# ACM 常用算法模板

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## 1 数据结构

#### 1.1 并查集

```
struct dsu {
2
      int n;
3
      vector<int> fa, sz;
4
      dsu(int _n) : n(_n), fa(n + 1), sz(n + 1, 1) { iota(fa.begin(), fa.end(), 0); }
5
      int find(int x) { return x == fa[x] ? x : fa[x] = find(fa[x]); }
6
      int merge(int x, int y) {
7
          int fax = find(x), fay = find(y);
8
          if (fax == fay) return 0; // 一个集合
9
          sz[fay] += sz[fax];
10
          return fa[fax] = fay; // 合并到哪个集合了
11
12
      int size(int x) { return sz[find(x)]; }
13 };
```

#### 1.2 树状数组

#### 1.2.1 一维

```
template <class T>
 2
   struct fenwick {
3
       int n;
 4
       vector<T> t;
 5
       fenwick(int _n) : n(_n), t(n + 1) {}
 6
       T query(int 1, int r) {
 7
           auto query = [&](int pos) {
 8
                T res = 0;
9
                while (pos) {
10
                    res += t[pos];
11
                    pos -= lowbit(pos);
12
                }
13
                return res;
14
           };
15
           return query(r) - query(l - 1);
16
17
       void add(int pos, T num) {
18
           while (pos <= n) {</pre>
19
                t[pos] += num;
20
                pos += lowbit(pos);
21
22
       }
23 };
```

#### 1.2.2 二维

```
template <class T>
struct Fenwick_tree_2 {
    Fenwick_tree_2(int n, int m) : n(n), m(m), tree(n + 1, vector<T>(m + 1)) {}

T query(int l1, int r1, int l2, int r2) {
    auto query = [&](int l, int r) {
    T res = 0;
    for (int i = 1; i; i -= lowbit(i))
```

```
8
                for (int j = r; j; j -= lowbit(j)) res += tree[i][j];
9
            return res;
10
         11
12
13
     void update(int x, int y, T num) {
         for (int i = x; i <= n; i += lowbit(i))</pre>
14
15
            for (int j = y; j \leftarrow m; j \leftarrow lowbit(j)) tree[i][j] += num;
16
17
  private:
18
     int n, m;
19
     vector<vector<T>> tree;
20 };
```

#### 1.2.3 三维

```
1
  template <class T>
2
   struct Fenwick_tree_3 {
3
       Fenwick_tree_3(int n, int m, int k)
 4
           : n(n), m(m), k(k), tree(n + 1, vector<vector<T>>(m + 1, vector<T>(k + 1))) {}
5
       T query(int a, int b, int c, int d, int e, int f) {
6
           auto query = [&](int x, int y, int z) {
7
               T res = 0;
8
               for (int i = x; i; i -= lowbit(i))
9
                   for (int j = y; j; j -= lowbit(j))
10
                        for (int p = z; p; p -= lowbit(p)) res += tree[i][j][p];
11
               return res;
12
           };
13
           T res = query(d, e, f);
14
           res -= query(a - 1, e, f) + query(d, b - 1, f) + query(d, e, c - 1);
15
           res += query(a - 1, b - 1, f) + query(a - 1, e, c - 1) + query(d, b - 1, c - 1);
           res -= query(a - 1, b - 1, c - 1);
16
17
           return res;
18
19
       void update(int x, int y, int z, T num) {
20
           for (int i = x; i <= n; i += lowbit(i))</pre>
21
               for (int j = y; j <= m; j += lowbit(j))</pre>
22
                   for (int p = z; p \leftarrow k; p += lowbit(p)) tree[i][j][p] += num;
23
24
  private:
25
       int n, m, k;
26
       vector<vector<T>>> tree;
27
  };
```

#### 1.3 线段树

```
template <class Data, class Num>
2
  struct Segment_Tree {
3
       inline void update(int l, int r, Num x) { update(1, l, r, x); }
      inline Data query(int 1, int r) { return query(1, 1, r); }
4
5
      Segment_Tree(vector<Data>& a) {
           n = a.size();
6
7
           tree.assign(n * 4 + 1, \{\});
8
           build(a, 1, 1, n);
9
      }
10 private:
```

```
11
       int n;
       struct Tree {
12
13
           int 1, r;
14
           Data data;
15
16
       vector<Tree> tree;
17
       inline void pushup(int pos) { tree[pos].data = tree[pos << 1].data + tree[pos << 1 | 1].data; }</pre>
18
       inline void pushdown(int pos) {
19
           tree[pos << 1].data = tree[pos << 1].data + tree[pos].data.lazytag;</pre>
20
           tree[pos << 1 | 1].data = tree[pos << 1 | 1].data + tree[pos].data.lazytag;</pre>
21
           tree[pos].data.lazytag = Num::zero();
22
23
       void build(vector<Data>& a, int pos, int 1, int r) {
24
           tree[pos].1 = 1;
25
           tree[pos].r = r;
26
           if (1 == r) {
27
               tree[pos].data = a[l - 1];
28
               return;
29
           }
30
           int mid = (tree[pos].1 + tree[pos].r) >> 1;
31
           build(a, pos << 1, 1, mid);</pre>
32
           build(a, pos << 1 | 1, mid + 1, r);
33
           pushup(pos);
34
35
       void update(int pos, int& 1, int& r, Num& x) {
36
           if (1 > tree[pos].r || r < tree[pos].1) return;</pre>
37
           if (1 <= tree[pos].1 && tree[pos].r <= r) {</pre>
38
               tree[pos].data = tree[pos].data + x;
39
               return;
40
           }
41
           pushdown(pos);
42
           update(pos << 1, l, r, x);
43
           update(pos << 1 | 1, 1, r, x);
44
           pushup(pos);
45
46
       Data query(int pos, int& 1, int& r) {
47
           if (1 > tree[pos].r || r < tree[pos].l) return Data::zero();</pre>
48
           if (1 <= tree[pos].1 && tree[pos].r <= r) return tree[pos].data;</pre>
49
           pushdown(pos);
50
           return query(pos << 1, 1, r) + query(pos << 1 | 1, 1, r);
51
       }
52 };
53
  struct Num {
54
       11 add;
55
       inline static Num zero() { return {0}; }
56
       inline Num operator+(Num b) { return {add + b.add}; }
57 };
58
  struct Data {
59
       11 sum, len;
60
       Num lazytag;
61
       inline static Data zero() { return {0, 0, Num::zero()}; }
62
       inline Data operator+(Num b) { return {sum + len * b.add, len, lazytag + b}; }
63
       inline Data operator+(Data b) { return {sum + b.sum, len + b.len, Num::zero()}; }
64
  };
```

#### 1.4 普通平衡树

#### 1.4.1 树状数组实现

需要预先处理出来所有可能的数。

```
int lowbit(int x) { return x & -x; }
2
 3
  template <typename T>
4
   struct treap {
 5
       int n, size;
 6
       vector<int> t;
 7
       vector<T> t2, S;
 8
       treap(const vector<T>& a) : S(a) {
9
           sort(S.begin(), S.end());
           S.erase(unique(S.begin(), S.end()), S.end());
10
11
           n = S.size();
12
           size = 0;
13
           t = vector<int>(n + 1);
14
           t2 = vector < T > (n + 1);
15
       }
16
       int pos(T x) { return lower_bound(S.begin(), S.end(), x) - S.begin() + 1; }
17
       int sum(int pos) {
           int res = 0;
18
19
           while (pos) {
20
               res += t[pos];
21
               pos -= lowbit(pos);
22
           }
23
           return res;
24
       }
25
26
       // 插入cnt个x
27
       void insert(T x, int cnt) {
28
           size += cnt;
29
           int i = pos(x);
30
           assert(i <= n && S[i - 1] == x);
31
           for (; i <= n; i += lowbit(i)) {</pre>
32
               t[i] += cnt;
33
               t2[i] += cnt * x;
34
           }
35
       }
36
37
       // 删除cnt个x
38
       void erase(T x, int cnt) {
39
           assert(cnt <= count(x));</pre>
40
           insert(x, -cnt);
41
       }
42
43
       // x的排名
44
       int rank(T x) {
45
           assert(count(x));
46
           return sum(pos(x) - 1) + 1;
47
48
       // 统计出现次数
49
50
       int count(T x) { return sum(pos(x)) - sum(pos(x) - 1); }
51
52
       // 第k小
53
       T kth(int k) {
```

```
54
           assert(0 < k && k <= size);
55
           int cnt = 0, x = 0;
56
           for (int i = __lg(n); i >= 0; i--) {
57
               x += 1 << i;
               if (x >= n \mid \mid cnt + t[x] >= k) x -= 1 << i;
58
59
               else cnt += t[x];
60
           }
61
           return S[x];
62
       }
63
64
       // 前k小的数之和
65
       T pre_sum(int k) {
66
           assert(0 < k && k <= size);
67
           int cnt = 0, x = 0;
68
           T res = 0;
69
           for (int i = __lg(n); i >= 0; i--) {
70
               x += 1 << i;
71
               if (x >= n \mid | cnt + t[x] >= k) x -= 1 << i;
72
               else {
73
                   cnt += t[x];
                   res += t2[x];
74
75
               }
76
           }
77
           return res + (k - cnt) * S[x];
78
       }
79
80
       // 小于x, 最大的数
81
       T prev(T x) { return kth(sum(pos(x) - 1)); }
82
83
       // 大于x, 最小的数
84
       T next(T x) { return kth(sum(pos(x)) + 1); }
85
  };
```

#### 1.5 可持久化线段树

```
1 constexpr int MAXN = 200000;
  vector<int> root(MAXN << 5);</pre>
3
  struct Persistent_seg {
4
       int n;
5
       struct Data {
6
           int ls, rs;
7
           int val;
8
9
       vector<Data> tree;
10
       Persistent_seg(int n, vector<int>& a) : n(n) { root[0] = build(1, n, a); }
11
       int build(int 1, int r, vector<int>& a) {
12
           if (1 == r) {
13
               tree.push_back({0, 0, a[1]});
14
               return tree.size() - 1;
15
           }
16
           int mid = 1 + r \gg 1;
17
           int ls = build(l, mid, a), rs = build(mid + 1, r, a);
18
           tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
19
           return tree.size() - 1;
20
21
       int update(int rt, const int& idx, const int& val, int l, int r) {
22
           if (1 == r) {
```

```
23
               tree.push_back({0, 0, tree[rt].val + val});
24
               return tree.size() - 1;
25
26
           int mid = 1 + r >> 1, ls = tree[rt].ls, rs = tree[rt].rs;
27
           if (idx <= mid) ls = update(ls, idx, val, l, mid);</pre>
28
           else rs = update(rs, idx, val, mid + 1, r);
29
           tree.push_back({ls, rs, tree[ls].val + tree[rs].val});
30
           return tree.size() - 1;
31
       }
32
       int query(int rt1, int rt2, int k, int 1, int r) {
33
           if (1 == r) return 1;
34
           int mid = 1 + r \gg 1;
35
           int lcnt = tree[tree[rt2].ls].val - tree[tree[rt1].ls].val;
36
           if (k <= lcnt) return query(tree[rt1].ls, tree[rt2].ls, k, l, mid);</pre>
37
           else return query(tree[rt1].rs, tree[rt2].rs, k - lcnt, mid + 1, r);
38
39
  };
```

