14                     for (int j = i + 1; j <= 63; j++)
15                         if (p[j] >> i & 1) p[j] ^= x;
16                     p[i] = x;
17                     rnk++;
18                     break;
19                 }
20             }
```



```

21     }
22 }
23
24 // 将另一个线性基插入此线性基中
25 void insert(basis other) {
26     for (int i = 0; i <= 63; i++) {
27         if (!other.p[i]) continue;
28         insert(other.p[i]);
29     }
30 }
31
32 // 最大异或值
33 ull max_basis() {
34     ull res = 0;
35     for (int i = 63; i >= 0; i--)
36         if ((res ^ p[i]) > res) res ^= p[i];
37     return res;
38 }
39 };

```

4.14 矩阵快速幂

```

1 constexpr ll mod = 2147493647;
2 struct Mat {
3     int n, m;
4     vector<vector<ll>> mat;
5     Mat(int n, int m) : n(n), m(m), mat(n, vector<ll>(m, 0)) {}
6     Mat(vector<vector<ll>> mat) : n(mat.size()), m(mat[0].size()), mat(mat) {}
7     Mat operator*(const Mat& other) {
8         assert(m == other.n);
9         Mat res(n, other.m);
10        for (int i = 0; i < res.n; i++)
11            for (int j = 0; j < res.m; j++)
12                for (int k = 0; k < m; k++)
13                    res.mat[i][j] =
14                        (res.mat[i][j] + mat[i][k] * other.mat[k][j] % mod) %
15                        mod;
16        return res;
17    }
18 };
19 Mat ksm(Mat a, ll b) {
20     assert(a.n == a.m);
21     Mat res(a.n, a.m);
22     for (int i = 0; i < res.n; i++) res.mat[i][i] = 1;
23     while (b) {
24         if (b & 1) res = res * a;
25         b >>= 1;
26         a = a * a;
27     }
28     return res;
29 }

```

5 计算几何

```

1 double eps = 1e-8;
2 const double PI = acos(-1);
3 using T = ll;
4
5 template <typename T>
6 int cmp(T a, T b) {
7     return a != b ? a < b ? -1 : 1 : 0;
8 }
9
10 int cmp(double a, double b) {
11     double c = a - b;
12     if (abs(c) < eps) return 0;
13     return c < 0 ? -1 : 1;
14 }
15
16 // 向量
17 struct vec {
18     T x, y;
19     vec(T _x = 0, T _y = 0) : x(_x), y(_y) {}
20
21     // 模
22     double length2() const { return x * x + y * y; }
23     double length() const { return sqrt(x * x + y * y); }
24
25     // 与x轴正方向的夹角
26     double angle() const {
27         double angle = atan2(y, x);
28         if (angle < 0) angle += 2 * PI;
29         return angle;
30     }
31
32     // 逆时针旋转
33     vec &rotate(const double &theta) {
34         double tmp = x;
35         x = x * cos(theta) - y * sin(theta);
36         y = y * cos(theta) + tmp * sin(theta);
37         return *this;
38     }
39
40     bool operator==(const vec &other) const {
41         return !cmp(x, other.x) && !cmp(y, other.y);
42     }
43     bool operator<(const vec &other) const {
44         int tmp = cmp(angle(), other.angle());
45         if (tmp) return tmp == -1 ? 0 : 1;
46         tmp = cmp(x, other.x);
47         return tmp == -1 ? 0 : 1;
48     }
49
50     vec operator+(const vec &other) const { return {x + other.x, y + other.y}; }

```