#### 1.6 st 表

```
auto lg = []() {
2
       array<int, 10000001> lg;
3
       lg[1] = 0;
4
       for (int i = 2; i \leftarrow 10000000; i++) lg[i] = lg[i >> 1] + 1;
5
       return lg;
6
  }();
  template <typename T>
   struct st {
8
9
       int n;
10
       vector<vector<T>> a;
11
       st(vector<T>& _a) : n(_a.size()) {
12
           a.assign(lg[n] + 1, vector<int>(n));
           for (int i = 0; i < n; i++) a[0][i] = _a[i];</pre>
13
14
           for (int j = 1; j <= lg[n]; j++)</pre>
15
                for (int i = 0; i + (1 << j) - 1 < n; i++)
16
                    a[j][i] = max(a[j - 1][i], a[j - 1][i + (1 << (j - 1))]);
17
18
       T query(int 1, int r) {
19
           int k = lg[r - l + 1];
20
           return max(a[k][1], a[k][r - (1 << k) + 1]);</pre>
21
       }
22 };
```

# 2 图论

存图

```
1
   struct Graph {
2
       int n;
3
       struct Edge {
 4
           int to, w;
5
6
       vector<vector<Edge>> graph;
7
       Graph(int _n) {
8
           n = _n;
9
           graph.assign(n + 1, vector<Edge>());
10
11
       void add(int u, int v, int w) { graph[u].push_back({v, w}); }
12 };
```

#### 2.1 最短路

#### 2.1.1 dijkstra

```
void dij(Graph& graph, vector<int>& dis, int t) {
2
       vector<int> visit(graph.n + 1, 0);
3
       priority_queue<pair<int, int>> que;
4
       dis[t] = 0;
5
       que.emplace(0, t);
6
       while (!que.empty()) {
7
           int u = que.top().second;
8
           que.pop();
9
           if (visit[u]) continue;
10
           visit[u] = 1;
11
           for (auto& [to, w] : graph.graph[u]) {
12
               if (dis[to] > dis[u] + w) {
13
                   dis[to] = dis[u] + w;
14
                   que.emplace(-dis[to], to);
15
               }
16
           }
17
       }
18 }
```

#### 2.2 树上问题

#### 2.2.1 最近公公祖先

倍增法

```
vector<int> dep;
  vector<array<int, 21>> fa;
3
  dep.assign(n + 1, 0);
4
  fa.assign(n + 1, array<int, 21>{});
5
  void binary_jump(int root) {
6
      function<void(int)> dfs = [&](int t) {
7
           dep[t] = dep[fa[t][0]] + 1;
8
           for (auto& [to] : graph[t]) {
9
               if (to == fa[t][0]) continue;
10
               fa[to][0] = t;
11
               dfs(to);
```

```
12
13
       };
14
       dfs(root);
15
       for (int j = 1; j <= 20; j++)
16
           for (int i = 1; i \leftarrow n; i++) fa[i][j] = fa[fa[i][j - 1]][j - 1];
17
   int lca(int x, int y) {
18
19
       if (dep[x] < dep[y]) swap(x, y);</pre>
20
       for (int i = 20; i >= 0; i--)
21
           if (dep[fa[x][i]] >= dep[y]) x = fa[x][i];
22
       if (x == y) return x;
23
       for (int i = 20; i >= 0; i--) {
24
           if (fa[x][i] != fa[y][i]) {
25
                x = fa[x][i];
26
                y = fa[y][i];
27
28
       return fa[x][0];
29
30 }
```

树剖

```
int lca(int x, int y) {
    while (top[x] != top[y]) {
        if (dep[top[x]] < dep[top[y]]) swap(x, y);
        x = fa[top[x]];
    }
    if (dep[x] < dep[y]) swap(x, y);
    return y;
}</pre>
```

#### 2.2.2 树链剖分

```
1 vector<int> fa, siz, dep, son, dfn, rnk, top;
2 fa.assign(n + 1, 0);
3 siz.assign(n + 1, 0);
4
  dep.assign(n + 1, 0);
5
  son.assign(n + 1, 0);
6
  dfn.assign(n + 1, 0);
7
  rnk.assign(n + 1, 0);
  top.assign(n + 1, 0);
8
   void hld(int root) {
9
10
       function < void(int) > dfs1 = [\&](int t) {
11
           dep[t] = dep[fa[t]] + 1;
12
           siz[t] = 1;
           for (auto& [to, w] : graph[t]) {
13
14
               if (to == fa[t]) continue;
15
               fa[to] = t;
16
               dfs1(to);
17
               if (siz[son[t]] < siz[to]) son[t] = to;</pre>
18
               siz[t] += siz[to];
19
           }
20
       };
21
       dfs1(root);
22
       int dfn_tail = 0;
23
       for (int i = 1; i <= n; i++) top[i] = i;</pre>
24
       function<void(int)> dfs2 = [&](int t) {
           dfn[t] = ++dfn_tail;
```

```
26
            rnk[dfn_tail] = t;
27
           if (!son[t]) return;
28
           top[son[t]] = top[t];
29
           dfs2(son[t]);
           for (auto& [to, w] : graph[t]) {
30
                if (to == fa[t] || to == son[t]) continue;
31
32
                dfs2(to);
33
34
       };
35
       dfs2(root);
36
```

#### 2.3 强连通分量

```
void tarjan(Graph& g1, Graph& g2) {
2
       int dfn_tail = 0, cnt = 0;
3
       vector < int > dfn(g1.n + 1, 0), low(g1.n + 1, 0), exist(g1.n + 1, 0), belong(g1.n + 1, 0);
4
       stack<int> sta;
 5
       function<void(int)> dfs = [&](int t) {
6
           dfn[t] = low[t] = ++dfn_tail;
7
           sta.push(t);
8
           exist[t] = 1;
9
           for (auto& [to] : g1.graph[t])
10
               if (!dfn[to]) {
11
                    dfs(to);
12
                    low[t] = min(low[t], low[to]);
13
               } else if (exist[to]) low[t] = min(low[t], dfn[to]);
14
           if (dfn[t] == low[t]) {
15
               cnt++;
16
               while (int temp = sta.top()) {
17
                    belong[temp] = cnt;
18
                    exist[temp] = 0;
19
                    sta.pop();
                    if (temp == t) break;
20
               }
21
22
           }
23
       };
24
       for (int i = 1; i <= g1.n; i++)</pre>
25
           if (!dfn[i]) dfs(i);
26
       g2 = Graph(cnt);
27
       for (int i = 1; i <= g1.n; i++) g2.w[belong[i]] += g1.w[i];</pre>
28
       for (int i = 1; i <= g1.n; i++)
29
           for (auto& [to] : g1.graph[i])
30
               if (belong[i] != belong[to]) g2.add(belong[i], belong[to]);
31
```

#### 2.4 拓扑排序

```
dis[i] = g.w[i]; // dp
9
10
         }
11
       while (!que.empty()) {
12
          int u = que.front();
13
          que.pop();
14
           for (auto& [to] : g.graph[u]) {
15
               in[to]--;
               dis[to] = max(dis[to], dis[u] + g.w[to]); // dp
16
17
              if (!in[to]) que.push(to);
18
          }
19
      }
20 }
```

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# 3 字符串

#### 3.1 kmp

```
auto kmp(string& s) {
2
       vector next(s.size(), -1);
3
       for (int i = 1, j = -1; i < s.size(); i++) {</pre>
4
           while (j \ge 0 \&\& s[i] != s[j + 1]) j = next[j];
5
           if (s[i] == s[j + 1]) j++;
           next[i] = j;
6
7
8
       // next 意为长度
9
       for (auto& i : next) i++;
10
       return next;
11 }
```

#### 3.2 哈希

```
1 constexpr int N = 1e6;
  int pow_base[N + 1][2];
  constexpr 11 mod[2] = {(int)2e9 + 11, (int)2e9 + 33}, base[2] = {(int)2e5 + 11, (int)2e5 + 33};
3
5
   struct Hash {
6
       int size;
 7
       vector<array<int, 2>> a;
8
       Hash() {}
       Hash(const string& s) {
9
10
           size = s.size();
11
           a.resize(size);
           a[0][0] = a[0][1] = s[0];
12
13
           for (int i = 1; i < size; i++) {</pre>
               a[i][0] = (a[i - 1][0] * base[0] + s[i]) % mod[0];
14
15
               a[i][1] = (a[i - 1][1] * base[1] + s[i]) % mod[1];
16
           }
17
18
       array<int, 2> get(int 1, int r) const {
19
           if (1 == 0) return a[r];
20
           auto getone = [&](bool f) {
               int x = (a[r][f] - 111 * a[1 - 1][f] * pow_base[r - 1 + 1][f]) % mod[f];
21
22
               if (x < 0) x += mod[f];
23
               return x;
24
25
           return {getone(0), getone(1)};
26
27
  };
28
29
  auto _ = []() {
30
       pow_base[0][0] = pow_base[0][1] = 1;
31
       for (int i = 1; i <= N; i++) {
32
           pow_base[i][0] = pow_base[i - 1][0] * base[0] % mod[0];
           pow_base[i][1] = pow_base[i - 1][1] * base[1] % mod[1];
33
34
35
       return true;
36 }();
```

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#### 3.3 manacher

```
1 auto manacher(const string& _s) {
2
       string s(_s.size() * 2 + 1, '$');
3
       for (int i = 0; i < _s.size(); i++) s[2 * i + 1] = _s[i];</pre>
4
       vector r(s.size(), 0);
5
       for (int i = 0, maxr = 0, mid = 0; i < s.size(); i++) {</pre>
6
           if (i < maxr) r[i] = min(r[mid * 2 - i], maxr - i);</pre>
7
           while (i - r[i] - 1 >= 0 && i + r[i] + 1 < s.size() && s[i - r[i] - 1] == s[i + r[i] + 1])
8
               ++r[i];
9
           if (i + r[i] > maxr) maxr = i + r[i], mid = i;
10
11
       return r;
12 }
```

# 4 数学

#### 4.1 扩展欧几里得

```
需保证 a,b>=0 x=x+k*dx, y=y-k*dy 若要求 x\geq p,\ k\geq \left\lceil\frac{p-x}{dx}\right\rceil 若要求 x\leq q,\ k\leq \left\lfloor\frac{q-x}{dx}\right\rfloor 若要求 y\geq p,\ k\leq \left\lfloor\frac{y-p}{dy}\right\rfloor 若要求 y\leq q,\ k\geq \left\lceil\frac{y-q}{dy}\right\rceil
```

```
int __exgcd(int a, int b, int& x, int& y) {
2
       if (!b) {
3
           x = 1;
           y = 0;
4
5
           return a;
6
7
       int g = __exgcd(b, a % b, y, x);
8
       y -= a / b * x;
9
       return g;
10
11
  array<int, 2> exgcd(int a, int b, int c) {
12
13
      int x, y;
       int g = \__exgcd(a, b, x, y);
14
15
       if (c % g) return {INT_MAX, INT_MAX};
16
       int dx = b / g;
17
      int dy = a / g;
       x = c / g % dx * x % dx;
18
19
       if (x < 0) x += dx;
20
       y = (c - a * x) / b;
21
       return {x, y};
22 }
```

#### 4.2 线性代数

#### 4.2.1 向量公约数

```
1 // 将这两个向量组转化为b.y=0的形式
2
  array<vec, 2> gcd(vec a, vec b) {
3
      while (b.y != 0) {
4
          int t = a.y / b.y;
5
          a = a - b * t;
6
          swap(a, b);
7
8
      return {a, b};
9
10
11
  array<vec, 2> gcd(array<vec, 2> g, vec a) {
12
      auto [b, c] = gcd(g[0], a);
13
      g[0] = b;
14
      g[1] = vec(gcd(g[1].x, c.x), 0);
15
      if (g[1].x != 0) g[0].x %= g[1].x;
16
      return g;
17 }
```

#### 4.3 线性筛法

```
1 constexpr int N = 10000000;
  array<int, N + 1> min_prime;
3
  vector<int> primes;
4
  bool ok = []() {
       for (int i = 2; i <= N; i++) {</pre>
5
6
           if (min_prime[i] == 0) {
7
               min_prime[i] = i;
               primes.push_back(i);
8
9
           }
10
           for (auto& j : primes) {
11
               if (j > min_prime[i] || j > N / i) break;
12
               min_prime[j * i] = j;
13
14
15
       return 1;
16 }();
```

#### 4.4 分解质因数

```
auto getprimes(int n) {
2
       vector<array<int, 2>> res;
3
       for (auto& i : primes) {
4
           if (i > n / i) break;
5
           if (n % i == 0) {
6
               res.push_back({i, 0});
7
               while (n % i == 0) {
8
                    n /= i;
9
                    res.back()[1]++;
10
               }
11
           }
12
13
       if (n > 1) res.push_back({n, 1});
14
       return res;
15 }
```

#### 4.5 pollard rho

```
using LL = __int128_t;
3
  random_device rd;
  mt19937 seed(rd());
5
  11 power(11 a, 11 b, 11 mod) {
6
7
      ll res = 1;
8
       while (b) {
9
           if (b & 1) res = (LL)res * a % mod;
10
           a = (LL)a * a % mod;
11
           b >>= 1;
12
13
       return res;
14
15
16 bool isprime(ll n) {
```

```
17
       static array primes{2, 3, 5, 7, 11, 13, 17, 19, 23};
18
       static unordered_map<11, bool> S;
19
       if (n < 2) return 0;</pre>
20
       if (S.count(n)) return S[n];
21
       11 d = n - 1, r = 0;
22
       while (!(d & 1)) {
23
           r++;
24
           d >>= 1;
25
       }
26
       for (auto& a : primes) {
27
           if (a == n) return S[n] = 1;
28
           11 x = power(a, d, n);
29
           if (x == 1 \mid \mid x == n - 1) continue;
30
           for (int i = 0; i < r - 1; i++) {</pre>
31
                x = (LL)x * x % n;
32
                if (x == n - 1) break;
33
34
           if (x != n - 1) return S[n] = 0;
35
36
       return S[n] = 1;
37
38
   11 pollard_rho(11 n) {
39
40
       11 s = 0, t = 0;
41
       11 c = seed() % (n - 1) + 1;
42
       ll val = 1;
       for (int goal = 1;; goal *= 2, s = t, val = 1) {
43
           for (int step = 1; step <= goal; step++) {</pre>
44
45
                t = ((LL)t * t + c) % n;
                val = (LL)val * abs(t - s) % n;
46
47
                if (step % 127 == 0) {
48
                    ll g = gcd(val, n);
                    if (g > 1) return g;
49
50
                }
51
           }
52
           11 g = gcd(val, n);
53
           if (g > 1) return g;
54
55
56
   auto getprimes(ll n) {
57
       unordered_set<11> S;
       auto get = [&](auto self, ll n) {
58
59
           if (n < 2) return;</pre>
           if (isprime(n)) {
60
61
                S.insert(n);
                return;
62
63
           }
64
           11 mx = pollard_rho(n);
           self(self, n / mx);
65
66
           self(self, mx);
67
       };
68
       get(get, n);
69
       return S;
70
```

#### 4.6 组合数

```
1 constexpr int N = 1e6;
2 array<modint, N + 1> fac, ifac;
4
  modint C(int n, int m) {
5
      if (m < 0 || m > n) return 0;
6
      if (n <= mod) return fac[n] * ifac[m] * ifac[n - m];</pre>
7
       // n >= mod 时需要这个
8
       return C(n % mod, m % mod) * C(n / mod, m / mod);
9
10
11 auto _ = []() {
12
       fac[0] = 1;
13
       for (int i = 1; i <= N; i++) fac[i] = fac[i - 1] * i;</pre>
      ifac[N] = fac[N].inv();
14
15
       for (int i = N - 1; i \ge 0; i--) ifac[i] = ifac[i + 1] * (i + 1);
16
       return true;
17 }();
```

#### 4.7 数论分块

求解形如  $\sum_{i=1}^{n} f(i)g(\lfloor \frac{n}{i} \rfloor)$  的合式  $s(n) = \sum_{i=1}^{n} f(i)$ 

```
modint sqrt_decomposition(int n) {
2
       auto s = [&](int x) { return x; };
3
       auto g = [&](int x) { return x; };
4
       modint res = 0;
5
       while (1 <= R) {
6
           int r = n / (n / 1);
7
           res = res + (s(r) - s(l - 1)) * g(n / 1);
8
           1 = r + 1;
9
10
       return res;
11 }
```

#### 4.8 积性函数

#### 4.8.1 定义

函数 f(n) 满足 f(1) = 1 且  $\forall x, y \in \mathbf{N}^*$ ,  $\gcd(x, y) = 1$  都有 f(xy) = f(x)f(y), 则 f(n) 为积性函数。 函数 f(n) 满足 f(1) = 1 且  $\forall x, y \in \mathbf{N}^*$  都有 f(xy) = f(x)f(y), 则 f(n) 为完全积性函数。

#### 4.8.2 例子

- 单位函数:  $\varepsilon(n) = [n = 1]$ 。(完全积性)
- 恒等函数:  $id_k(n) = n^k$ 。(完全积性)
- 常数函数: 1(n) = 1。(完全积性)
- 除数函数:  $\sigma_k(n) = \sum_{d|n} d^k$ 。  $\sigma_0(n)$  通常简记作 d(n) 或  $\tau(n)$ ,  $\sigma_1(n)$  通常简记作  $\sigma(n)$ 。
- 欧拉函数:  $\varphi(n) = \sum_{i=1}^{n} [\gcd(i, n) = 1]$ .

• 莫比乌斯函数:  $\mu(n) = \begin{cases} 1 & n=1 \\ 0 & \exists d>1, d^2 \mid n \text{, 其中 } \omega(n) \text{ 表示 } n \text{ 的本质不同质因子个数,它是} \\ (-1)^{\omega(n)} & \text{otherwise} \end{cases}$ 

#### 4.9 狄利克雷卷积

对于两个数论函数 f(x) 和 g(x),则它们的狄利克雷卷积得到的结果 h(x) 定义为:  $h(x)=\sum_{d|x}f(d)g\left(\frac{x}{d}\right)=\sum_{ab=x}f(a)g(b)$  可以简记为: h=f\*g。

#### 4.9.1 性质

**交換律:** f \* g = g \* f。

**结合律:** (f \* g) \* h = f \* (g \* h)。

**分配律:** (f+g)\*h = f\*h+g\*h。

等式的性质: f = g 的充要条件是 f \* h = g \* h,其中数论函数 h(x) 要满足  $h(1) \neq 0$ 。

#### 4.9.2 例子

- $\varepsilon = \mu * 1 \iff \varepsilon(n) = \sum_{d|n} \mu(d)$
- $id = \varphi * 1 \iff id(n) = \sum_{d|n} \varphi(d)$
- $d=1*1 \iff d(n)=\sum_{d|n}1$
- $\sigma = id *1 \iff \sigma(n) = \sum_{d|n} d$
- $\varphi = \mu * id \iff \varphi(n) = \sum_{d|n} d \cdot \mu(\frac{n}{d})$

#### 4.10 欧拉函数

#### 4.11 莫比乌斯反演

#### 4.11.1 莫比乌斯函数性质

$$\bullet \quad \textstyle \sum_{d|n} \mu(d) = \begin{cases} 1 & n=1 \\ 0 & n \neq 1 \end{cases}, \quad \text{RF} \, \textstyle \sum_{d|n} \mu(d) = \varepsilon(n) \,, \quad \mu*1 = \varepsilon$$

• 
$$[\gcd(i,j) = 1] = \sum_{d | \gcd(i,j)} \mu(d)$$

```
constexpr int N = 1e6;
  array<int, N + 1> miu;
3 array<bool, N + 1> ispr;
5
  auto _ = []() {
6
      miu.fill(1);
7
      ispr.fill(1);
       for (int i = 2; i <= N; i++) {
9
           if (!ispr[i]) continue;
10
           miu[i] = -1;
11
           for (int j = 2 * i; j <= N; j += i) {
12
               ispr[j] = 0;
13
               if ((j / i) % i == 0) miu[j] = 0;
               else miu[j] *= -1;
14
15
           }
16
17
       return true;
18 }();
```

#### 4.11.2 莫比乌斯变换/反演

```
f(n) = \sum_{d|n} g(d),那么有 g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) = \sum_{n|d} \mu(\frac{d}{n}) f(d)。
用狄利克雷卷积表示则为 f = g * 1,有 g = f * \mu。
f \to g 称为莫比乌斯反演,g \to f 称为莫比乌斯反演。
```

#### 4.12 杜教筛

杜教筛被用于处理一类数论函数的前缀和问题。对于数论函数 f,杜教筛可以在低于线性时间的复杂 度內计算  $S(n) = \sum_{i=1}^{n} f(i)$ 。

$$S(n) = \frac{\sum_{i=1}^{n} (f * g)(i) - \sum_{i=2}^{n} g(i)S\left(\left\lfloor \frac{n}{i} \right\rfloor\right)}{q(1)}$$

可以构造恰当的数论函数 g 使得:

- 可以快速计算  $\sum_{i=1}^{n} (f * g)(i)$ 。
- 可以快速计算 g 的单点值,用数论分块求解  $\sum_{i=2}^{n} g(i) S\left(\left|\frac{n}{i}\right|\right)$ 。

#### 4.12.1 示例

```
1 | 11 | sum_phi(11 n) {
       if (n <= N) return sp[n];</pre>
3
       if (sp2.count(n)) return sp2[n];
4
       11 \text{ res} = 0, 1 = 2;
5
       while (1 <= n) {
6
           11 r = n / (n / 1);
7
            res = res + (r - 1 + 1) * sum_phi(n / 1);
8
9
10
       return sp2[n] = (ll)n * (n + 1) / 2 - res;
11 }
12
13 ll sum_miu(ll n) {
    if (n <= N) return sm[n];</pre>
```