```

51 vec operator-() const { return {-x, -y}; }
52 vec operator-(const vec &other) const { return -other + (*this); }
53 vec operator*(const T &other) const { return {x * other, y * other}; }
54 vec operator/(const T &other) const { return {x / other, y / other}; }
55 T operator*(const vec &other) const { return x * other.x + y * other.y; }
56
57 // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
58 // 等于0共线, 可能同向或反向, 结果绝对值表示 a b 形成的平行四边形的面积
59 T operator^(const vec &other) const { return x * other.y - y * other.x; }
60
61 friend istream &operator>>(istream &input, vec &data) {
62     input >> data.x >> data.y;
63     return input;
64 }
65 friend ostream &operator<<(ostream &output, const vec &data) {
66     output << fixed << setprecision(6);
67     output << data.x << " " << data.y;
68     return output;
69 }
70 };
71
72 bool xycmp(const vec &a, const vec &b) {
73     int tmp = cmp(a.x, b.x);
74     if (tmp) return tmp == -1 ? 0 : 1;
75     tmp = cmp(a.y, b.y);
76     return tmp == -1 ? 0 : 1;
77 }
78
79 T cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
80
81 // 两点间的距离
82 T distance(const vec &a, const vec &b) {
83     return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y);
84 }
85
86 // 两向量夹角
87 double angle(const vec &a, const vec &b) {
88     double theta = abs(a.angle() - b.angle());
89     if (theta > PI) theta = 2 * PI - theta;
90     return theta;
91 }
92
93 // 判断点是否在凸包内
94 bool in_polygon(const vec &a, vector<vec> &p) {
95     int n = p.size();
96     if (n == 1) return a == p[0];
97     if (cross(a, p[1], p[0]) > 0 || cross(p.back(), a, p[0]) > 0) return 0;
98     auto cmp = [&p](vec &x, const vec &y) { return ((x - p[0]) ^ y) >= 0; };
99     int i = lower_bound(p.begin() + 2, p.end(), a - p[0], cmp) - p.begin() - 1;
100    return cross(p[(i + 1) % n], a, p[i]) >= 0;
101 }
102

```

```
103 // 多边形的面积
104 double polygon_area(vector<vec> &p) {
105     T area = 0;
106     for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];
107     area += p.back() ^ p[0];
108     return abs(area / 2.0);
109 }
110
111 // 多边形的周长
112 double polygon_length(vector<vec> &p) {
113     double length = 0;
114     for (int i = 1; i < p.size(); i++) length += (p[i - 1] - p[i]).length();
115     length += (p.back() - p[0]).length();
116     return length;
117 }
118
119 // 以整点为顶点的线段上的整点个数
120 T count(const vec &a, const vec &b) {
121     vec c = a - b;
122     return gcd(abs(c.x), abs(c.y)) + 1;
123 }
124
125 // 以整点为顶点的多边形边上整点个数
126 T count(vector<vec> &p) {
127     T cnt = 0;
128     for (int i = 1; i < p.size(); i++) cnt += count(p[i - 1], p[i]);
129     cnt += count(p.back(), p[0]);
130     return cnt - p.size();
131 }
132
133 // 凸包直径的两个端点
134 auto polygon_dia(vector<vec> &p) {
135     int n = p.size();
136     array<vec, 2> res{};
137     if (n == 1) return res;
138     if (n == 2) return res = {p[0], p[1]};
139     T mx = 0;
140     for (int i = 0, j = 2; i < n; i++) {
141         while (abs(cross(p[i], p[(i + 1) % n], p[j])) <=
142             abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n])))
143             j = (j + 1) % n;
144         if (T tmp = distance(p[i], p[j]); tmp > mx) {
145             mx = tmp;
146             res = {p[i], p[j]};
147         }
148         if (T tmp = distance(p[(i + 1) % n], p[j]); tmp > mx) {
149             mx = tmp;
150             res = {p[(i + 1) % n], p[j]};
151         }
152     }
153     return res;
154 }
```