#### 4.13 多项式

```
#define countr_zero(n) __builtin_ctz(n)
  constexpr int N = 1e6;
   array<int, N + 1> inv;
 3
 4
 5
   int power(int a, int b) {
 6
       int res = 1;
 7
       while (b) {
 8
           if (b & 1) res = 1ll * res * a % mod;
9
           a = 111 * a * a % mod;
10
           b >>= 1;
11
       }
12
       return res;
13
14
15 namespace NFTS {
16 \mid int g = 3;
   vector<int> rev, roots{0, 1};
17
18
   void dft(vector<int> &a) {
19
       int n = a.size();
20
       if (rev.size() != n) {
21
           int k = countr_zero(n) - 1;
22
           rev.resize(n);
23
           for (int i = 0; i < n; ++i) rev[i] = rev[i >> 1] >> 1 | (i & 1) << k;
24
25
       if (roots.size() < n) {</pre>
26
           int k = countr_zero(roots.size());
27
           roots.resize(n);
28
           while ((1 << k) < n) {
29
                int e = power(g, (mod - 1) >> (k + 1));
                for (int i = 1 << (k - 1); i < (1 << k); ++i) {
30
31
                    roots[2 * i] = roots[i];
                    roots[2 * i + 1] = 1ll * roots[i] * e % mod;
32
33
                }
34
                ++k;
35
           }
36
37
       for (int i = 0; i < n; ++i)</pre>
38
           if (rev[i] < i) swap(a[i], a[rev[i]]);</pre>
39
       for (int k = 1; k < n; k *= 2) {
           for (int i = 0; i < n; i += 2 * k) {
40
41
                for (int j = 0; j < k; ++j) {</pre>
42
                    int u = a[i + j];
43
                    int v = 111 * a[i + j + k] * roots[k + j] % mod;
44
                    int x = u + v, y = u - v;
45
                    if (x >= mod) x -= mod;
```

```
46
                     if (y < 0) y += mod;
47
                     a[i + j] = x;
48
                     a[i + j + k] = y;
49
                }
50
            }
51
52
53
   void idft(vector<int> &a) {
54
        int n = a.size();
55
        reverse(a.begin() + 1, a.end());
56
        dft(a);
        int inv_n = power(n, mod - 2);
57
58
        for (int i = 0; i < n; ++i) a[i] = 1ll * a[i] * inv_n % mod;</pre>
59
60
      // namespace NFTS
61
62
    struct poly {
63
        poly &format() {
64
            while (!a.empty() && a.back() == 0) a.pop_back();
65
            return *this;
66
67
        poly &reverse() {
68
            ::reverse(a.begin(), a.end());
69
            return *this;
70
71
        vector<int> a;
72
        poly() {}
73
        poly(int x) {
74
            if (x) a = \{x\};
75
76
        poly(const vector<int> &_a) : a(_a) {}
77
        int size() const { return a.size(); }
78
        int &operator[](int id) { return a[id]; }
79
        int at(int id) const {
80
            if (id < 0 || id >= (int)a.size()) return 0;
81
            return a[id];
82
83
        poly operator-() const {
84
            auto A = *this;
            for (auto &x : A.a) x = (x == 0 ? 0 : mod - x);
85
86
            return A;
87
88
        poly mulXn(int n) const {
89
            auto b = a;
90
            b.insert(b.begin(), n, 0);
91
            return poly(b);
92
93
        poly modXn(int n) const {
            if (n > size()) return *this;
94
95
            return poly({a.begin(), a.begin() + n});
96
97
        poly divXn(int n) const {
            if (size() <= n) return poly();</pre>
98
99
            return poly({a.begin() + n, a.end()});
100
101
        poly &operator+=(const poly &rhs) {
102
            if (size() < rhs.size()) a.resize(rhs.size());</pre>
103
            for (int i = 0; i < rhs.size(); ++i)</pre>
```

```
104
                if ((a[i] += rhs.a[i]) >= mod) a[i] -= mod;
            return *this;
105
106
107
        poly &operator-=(const poly &rhs) {
108
            if (size() < rhs.size()) a.resize(rhs.size());</pre>
            for (int i = 0; i < rhs.size(); ++i)</pre>
109
                if ((a[i] -= rhs.a[i]) < 0) a[i] += mod;</pre>
110
111
            return *this;
112
        }
113
        poly &operator*=(poly rhs) {
114
            int n = size(), m = rhs.size(), tot = max(1, n + m - 1);
            int sz = 1 << __lg(tot * 2 - 1);</pre>
115
116
            a.resize(sz);
117
            rhs.a.resize(sz);
118
            NFTS::dft(a);
119
            NFTS::dft(rhs.a);
120
            for (int i = 0; i < sz; ++i) a[i] = 1ll * a[i] * rhs.a[i] % mod;</pre>
121
            NFTS::idft(a);
122
            return *this;
123
        }
        poly &operator/=(poly rhs) {
124
125
            int n = size(), m = rhs.size();
126
            if (n < m) return (*this) = poly();</pre>
127
            reverse();
128
            rhs.reverse();
129
            (*this) *= rhs.inv(n - m + 1);
130
            a.resize(n - m + 1);
131
            reverse();
132
            return *this;
133
134
        poly &operator%=(poly rhs) { return (*this) -= (*this) / rhs * rhs; }
135
        poly operator+(const poly &rhs) const { return poly(*this) += rhs; }
136
        poly operator-(const poly &rhs) const { return poly(*this) -= rhs; }
137
        poly operator*(poly rhs) const { return poly(*this) *= rhs; }
        poly operator/(poly rhs) const { return poly(*this) /= rhs; }
138
        poly operator%(poly rhs) const { return poly(*this) %= rhs; }
139
        poly powModPoly(int n, poly p) {
140
141
            poly r(1), x(*this);
142
            while (n) {
                if (n & 1) (r *= x) %= p;
143
144
                 (x *= x) %= p;
145
                n >>= 1;
146
            }
147
            return r;
148
149
        int inner(const poly &rhs) {
150
            int r = 0, n = min(size(), rhs.size());
151
            for (int i = 0; i < n; ++i) r = (r + 111 * a[i] * rhs.a[i]) % mod;
            return r;
152
153
        poly derivation() const {
154
155
            if (a.empty()) return poly();
156
            int n = size();
157
            vector<int> r(n - 1);
158
            for (int i = 1; i < n; ++i) r[i - 1] = 1ll * a[i] * i % mod;
159
            return poly(r);
160
161
        poly integral() const {
```

```
162
            if (a.empty()) return poly();
163
            int n = size();
164
            vector < int > r(n + 1);
            for (int i = 0; i < n; ++i) r[i + 1] = 1ll * a[i] * ::inv[i + 1] % mod;</pre>
165
166
            return poly(r);
167
        }
        poly inv(int n) const {
168
169
            assert(a[0] != 0);
170
            poly x(power(a[0], mod - 2));
171
            int k = 1;
            while (k < n) {
172
173
                k *= 2;
174
                x *= (poly(2) - modXn(k) * x).modXn(k);
175
            }
176
            return x.modXn(n);
177
        }
178
        // 需要保证首项为 1
179
        poly log(int n) const { return (derivation() * inv(n)).integral().modXn(n); }
180
        // 需要保证首项为 0
181
        poly exp(int n) const {
182
            poly x(1);
183
            int k = 1;
            while (k < n) {
184
185
                k *= 2;
186
                x = (x * (poly(1) - x.log(k) + modXn(k))).modXn(k);
187
            }
188
            return x.modXn(n);
189
190
        // 需要保证首项为 1, 开任意次方可以先 1n 再 exp 实现。
191
        poly sqrt(int n) const {
192
            poly x(1);
193
            int k = 1;
194
            while (k < n) {
195
                k *= 2;
                x += modXn(k) * x.inv(k);
196
                x = x.modXn(k) * inv2;
197
198
            }
199
            return x.modXn(n);
200
        }
        // 减法卷积, 也称转置卷积 {\rm MULT}(F(x),G(x))=\sum_{i\ge0}(\sum_{j\ge
201
202
        // 0}f_{i+j}g_j)x^i
        poly mulT(poly rhs) const {
203
            if (rhs.size() == 0) return poly();
204
205
            int n = rhs.size();
206
            ::reverse(rhs.a.begin(), rhs.a.end());
207
            return ((*this) * rhs).divXn(n - 1);
208
209
        int eval(int x) {
            int r = 0, t = 1;
210
211
            for (int i = 0, n = size(); i < n; ++i) {</pre>
212
                r = (r + 111 * a[i] * t) % mod;
213
                t = 111 * t * x % mod;
214
            }
215
            return r;
216
        // 多点求值新科技: https://jkloverdcoi.github.io/2020/08/04/转置原理及其应用/
217
218
        // 模板例题: https://www.luogu.com.cn/problem/P5050
219
        auto evals(vector<int> &x) const {
```

```
220
            if (size() == 0) return vector(x.size(), 0);
221
            int n = x.size();
            vector ans(n, 0);
222
223
            vector<poly> g(4 * n);
            auto build = [&](auto self, int 1, int r, int p) -> void {
224
225
                if (r - 1 == 1) {
226
                    g[p] = poly(\{1, x[1] ? mod - x[1] : 0\});
227
                } else {
228
                    int m = (1 + r) / 2;
                    self(self, 1, m, 2 * p);
229
                    self(self, m, r, 2 * p + 1);
230
231
                    g[p] = g[2 * p] * g[2 * p + 1];
232
                }
233
            };
234
            build(build, 0, n, 1);
235
            auto solve = [&](auto self, int l, int r, int p, poly f) -> void {
                if (r - 1 == 1) {
236
237
                    ans[1] = f[0];
238
                } else {
239
                    int m = (1 + r) / 2;
                    self(self, 1, m, 2 * p, f.mulT(g[2 * p + 1]).modXn(m - 1));
240
241
                    self(self, m, r, 2 * p + 1, f.mulT(g[2 * p]).modXn(r - m));
                }
242
243
            };
244
            solve(solve, 0, n, 1, mulT(g[1].inv(size())).modXn(n));
245
            return ans;
246
   }; // 全家桶测试: https://www.luogu.com.cn/training/3015#information
247
248
249
   auto _ = []() {
250
        inv[0] = inv[1] = 1;
251
        for (int i = 2; i < inv.size(); i++) inv[i] = 111 * (mod - mod / i) * inv[mod % i] % mod;</pre>
252
        return true;
253 }();
```

#### 4.14 盒子与球

n 个球, m 个盒

球同	盒同	可空	公式
<b>√</b>	<b>√</b>	<b>√</b>	$f_{n,m}=f_{n,m-1}+f_{n-m,m}$ 或 $[x^n]e^{\sum\limits_{i=1}^{m}\sum\limits_{j=1}^{\infty}rac{x^{ij}}{j}}$
<b>√</b>	<b>√</b>	×	$f_{n-m,m}$
×	<b>√</b>	✓	$f_{n-m,m}$ $\sum_{i=1}^m g_{n,i}$ 或 $\sum_{i=1}^m \sum_{j=0}^i rac{j^n}{j!} rac{(-1)^{i-j}}{(i-j)!}$
×	<b>√</b>	×	$g_{n,m} = g_{n-1,m-1} + m * g_{n-1,m}$ 或 $\frac{1}{m!} \sum_{i=0}^m (-1)^i {m \choose i} (m-i)^n$
<b>√</b>	×	<b>√</b>	$C_{n+m-1}^{m-1}$
<b>√</b>	×		$C_{n-1}^{m-1}$
×	×	✓	$m^n$
×	×	×	$m!*g_{n,m}$ 或 $\sum\limits_{i=0}^{m}(-1)^iinom{m}{i}(m-i)^n$

#### 4.14.1 球同, 盒同, 可空

```
int solve(int n, int m) {
    vector a(n + 1, 0);
    for (int i = 1; i <= m; i++)
        for (int j = i, k = 1; j <= n; j += i, k++) a[j] = (a[j] + inv[k]) % mod;
    auto p = poly(a).exp(n + 1);
    return (p.a[n] + mod) % mod;
}</pre>
```

若要求不超过 k 个,答案为  $[x^ny^m]\prod\limits_{i=0}^k\left(\sum\limits_{j=0}^mx^{ij}y^j\right)$ 。

#### 4.14.2 球不同, 盒同, 可空

```
int solve(int n, int m) {
2
       vector a(n + 1, 0);
3
       vector b(n + 1, 0);
       for (int i = 0; i <= n; i++) {</pre>
4
5
           a[i] = ifac[i];
6
           if (i & 1) a[i] = -a[i];
7
           b[i] = 1ll * power(i, n) * ifac[i] % mod;
8
9
       auto p = poly(a) * poly(b);
10
       int ans = 0;
11
       for (int i = 1; i \leftarrow min(n, m); i++) ans = (ans + p.a[i]) % mod;
12
       return (ans + mod) % mod;
13 }
```

若要求不超过 k 个,答案为  $m! \cdot [x^n y^m] \prod_{i=0}^k \left( \sum_{j=0}^n \frac{1}{i!^j} x^{ij} y^j \right)$ 。

#### 4.14.3 球同, 盒不同, 可空

若要求不超过 k 个,答案为  $\left[x^n\right] \left(\sum_{i=0}^k x^i\right)^m = \left[x^n\right] \frac{(x^{k+1}-1)^m}{(x-1)^m}$ 。 也可以考虑容斥,令 f(i) 表示至少有 i 个盒子装了 > k 个球方案数,  $f(i) = {m \choose i} {n-(k+1)i+m-1 \choose m-1}$ 。 总方案数则为  $\sum_{i=0}^m (-1)^i f(i)$ 。

#### 4.14.4 球同,盒不同,不可空

若要求不超过 k 个,答案为  $\left[x^n\right] \left(\sum_{i=1}^k x^i\right)^m = \left[x^n\right] \frac{(x^k-1)^m x^m}{(x-1)^m}$ 。 也可以考虑容斥,令 f(i) 表示至少有 i 个盒子装了 > k 个球方案数,  $f(i) = {m \choose i} {n-ki-1 \choose m-1}$ 。 总方案数则为  $\sum_{i=0}^m (-1)^i f(i)$ 。

#### 4.14.5 球不同,盒不同,可空

若要求不超过 k 个,答案为  $m! \cdot [x^n] \left(\sum_{i=0}^k \frac{1}{i!} x^i\right)^m$ 。

#### 4.14.6 球不同, 盒不同, 不可空

若要求不超过 k 个,答案为  $m! \cdot [x^n] \left(\sum_{i=1}^k \frac{1}{i!} x^i\right)^m$ 。

#### 4.15 线性基

```
// 线性基
  struct basis {
      int rnk = 0;
      array<ull, 64> p{};
5
      // 将x插入此线性基中
7
      void insert(ull x) {
          for (int i = 63; i >= 0; i--) {
9
              if (!(x >> i & 1)) continue;
10
              if (p[i]) x ^= p[i];
              else {
12
                  p[i] = x;
13
                  rnk++;
14
                  break;
15
              }
16
          }
17
18
19
      // 将另一个线性基插入此线性基中
20
      void insert(basis other) {
21
          for (int i = 0; i <= 63; i++) {
22
              if (!other.p[i]) continue;
23
              insert(other.p[i]);
24
          }
25
26
27
      // 最大异或值
```

```
28     ull max_basis() {
29         ull res = 0;
30         for (int i = 63; i >= 0; i--)
31             if ((res ^ p[i]) > res) res ^= p[i];
32             return res;
33         }
34     };
```

#### 4.16 矩阵快速幂

```
constexpr 11 mod = 2147493647;
 2
   struct Mat {
3
       int n, m;
 4
       vector<vector<ll>>> mat;
5
       Mat(int n, int m) : n(n), m(n), mat(n, vector<ll>(m, 0)) {}
6
       Mat(vector<vector<ll>>> mat) : n(mat.size()), m(mat[0].size()), mat(mat) {}
 7
       Mat operator*(const Mat& other) {
8
           assert(m == other.n);
9
           Mat res(n, other.m);
10
           for (int i = 0; i < res.n; i++)</pre>
                for (int j = 0; j < res.m; j++)</pre>
11
12
                    for (int k = 0; k < m; k++)
13
                        res.mat[i][j] = (res.mat[i][j] + mat[i][k] * other.mat[k][j] \% mod) \% mod;
14
           return res;
15
       }
16 };
17
  Mat ksm(Mat a, ll b) {
18
       assert(a.n == a.m);
19
       Mat res(a.n, a.m);
20
       for (int i = 0; i < res.n; i++) res.mat[i][i] = 1;</pre>
21
       while (b) {
22
           if (b & 1) res = res * a;
23
           b >>= 1;
24
           a = a * a;
25
26
       return res;
27 }
```

## 5 计算几何

#### 5.1 整数

```
constexpr double inf = 1e100;
2
3
  // 向量
4
  struct vec {
5
       static bool cmp(const vec &a, const vec &b) { return tie(a.x, a.y) < tie(b.x, b.y); }</pre>
6
7
       11 x, y;
8
       vec() : x(0), y(0) \{ \}
9
       vec(11 _x, 11 _y) : x(_x), y(_y) {}
10
11
       // 模
12
       11 len2() const { return x * x + y * y; }
13
       double len() const { return sqrt(x * x + y * y); }
14
15
       // 是否在上半轴
16
       bool up() const { return y > 0 \mid | y == 0 && x >= 0; }
17
18
       bool operator==(const vec &b) const { return tie(x, y) == tie(b.x, b.y); }
19
       // 极角排序
20
       bool operator<(const vec &b) const {</pre>
21
           if (up() != b.up()) return up() > b.up();
22
           11 tmp = (*this) ^ b;
23
           return tmp ? tmp > 0 : cmp(*this, b);
24
       }
25
26
       vec operator+(const vec &b) const { return {x + b.x, y + b.y}; }
27
       vec operator-() const { return {-x, -y}; }
28
       vec operator-(const vec &b) const { return -b + (*this); }
29
       vec operator*(11 b) const { return {x * b, y * b}; }
30
       11 operator*(const vec &b) const { return x * b.x + y * b.y; }
31
       // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
32
33
       // 等于0共线,可能同向或反向,结果绝对值表示 a b 形成的平行四边行的面积
34
       11 operator^(const vec &b) const { return x * b.y - y * b.x; }
35
36
       friend istream &operator>>(istream &in, vec &data) {
37
           in >> data.x >> data.y;
38
           return in;
39
       }
40
       friend ostream &operator<<(ostream &out, const vec &data) {</pre>
41
           out << fixed << setprecision(6);</pre>
           out << data.x << " " << data.y;
42
43
           return out;
44
45
  };
46
47
  11 cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
48
49
  // 多边形的面积a
50
  double polygon_area(vector<vec> &p) {
51
       11 area = 0;
52
       for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];</pre>
       area += p.back() ^ p[0];
53
54
       return abs(area / 2.0);
```

```
55 }
56
57
   // 多边形的周长
58
   double polygon_len(vector<vec> &p) {
59
       double len = 0;
       for (int i = 1; i < p.size(); i++) len += (p[i - 1] - p[i]).len();</pre>
60
       len += (p.back() - p[0]).len();
61
62
       return len;
63
64
65
   // 以整点为顶点的线段上的整点个数
66
   11 count(const vec &a, const vec &b) {
67
       vec c = a - b;
68
       return gcd(abs(c.x), abs(c.y)) + 1;
69
70
71
   // 以整点为顶点的多边形边上整点个数
72
   11 count(vector<vec> &p) {
73
       11 cnt = 0;
74
       for (int i = 1; i < p.size(); i++) cnt += count(p[i - 1], p[i]);</pre>
75
       cnt += count(p.back(), p[0]);
76
       return cnt - p.size();
77
   }
78
79
   // 判断点是否在凸包内, 凸包必须为逆时针顺序
80
   bool in_polygon(const vec &a, vector<vec> &p) {
81
       int n = p.size();
82
       if (n == 0) return 0;
83
       if (n == 1) return a == p[0];
84
       if (n == 2) return cross(a, p[1], p[0]) == 0 && <math>(p[0] - a) * (p[1] - a) <= 0;
85
       if (cross(a, p[1], p[0]) > 0 || cross(p.back(), a, p[0]) > 0) return 0;
86
       auto cmp = [\&](vec \&x, const vec \&y) { return ((x - p[0]) ^ y) >= 0; };
87
       int i = lower_bound(p.begin() + 2, p.end() - 1, a - p[0], cmp) - p.begin() - 1;
88
       return cross(p[(i + 1) % n], a, p[i]) >= 0;
89
90
91
   // 凸包直径的两个端点
92
   auto polygon_dia(vector<vec> &p) {
93
       int n = p.size();
94
       array<vec, 2> res{};
95
       if (n == 1) return res;
       if (n == 2) return res = {p[0], p[1]};
96
97
       11 mx = 0;
98
       for (int i = 0, j = 2; i < n; i++) {
99
            while (abs(cross(p[i], p[(i + 1) % n], p[j])) <=</pre>
                   abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n])))
100
101
                j = (j + 1) \% n;
102
            11 tmp = (p[i] - p[j]).len2();
            if (tmp > mx) {
103
104
                mx = tmp;
105
                res = {p[i], p[j]};
106
107
           tmp = (p[(i + 1) % n] - p[j]).len2();
108
            if (tmp > mx) {
109
                mx = tmp;
110
                res = {p[(i + 1) % n], p[j]};
111
           }
112
```

```
113
        return res;
114
115
116
    // 凸包
117
   auto convex_hull(vector<vec> &p) {
118
        sort(p.begin(), p.end(), vec::cmp);
        int n = p.size();
119
120
        vector sta(n + 1, 0);
121
        vector v(n, false);
122
        int tp = -1;
123
        sta[++tp] = 0;
124
        auto update = [&](int lim, int i) {
125
            while (tp > lim && cross(p[i], p[sta[tp]], p[sta[tp - 1]]) >= 0) v[sta[tp--]] = 0;
126
            sta[++tp] = i;
127
            v[i] = 1;
128
        };
129
        for (int i = 1; i < n; i++) update(0, i);</pre>
130
        int cnt = tp;
131
        for (int i = n - 1; i >= 0; i--) {
            if (v[i]) continue;
132
133
            update(cnt, i);
134
        vector<vec> res(tp);
135
136
        for (int i = 0; i < tp; i++) res[i] = p[sta[i]];</pre>
137
        return res;
138
139
    // 闵可夫斯基和,两个点集的和构成一个凸包
140
    auto minkowski(vector<vec> &a, vector<vec> &b) {
141
142
        rotate(a.begin(), min_element(a.begin(), a.end(), vec::cmp), a.end());
143
        rotate(b.begin(), min_element(b.begin(), b.end(), vec::cmp), b.end());
        int n = a.size(), m = b.size();
144
145
        vector<vec> c{a[0] + b[0]};
146
        c.reserve(n + m);
        int i = 0, j = 0;
147
148
        while (i < n && j < m) {
            vec x = a[(i + 1) \% n] - a[i];
149
150
            vec y = b[(j + 1) \% m] - b[j];
151
            c.push_back(c.back() + ((x ^ y) >= 0 ? (i++, x) : (j++, y)));
152
153
        while (i + 1 < n) {
            c.push_back(c.back() + a[(i + 1) % n] - a[i]);
154
155
            i++;
156
157
        while (j + 1 < m) {
            c.push_back(c.back() + b[(j + 1) \% m] - b[j]);
158
159
            j++;
160
        }
161
        return c;
162
163
164
    // 过凸多边形外一点求凸多边形的切线,返回切点下标
   auto tangent(const vec &a, vector<vec> &p) {
165
166
        int n = p.size();
167
        int l = -1, r = -1;
        for (int i = 0; i < n; i++) {</pre>
168
            ll tmp1 = cross(p[i], p[(i - 1 + n) % n], a);
169
            ll tmp2 = cross(p[i], p[(i + 1) % n], a);
170
```

```
171
           if (1 == -1 \&\& tmp1 <= 0 \&\& tmp2 <= 0) 1 = i;
172
           else if (r == -1 \&\& tmp1 >= 0 \&\& tmp2 >= 0) r = i;
173
       return array{1, r};
174
175
176
   // 直线
177
178
   struct line {
179
       vec p, d;
180
       line() : p(vec()), d(vec()) {}
       line(const vec \&_p, const vec \&_d) : p(_p), d(_d) {}
181
182
183
   // 点到直线距离
184
185
   double dis(const vec &a, const line &b) { return abs((b.p - a) ^ (b.p + b.d - a)) / b.d.len(); }
186
187
    // 点在直线哪边,大于0在左边,等于0在线上,小于0在右边
   11 side_line(const vec &a, const line &b) { return b.d ^ (a - b.p); }
188
189
   // 两直线是否垂直
190
   bool perpen(const line &a, const line &b) { return a.d * b.d == 0; }
191
192
193
   // 两直线是否平行
194
   bool parallel(const line &a, const line &b) { return (a.d ^ b.d) == 0; }
195
196
   // 点的垂线是否与线段有交点
197
   bool perpen(const vec &a, const line &b) {
198
       vec p(-b.d.y, b.d.x);
       bool cross1 = (p ^ (b.p - a)) > 0;
199
       bool cross2 = (p ^ (b.p + b.d - a)) > 0;
200
201
       return cross1 != cross2;
202
   }
203
204
   // 点到线段距离
205
   double dis_seg(const vec &a, const line &b) {
206
       if (perpen(a, b)) return dis(a, b);
207
       return min((b.p - a).len(), (b.p + b.d - a).len());
208
209
   // 点到凸包距离
210
211
   double dis(const vec &a, vector<vec> &p) {
       double res = inf;
212
213
       for (int i = 1; i < p.size(); i++) res = min(dis_seg(a, line(p[i - 1], p[i] - p[i - 1])), res);</pre>
       res = min(dis_seg(a, line(p.back(), p[0] - p.back())), res);
214
215
       return res;
216
217
218 // 两直线交点
219 vec intersection(11 A, 11 B, 11 C, 11 D, 11 E, 11 F) {
220
       return {(B * F - C * E) / (A * E - B * D), (C * D - A * F) / (A * E - B * D)};
221
   }
222
223 // 两直线交点
224
   vec intersection(const line &a, const line &b) {
225
       return intersection(a.d.y, -a.d.x, a.d.x * a.p.y - a.d.y * a.p.x, b.d.y, -b.d.x,
226
                           b.d.x * b.p.y - b.d.y * b.p.x);
227
   }
```