```

155
156 // 凸包
157 auto convex_hull(vector<vec> &p) {
158     sort(p.begin(), p.end(), xycmp);
159     int n = p.size();
160     vector sta(n + 1, 0);
161     vector v(n, false);
162     int tp = -1;
163     sta[++tp] = 0;
164     auto update = [&](int lim, int i) {
165         while (tp > lim &&
166             ((p[sta[tp]] - p[sta[tp - 1]]) ^ (p[i] - p[sta[tp]])) <= 0)
167             v[sta[tp--]] = 0;
168         sta[++tp] = i;
169         v[i] = 1;
170     };
171     for (int i = 1; i < n; i++) update(0, i);
172     int cnt = tp;
173     for (int i = n - 1; i >= 0; i--) {
174         if (v[i]) continue;
175         update(cnt, i);
176     }
177     vector<vec> res(tp);
178     for (int i = 0; i < tp; i++) res[i] = p[sta[i]];
179     return res;
180 }
181
182 // 闵可夫斯基和 两个点集的和构成一个凸包
183 auto minkowski(vector<vec> &a, vector<vec> &b) {
184     rotate(a.begin(), min_element(a.begin(), a.end(), xycmp), a.end());
185     rotate(b.begin(), min_element(b.begin(), b.end(), xycmp), b.end());
186     int n = a.size(), m = b.size();
187     vector<vec> c{a[0] + b[0]};
188     c.reserve(n + m);
189     int i = 0, j = 0;
190     while (i < n && j < m) {
191         vec x = a[(i + 1) % n] - a[i];
192         vec y = b[(j + 1) % m] - b[j];
193         c.push_back(c.back() + ((x ^ y) >= 0 ? (i++, x) : (j++, y)));
194     }
195     while (i + 1 < n) {
196         c.push_back(c.back() + a[(i + 1) % n] - a[i]);
197         i++;
198     }
199     while (j + 1 < m) {
200         c.push_back(c.back() + b[(j + 1) % m] - b[j]);
201         j++;
202     }
203     return c;
204 }
205
206 // 过凸多边形外一点求凸多边形的切线，返回切点下标

```

```

207 auto tangent(const vec &a, vector<vec> &p) {
208     int n = p.size();
209     int l = -1, r = -1;
210     for (int i = 0; i < n; i++) {
211         T tmp1 = cross(p[i], p[(i - 1 + n) % n], a);
212         T tmp2 = cross(p[i], p[(i + 1) % n], a);
213         if (l == -1 && tmp1 <= 0 && tmp2 <= 0) l = i;
214         else if (r == -1 && tmp1 >= 0 && tmp2 >= 0) r = i;
215     }
216     return array{l, r};
217 }
218
219 // 直线
220 struct line {
221     vec point, direction;
222     line(const vec &p = vec(), const vec &d = vec()) : point(p), direction(d) {}
223 };
224
225 // 点到直线距离
226 double distance(const vec &a, const line &b) {
227     return abs((b.point - a) ^ (b.point + b.direction - a)) /
228         b.direction.length();
229 }
230
231 // 判断点在直线哪边, 大于0在左边, 等于0在线上, 小于0在右边
232 T side_line(const vec &a, const line &b) { return b.direction ^ (a - b.point); }
233
234 // 两直线是否垂直
235 bool perpendicular(const line &a, const line &b) {
236     return !cmp(a.direction * b.direction, 0);
237 }
238
239 // 两直线是否平行
240 bool parallel(const line &a, const line &b) {
241     return !cmp(a.direction ^ b.direction, 0);
242 }
243
244 // 点的垂线是否与线段有交点
245 bool perpendicular(const vec &a, const line &b) {
246     vec perpen(-b.direction.y, b.direction.x);
247     bool cross1 = (perpen ^ (b.point - a)) > 0;
248     bool cross2 = (perpen ^ (b.point + b.direction - a)) > 0;
249     return cross1 != cross2;
250 }
251
252 // 点到线段距离
253 double distance_seg(const vec &a, const line &b) {
254     if (perpendicular(a, b)) return distance(a, b);
255     return min(distance(a, b.point), distance(a, b.point + b.direction));
256 }
257
258 // 两直线交点

```

```

259 vec intersection(T A, T B, T C, T D, T E, T F) {
260     return {(B * F - C * E) / (A * E - B * D),
261             (C * D - A * F) / (A * E - B * D)};
262 }
263
264 // 两直线交点
265 vec intersection(const line &a, const line &b) {
266     return intersection(a.direction.y, -a.direction.x,
267                         a.direction.x * a.point.y - a.direction.y * a.point.x,
268                         b.direction.y, -b.direction.x,
269                         b.direction.x * b.point.y - b.direction.y * b.point.x);
270 }
271
272 struct circle {
273     vec o;
274     double r;
275     circle(const vec &o, T _r) : o(_o), r(_r){};
276     // 点与圆的关系 -1在圆内, 0在圆上, 1在圆外
277     int relation(const vec &other) const {
278         double len = (other - o).length();
279         return cmp(len, r);
280     }
281     double area() { return PI * r * r; }
282 };
283
284 // 圆与直线交点
285 auto intersection(const circle &c, const line &l) {
286     double d = distance(c.o, l);
287     vector<vec> res;
288     double len = l.direction.length();
289     vec mid = l.point + l.direction * ((c.o - l.point) * l.direction / len);
290     if (!cmp(d, c.r)) res.push_back(mid);
291     else if (d < c.r) {
292         d = sqrt(c.r * c.r - d * d) / len;
293         res.push_back(mid + l.direction * d);
294         res.push_back(mid - l.direction * d);
295     }
296     return res;
297 }
298
299 // oab三角形与圆相交的面积
300 double area(const circle &c, const vec &a, const vec &b) {
301     vec oa = a - c.o, ob = b - c.o;
302     T cab = oa ^ ob;
303     if (!cmp(cab, 0)) return 0;
304     if (c.relation(a) != 1 && c.relation(b) != 1) return cab / 2.0;
305     vec ba = a - b, bo = -ob;
306     vec ab = -ba, ao = -oa;
307     auto r = c.r;
308     double ang;
309     double loa = oa.length(), lob = ob.length(), lab = ab.length();
310     double x =

```

```

311     (ba * bo + sqrt(r * r * lab * lab - (ba ^ bo) * (ba ^ bo))) / lab;
312     double y =
313         (ab * ao + sqrt(r * r * lab * lab - (ab ^ ao) * (ab ^ ao))) / lab;
314     if (cmp(lob, r) == -1 && cmp(loa, r) != -1) {
315         ang = cab * (1 - x / lab) / (r * loa);
316         ang = min(max((double)-1, ang), (double)1);
317         return (asin(ang) * r * r + cab * x / lab) / 2;
318     }
319     if (cmp(lob, r) != -1 && cmp(loa, r) == -1) {
320         ang = cab * (1 - y / lab) / (r * lob);
321         ang = min(max((double)-1, ang), (double)1);
322         return (asin(ang) * r * r + cab * y / lab) / 2;
323     }
324     if (cmp(abs(cab), r * lab) != -1 || cmp(ab * ao, 0) != 1 ||
325         cmp(ba * bo, 0) != 1) {
326         ang = cab / (loa * lob);
327         ang = min(max((double)-1, ang), (double)1);
328         double tmp = -asin(ang);
329         if (cmp(oa * ob, 0) == -1)
330             if (cmp(cab, 0) == -1) tmp -= PI;
331             else tmp += PI;
332         else tmp = -tmp;
333         return tmp * r * r / 2;
334     }
335     ang = cab * (1 - x / lab) / (r * loa);
336     ang = min(max((double)-1, ang), (double)1);
337     double ang2 = cab * (1 - y / lab) / (r * lob);
338     ang2 = min(max((double)-1, ang2), (double)1);
339     return ((asin(ang) + asin(ang2)) * r * r + cab * ((x + y) / lab - 1)) / 2;
340 }
341
342 // 多边形与圆相交的面积
343 double area(vector<vec> &p, circle c) {
344     double res = 0;
345     for (int i = 1; i < p.size(); i++) res += area(c, p[i - 1], p[i]);
346     res += area(c, p.back(), p[0]);
347     return abs(res);
348 }

```

5.1 扫描线

```

1 #define ls (pos << 1)
2 #define rs (ls | 1)
3 #define mid ((tree[pos].l + tree[pos].r) >> 1)
4 struct Rectangle {
5     ll x_l, y_l, x_r, y_r;
6 };
7 ll area(vector<Rectangle>& rec) {
8     struct Line {
9         ll x, y_up, y_down;
10        int pd;

```