三维

```
1
  // 向量
2
  struct vec3 {
3
       static bool cmp(const vec3 &a, const vec3 &b) {
 4
           return tie(a.x, a.y, a.z) < tie(b.x, b.y, b.z);</pre>
5
6
7
       11 x, y, z;
 8
       vec3() : x(0), y(0), z(0) {}
9
       vec3(11 _x, 11 _y, 11 _z) : x(_x), y(_y), z(_z) {}
10
11
       // 模
12
       11 len2() const { return x * x + y * y + z * z; }
13
       double len() const { return hypot(x, y, z); }
14
       bool operator==(const vec3 &b) const { return tie(x, y, z) == tie(b.x, b.y, b.z); }
15
16
       bool operator!=(const vec3 &b) const { return !(*this == b); }
17
18
       vec3 operator+(const vec3 &b) const { return \{x + b.x, y + b.y, z + b.z\}; }
19
       vec3 operator-() const { return {-x, -y, -z}; }
20
       vec3 operator-(const vec3 &b) const { return -b + (*this); }
21
       vec3 operator*(ll b) const { return \{b * x, b * y, b * z\}; \}
22
       11 operator*(const vec3 &b) const { return x * b.x + y * b.y + z * b.z; }
23
24
       // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
25
       // 等于O共线,可能同向或反向,结果绝对值表示 a b 形成的平行四边行的面积
26
       vec3 operator^(const vec3 &b) const {
27
           return {y * b.z - z * b.y, z * b.x - x * b.z, x * b.y - y * b.x};
28
       }
29
30
       friend istream &operator>>(istream &in, vec3 &data) {
31
           in >> data.x >> data.y >> data.z;
32
           return in;
33
       }
34
       friend ostream &operator<<(ostream &out, const vec3 &data) {</pre>
35
           out << fixed << setprecision(6);</pre>
           out << data.x << " " << data.y << " " << data.z;
36
37
           return out;
38
39
  };
40
41
   struct line3 {
42
       vec3 p, d;
43
44
       line3(const vec3 &a, const vec3 &b) : p(a), d(b - a) {}
45
  };
46
47
   struct plane {
48
       vec3 p, d;
49
       plane() {}
50
       plane(const vec3 &a, const vec3 &b, const vec3 &c) : p(a) {
51
           d = (b - a) ^ (c - a);
           assert(d != vec3());
52
53
       }
54
  };
55
56 // 线面是否垂直
57 bool perpen(const line3 &a, const plane &b) { return (a.d ^ b.d) == vec3(); }
```

#### 5.2 浮点数

```
1
  using lf = double;
3
  constexpr lf eps = 1e-8;
  constexpr lf inf = 1e100;
4
  const lf PI = acos(-1);
5
6
 7
   int sgn(lf a, lf b) {
8
      lf c = a - b;
9
       return c < -eps ? -1 : c < eps ? 0 : 1;
10
11
12
  // 向量
  struct vec {
13
14
       static bool cmp(const vec &a, const vec &b) {
15
           return sgn(a.x, b.x) ? a.x < b.x : sgn(a.y, b.y) < 0;
16
17
       If x, y;
18
19
       vec() : x(0), y(0) \{ \}
       vec(lf _x, lf _y) : x(_x), y(_y) {}
20
21
22
       // 模
       1f len2() const { return x * x + y * y; }
23
24
       lf len() const { return sqrt(x * x + y * y); }
25
26
       // 与×轴正方向的夹角
27
       lf angle() const {
           If angle = atan2(y, x);
28
29
           if (angle < 0) angle += 2 * PI;</pre>
30
           return angle;
31
       }
32
       // 逆时针旋转
33
34
       vec rotate(const 1f &theta) const {
35
           return \{x * cos(theta) - y * sin(theta), y * cos(theta) + x * sin(theta)\};
36
37
38
       vec e() const {
39
           lf tmp = len();
40
           return {x / tmp, y / tmp};
41
42
       // 是否在上半轴
43
44
       bool up() const { return sgn(y, 0) > 0 \mid | sgn(y, 0) == 0 && sgn(x, 0) >= 0; }
45
       bool operator==(const vec &other) const { return sgn(x, other.x) == 0 && sgn(y, other.y) == 0; }
46
47
       // 极角排序
       bool operator<(const vec &b) const {</pre>
48
49
           if (up() != b.up()) return up() > b.up();
           If tmp = (*this) ^ b;
50
           return sgn(tmp, 0) ? tmp > 0 : cmp(*this, b);
```

```
52
53
54
       vec operator+(const vec &b) const { return {x + b.x, y + b.y}; }
55
       vec operator-() const { return {-x, -y}; }
56
       vec operator-(const vec &b) const { return -b + (*this); }
57
       vec operator*(lf b) const { return {x * b, y * b}; }
       vec operator/(lf b) const { return {x / b, y / b}; }
58
59
       lf operator*(const vec &b) const { return x * b.x + y * b.y; }
60
61
       // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
62
       // 等于0共线,可能同向或反向,结果绝对值表示 a b 形成的平行四边行的面积
       lf operator^(const vec &b) const { return x * b.y - y * b.x; }
63
64
65
       friend istream &operator>>(istream &in, vec &data) {
66
            in >> data.x >> data.y;
67
            return in;
68
69
       friend ostream &operator<<(ostream &out, const vec &data) {</pre>
70
            out << fixed << setprecision(6);</pre>
            out << data.x << " " << data.y;
71
72
            return out;
73
74
   };
75
76 | 1f cross(const vec &a, const vec &b, const vec &c) { return (a - c) ^ (b - c); }
77
   lf angle(const vec &a, const vec &b) { return atan2(abs(a ^ b), a * b); }
78
79
   // 多边形的面积
80
81
   lf polygon_area(vector<vec> &p) {
82
       If area = 0;
       for (int i = 1; i < p.size(); i++) area += p[i - 1] ^ p[i];</pre>
83
84
       area += p.back() ^ p[0];
85
       return abs(area / 2.0);
86
87
   // 多边形的周长
88
89
   lf polygon_len(vector<vec> &p) {
90
       lf len = 0;
       for (int i = 1; i < p.size(); i++) len += (p[i - 1] - p[i]).len();</pre>
91
92
       len += (p.back() - p[0]).len();
93
       return len;
94
95
96
   // 判断点是否在凸包内, 凸包必须为逆时针顺序
   bool in_polygon(const vec &a, vector<vec> &p) {
97
98
       int n = p.size();
99
       if (n == 0) return 0;
       if (n == 1) return a == p[0];
100
       if (n == 2) return sgn(cross(a, p[1], p[0]), 0) == 0 && <math>sgn((p[0] - a) * (p[1] - a), 0) <= 0;
101
       if (sgn(cross(a, p[1], p[0]), 0) > 0 \mid | sgn(cross(p.back(), a, p[0]), 0) > 0) return 0;
102
103
       auto cmp = [\&](\text{vec }\&x, \text{const vec }\&y) \{ \text{return sgn}((x - p[0]) ^ y, 0) >= 0; \};
       int i = lower_bound(p.begin() + 2, p.end() - 1, a - p[0], cmp) - p.begin() - 1;
104
105
       return sgn(cross(p[(i + 1) % n], a, p[i]), 0) >= 0;
106 }
107
   // 凸包直径的两个端点
109 auto polygon_dia(vector<vec> &p) {
```

```
110
        int n = p.size();
        array<vec, 2> res{};
111
112
        if (n == 1) return res;
113
        if (n == 2) return res = {p[0], p[1]};
114
        If mx = 0;
        for (int i = 0, j = 2; i < n; i++) {
115
            while (sgn(abs(cross(p[i], p[(i + 1) % n], p[j])),
116
117
                        abs(cross(p[i], p[(i + 1) % n], p[(j + 1) % n]))) <= 0)
                j = (j + 1) \% n;
118
119
            lf tmp = (p[i] - p[j]).len();
            if (tmp > mx) {
120
121
                mx = tmp;
122
                res = {p[i], p[j]};
123
            }
124
            tmp = (p[(i + 1) % n] - p[j]).len();
125
            if (tmp > mx) {
126
                res = {p[(i + 1) % n], p[j]};
127
128
            }
129
        }
130
        return res;
131
132
133
   // 凸包
   auto convex_hull(vector<vec> &p) {
134
135
        sort(p.begin(), p.end(), vec::cmp);
136
        int n = p.size();
137
        vector sta(n + 1, 0);
138
        vector v(n, false);
        int tp = -1;
139
140
        sta[++tp] = 0;
        auto update = [&](int lim, int i) {
141
            while (tp > \lim \&  sgn(cross(p[i], p[sta[tp]], p[sta[tp - 1]]), 0) >= 0) v[sta[tp--]] = 0;
142
143
            sta[++tp] = i;
            v[i] = 1;
144
145
        };
        for (int i = 1; i < n; i++) update(0, i);</pre>
146
147
        int cnt = tp;
        for (int i = n - 1; i >= 0; i--) {
148
            if (v[i]) continue;
149
150
            update(cnt, i);
151
        vector<vec> res(tp);
152
        for (int i = 0; i < tp; i++) res[i] = p[sta[i]];</pre>
153
154
        return res;
155
156
    // 闵可夫斯基和,两个点集的和构成一个凸包
157
    auto minkowski(vector<vec> &a, vector<vec> &b) {
158
        rotate(a.begin(), min_element(a.begin(), a.end(), vec::cmp), a.end());
159
        rotate(b.begin(), min_element(b.begin(), b.end(), vec::cmp), b.end());
160
161
        int n = a.size(), m = b.size();
        vector<vec> c{a[0] + b[0]};
162
163
        c.reserve(n + m);
164
        int i = 0, j = 0;
        while (i < n \&\& j < m) {
165
166
            vec x = a[(i + 1) \% n] - a[i];
            vec y = b[(j + 1) \% m] - b[j];
167
```

```
168
           c.push_back(c.back() + (sgn(x ^ y, 0) >= 0 ? (i++, x) : (j++, y)));
169
       }
170
       while (i + 1 < n) {
           c.push_back(c.back() + a[(i + 1) % n] - a[i]);
171
172
           i++;
173
       }
       while (j + 1 < m) {
174
175
           c.push_back(c.back() + b[(j + 1) \% m] - b[j]);
176
           j++;
177
178
       return c;
179
180
    // 过凸多边形外一点求凸多边形的切线, 返回切点下标
181
182
   auto tangent(const vec &a, vector<vec> &p) {
183
       int n = p.size();
184
       int l = -1, r = -1;
       for (int i = 0; i < n; i++) {</pre>
185
186
           lf tmp1 = cross(p[i], p[(i - 1 + n) % n], a);
187
           lf tmp2 = cross(p[i], p[(i + 1) % n], a);
           if (1 == -1 \&\& sgn(tmp1, 0) <= 0 \&\& sgn(tmp2, 0) <= 0) 1 = i;
188
189
           else if (r == -1 \&\& sgn(tmp1, 0) >= 0 \&\& sgn(tmp2, 0) >= 0) r = i;
190
191
       return array{1, r};
192
193
   // 直线
194
   struct line {
195
196
       vec p, d;
197
       line() : p(vec()), d(vec()) {}
198
       line(const vec \&_p, const vec \&_d) : p(_p), d(_d) {}
199
   };
200
201
   // 点到直线距离
202 If dis(const vec &a, const line &b) { return abs((b.p - a) ^ (b.p + b.d - a)) / b.d.len(); }
203
204
   // 点在直线哪边,大于0在左边,等于0在线上,小于0在右边
205 int side_line(const vec &a, const line &b) { return sgn(b.d ^ (a - b.p), 0); }
206
   // 两直线是否垂直
207
208
   bool perpen(const line &a, const line &b) { return sgn(a.d * b.d, 0) == 0; }
209
210
   // 两直线是否平行
211 bool parallel(const line &a, const line &b) { return sgn(a.d ^ b.d, 0) == 0; }
212
213
   // 点的垂线是否与线段有交点
214
   bool perpen(const vec &a, const line &b) {
215
       vec p(-b.d.y, b.d.x);
216
       bool cross1 = sgn(p \land (b.p - a), 0) > 0;
217
       bool cross2 = sgn(p ^(b.p + b.d - a), 0) > 0;
218
       return cross1 != cross2;
219
220
   // 点到线段距离
221
222
   lf dis_seg(const vec &a, const line &b) {
223
       if (perpen(a, b)) return dis(a, b);
224
       return min((b.p - a).len(), (b.p + b.d - a).len());
225 }
```

```
226
227
    // 点到凸包距离
228
   lf dis(const vec &a, vector<vec> &p) {
229
       lf res = inf;
       for (int i = 1; i < p.size(); i++) res = min(dis_seg(a, line(p[i - 1], p[i] - p[i - 1])), res);</pre>
230
231
       res = min(dis_seg(a, line(p.back(), p[0] - p.back())), res);
232
       return res;
233
234
235
   // 两直线交点
236
   vec intersection(lf A, lf B, lf C, lf D, lf E, lf F) {
       return {(B * F - C * E) / (A * E - B * D), (C * D - A * F) / (A * E - B * D)};
237
238
239
240
   // 两直线交点
241
   vec intersection(const line &a, const line &b) {
242
        return intersection(a.d.y, -a.d.x, a.d.x * a.p.y - a.d.y * a.p.x, b.d.y, -b.d.x,
243
                            b.d.x * b.p.y - b.d.y * b.p.x);
244
245
246
   struct circle {
247
       vec o;
       lf r;
248
249
       circle(const vec &_o, lf _r) : o(_o), r(_r){};
250
251
       // 点与圆的关系 -1在圆内, 0在圆上, 1在圆外
252
       int relation(const vec &a) const { return sgn((a - o).len(), r); }
253
254
       // 圆与圆的关系 -3包含, -2内切, -1相交, 0外切, 1相离
255
       int relation(const circle &a) const {
256
            lf 1 = (a.o - o).len();
257
            if (sgn(1, abs(r - a.r)) < 0) return -3;</pre>
258
            if (sgn(1, abs(r - a.r)) == 0) return -2;
259
            if (sgn(1, abs(r + a.r)) < 0) return -1;
            if (sgn(1, abs(r + a.r)) == 0) return 0;
260
261
            return 1;
262
263
264
       lf area() { return PI * r * r; }
265
266
   // 圆与直线交点
267
   auto intersection(const circle &c, const line &l) {
268
269
       lf d = dis(c.o, 1);
270
       vector<vec> res;
271
       vec mid = 1.p + 1.d.e() * ((c.o - 1.p) * 1.d / 1.d.len());
272
       if (sgn(d, c.r) == 0) res.push_back(mid);
273
       else if (sgn(d, c.r) < 0) {
            d = sqrt(c.r * c.r - d * d);
274
275
            res.push_back(mid + 1.d.e() * d);
276
            res.push_back(mid - 1.d.e() * d);
277
278
       return res;
279
280
   // oab三角形与圆相交的面积
281
282 If area(const circle &c, const vec &a, const vec &b) {
       if (sgn(cross(a, b, c.o), 0) == 0) return 0;
283
```

```
284
        vector<vec> p;
285
        p.push_back(a);
286
        line l(a, b - a);
287
        auto tmp = intersection(c, 1);
288
        if (tmp.size() == 2) {
289
            for (auto &i : tmp)
                if (sgn((a - i) * (b - i), 0) < 0) p.push_back(i);</pre>
290
291
292
        p.push_back(b);
        if (p.size() == 4 \& sgn((p[0] - p[1]) * (p[2] - p[1]), 0) > 0) swap(p[1], p[2]);
293
294
        If res = 0;
295
        for (int i = 1; i < p.size(); i++)</pre>
296
            if (c.relation(p[i - 1]) == 1 || c.relation(p[i]) == 1) {
297
                If ang = angle(p[i - 1] - c.o, p[i] - c.o);
298
                res += c.r * c.r * ang / 2;
299
            } else res += abs(cross(p[i - 1], p[i], c.o)) / 2.0;
300
        return res;
301
302
    // 多边形与圆相交的面积
303
304 | 1f area(vector<vec> &p, circle c) {
305
        If res = 0;
        for (int i = 0; i < p.size(); i++) {</pre>
306
307
            int j = i + 1 == p.size() ? 0 : i + 1;
308
            if (sgn(cross(p[i], p[j], c.o), 0) <= 0) res += area(c, p[i], p[j]);</pre>
309
            else res -= area(c, p[i], p[j]);
310
311
        return abs(res);
312 }
```