```

11     };
12     vector<Line> line(rec.size() * 2);
13     vector<ll> y_set(rec.size() * 2);
14     for (int i = 0; i < rec.size(); i++) {
15         y_set[i * 2] = rec[i].y_l;
16         y_set[i * 2 + 1] = rec[i].y_r;
17         line[i * 2] = {rec[i].x_l, rec[i].y_r, rec[i].y_l, 1};
18         line[i * 2 + 1] = {rec[i].x_r, rec[i].y_r, rec[i].y_l, -1};
19     }
20     sort(y_set.begin(), y_set.end());
21     y_set.erase(unique(y_set.begin(), y_set.end()), y_set.end());
22     sort(line.begin(), line.end(), [](Line a, Line b) { return a.x < b.x; });
23     struct Data {
24         int l, r;
25         ll len, cnt, raw_len;
26     };
27     vector<Data> tree(4 * y_set.size());
28     function<void(int, int, int)> build = [&](int pos, int l, int r) {
29         tree[pos].l = l;
30         tree[pos].r = r;
31         if (l == r) {
32             tree[pos].raw_len = y_set[r + 1] - y_set[l];
33             tree[pos].cnt = tree[pos].len = 0;
34             return;
35         }
36         build(ls, l, mid);
37         build(rs, mid + 1, r);
38         tree[pos].raw_len = tree[ls].raw_len + tree[rs].raw_len;
39     };
40     function<void(int, int, int, int)> update = [&](int pos, int l, int r,
41                                                     int num) {
42         if (l <= tree[pos].l && tree[pos].r <= r) {
43             tree[pos].cnt += num;
44             tree[pos].len = tree[pos].cnt ? tree[pos].raw_len
45                             : tree[pos].l == tree[pos].r
46                             ? 0
47                             : tree[ls].len + tree[rs].len;
48             return;
49         }
50         if (l <= mid) update(ls, l, r, num);
51         if (r > mid) update(rs, l, r, num);
52         tree[pos].len =
53             tree[pos].cnt ? tree[pos].raw_len : tree[ls].len + tree[rs].len;
54     };
55     build(1, 0, y_set.size() - 2);
56     auto find_pos = [&](ll num) {
57         return lower_bound(y_set.begin(), y_set.end(), num) - y_set.begin();
58     };
59     ll res = 0;
60     for (int i = 0; i < line.size() - 1; i++) {
61         update(1, find_pos(line[i].y_down), find_pos(line[i].y_up) - 1,
62               line[i].pd);

```

```
63     res += (line[i + 1].x - line[i].x) * tree[1].len;  
64 }  
65 return res;  
66 }
```

6 杂项

6.1 高精度

```

1 struct bignum {
2     string num;
3
4     bignum() : num("0") {}
5     bignum(const string& num) : num(num) {
6         reverse(this->num.begin(), this->num.end());
7     }
8     bignum(ll num) : num(to_string(num)) {
9         reverse(this->num.begin(), this->num.end());
10    }
11
12    bignum operator+(const bignum& other) {
13        bignum res;
14        res.num.pop_back();
15        res.num.reserve(max(num.size(), other.num.size()) + 1);
16        for (int i = 0, j = 0, x; i < num.size() || i < other.num.size() || j;
17             i++) {
18            x = j;
19            j = 0;
20            if (i < num.size()) x += num[i] - '0';
21            if (i < other.num.size()) x += other.num[i] - '0';
22            if (x >= 10) j = 1, x -= 10;
23            res.num.push_back(x + '0');
24        }
25        res.num.capacity();
26        return res;
27    }
28
29    bignum operator*(const bignum& other) {
30        vector<int> res(num.size() + other.num.size() - 1, 0);
31        for (int i = 0; i < num.size(); i++)
32            for (int j = 0; j < other.num.size(); j++)
33                res[i + j] += (num[i] - '0') * (other.num[j] - '0');
34        int g = 0;
35        for (int i = 0; i < res.size(); i++) {
36            res[i] += g;
37            g = res[i] / 10;
38            res[i] %= 10;
39        }
40        while (g) {
41            res.push_back(g % 10);
42            g /= 10;
43        }
44        int lim = res.size();
45        while (lim > 1 && res[lim - 1] == 0) lim--;
46        bignum res2;
47        res2.num.resize(lim);
48        for (int i = 0; i < lim; i++) res2.num[i] = res[i] + '0';

```

```

49     return res2;
50 }
51
52 bool operator<(const bignum& other) {
53     if (num.size() == other.num.size())
54         for (int i = num.size() - 1; i >= 0; i--)
55             if (num[i] == other.num[i]) continue;
56             else return num[i] < other.num[i];
57     return num.size() < other.num.size();
58 }
59
60 friend istream& operator>>(istream& in, bignum& a) {
61     in >> a.num;
62     reverse(a.num.begin(), a.num.end());
63     return in;
64 }
65 friend ostream& operator<<(ostream& out, bignum a) {
66     reverse(a.num.begin(), a.num.end());
67     return out << a.num;
68 }
69 };

```

6.2 模运算

```

1 struct modint {
2     int x;
3     modint(ll _x = 0) : x(_x % mod) {}
4     modint inv() const { return power(*this, mod - 2); }
5     modint operator+(const modint& b) { return {x + b.x}; }
6     modint operator-() const { return {-x}; }
7     modint operator-(const modint& b) { return {-b + *this}; }
8     modint operator*(const modint& b) { return {(ll)x * b.x}; }
9     modint operator/(const modint& b) { return *this * b.inv(); }
10    friend istream& operator>>(istream& is, modint& other) {
11        ll _x;
12        is >> _x;
13        other = modint(_x);
14        return is;
15    }
16    friend ostream& operator<<(ostream& os, modint other) {
17        other.x = (other.x + mod) % mod;
18        return os << other.x;
19    }
20 };

```

6.3 分数

```

1 struct frac {
2     ll a, b;
3     frac() : a(0), b(1) {}

```

```

4   frac(ll _a, ll _b) : a(_a), b(_b) {
5       assert(b);
6       if (a) {
7           int tmp = gcd(a, b);
8           a /= tmp;
9           b /= tmp;
10      } else *this = frac();
11  }
12  frac operator+(const frac& other) {
13      return frac(a * other.b + other.a * b, b * other.b);
14  }
15  frac operator-() const {
16      frac res = *this;
17      res.a = -res.a;
18      return res;
19  }
20  frac operator-(const frac& other) const { return -other + *this; }
21  frac operator*(const frac& other) const {
22      return frac(a * other.a, b * other.b);
23  }
24  frac operator/(const frac& other) const {
25      assert(other.a);
26      return *this * frac(other.b, other.a);
27  }
28  bool operator<(const frac& other) const { return (*this - other).a < 0; }
29  bool operator<=(const frac& other) const { return (*this - other).a <= 0; }
30  bool operator>=(const frac& other) const { return (*this - other).a >= 0; }
31  bool operator>(const frac& other) const { return (*this - other).a > 0; }
32  bool operator==(const frac& other) const {
33      return a == other.a && b == other.b;
34  }
35  bool operator!=(const frac& other) const { return !(*this == other); }
36 };

```

6.4 表达式求值

```

1  // 格式化表达式
2  string format(const string& s1) {
3      stringstream ss(s1);
4      string s2;
5      char ch;
6      while ((ch = ss.get()) != EOF) {
7          if (ch == ' ') continue;
8          if (isdigit(ch)) s2 += ch;
9          else {
10             if (s2.back() != ' ') s2 += ' ';
11             s2 += ch;
12             s2 += ' ';
13         }
14     }
15     return s2;

```