#### 三维

```
constexpr lf eps = 1e-8;
2
3
  int sgn(lf a, lf b) {
4
        1f c = a - b; 
5
       return c < -eps ? -1 : c < eps ? 0 : 1;
6
  }
7
8
   // 向量
   struct vec3 {
9
10
       If x, y, z;
11
       vec3() : x(0), y(0), z(0) {}
       vec3(1f _x, 1f _y, 1f _z) : x(_x), y(_y), z(_z) {}
12
13
       // 模
14
       lf len2() const { return x * x + y * y + z * z; }
15
16
       lf len() const { return hypot(x, y, z); }
17
18
       bool operator==(const vec3 &b) const {
19
           return sgn(x, b.x) == 0 && sgn(y, b.y) == 0 && sgn(z, b.z) == 0;
20
21
       bool operator!=(const vec3 &b) const { return !(*this == b); }
22
23
       vec3 operator+(const vec3 &b) const { return \{x + b.x, y + b.y, z + b.z\}; }
24
       vec3 operator-() const { return {-x, -y, -z}; }
25
       vec3 operator-(const vec3 &b) const { return -b + (*this); }
26
       vec3 operator*(lf b) const { return {b * x, b * y, b * z}; }
27
       lf operator*(const vec3 &b) const { return x * b.x + y * b.y + z * b.z; }
```

```
28
29
       // 叉积 结果大于0, a到b为逆时针, 小于0, a到b顺时针,
30
       // 等于0共线,可能同向或反向,结果绝对值表示 a b 形成的平行四边行的面积
31
       vec3 operator^(const vec3 &b) const {
           return \{y * b.z - z * b.y, z * b.x - x * b.z, x * b.y - y * b.x\};
32
33
34
35
       friend istream &operator>>(istream &in, vec3 &a) {
36
           in >> a.x >> a.y >> a.z;
37
           return in;
38
       }
       friend ostream &operator<<(ostream &out, const vec3 &a) {</pre>
39
40
           out << fixed << setprecision(6);</pre>
           out << a.x << " " << a.y << " " << a.z;
41
42
           return out;
43
44
  };
45
46
  struct line3 {
47
      vec3 p, d;
48
       line3() {}
49
       line3(const vec3 &a, const vec3 &b) : p(a), d(b - a) {}
50
  };
51
52
  struct plane {
53
       vec3 p, d;
54
       plane() {}
       plane(const vec3 &a, const vec3 &b, const vec3 &c) : p(a) {
55
56
          d = (b - a) ^ (c - a);
           assert(d != vec3());
57
58
59
  };
60
  // 线面是否垂直
62 bool perpen(const line3 &a, const plane &b) { return (a.d ^ b.d) == vec3(); }
63
64
  // 线面是否平行
65 bool parallel(const line3 &a, const plane &b) { return sgn(a.d * b.d, 0) == 0; }
66
67
  // 线面交点
68
  vec3 intersection(const line3 &a, const plane &b) {
69
       assert(!parallel(a, b));
70
       double t = (b.p - a.p) * b.d / (a.d * b.d);
71
       return a.p + a.d * t;
72
```

## 5.3 扫描线

```
10
           int pd;
11
       };
       vector<Line> line(rec.size() * 2);
12
13
       vector<ll> y_set(rec.size() * 2);
14
       for (int i = 0; i < rec.size(); i++) {</pre>
           y_set[i * 2] = rec[i].y_l;
15
16
           y_set[i * 2 + 1] = rec[i].y_r;
17
           line[i * 2] = {rec[i].x_l, rec[i].y_r, rec[i].y_l, 1};
           line[i * 2 + 1] = {rec[i].x_r, rec[i].y_r, rec[i].y_l, -1};
18
19
       }
20
       sort(y_set.begin(), y_set.end());
21
       y_set.erase(unique(y_set.begin(), y_set.end()), y_set.end());
22
       sort(line.begin(), line.end(), [](Line a, Line b) { return a.x < b.x; });</pre>
23
       struct Data {
24
           int 1, r;
25
           11 len, cnt, raw_len;
26
27
       vector<Data> tree(4 * y_set.size());
28
       function<void(int, int, int)> build = [&](int pos, int 1, int r) {
29
           tree[pos].1 = 1;
30
           tree[pos].r = r;
31
           if (1 == r) {
32
               tree[pos].raw_len = y_set[r + 1] - y_set[l];
33
               tree[pos].cnt = tree[pos].len = 0;
34
               return;
35
           build(ls, 1, mid);
36
37
           build(rs, mid + 1, r);
38
           tree[pos].raw_len = tree[ls].raw_len + tree[rs].raw_len;
39
       };
40
       function<void(int, int, int, int)> update = [&](int pos, int 1, int r, int num) {
41
           if (1 <= tree[pos].1 && tree[pos].r <= r) {</pre>
42
               tree[pos].cnt += num;
               tree[pos].len = tree[pos].cnt
43
                                                                ? tree[pos].raw_len
                                 : tree[pos].1 == tree[pos].r ? 0
44
45
                                                                : tree[ls].len + tree[rs].len;
46
               return;
47
           }
48
           if (1 <= mid) update(ls, 1, r, num);</pre>
49
           if (r > mid) update(rs, l, r, num);
50
           tree[pos].len = tree[pos].cnt ? tree[pos].raw_len : tree[ls].len + tree[rs].len;
51
       };
52
       build(1, 0, y_set.size() - 2);
53
       auto find_pos = [&](ll num) {
54
           return lower_bound(y_set.begin(), y_set.end(), num) - y_set.begin();
55
       };
56
       11 \text{ res} = 0;
57
       for (int i = 0; i < line.size() - 1; i++) {</pre>
           update(1, find_pos(line[i].y_down), find_pos(line[i].y_up) - 1, line[i].pd);
58
           res += (line[i + 1].x - line[i].x) * tree[1].len;
59
60
61
       return res;
62 }
```