```
16 }
17
18 // 中缀表达式转后缀表达式
19 string convert(const string& s1) {
20     unordered_map<char, int> rank{
21         {'+', 2}, {'-', 2}, {'*', 1}, {'/', 1}, {'^', 0}};
22     stringstream ss(s1);
23     string s2, temp;
24     stack<char> op;
25     while (ss >> temp) {
26         if (isdigit(temp[0])) s2 += temp + ' ';
27         else if (temp[0] == '(') op.push('(');
28         else if (temp[0] == ')') {
29             while (op.top() != '(') {
30                 s2 += op.top();
31                 s2 += ' ';
32                 op.pop();
33             }
34             op.pop();
35         } else {
36             while (!op.empty() && op.top() != '(' &&
37                 (temp[0] != '^' && rank[op.top()] <= rank[temp[0]] ||
38                 rank[op.top()] < rank[temp[0]])) {
39                 s2 += op.top();
40                 s2 += ' ';
41                 op.pop();
42             }
43             op.push(temp[0]);
44         }
45     }
46     while (!op.empty()) {
47         s2 += op.top();
48         s2 += ' ';
49         op.pop();
50     }
51     return s2;
52 }
53
54 // 计算后缀表达式
55 int calc(const string& s) {
56     stack<int> num;
57     stringstream ss(s);
58     string temp;
59     while (ss >> temp) {
60         if (isdigit(temp[0])) num.push(stoi(temp));
61         else {
62             int b = num.top();
63             num.pop();
64             int a = num.top();
65             num.pop();
66             if (temp[0] == '+') a += b;
67             else if (temp[0] == '-') a -= b;
```

```

68         else if (temp[0] == '*') a *= b;
69         else if (temp[0] == '/') a /= b;
70         else if (temp[0] == '^') a = ksm(a, b);
71         num.push(a);
72     }
73 }
74 return num.top();
75 }

```

6.5 日期

```

1 int month[] = {0, 31, 28, 31, 30, 31, 30, 31, 31, 30, 31, 30, 31};
2 int pre[13];
3 vector<int> leap;
4 struct Date {
5     int y, m, d;
6     bool operator<(const Date& other) const {
7         return array<int, 3>{y, m, d} <
8             array<int, 3>{other.y, other.m, other.d};
9     }
10    Date(const string& s) {
11        stringstream ss(s);
12        char ch;
13        ss >> y >> ch >> m >> ch >> d;
14    }
15    int dis() const {
16        int yd = (y - 1) * 365 +
17            (upper_bound(leap.begin(), leap.end(), y - 1) - leap.begin());
18        int md =
19            pre[m - 1] + (m > 2 && (y % 4 == 0 && y % 100 || y % 400 == 0));
20        return yd + md + d;
21    }
22    int dis(const Date& other) const { return other.dis() - dis(); }
23 };
24 for (int i = 1; i <= 12; i++) pre[i] = pre[i - 1] + month[2];
25 for (int i = 1; i <= 1000000; i++)
26     if (i % 4 == 0 && i % 100 || i % 400 == 0) leap.push_back(i);

```

6.6 对拍

linux/Mac

```

1 g++ a.cpp -o program/a -O2 -std=c++17
2 g++ b.cpp -o program/b -O2 -std=c++17
3 g++ suiji.cpp -o program/suiji -O2 -std=c++17
4
5 cnt=0
6
7 while true; do
8     let cnt++
9     echo TEST:$cnt

```

```

10
11 ./program/suiji > in
12 ./program/a < in > out.a
13 ./program/b < in > out.b
14
15 diff out.a out.b
16 if [ $? -ne 0 ];then break;fi
17 done

```

windows

```

1 @echo off
2
3 g++ a.cpp -o program/a -O2 -std=c++17
4 g++ b.cpp -o program/b -O2 -std=c++17
5 g++ suiji.cpp -o program/suiji -O2 -std=c++17
6
7 set cnt=0
8
9 :again
10 set /a cnt=cnt+1
11 echo TEST:%cnt%
12 .\program\suiji > in
13 .\program\a < in > out.a
14 .\program\b < in > out.b
15
16 fc output.a output.b
17 if not errorlevel 1 goto again

```

6.7 编译常用选项

```
1 -Wall -Woverflow -Wextra -Wpedantic -Wfloat-equal -Wshadow -fsanitize=address,undefined
```

6.8 开栈

不同的编译器可能命令不一样

```

1 -Wl,--stack=0x10000000
2 -Wl,-stack_size -Wl,0x10000000
3 -Wl,-z,stack-size=0x10000000

```

6.9 clang-format

```

1 BasedOnStyle: Google
2 IndentWidth: 4
3 ColumnLimit: 80
4 AllowShortIfStatementsOnASingleLine: AllIfsAndElse
5 AccessModifierOffset: -4
6 EmptyLineBeforeAccessModifier: Leave
7 RemoveBracesLLVM: true

```