# 6 杂项

## 6.1 快读

```
namespace IO {
   constexpr int N = (1 << 20) + 1;
 3
   char Buffer[N];
   int p = N;
 5
 6
   char& get() {
 7
       if (p == N) {
 8
           fread(Buffer, 1, N, stdin);
9
           p = 0;
10
11
       return Buffer[p++];
12
13
   template <typename T = int>
14
15
   T read() {
16
       T x = 0;
       bool f = 1;
17
18
       char c = get();
19
       while (!isdigit(c)) {
20
           f = c != '-';
21
           c = get();
22
23
       while (isdigit(c)) {
           x = x * 10 + c - '0';
24
25
           c = get();
26
27
       return f ? x : -x;
28
29
  } // namespace IO
  using IO::read;
```

#### 6.2 高精度

```
struct bignum {
2
       string num;
3
4
       bignum() : num("0") {}
5
       bignum(const string& num) : num(num) { reverse(this->num.begin(), this->num.end()); }
6
       bignum(11 num) : num(to_string(num)) { reverse(this->num.begin(), this->num.end()); }
7
8
       bignum operator+(const bignum& other) {
9
           bignum res;
10
           res.num.pop_back();
11
           res.num.reserve(max(num.size(), other.num.size()) + 1);
12
           for (int i = 0, j = 0, x; i < num.size() || i < other.num.size() || j; i++) {
13
               x = j;
               j = 0;
14
15
               if (i < num.size()) x += num[i] - '0';</pre>
16
               if (i < other.num.size()) x += other.num[i] - '0';</pre>
17
               if (x >= 10) j = 1, x -= 10;
18
               res.num.push_back(x + '0');
19
           }
20
           res.num.capacity();
```

```
21
            return res;
22
       }
23
24
       bignum operator*(const bignum& other) {
25
            vector<int> res(num.size() + other.num.size() - 1, 0);
            for (int i = 0; i < num.size(); i++)</pre>
26
27
                for (int j = 0; j < other.num.size(); j++)</pre>
28
                     res[i + j] += (num[i] - '0') * (other.num[j] - '0');
29
            int g = 0;
30
            for (int i = 0; i < res.size(); i++) {</pre>
31
                res[i] += g;
32
                g = res[i] / 10;
33
                res[i] %= 10;
34
            }
35
            while (g) {
36
                res.push_back(g % 10);
37
                g /= 10;
38
            }
39
            int lim = res.size();
            while (lim > 1 && res[lim - 1] == 0) lim--;
40
41
            bignum res2;
42
            res2.num.resize(lim);
            for (int i = 0; i < lim; i++) res2.num[i] = res[i] + '0';</pre>
43
44
45
       }
46
47
       bool operator<(const bignum& other) {</pre>
            if (num.size() == other.num.size())
48
49
                for (int i = num.size() - 1; i >= 0; i--)
50
                    if (num[i] == other.num[i]) continue;
51
                    else return num[i] < other.num[i];</pre>
52
            return num.size() < other.num.size();</pre>
53
       }
54
55
       friend istream& operator>>(istream& in, bignum& a) {
56
            in >> a.num;
57
            reverse(a.num.begin(), a.num.end());
58
            return in;
59
60
       friend ostream& operator<<(ostream& out, bignum a) {</pre>
61
            reverse(a.num.begin(), a.num.end());
62
            return out << a.num;</pre>
63
64
```

#### 6.3 离散化

```
template <typename T>
  struct Hash {
2
3
      vector<int> S;
4
      vector<T> a;
5
      Hash(const vector<int>& b) : S(b) {
6
           sort(S.begin(), S.end());
7
           S.erase(unique(S.begin(), S.end()), S.end());
8
           a = vector<T>(S.size());
9
10
      T& operator[](int i) const {
```

```
auto pos = lower_bound(S.begin(), S.end(), i) - S.begin();
assert(pos != S.size() && S[pos] == i);
return a[pos];
}
```

#### 6.4 模运算

```
constexpr int mod = 998244353;
 1
 2
 3
   template <typename T>
 4
   T power(T a, int b) {
5
       T res = 1;
 6
       while (b) {
 7
           if (b & 1) res = res * a;
           a = a * a;
 8
9
           b >>= 1;
10
11
       return res;
12
  }
13
14
   struct modint {
15
       int x;
16
       modint(int _x = 0) : x(_x) {
17
           if (x < 0) x += mod;
           else if (x >= mod) x -= mod;
18
19
       modint inv() const { return power(*this, mod - 2); }
20
21
       modint operator+(const modint& b) { return x + b.x; }
22
       modint operator-() const { return {-x}; }
23
       modint operator-(const modint& b) { return -b + *this; }
24
       modint operator*(const modint& b) { return int((ll)x * b.x % mod); }
25
       modint operator/(const modint& b) { return *this * b.inv(); }
26
       friend istream& operator>>(istream& is, modint& other) {
27
           11 _x;
28
           is \rightarrow _x;
29
           other = modint(_x);
30
           return is;
31
32
       friend ostream& operator<<(ostream& os, modint other) { return os << other.x; }</pre>
33 };
```

#### 6.5 分数

```
struct frac {
1
2
       11 a, b;
3
       frac() : a(0), b(1) {}
4
       frac(ll _a, ll _b) : a(_a), b(_b) {
5
           assert(b);
6
           if (a) {
7
               int tmp = gcd(a, b);
8
               a /= tmp;
9
               b /= tmp;
10
           } else *this = frac();
11
       frac operator+(const frac& other) { return frac(a * other.b + other.a * b, b * other.b); }
```

```
13
       frac operator-() const {
           frac res = *this;
14
15
           res.a = -res.a;
16
           return res;
17
       frac operator-(const frac& other) const { return -other + *this; }
18
19
       frac operator*(const frac& other) const { return frac(a * other.a, b * other.b); }
20
       frac operator/(const frac& other) const {
21
           assert(other.a);
22
           return *this * frac(other.b, other.a);
23
       }
24
       bool operator<(const frac& other) const { return (*this - other).a < 0; }</pre>
25
       bool operator<=(const frac& other) const { return (*this - other).a <= 0; }</pre>
26
       bool operator>=(const frac& other) const { return (*this - other).a >= 0; }
27
       bool operator>(const frac& other) const { return (*this - other).a > 0; }
28
       bool operator==(const frac& other) const { return a == other.a && b == other.b; }
29
       bool operator!=(const frac& other) const { return !(*this == other); }
30 };
```

#### 6.6 表达式求值

```
1
   // 格式化表达式
  string format(const string& s1) {
3
       stringstream ss(s1);
4
       string s2;
5
       char ch;
 6
       while ((ch = ss.get()) != EOF) {
           if (ch == ' ') continue;
7
8
           if (isdigit(ch)) s2 += ch;
           else {
9
               if (s2.back() != ' ') s2 += ' ';
10
               s2 += ch;
11
               s2 += ' ';
12
13
           }
14
       }
15
       return s2;
16 }
17
18
   // 中缀表达式转后缀表达式
19
  string convert(const string& s1) {
20
       unordered_map<char, int> rank{{'+', 2}, {'-', 2}, {'*', 1}, {'/', 1}, {'^', 0}};
21
       stringstream ss(s1);
22
       string s2, temp;
23
       stack<char> op;
24
       while (ss >> temp) {
25
           if (isdigit(temp[0])) s2 += temp + ' ';
26
           else if (temp[0] == '(') op.push('(');
27
           else if (temp[0] == ')') {
28
               while (op.top() != '(') {
29
                   s2 += op.top();
                   s2 += ' ';
30
31
                   op.pop();
32
               }
33
               op.pop();
34
           } else {
35
               while (!op.empty() && op.top() != '(' &&
36
                       (temp[0] != '^' && rank[op.top()] <= rank[temp[0]] ||</pre>
```

```
37
                        rank[op.top()] < rank[temp[0]])) {</pre>
38
                    s2 += op.top();
                    s2 += ' ';
39
40
                    op.pop();
41
42
                op.push(temp[0]);
43
           }
44
45
       while (!op.empty()) {
46
           s2 += op.top();
           s2 += ' ';
47
48
           op.pop();
49
50
       return s2;
51
52
53
   // 计算后缀表达式
54
   int calc(const string& s) {
55
       stack<int> num;
56
       stringstream ss(s);
57
       string temp;
58
       while (ss >> temp) {
           if (isdigit(temp[0])) num.push(stoi(temp));
59
60
61
                int b = num.top();
62
                num.pop();
63
                int a = num.top();
64
                num.pop();
65
                if (temp[0] == '+') a += b;
                else if (temp[0] == '-') a -= b;
66
                else if (temp[0] == '*') a *= b;
67
68
                else if (temp[0] == '/') a /= b;
                else if (temp[0] == '^') a = ksm(a, b);
69
70
                num.push(a);
71
72
73
       return num.top();
74 }
```

#### 6.7 日期

```
1 int month[] = {0, 31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31};
2
  int pre[13];
3
   vector<int> leap;
 4
   struct Date {
5
       int y, m, d;
6
       bool operator<(const Date& other) const {</pre>
7
            return array<int, 3>{y, m, d} < array<int, 3>{other.y, other.m, other.d};
8
9
       Date(const string& s) {
10
            stringstream ss(s);
11
            char ch;
12
            ss \Rightarrow y \Rightarrow ch \Rightarrow m \Rightarrow ch \Rightarrow d;
13
       }
14
       int dis() const {
            int yd = (y - 1) * 365 + (upper_bound(leap.begin(), leap.end(), y - 1) - leap.begin());
15
16
            int md = pre[m - 1] + (m > 2 && (y % 4 == 0 && y % 100 || y % 400 == 0));
```

```
17          return yd + md + d;
18          }
19          int dis(const Date& other) const { return other.dis() - dis(); }
20     };
21     for (int i = 1; i <= 12; i++) pre[i] = pre[i - 1] + month[2];
22     for (int i = 1; i <= 1000000; i++)
23     if (i % 4 == 0 && i % 100 || i % 400 == 0) leap.push_back(i);</pre>
```

#### 6.8 builtin 函数

如果是 long long 型,记得函数后多加个 ll。

- ctz, 从最低位连续的 0 的个数, 如果传入 0 则行为未定义。
- clz, 从最高位连续的 0 的个数, 如果传入 0 则行为未定义。
- popcount, 二进制 1 的个数。
- parity, 二进制 1 的个数奇偶性。

#### 6.9 对拍

linux/Mac

```
#!/bin/bash
2
  g++ $1 -o a -02
3
  g++ $2 -o b -02
  g++ random.cpp -o random -02
6
7
  cnt=0
8
  while true; do
9
      let cnt++
10
       echo TEST:$cnt
11
       ./random > in
12
       ./a < in > out.a
13
       ./b < in > out.b
14
      if ! diff out.a out.b; then break; fi
15 done
```

windows

```
@echo off
 1
2
3 g++ %1 -o a -02
4
  g++ %2 -o b -02
  g++ random.cpp -o random -02
5
 7
   set cnt=0
8
9
   :again
10
       set /a cnt=cnt+1
11
       echo TEST:%cnt%
12
       .\random > in
        \cdot \arraycolor{a} < in > out.a
13
14
       .\b < in \rightarrow out.b
       fc out.a out.b > nul
15
16 if not errorlevel 1 goto again
```

# 6.10 编译常用选项

-Wall -Woverflow -Wextra -Wpedantic -Wfloat-equal -Wshadow -fsanitize=address,undefined

# 6.11 开栈

不同的系统/编译器可能命令不一样

```
ulimit -s
-Wl,--stack=0x10000000
-Wl,-stack_size -Wl,0x10000000
-Wl,-z,stack-size=0x10000000
```

# 6.12 clang-format

转储配置

```
clang-format -style=Google -dump-config > ./.clang-format
```

## $. \\ clang-format$

```
BasedOnStyle: Google
IndentWidth: 4
AllowShortIfStatementsOnASingleLine: AllIfsAndElse
ColumnLimit: 100
```